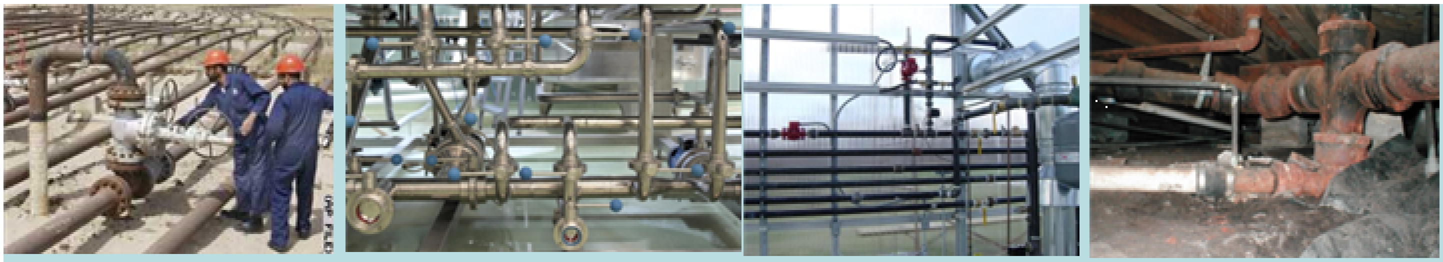
Lecture19



CSE 421
Introduction to Algorithms

Lecture 19
Winter 2024
Network Flow, Part 3

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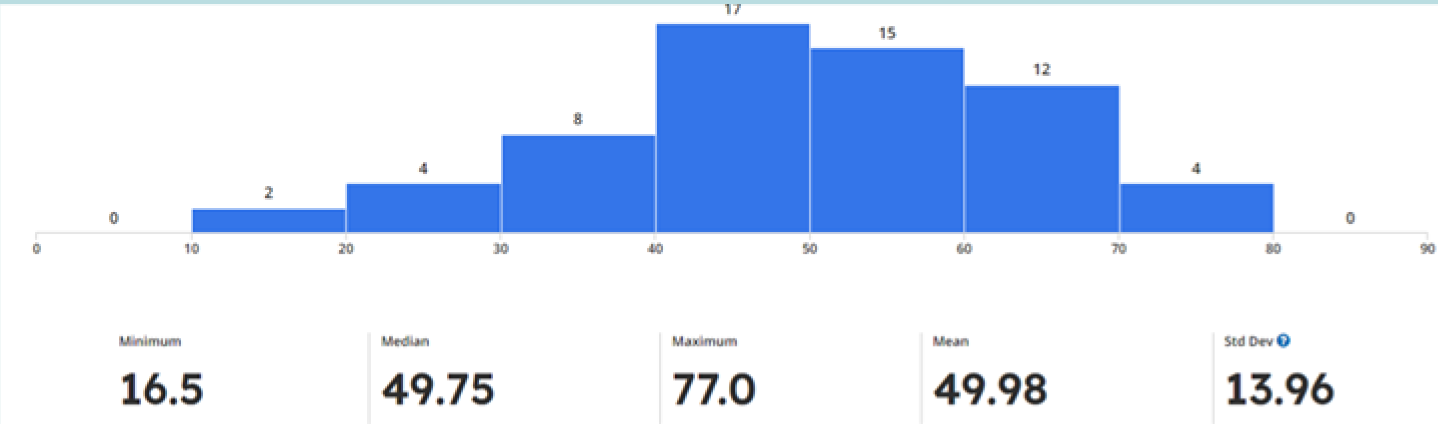
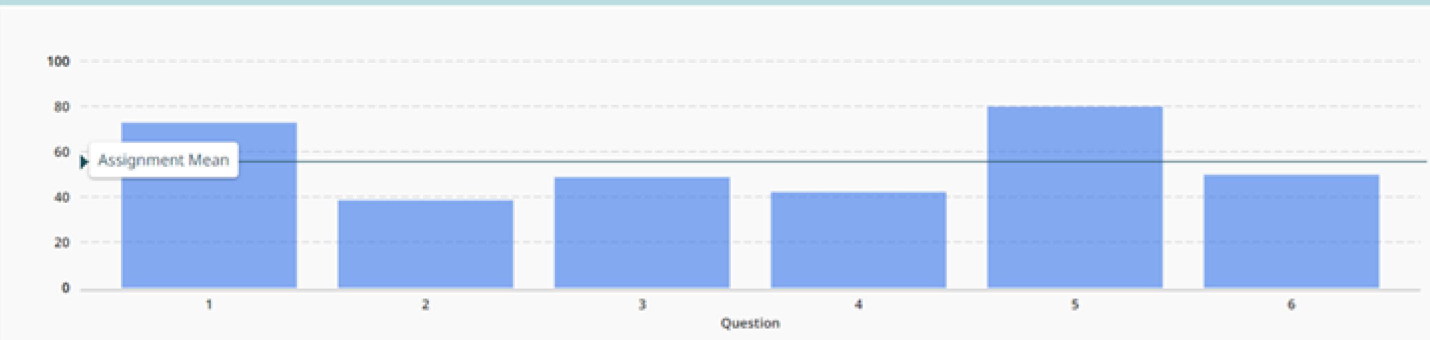
Midterm


17
15
12
8
4
2
0
4
0
0
10
20
30
40
50
60
70
80
90
Minimum
16.5
Median
49.75
Maximum
77.0
Mean
49.98
Std Dev
13.96


100
80
60
40
20
0
Assignment Mean
1
2
3
4
5
6
Question

Outline

- ~~Network flow definitions~~
- ~~Flow examples~~
- ~~Augmenting Paths~~
- ~~Residual Graph~~
- ~~Ford Fulkerson Algorithm~~
- ~~Cuts~~
- ~~Maxflow-MinCut Theorem~~
- ~~Worst Case Runtime for FF~~
- Improving Runtime bounds
  - Capacity Scaling
  - Fully Polynomial Time Algorithms
- Applications of Network Flow

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Ford-Fulkerson Algorithm (1956)

How many iterations?

while not done

- Construct residual graph G_R
- Find an s-t path P in G_R with capacity $b > 0$
- Add b units of flow along path P in G

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Ford Fulkerson Runtime

- Cost per phase X number of phases
- Phases
  - Capacity leaving source: C
  - Add at least one unit per phase
- Cost per phase
  - Build residual graph: O(m)
  - Find s-t path in residual: O(m)

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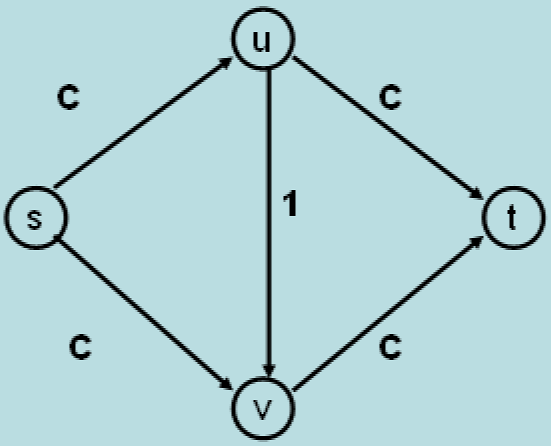
Performance

- The worst case performance of the Ford-Fulkerson algorithm O(Cm)



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Polynomial Time Algorithms

- Input of size n, runtime $T(n) = O(n^k)$
- Input size measures
  - Bits of input
  - Number of data items
- Maximum item magnitude C
  - $O(Cn^k)$: Exponential
  - $O(n^k \log C)$: Polynomial
  - $O(n^k)$: Fully polynomial

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Better methods of finding augmenting paths

- Find the maximum capacity augmenting path
  - $O(m^2 \log(C))$ time algorithm for network flow
- Find the shortest augmenting path
  - $O(m^2 n)$ time algorithm for network flow
- Find a blocking flow in the residual graph
  - $O(mn \log n)$ time algorithm for network flow

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Capacity Scaling Algorithm

- Choose $\Delta = 2^k$ such that all edges in G_R have capacity less than 2Δ

while $\Delta \geq 1$

while there is a path P in G_R with capacity Δ

Add Δ units of flow along path P in G

Update G_R

$\Delta = \Delta / 2$

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Edmonds-Karp: Easier analysis than Max Capacity First

Analysis

Runtime $O(m^2 \log C)$

- If capacities are integers, then graph is disconnected when Δ = ½
- If largest edge capacity is C, then there are at most log C outer phases
- At the start of each outer phase, the flow is within 2mΔ of the maximum
  - So there are at most 2m inner phases for each Δ

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Shortest Augmenting Path

- Find augmenting paths by BFS

for k = 1 to n

  while there is a path P in G_R of length k and capacity b > 0

    Add b units of flow along path P in G

    Update G_R

- Need to show:
  - The length of the shortest augmenting path is non-decreasing
  - Each while loop finds at most m paths

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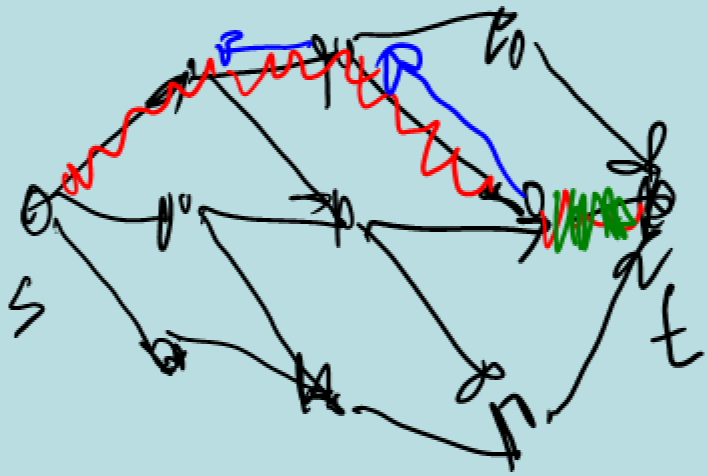
Analysis

- Augmenting along shortest path from s to t does not decrease distance from s to t



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Analysis

- The distance from s to t must increase in G_R after m augmentations by shortest paths

$O(nm^2)$

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Improving the shortest augmenting path algorithm

- Find a blocking flow in one phase to increase the length of augmenting paths
  - Dinitz (Ефим Абрамович Диниц) Algorithm
  - $O(n^2m)$
- Dynamic Trees to decrease cost per augmentation
  - O(nm log n)

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APPLICATIONS OF NETWORK FLOW

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Problem Reduction

- Reduce Problem A to Problem B
  - Convert an instance of Problem A to an instance of Problem B
  - Use a solution of Problem B to get a solution to Problem A
- Practical
  - Use a program for Problem B to solve Problem A
- Theoretical
  - Show that Problem B is at least as hard as Problem A

Problem Reduction Examples

- Reduce the problem of finding the Maximum of a set of integers to finding the Minimum of a set of integers

Find the maximum of: 8, -3, 2, 12, 1, -6

-8 3 -2 -12 -1 6

-12

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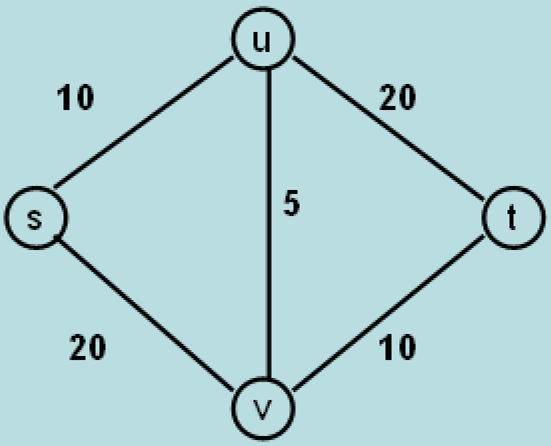
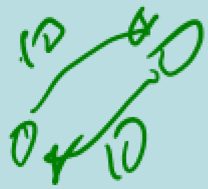
Construct an equivalent minimization problem

Undirected Network Flow

- Undirected graph with edge capacities
- Flow may go either direction along the edges (subject to the capacity constraints)





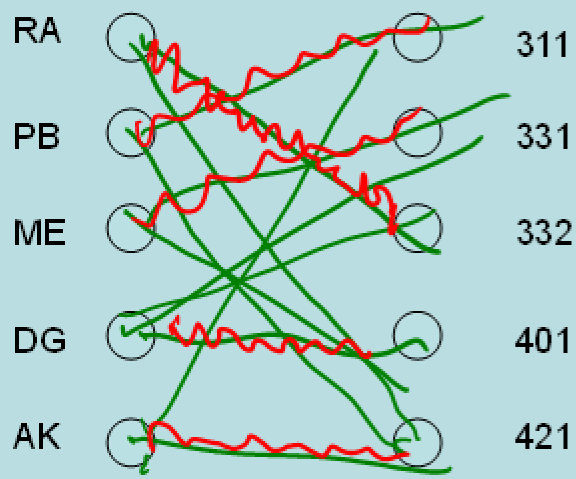
Construct an equivalent flow problem

Bipartite Matching

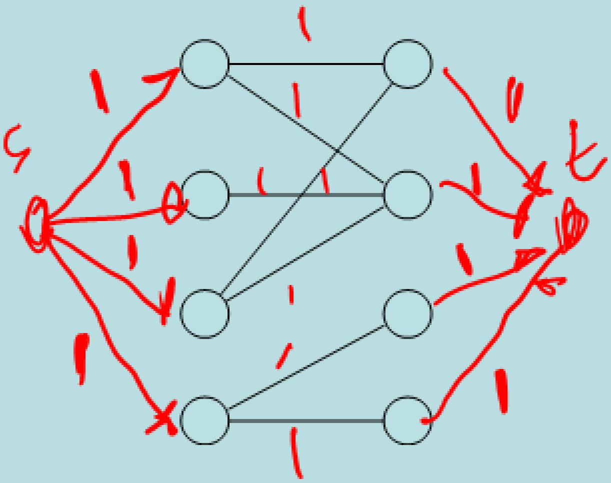
- A graph G=(V,E) is bipartite if the vertices can be partitioned into disjoints sets X,Y
- A matching M is a subset of the edges that does not share any vertices
- Find a matching as large as possible

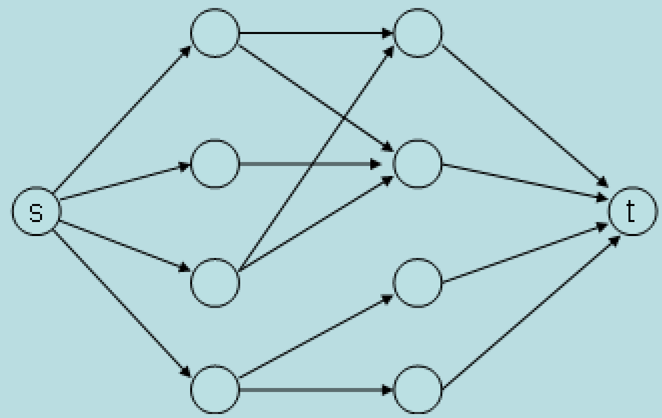
Application

- A collection of teachers
- A collection of courses
- And a graph showing which teachers can teach which courses

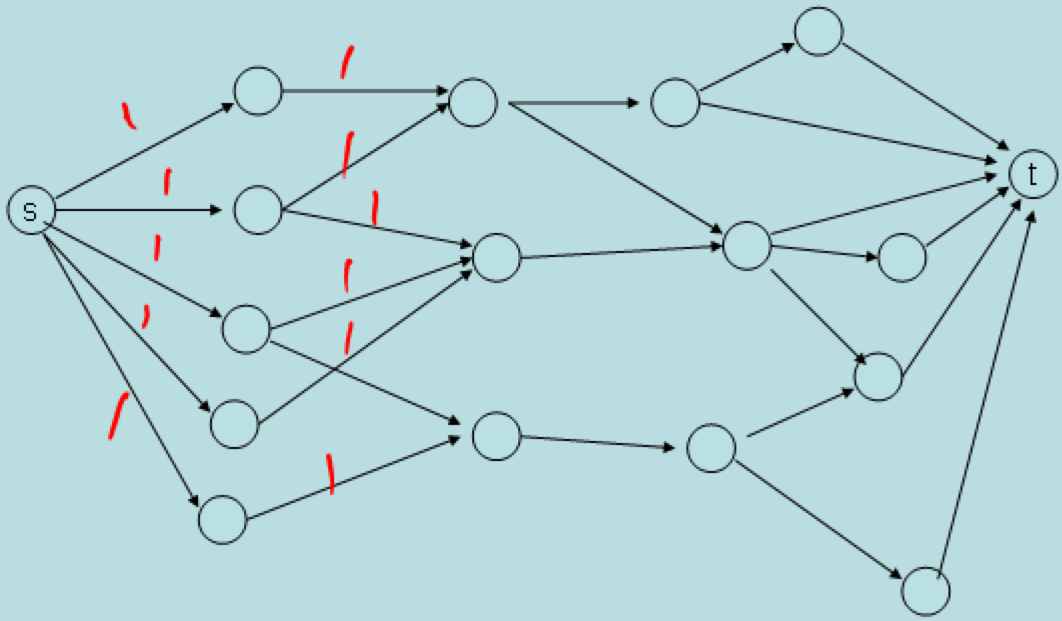


Converting Matching to Network Flow





Finding edge disjoint paths



Construct a maximum cardinality set of edge disjoint paths

Multi-source network flow

- Multi-source network flow
  - Sources $s_1, s_2, \ldots, s_k$
  - Sinks $t_1, t_2, \ldots, t_j$
- Solve with Single source network flow

Resource Allocation: Assignment of reviewers

- A set of papers $P_1, \ldots, P_n$
- A set of reviewers $R_1, \ldots, R_m$
- Paper P_i requires A_i reviewers
- Reviewer R_j can review B_j papers
- For each reviewer R_j, there is a list of paper $L_{j1}, \ldots, L_{jk}$ that R_j is qualified to review

Resource Allocation: Illegal Campaign Donations

- Candidates C_i, . . ., C_n
  - Donate b_i to C_i
- With a little help from your friends
  - Friends F_1, . . ., F_m
  - F_i can give a_{ij} to candidate C_j
  - You can give at most M_i to F_i