CSE 431 Spring 2006
Assignment #4

Due: Friday, April 28, 2006

Reading assignment: Finish reading Chapter 5 of Sipser's text. (You may also want to skim section 6.3 of the text.)

Problems:

1. (10 points) Show that A is decidable if and only if $A \leq_m 0^*1^*$.
2. (10 points) Show that there is a undecidable language contained in 1^*.
3. Which of the following problems are decidable? Justify each answer:
   (a) (10 points) Given Turing machines M and N, is $L(N)$ the complement of $L(M)$?
   (b) (10 points) Given a Turing machine M, integers a and b, and input x, does M run for more than $a|x|^2 + b$ steps on input x?
   (c) (20 points) Given a program P written in Java, or C, or (insert your favorite programming language) that does not read any input but is executed with no bound on the size of integers, does P ever attempt to index an array outside its allocated array bounds.
4. (15 points) Sipser's text: 1st edition problem 5.19; 2nd edition problem 5.21.
5. (Bonus) Show that the following problem is undecidable: Given a Turing machine M and integers a and b, does there exist an input x on which M runs for more than $a|x|^2 + b$ steps on input x?
6. (Bonus) Rice's Theorem shows that for every 'non-trivial' property $\mathcal{P}$ of languages,

$$\mathcal{P}_{TM} = \{\langle M \rangle \mid L(M) \text{ has property } \mathcal{P}\}$$

   is undecidable where by '$\mathcal{P}$ is non-trivial' we mean that $\mathcal{P}$ contains some but not all Turing-recognizable languages. Some of these $\mathcal{P}_{TM}$ are not only undecidable, they are also not Turing-recognizable:

   Show that if there is some *infinite* Turing-recognizable language L that has property $\mathcal{P}$ but none of the finite subsets of L have property $\mathcal{P}$ then $\mathcal{P}_{TM}$ is not Turing recognizable.