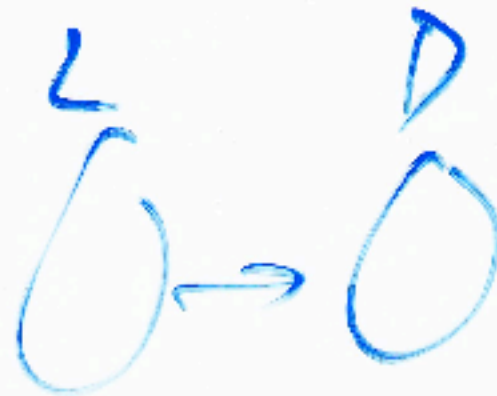


Lecture 15

Some notes on HW #4

\exists decidable D
 \forall decidable L

$$\underline{L \leq_m D}$$



$x_0 \notin D$

$x_1 \in D$

$f(w)$: runs decider for L
~~returns decision for L~~

if γ output x_1

else x_0

$A \leq_m B \iff \bar{A} \leq_m \bar{B}$
If A not T. REC.
Then B " "

EQ_{TM} neither T. REC nor Co-T. REC.

\overline{A} prog \leq_m Uselect prog

$\langle P, w \rangle \rightarrow \langle P', l \rangle$

P' ~~is~~ ignores its input

runs P on w

$l =$ line# of "accept" in P

$P'(x)$
 ~~\equiv~~ $[P$

ans = $P(w)$

if ans = 1 then return 1 l
else return 0

More on P vs NP

$$P = \bigcup_k \text{TIME}(n^k)$$

$n \log n = O(n^2)$

Given $G, a, b \exists \text{ path } a \rightarrow b$

Sorting is x the 17^{th} largest element in list ...

given matrices A, B i, j

$$\text{is } (A \cdot B)_{ij} = c$$

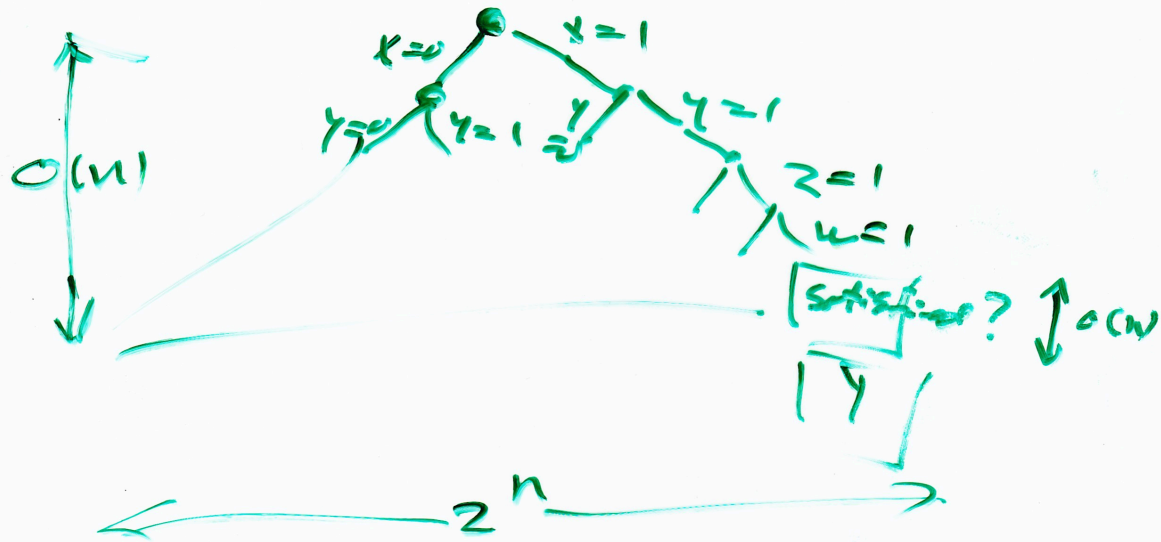
Shortest path

$G, a, b, k \exists \text{ path } a \rightarrow b$
of length $\leq k$

CFL recognition
 $O(n^3)$

SAT

$$(x \vee y) \wedge (\bar{y} \vee z \vee w) \wedge (x \vee y)$$



$$NP = \bigcup_{k \geq 1} \text{Nondet-Time}(n^k)$$

