


CSE/EE 461 – Lecture 10



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Last Time



- We finished up the Network layer
 - Internetworks (IP)
 - Routing (DV/RIP, LS/OSPF)
 - Scalable addressing/routing (BGP, CIDR)
 - Routers

Application
Presentation
Session
Transport
Network
Data Link
Physical

This Time

- We begin on the Transport layer
- Focus
 - How do we send information reliably?
- Topics
 - The Transport layer
 - Acknowledgements and retransmissions (ARQ)
 - End-to-End argument (E2E)

Application
Presentation
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The Transport Layer

- Builds on the services of the Network layer
- Communication between processes running on hosts
 - Naming/Addressing
- Stronger guarantees of message delivery
 - Reliability

Example – Common Properties

TCP

- Guaranteed delivery
- In-order delivery
- Single delivery
- Arbitrarily long messages
- Synchronization
- Flow control
- Multiple processes

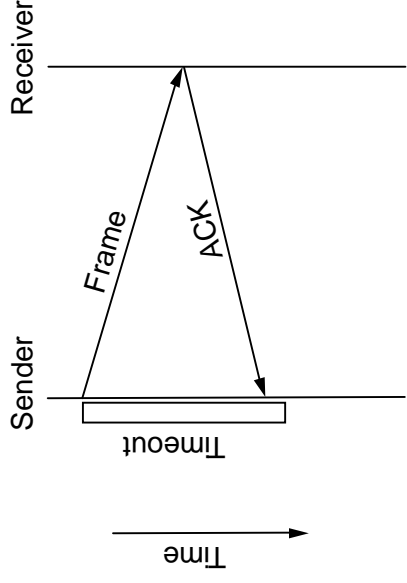
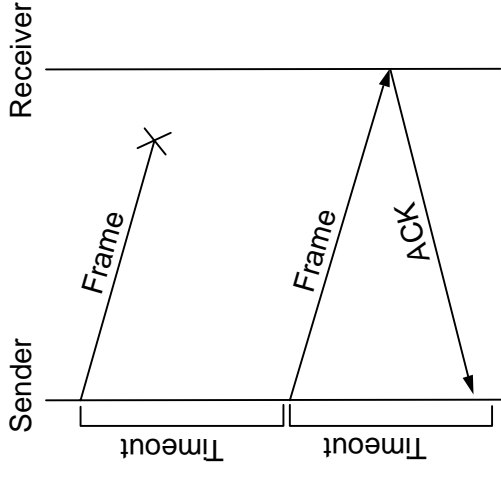
IP

- Lost packets
- Reordered packets
- Duplicate packets
- Limited size packets

Internet Transport Protocols

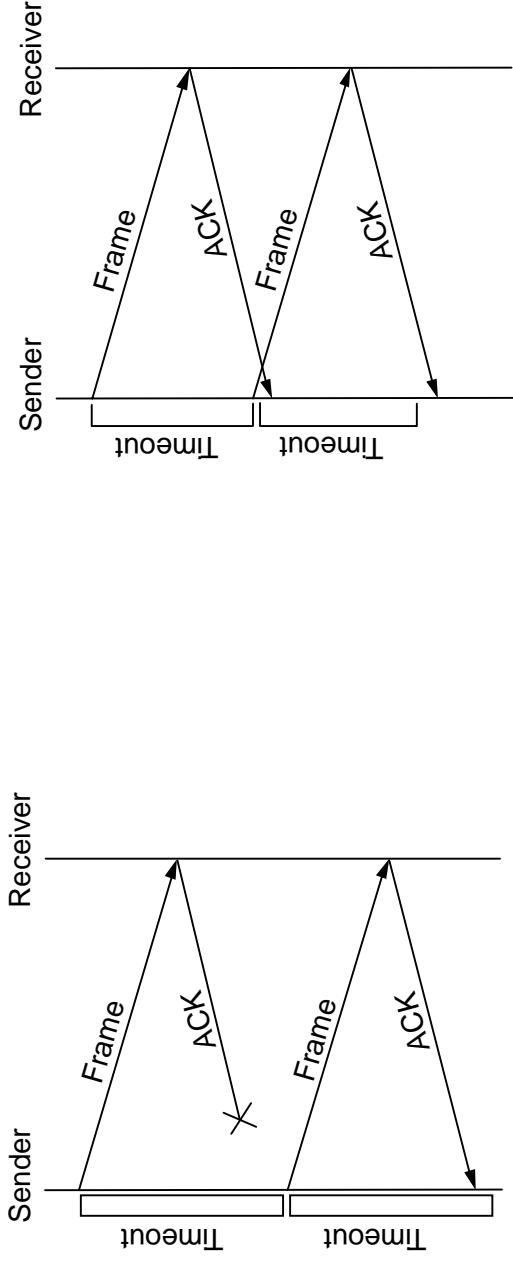
- UDP
 - Datagram abstraction between processes
 - With error detection
- TCP
 - Bytestream abstraction between processes
 - With reliability
 - Plus congestion control (later!)

Automatic Repeat Request (ARQ)



- Packets can be corrupted or lost. How do we add reliability?
- Acknowledgments (ACKs) and retransmissions after a timeout
- ARQ is generic name for protocols based on this strategy

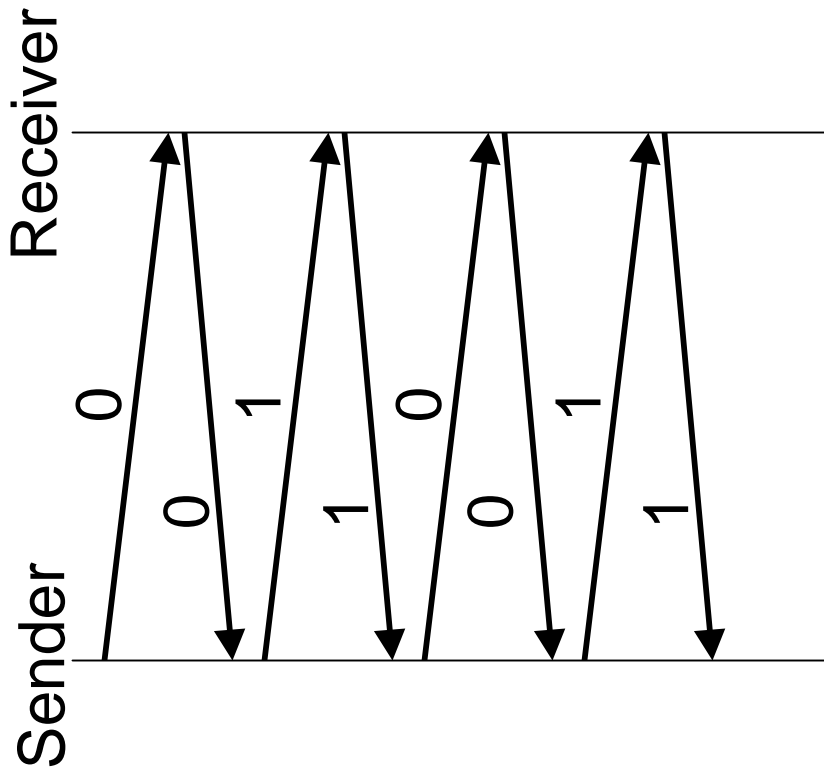
The Need for Sequence Numbers



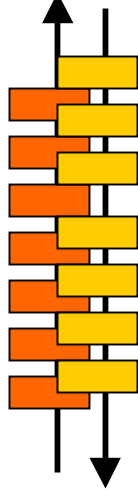
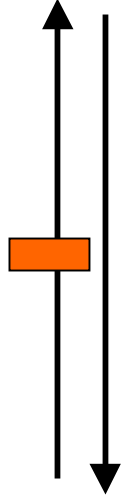
- In the case of ACK loss (or poor choice of timeout) the receiver can't distinguish this message from the next
 - Number packets; here, a single bit will do

Stop-and-Wait

- Only one outstanding packet at a time
- Also called alternating bit protocol

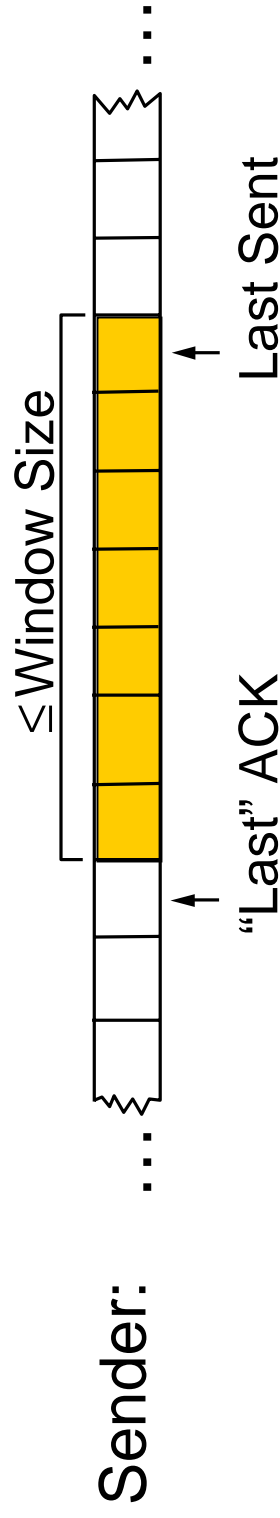


Limitation of Stop-and-Wait



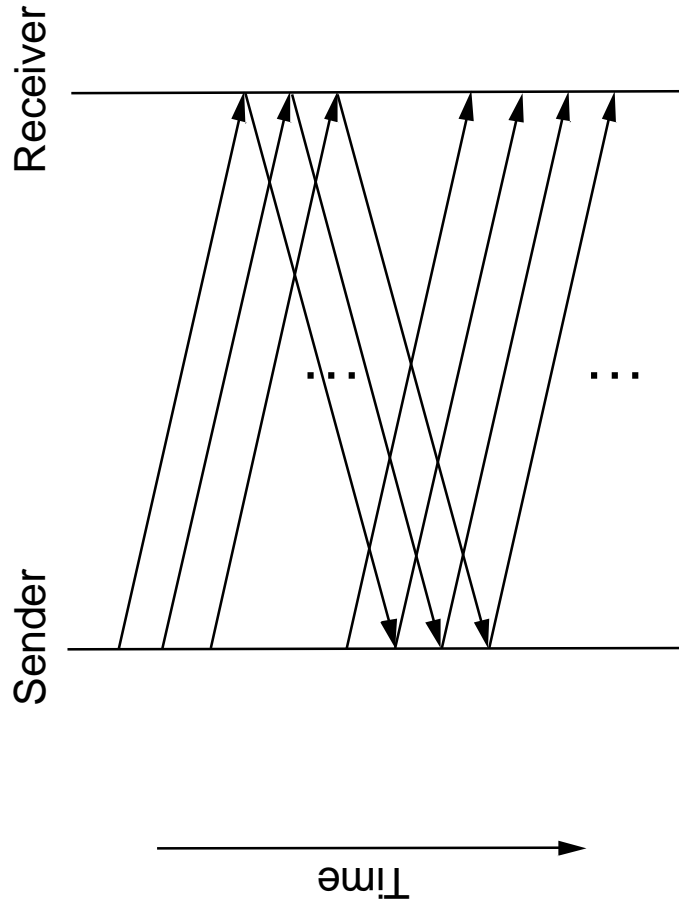
- Lousy performance if wire time \ll prop. delay
 - How bad? You do the math
- Want to utilize all available bandwidth
 - Need to keep more data “in flight”
 - How much? Bandwidth-delay product
- Leads to Sliding Window Protocol

Sliding Window – Sender

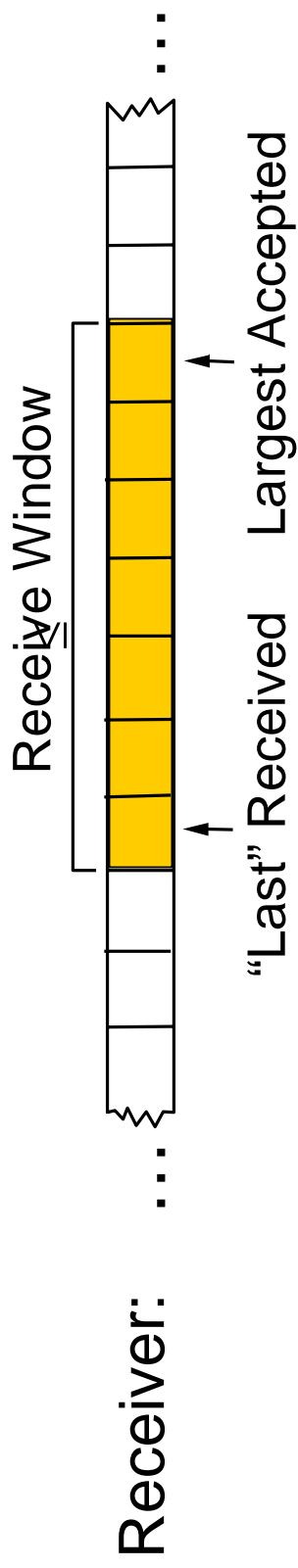


- Window bounds outstanding data
 - Implies need for buffering at sender
- “Last” ACK applies to in-order data
- Sender maintains timers too
 - Go-Back-N: one timer, send all unacknowledged on timeout
 - Selective Repeat: timer per packet, resend as needed

Sliding Window – Timeline



Sliding Window – Receiver



- Receiver buffers too:
 - data may arrive out-of-order
 - or faster than can be consumed (flow control)
- Receiver ACK choices:
 - Individual, Cumulative (TCP), Selective (newer TCP), Negative

Sliding Window Functions



- Sliding window is a mechanism
- It supports multiple functions:
 - Reliable delivery
 - In-order delivery
 - Flow control

Which layer provides Reliability?

- We've been talking about the Transport layer but ...
- ARQ is used by some link layers
 - Acknowledgements in 802.11
- Error detection/correction codes boost reliability
 - Ethernet CRC, IP header checksum, etc.
- Where is the "right" place in the protocol stack?

End-to-End Argument

- Key design principle applied in the Internet
- Reliability is needed end-to-end and can't be replaced by lower layer mechanisms. So put it end-to-end; use lower mechanisms to improve performance as needed.
- TCP provides reliable delivery
 - Checksums packet data as well
- Lower layers keep their residual error rate is low
 - CRC enough for Ethernet; wireless links more problematic

Key Concepts

- Transport layer allows processes to communicate with stronger guarantees, e.g., reliability
- Basic reliability is provided by ARQ mechanisms
 - Stop-and-Wait through Sliding Window
- End-to-End principle guides placement of functions
- Coming Next: Connections and Congestion Control
- Read Keshav 12.4 and Ch 13, esp. 13.4