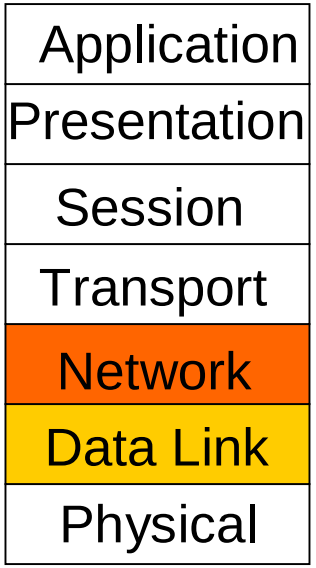
CSE/EE 461 – Lecture 6

David Wetherall
djw@cs.washington.edu

Last Time

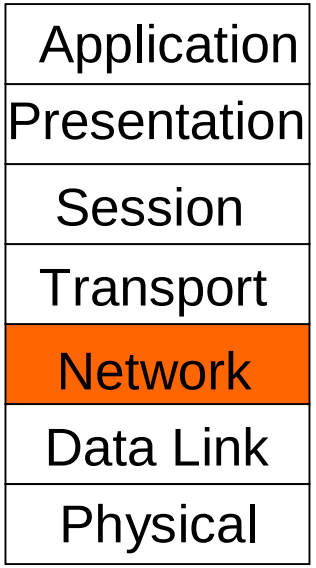
- Focus:
  - What to do when one shared LAN isn't big enough?
- Interconnecting LANs
  - Bridges and LAN switches
- A preview of the Network layer



This Time

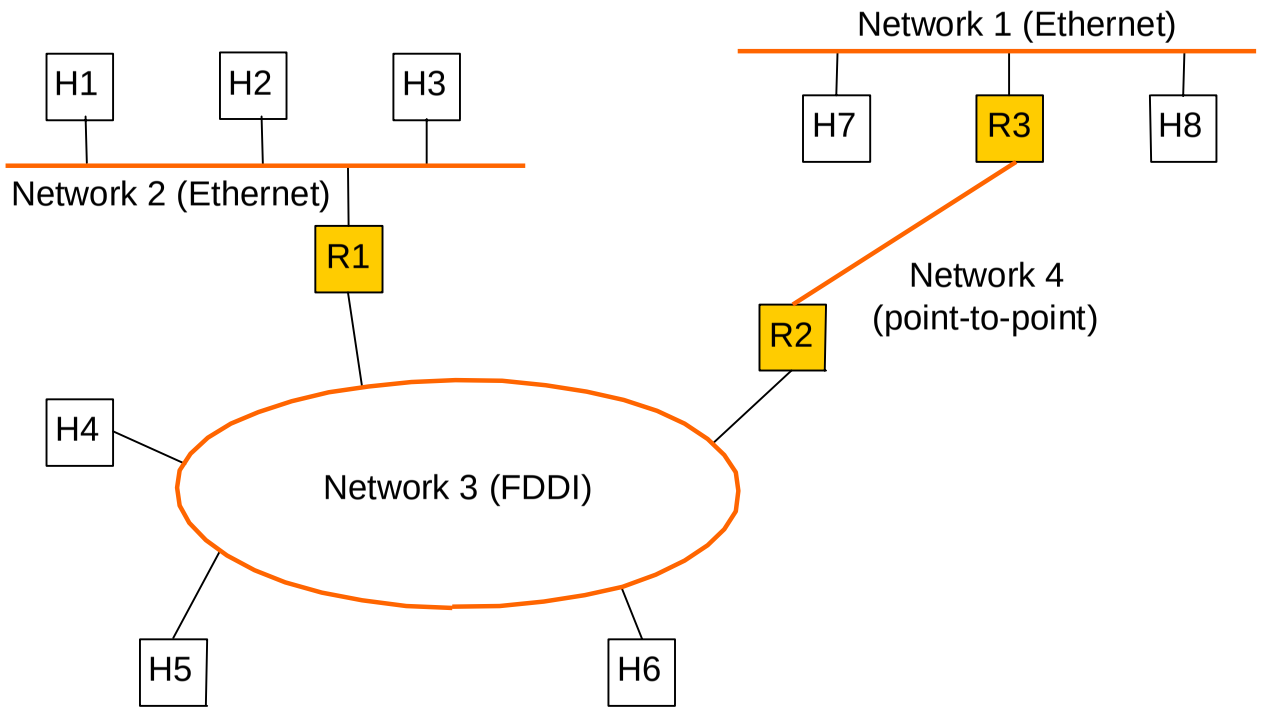
- Focus:
  - How do we build large networks?

- Introduction to the Network layer
  - Internetworks
  - Service models
  - IP (addressing, routing)
  - IP packet format



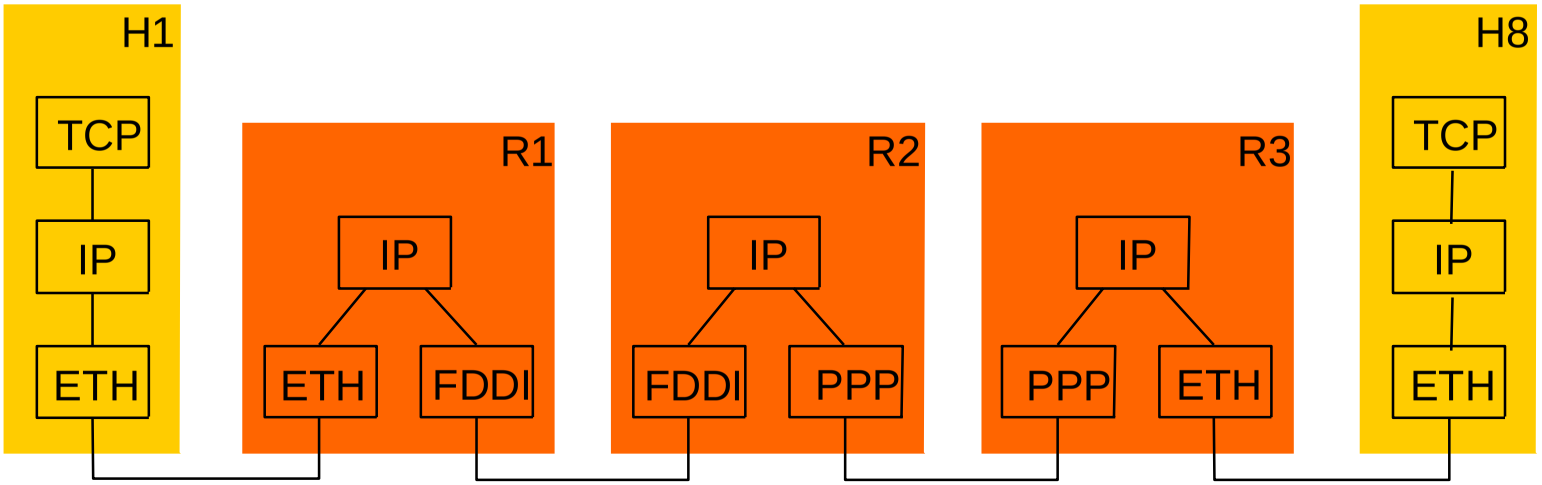
Internetworks

- Set of interconnected networks, e.g., the Internet
  - Scale and heterogeneity



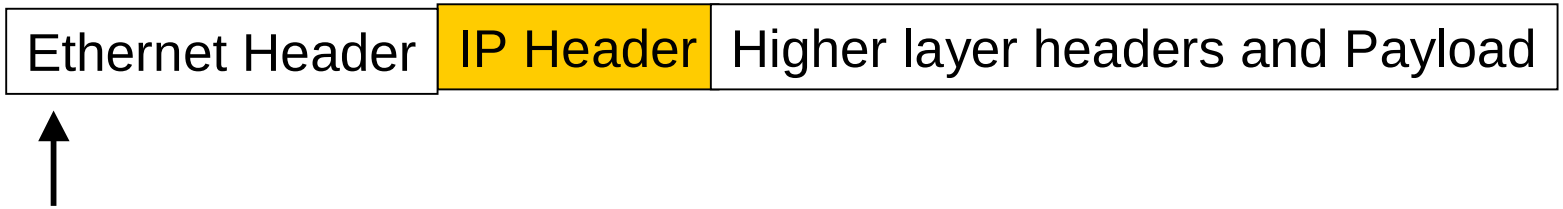
In terms of protocol stacks

- Routers are network level gateways
  - IP is the network layer protocol used in the Internet
  - Packet is the term for network layer PDUs



In terms of packet formats

- View of a packet on the wire on network 1 or 2
- Routers work with IP header, not higher
  - Higher would be a "layer violation"
- Routers strip and add link layer headers



Front of packet to left (and uppermost)

Network Service Models

- Datagram delivery: postal service
  - Also connectionless, best-effort or unreliable service
  - Network can't guarantee delivery of the packet
  - Each packet from a host is routed independently
  - Example: IP

- Virtual circuit models: telephone
  - Also connection-oriented service
  - Signaling: connection establishment, data transfer, teardown
  - All packets from a host are routed the same way (router state)
  - Example: ATM, Frame Relay, X.25

Datagrams or Virtual Circuits?

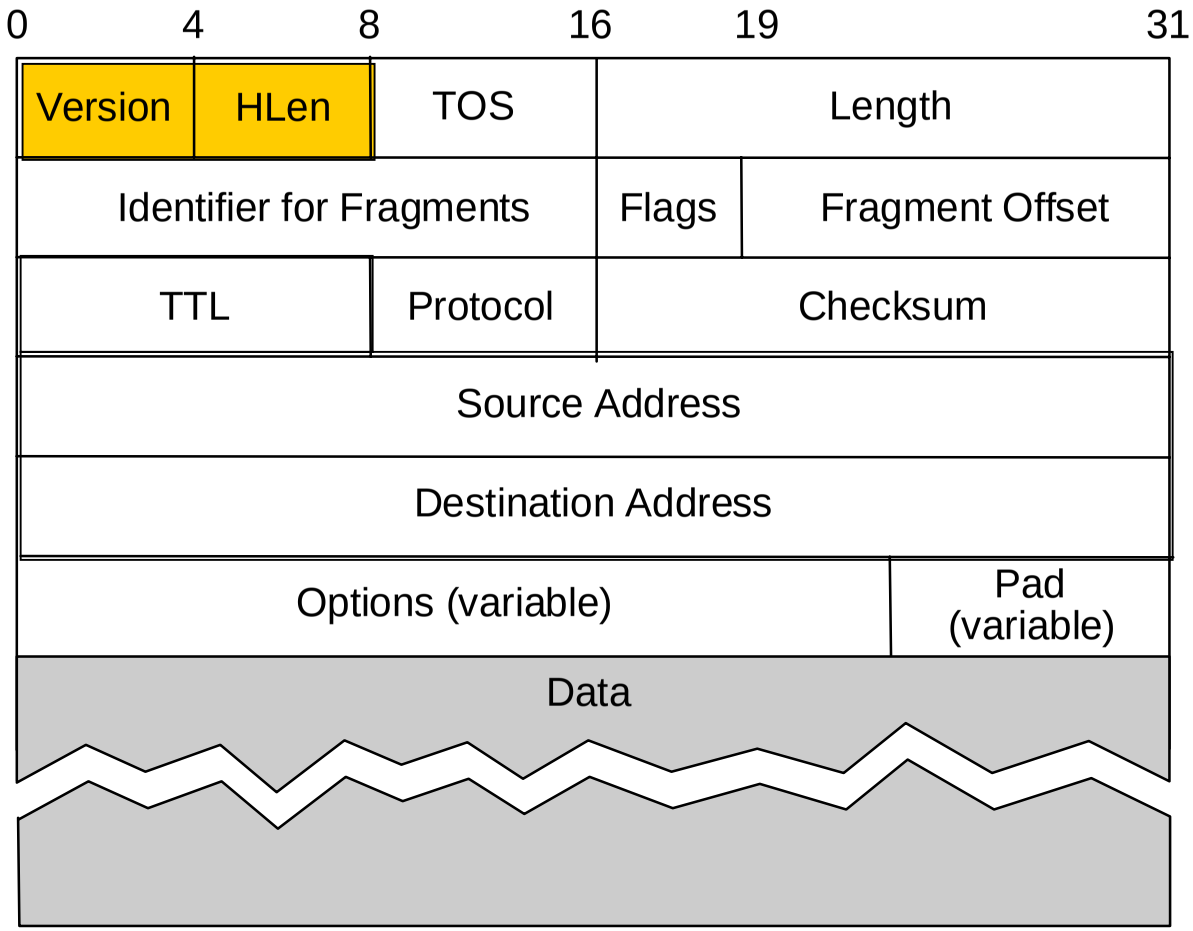
- Pros and Cons?
  - Simplicity/robustness versus stronger resource allocation
  - Examples …
- We return to these tradeoffs later
  - Quality of Service (QOS)
  - These issues at the heart of current Internet evolution
  - Intserv (connection oriented) vs Diffserv ("connectionless")

Internet Protocol (IP)

- Network layer provides end-to-end delivery (routing)
  - Between networks rather than a single network
- IP defines a best effort service without error detection
  - Can be loss, reordering, even duplication!
  - Currently IPv4 (IP version 4), IPv6 on the way
- Global, hierarchical addresses, not flat addresses
  - 32 bits in IPv4 not enough; 128 bits in IPv6
- Routing protocols (RIP, OSPF, BGP) run between routers to maintain routing table (FIB)
  - Which way to forward depending on IP address

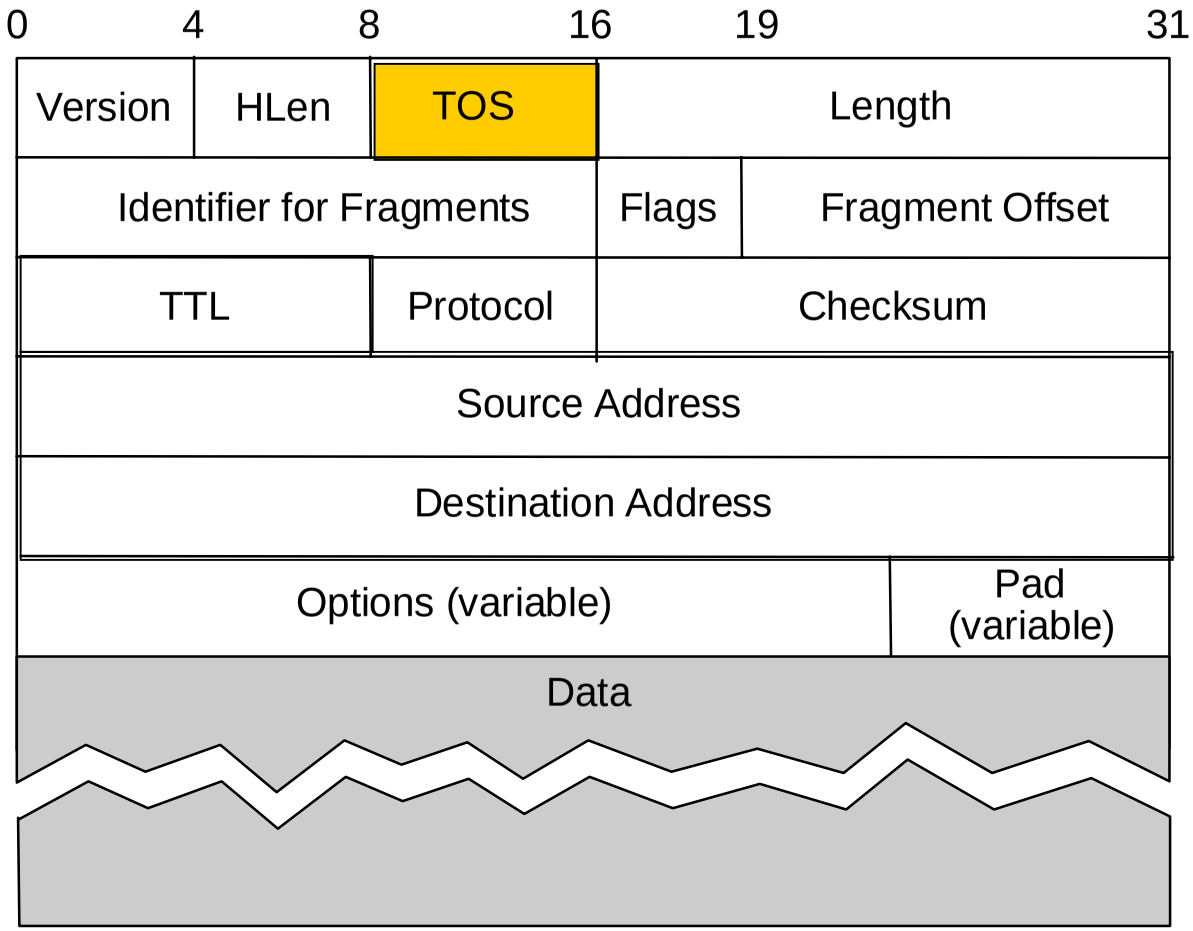
IPv4 Packet Format

- Version is 4
- Header length is number of 32 bit words
- Limits size of options



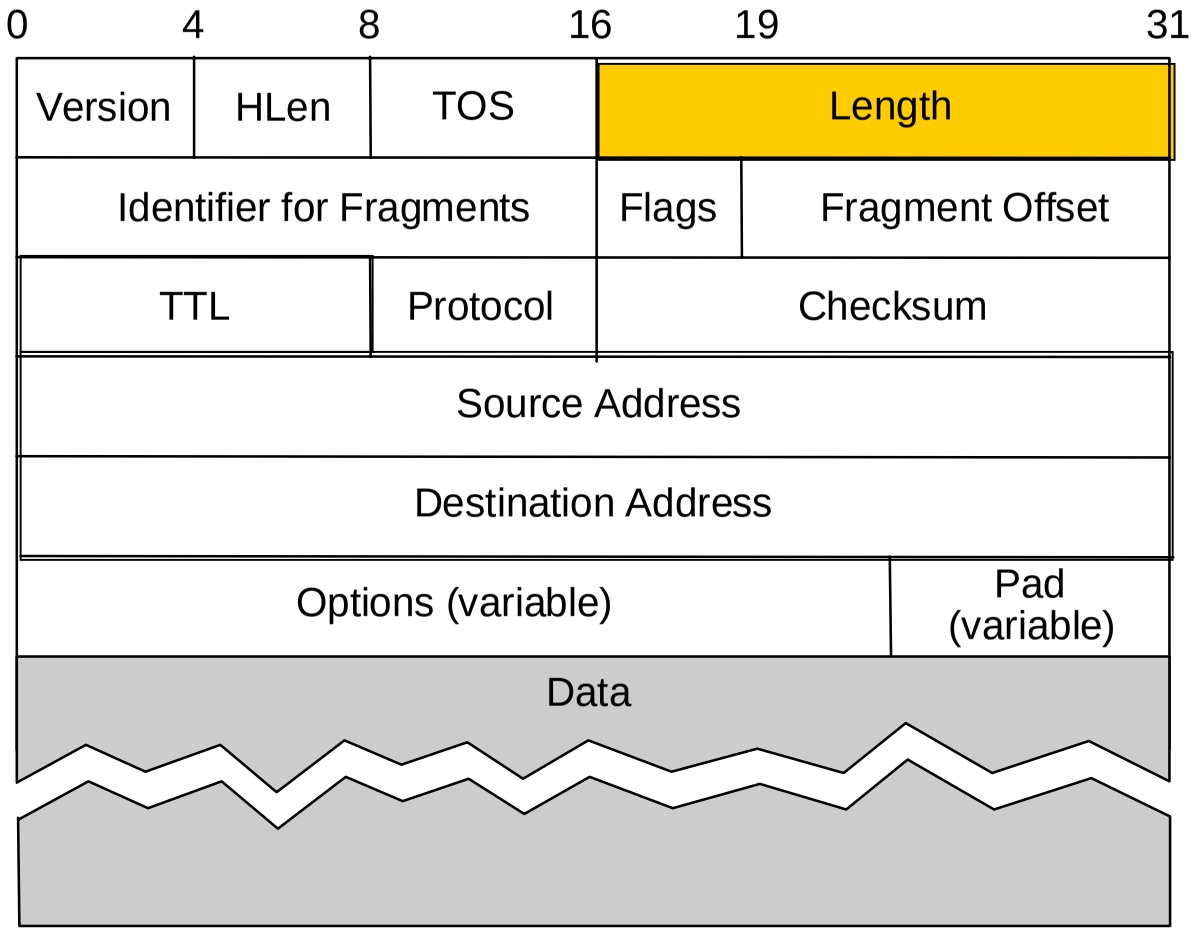
IPv4 Header Fields …

- Type of Service
- Abstract notion of kind of service, never really worked out
  - Routers ignore
- Field will be reused for Diffserv



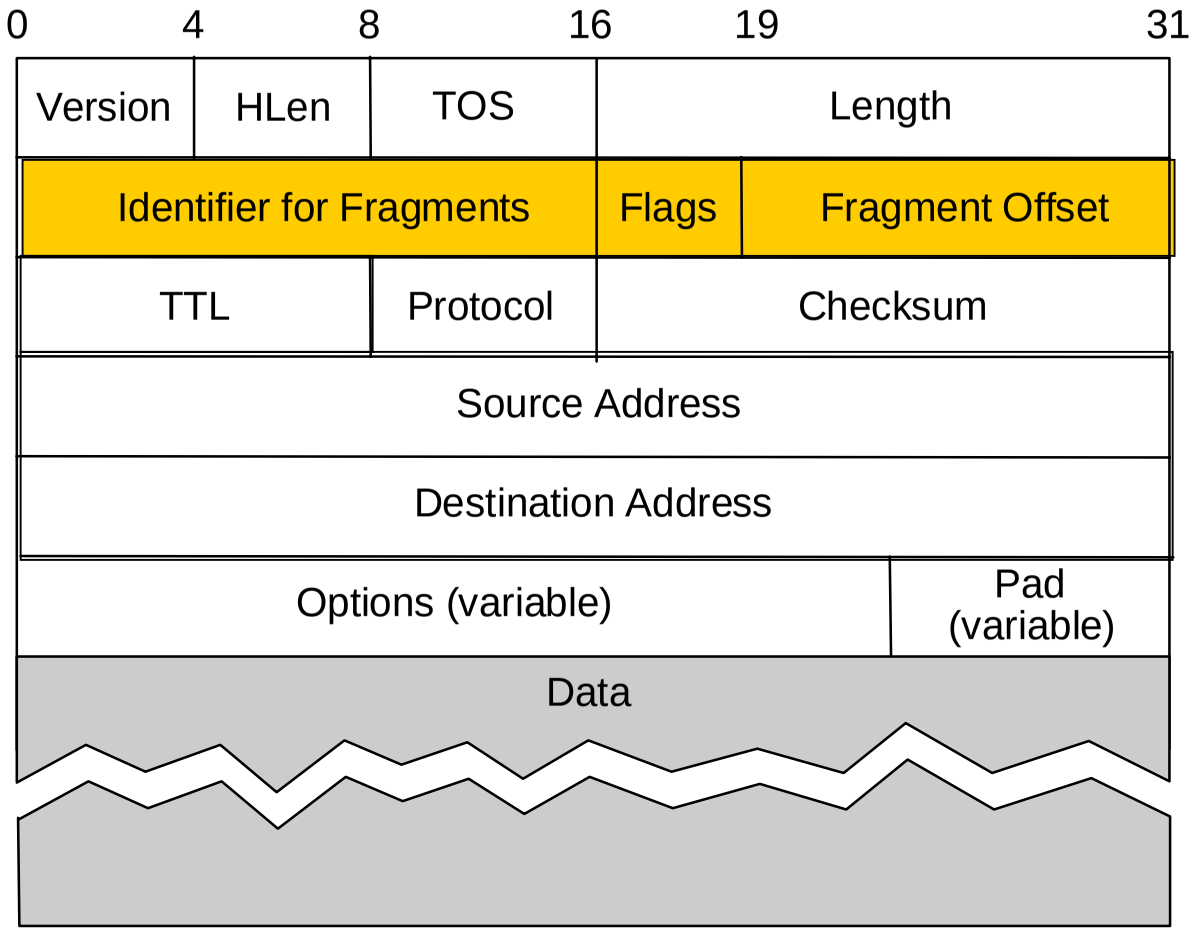
IPv4 Header Fields …

- Length of packet
- Min 20 bytes, max 65K bytes (limit to packet size)



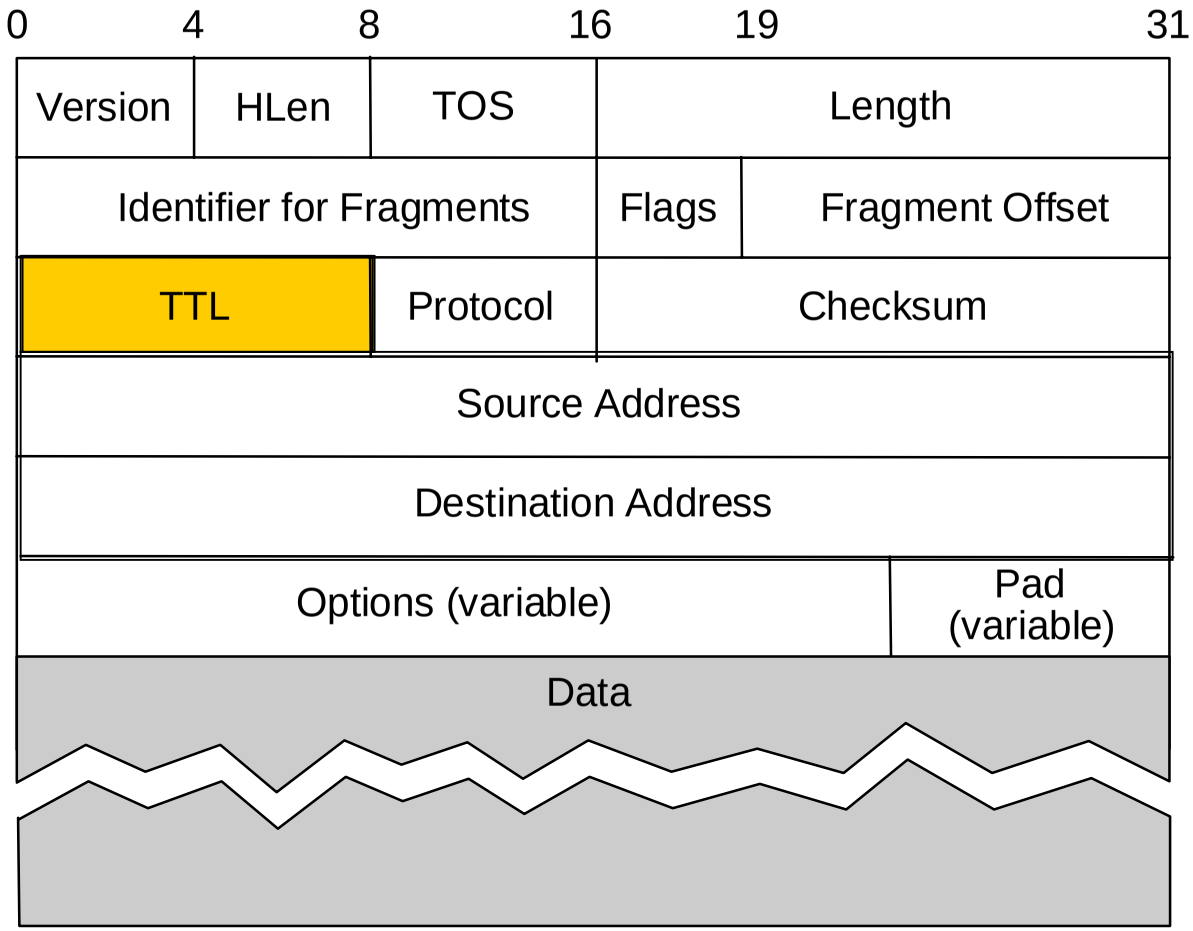
IPv4 Header Fields …

- Fragment fields
- Different LANs have different frame size limits
- May need to break large packet into smaller fragments



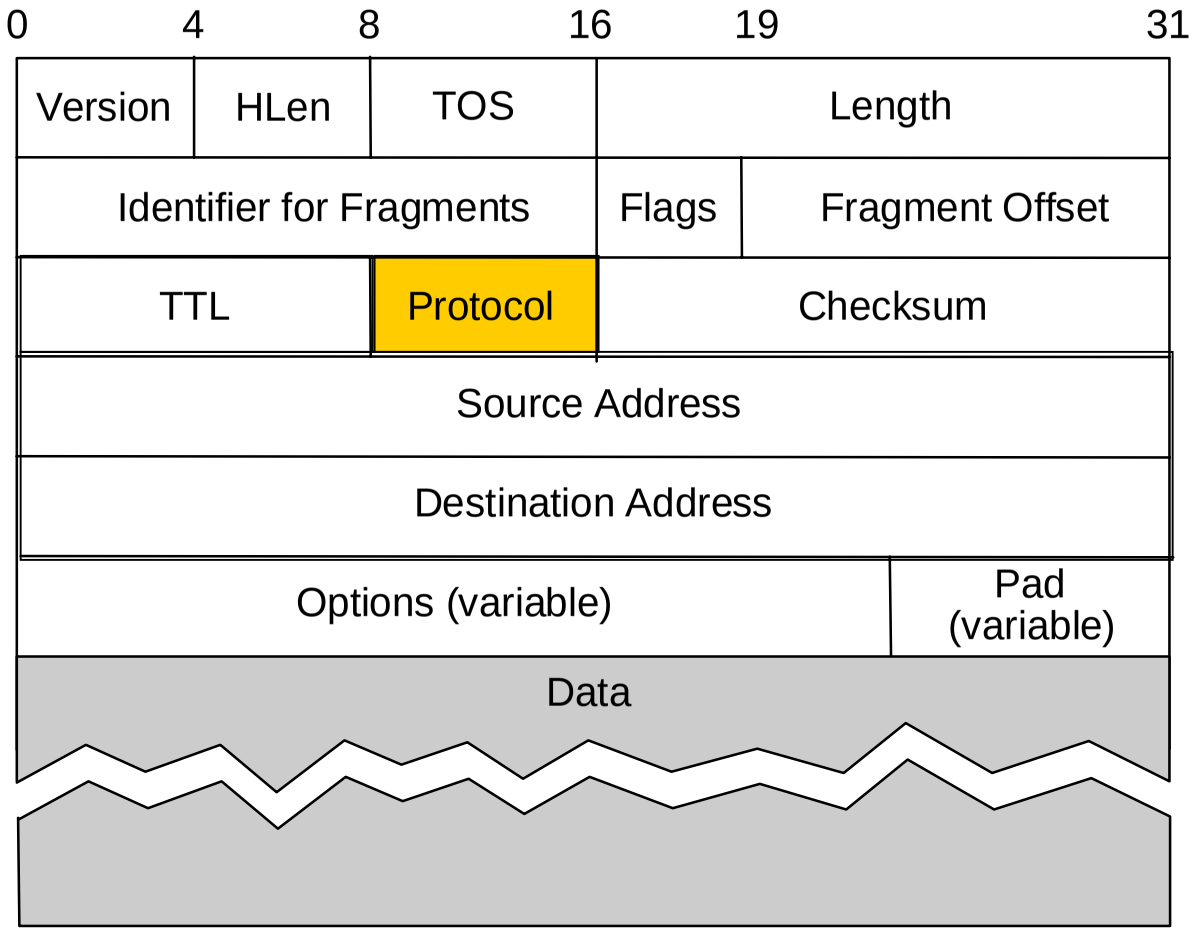
IPv4 Header Fields …

- Time To Live
- Decremented by router and packet discarded if = 0
- Prevents immortal packets



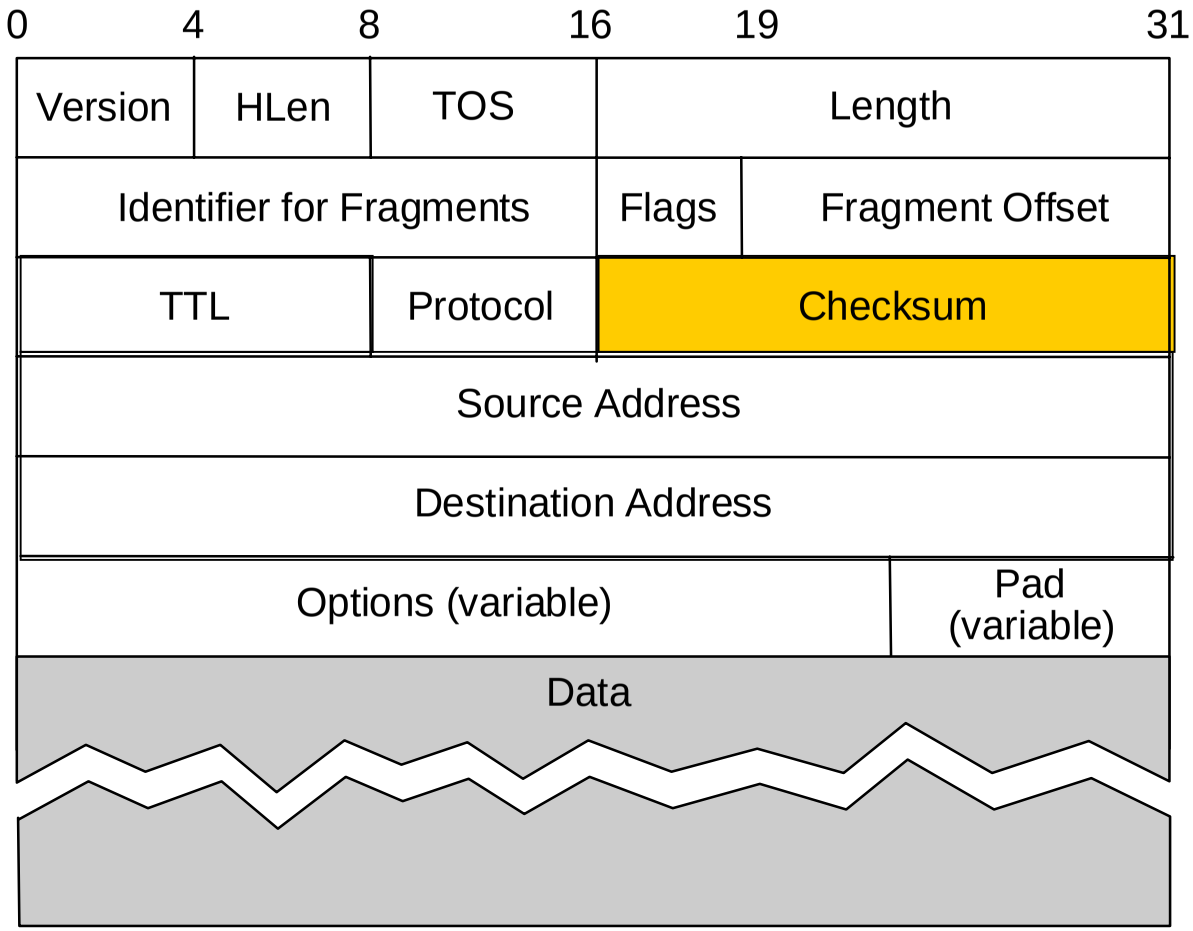
IPv4 Header Fields …

- Identifies higher layer protocol
  - E.g., TCP, UDP



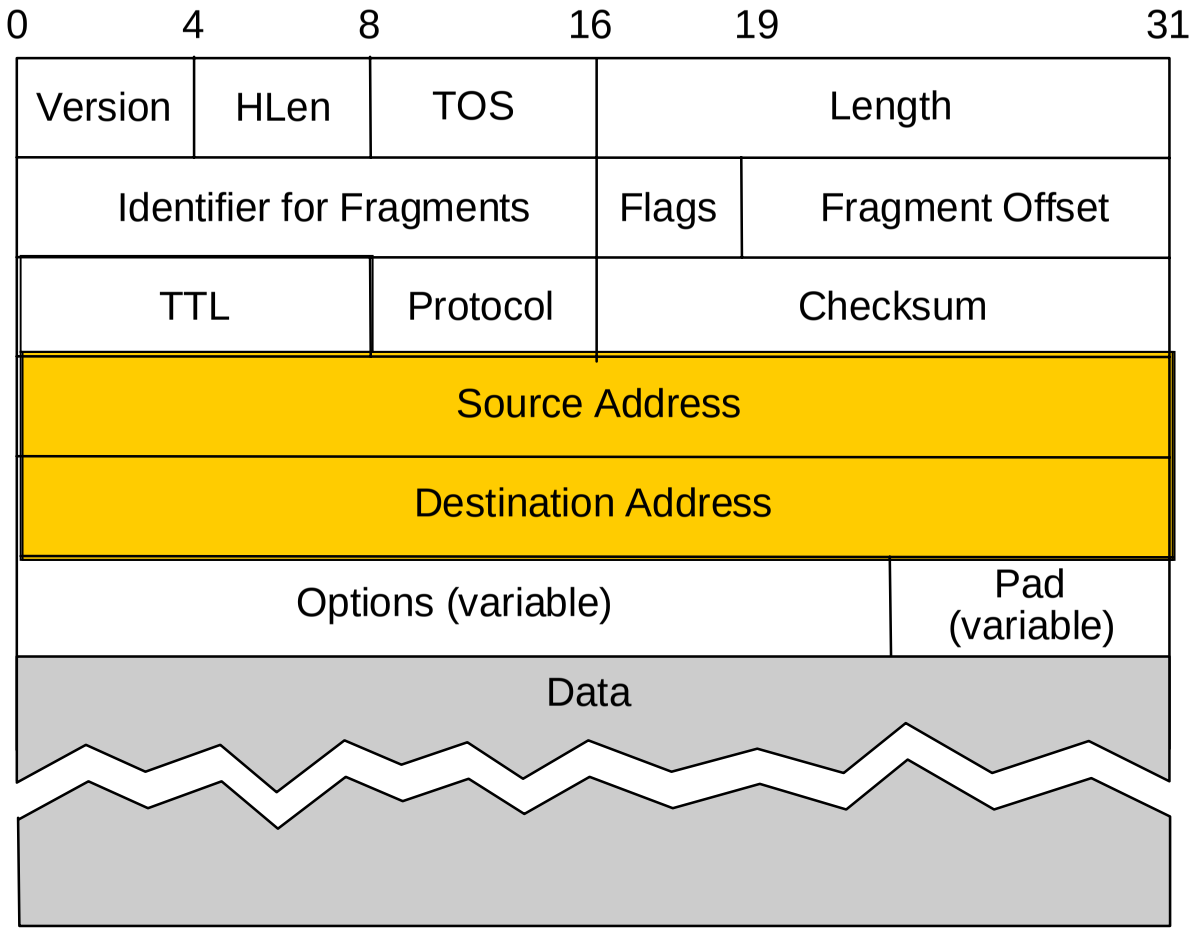
IPv4 Header Fields …

- Header checksum
- Recalculated by routers (TTL drops)
- Doesn't cover data
- Disappears for IPv6



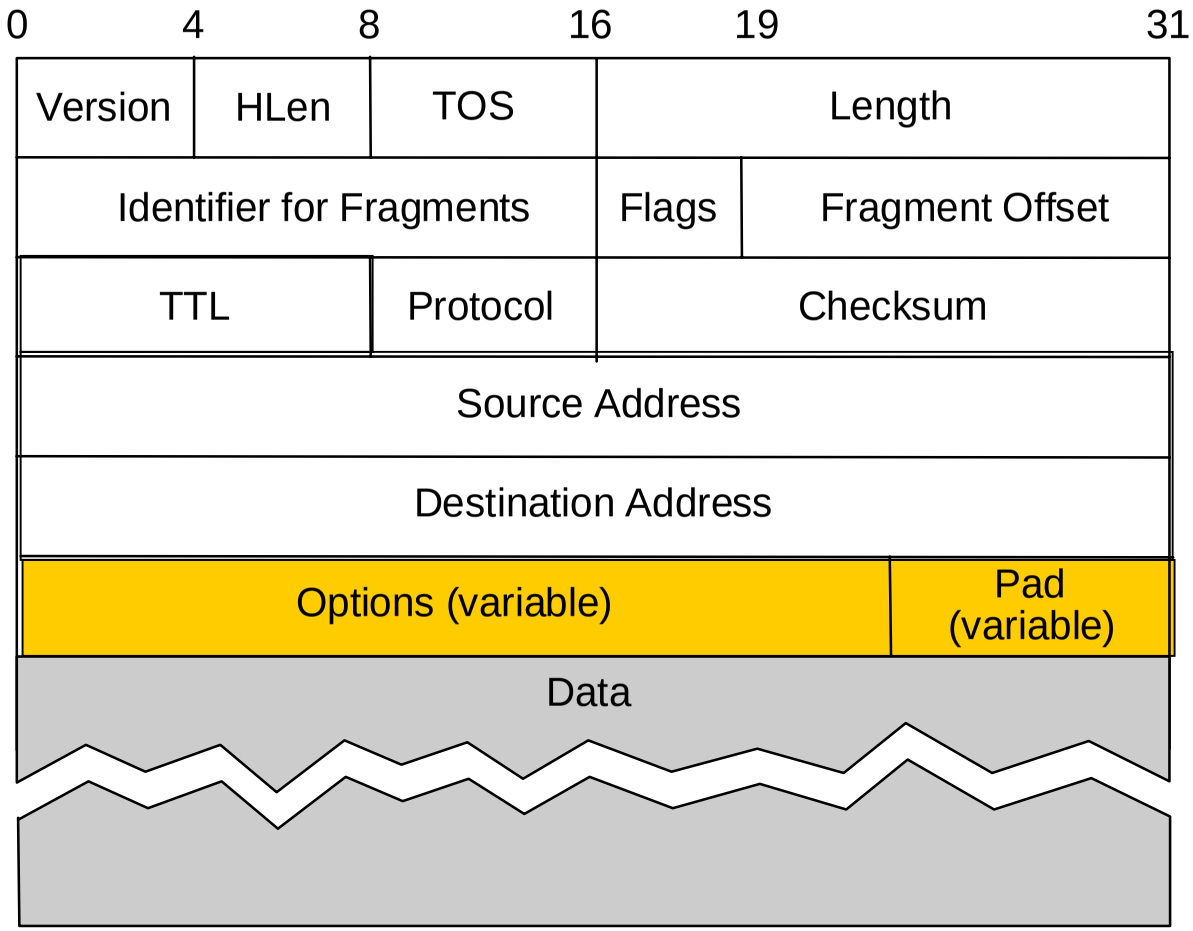
IPv4 Header Fields …

- Source/destination addresses
  - Not Ethernet
- Unchanged by routers
- Not authenticated by default



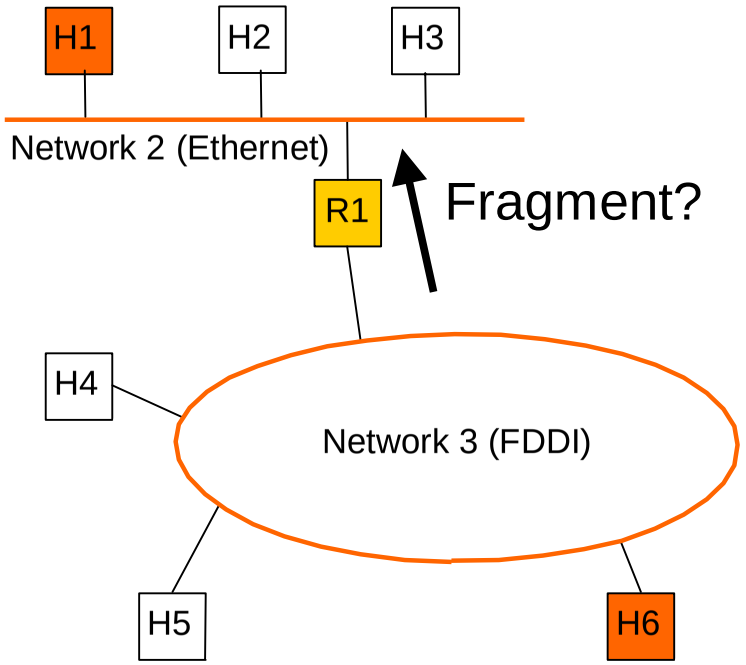
IPv4 Header Fields …

- IP options indicate special handling
  - Timestamps
  - "Source" routes
- Rarely used …



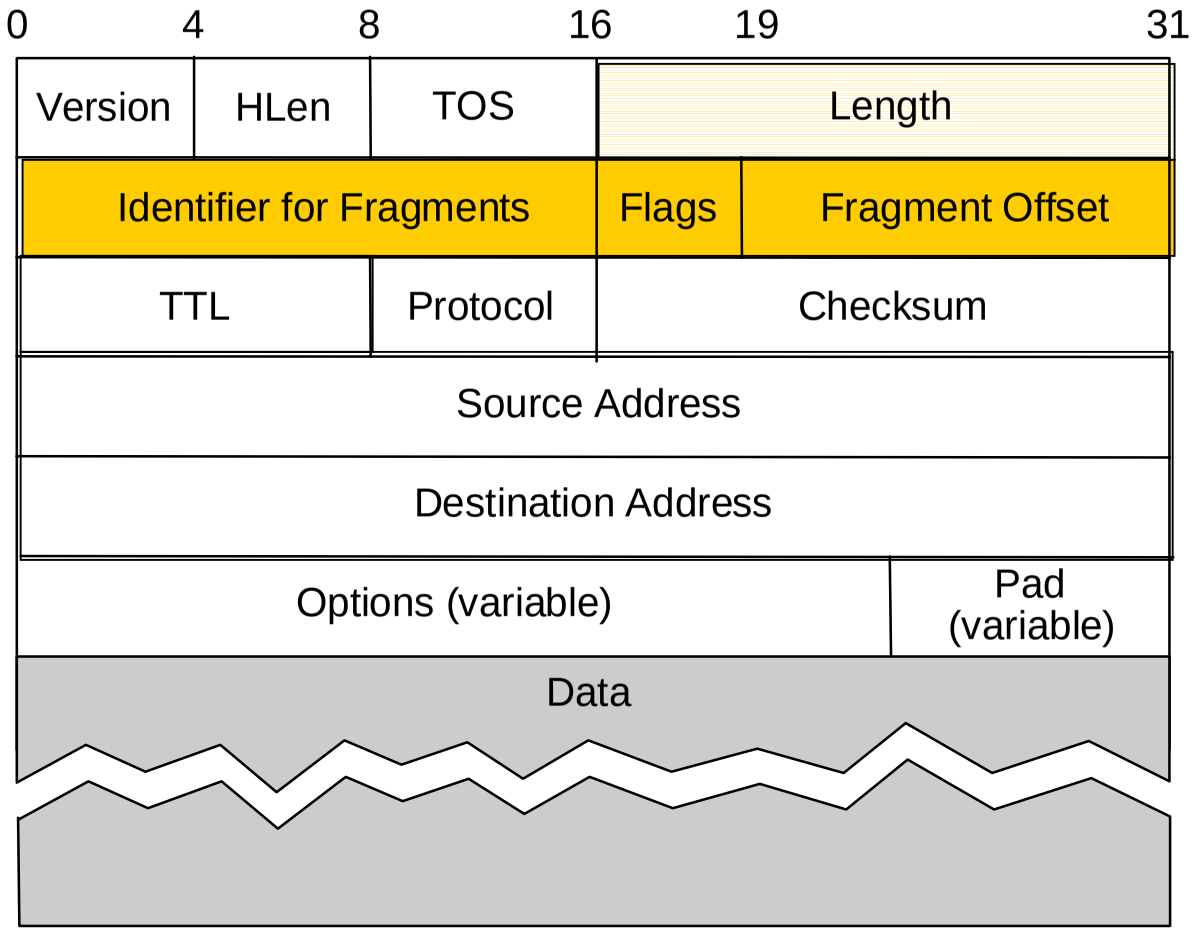
Fragmentation Issue

- Different networks may have different frame limits (MTUs)
  - Ethernet 1.5K, FDDI 4.5K
- Don't know if packet will be too big for path beforehand
  - IPv4: fragment on demand and reassemble at destination
  - IPv6: network returns error message so host can learn limit



Fragment Fields

- Fragments of one packet identified by (source, dest, frag id) triple
  - Make unique
- Offset gives start, length changed
- Flags are More Fragments (MF) Don't Fragment (DF)



Fragment Considerations

- Relating fragments to original datagram provides:
  - Tolerance of loss, reordering and duplication
  - Ability to fragment fragments

- Consequences of fragmentation:
  - Loss of any fragments causes loss of entire packet
  - Need to time-out reassembly when any fragments lost

Path MTU Discovery

- Path MTU is the smallest MTU along path
  - Packets less than this size don't get fragmented
- Fragmentation is a burden for routers
  - We already avoid reassembling at routers
  - Avoid fragmentation too by having hosts learn path MTUs
- Hosts send packets, routers return error if too large
  - Hosts discover limits, can fragment at source
  - Reassembly at destination as before
- Learned lesson from IPv4, streamlined in IPv6

Key Concepts

- Network layer provides end-to-end data delivery across an internetwork, not just a LAN
  - Issues of scale and heterogeneity
  - Datagram and virtual circuit service models
  - IP is the network layer protocol of the Internet

- Up next: More detailed look at routing and addressing