

## CSE/EE 461 – Lecture 2

### Protocols and Layering

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### Last Time ...

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- Networks are used to share distributed resources
  - Key problems revolve around effective resource sharing
- Statistical multiplexing
  - It's well-suited to data communications

## This Lecture

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1. A top-down look at the Internet
2. Mechanics of protocols and layering
3. The OSI/Internet models

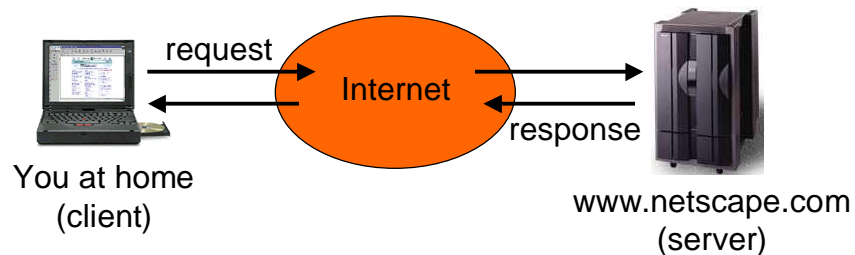
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L2.3

## 1. A Brief Tour of the Internet

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- What happens when you “click” on a web link?



- This is the view from 10,000 ft ...

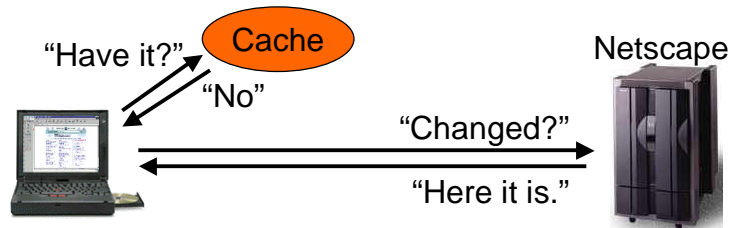
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## 9,000 ft: Scalability

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- Caching improves scalability



- We cut down on transfers:
  - Check cache (local or proxy) for a copy
  - Check with server for a new version

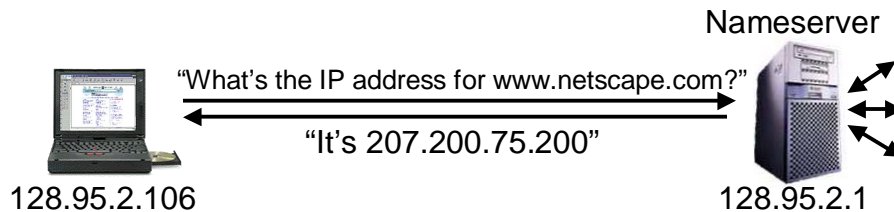
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L2.5

## 8,000 ft: Naming (DNS)

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- Map domain names to IP network addresses



- All messages are sent using IP addresses
  - So we have to translate names to addresses first
  - But we cache translations to avoid next time

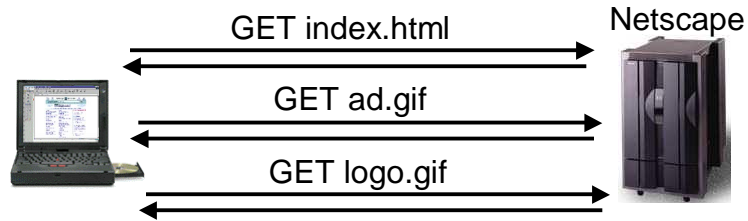
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## 7,000 ft: Sessions (HTTP)

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- A single web page can be multiple “objects”



- Fetch each “object”
  - either sequentially or in parallel

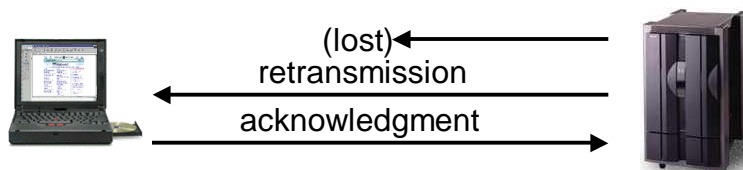
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## 6,000 ft: Reliability (TCP)

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- Messages can get lost



- We acknowledge successful receipt and detect and retransmit lost messages (e.g., timeouts)

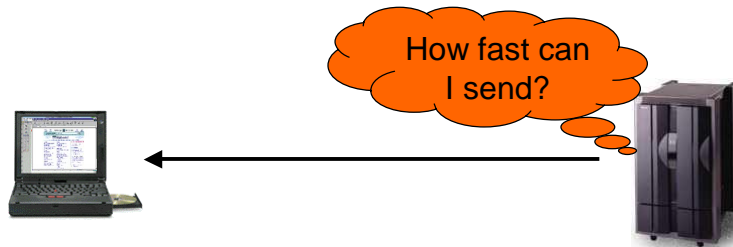
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## 5,000 ft: Congestion (TCP)

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- Need to allocate bandwidth between users



- Senders balance available and required bandwidths by probing network path and observing the response

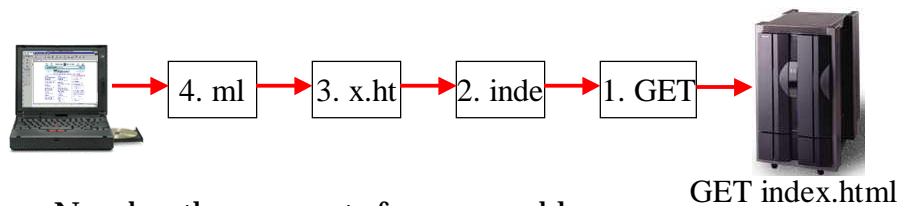
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## 4,000 ft: Packets (TCP/IP)

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- Long messages are broken into packets
  - Maximum Ethernet packet is 1.5 Kbytes
  - Typical web page is 10 Kbytes



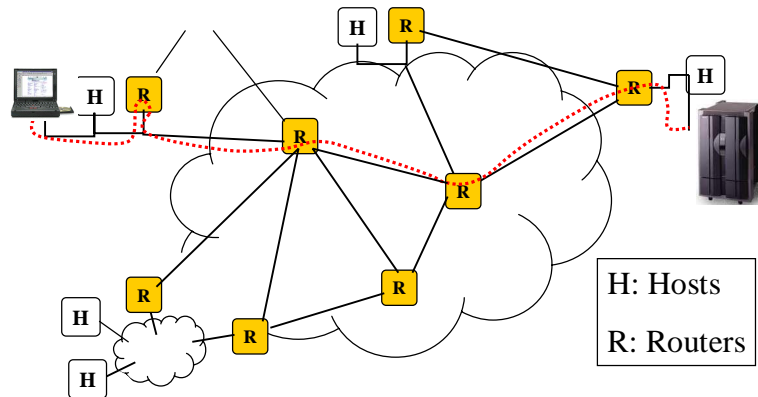
- Number the segments for reassembly

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## 3,000 ft: Routing (IP)

- Packets are directed through many routers

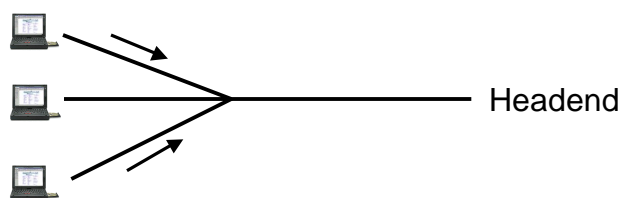


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## 2,000 ft: Multi-access (e.g., Cable)

- May need to share links with other senders



- Poll headend to receive a timeslot to send upstream
  - Headend controls all downstream transmissions
  - A lower level of addressing is used ...

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## 1,000 ft: Framing/Modulation

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- Protect, delimit and modulate payload as a signal

Sync / Unique	Header	Payload w/ error correcting code
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- E.g. for cable, take payload, add error protection (Reed-Solomon), header and framing, then turn into a signal
  - Modulate data to assigned channel and time (upstream)
  - Downstream, 6 MHz (~30 Mbps), Upstream ~2 MHz (~3 Mbps)

## 2. Protocols and Layering

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- We need abstractions to handle complexity
- A protocol is an agreement dictating the form and function of data exchanged between parties to effect communication
- Two parts:
  - Syntax: where the bits go
  - Semantics: what they mean, what to do with them
- Examples:
  - IP, the Internet protocol
  - TCP and HTTP, for the Web

## Protocol Standards

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- Different functions require different protocols
- Thus there are many protocol standards
  - E.g., IP, TCP, UDP, HTTP, DNS, FTP, SMTP, NNTP, ARP, Ethernet/802.3, 802.11, RIP, OSPF, 802.1D, NFS, ICMP, IGMP, DVMRP, IPSEC, PIM-SM, BGP, ...
- Organizations: IETF, IEEE, ITU
- IETF ([www.ietf.org](http://www.ietf.org)) specifies Internet-related protocols
  - RFCs (Requests for Comments)
  - “We reject kings, presidents and voting. We believe in rough consensus and running code.” – Dave Clark.

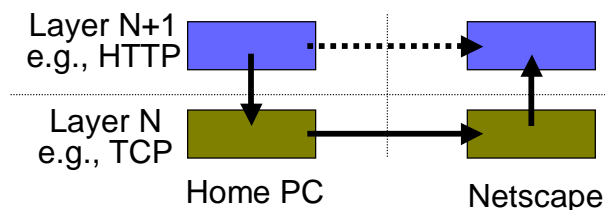
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## Layering and Protocol Stacks

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- Layering is how we combine protocols
  - Higher level protocols build on services provided by lower levels
  - Peer layers communicate with each other

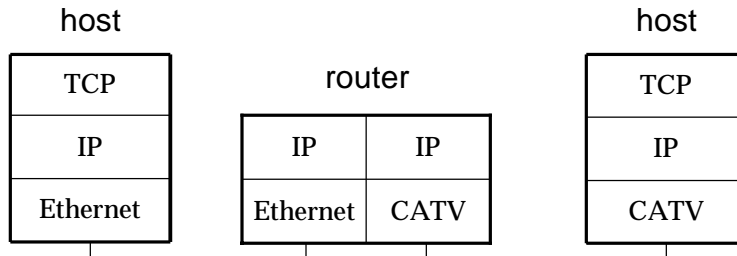


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## Example – Layering at work



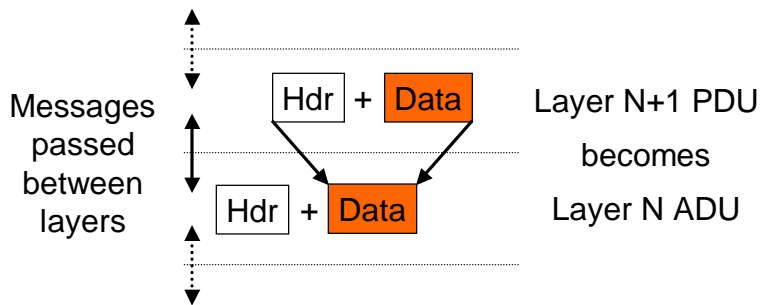
- We can connect different systems

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## Layering Mechanics

- Encapsulation and decapsulation



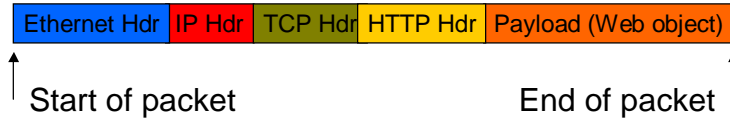
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## A Packet on the Wire

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- Starts looking like an onion!



- This isn't entirely accurate
  - ignores segmentation and reassembly, Ethernet trailers, etc.
- But you can see that layering adds overhead

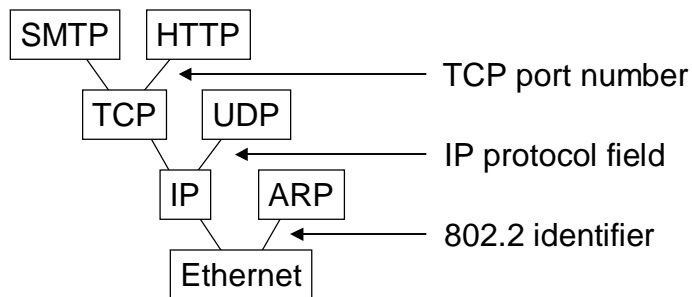
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## More Layering Mechanics

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- Multiplexing and demultiplexing in a protocol graph



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### 3. OSI/Internet Protocol Stacks

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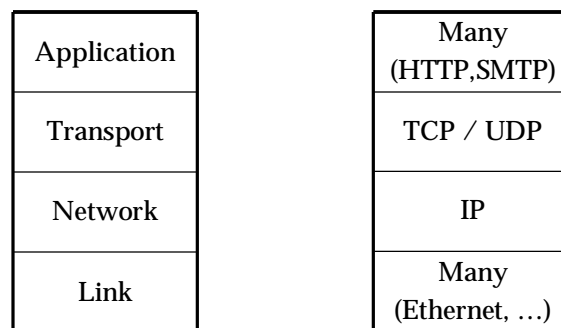
Key Question: What functionality goes in which protocol?

The “End to End Argument” (Reed, Saltzer, Clark, 1984):

- *Functionality should be implemented at a lower layer only if it can be correctly and completely implemented. (Sometimes an incomplete implementation can be useful as a performance optimization.)*
- Tends to push functions to the endpoints, which has aided the transparency and extensibility of the Internet.

### Internet Protocol Framework

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Model

Protocols

## OSI “Seven Layer” Reference Model

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- Seven Layers:

Application
Presentation
Session
Transport
Network
Link
Physical

Their functions:

- Your call
- Encode/decode messages
- Manage connections
- Reliability, congestion control
- Routing
- Framing, multiple access
- Symbol coding, modulation

## Key Concepts

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- Protocol layers are the modularity that is used in networks to handle complexity
- The Internet/OSI models give us a roadmap of what kind of function belongs at what layer