

## CSE/EE 461 – Lecture 20

### Congestion Avoidance

---

David Wetherall  
djw@cs.washington.edu

### Last Time ...

---

- Introduction to Quality of Service
- Focus
  - What transports do applications need?
- Topics
  - Real-time versus Elastic applications
  - Adapting to variable delay
  - Token buckets as bandwidth descriptors

Application
Presentation
Session
Transport
Network
Data Link
Physical

## This Lecture

---

- Congestion Avoidance
- Focus
  - How to we avoid congestion?
- Topics
  - Random Early Detection (RED) gateways
  - Explicit Congestion Notification (ECN)

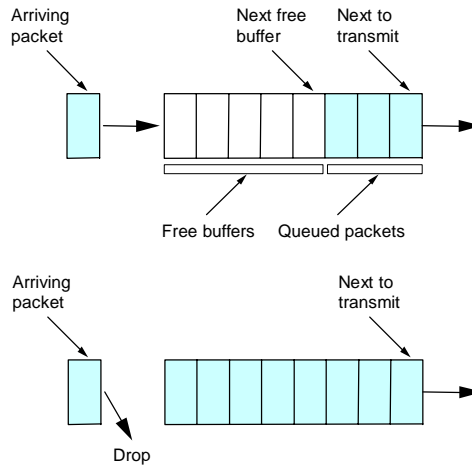
Application
Presentation
Session
Transport
Network
Data Link
Physical

## Why Congestion Avoidance?

---

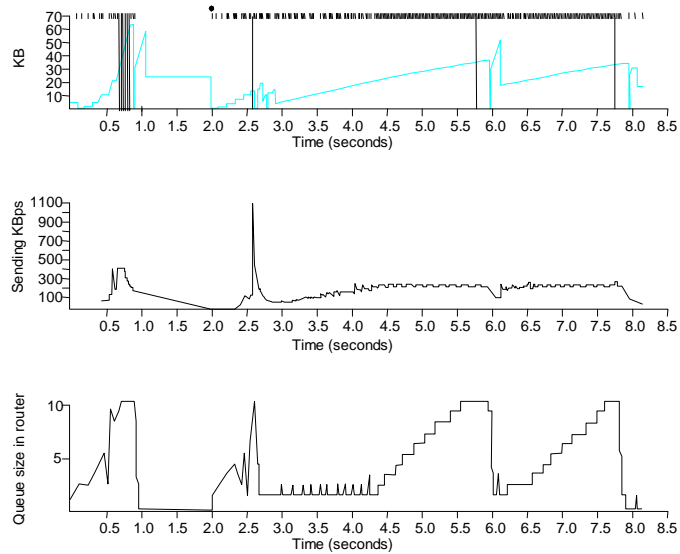
- TCP causes congestion as it probes for the available bandwidth and then recovers from it after the fact
  - Leads to loss, delay and bandwidth fluctuations (Yuck!)
  - We want congestion avoidance, not congestion control
- Congestion avoidance mechanisms
  - Aim to detect incipient congestion, before loss. So monitor queues to see that they absorb bursts, but not build steadily

# FIFO with Tail Drop



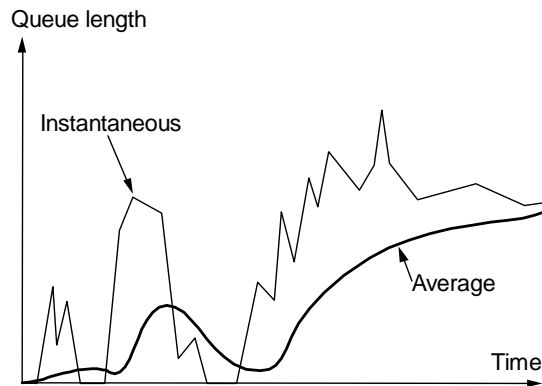
djw // CSE/EE 461, Autumn 2002

L20.5



## Incipient Congestion at a Router

- Sustained overload causes queue to build and overflow

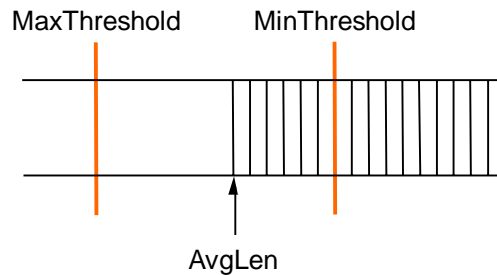


djw // CSE/EE 461, Autumn 2002

L20.7

## Random Early Detection (RED)

- Common approach is to have routers monitor average queue and send “early” signal to source when it builds by probabilistically dropping a packet



- Paradox: early loss can improve performance!

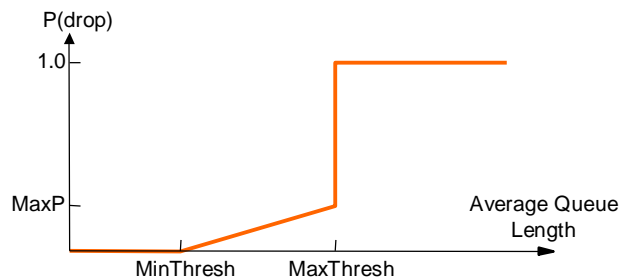
djw // CSE/EE 461, Autumn 2002

L20.8

## Red Drop Curve

---

- Start dropping a fraction of the traffic as queue builds
  - Expected drops proportional to bandwidth usage
  - When queue is too high, revert to drop tail
  - Nice theory, difficult to set parameters in practice



djw // CSE/EE 461, Autumn 2002

L20.9

## Explicit Congestion Notification (ECN)

---

- Why drop packets to signal congestion?
  - Drops are a robust signal, but there are other means ...
  - We need to be careful though: no extra packets
- ECN signals congestion with a bit in the IP header
- Receiver returns indication to the sender, who slows
  - Need to signal this reliably or we risk instability
- RED actually works by “marking” packets
  - Mark can be a drop or ECN signal if hosts understand ECN
  - Supports congestion avoidance without loss

djw // CSE/EE 461, Autumn 2002

L20.10

## Aside: TCP Vegas (Peterson '94)

---

- RED needs router upgrades but no host upgrades
- Instead, can we upgrade host but not router?
  
- TCP Vegas looks at the difference between cwnd (the amount of outstanding data in the network) and that acknowledged from the other side in the last interval
  - Excess must be buffered in the network at router queues
  - Vegas slows down when it believes there is a queue and otherwise increases to use the available bandwidth

## Key Concepts

---

- We want to avoid congestion rather than control it after it has occurred
  - Think of in terms of the queues at routers
  
- Random early packet drops, rather than tail drop, can have unintuitive advantages
  - Signal congestion early, before we're forced to drop repeatedly
  
- ECN signals congestion using bit in the IP header
  - No loss and no extra packets at overloaded times