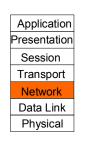
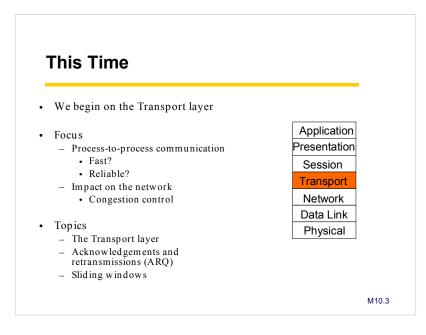
CSE 461 – Module 10

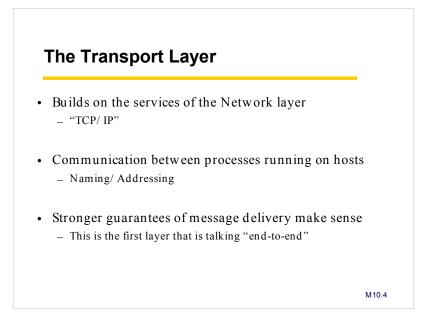
Introduction to the Transport Layer

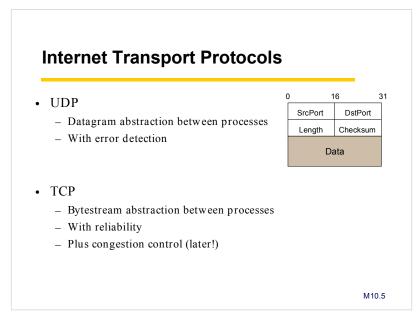
Last Time

- We finished up the Network layer – Internetworks (IP)
 - Routing (DV/ RIP, LS/ OSPF, BGP)
- It was all about routing: how to provide end-to-end delivery of packets.









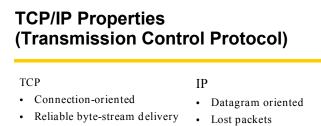
UDP/IP Properties (User Datagram Protocol)

UDP

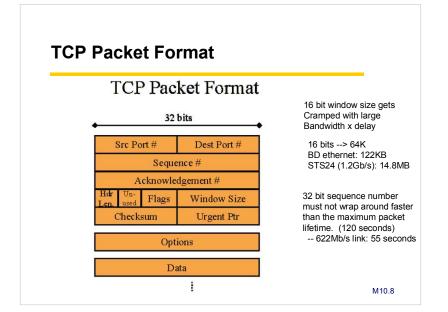
- Datagram oriented
- Lost packets
- Reordered packets
- Duplicate packets
- Limited size packets

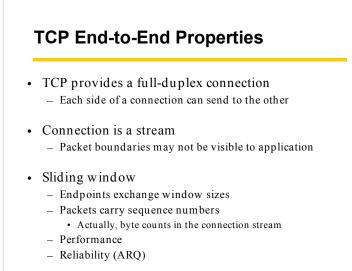
IP

- Datagram oriented
- Lost packets
- Reordered packets
- Duplicate packets
 - Limited size packets



- In-order delivery
- Single delivery
- Arbitrarily long messages
- Synchronization •
- Flow control ٠
- Congestion control
- Reordered packets •
- Duplicate packets
- Limited size packets

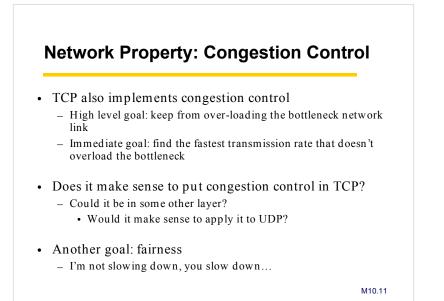




M10.9

End-to-end Properties

- Performance
 - Sliding Window
 - Try to enable sender to put bandwidth *x* delay product bytes on the wire
- Reliability
 - Lost packets?
 - Sliding window performs flow control
 - Sliding window performs ARQ (Automatic Repeat Request)
 - Duplicate / out-of-order packets?
 - Sliding window receive (re-order) buffer



ТСР	UDP
Reliable	Unreliable
Stream-oriented	Packet-oriented
Connection	Connectionless



- Stream- vs. packet-oriented
 - Visible packet boundaries can act as "end of record" indicators to application
 - In a stream, if the application wants the notion of "records", it must embed them in the data
 - Example: lines in a text file
 - Since TCP doesn't know about app record boundaries, reading records can be cumbersome
 - Each read() operation returns whatever data happens to have arrived in the stream to this point

M10.13

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