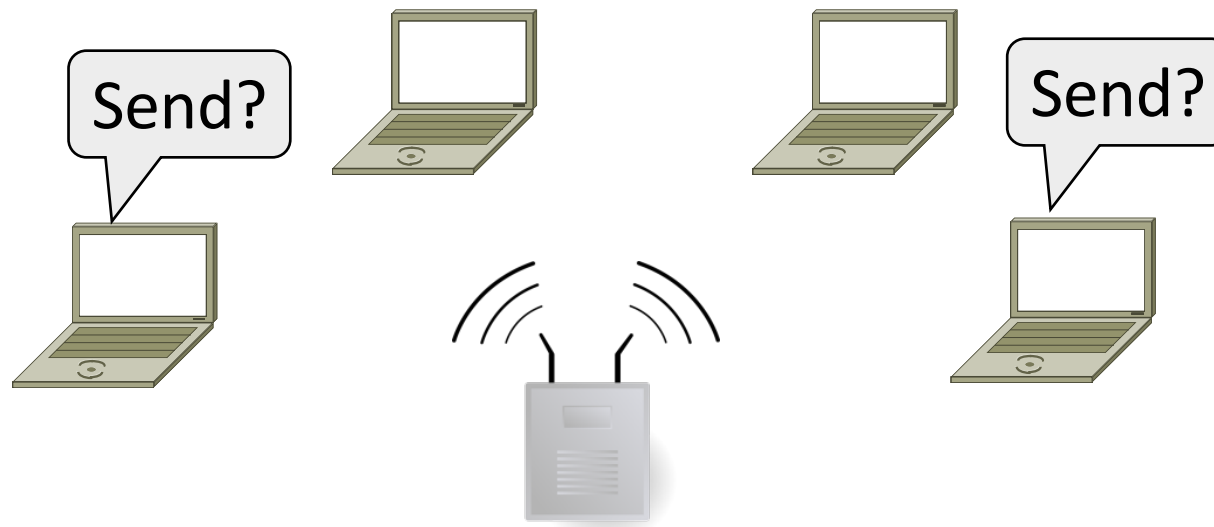


Wireless MACs

- How do wireless nodes share a single link? (Yes, this is WiFi!)
 - Build on our simple, wired model



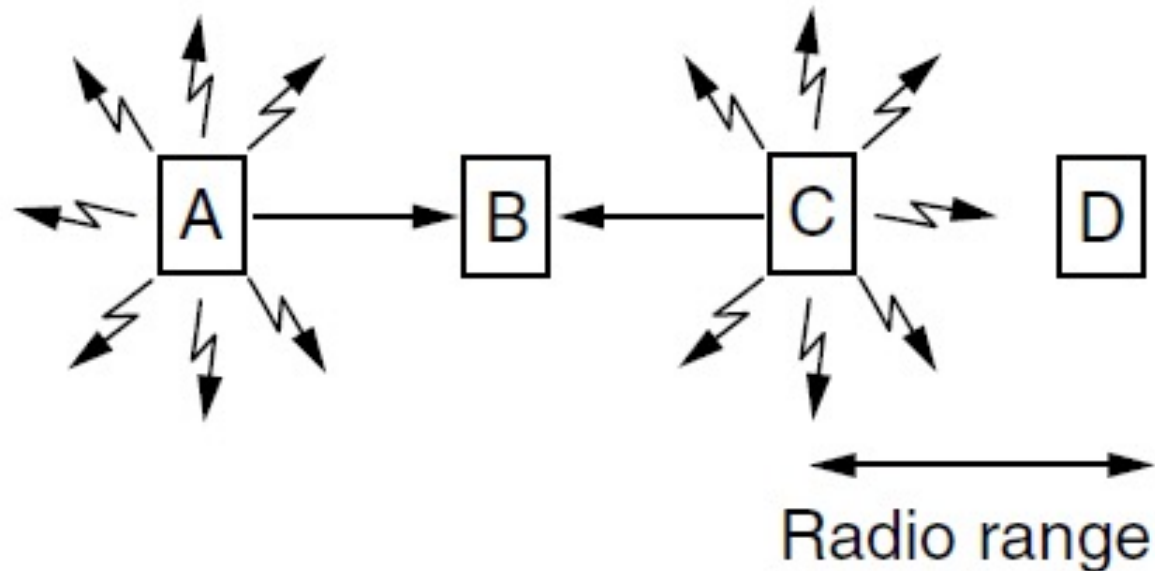
Wireless Complications

- Wireless is more complicated than wired (surprise!)
 1. Media is infinite – can't reliably Carrier Sense
 2. Nodes usually can't hear while sending – can't Collision Detect



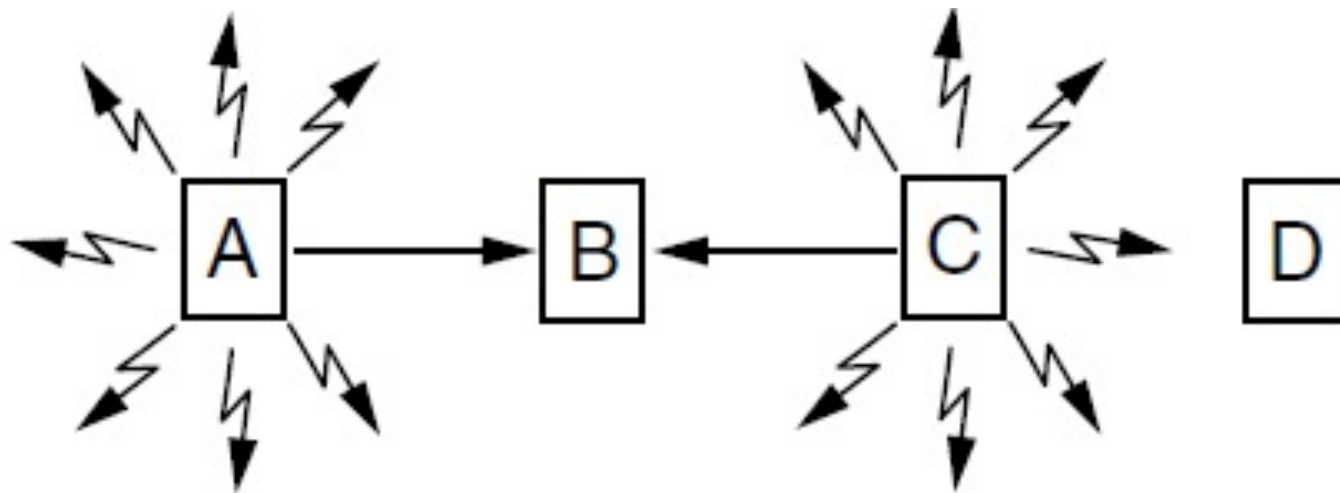
No CS: Different Coverage Areas

- Wireless signal is broadcast and received nearby, where there is sufficient SNR



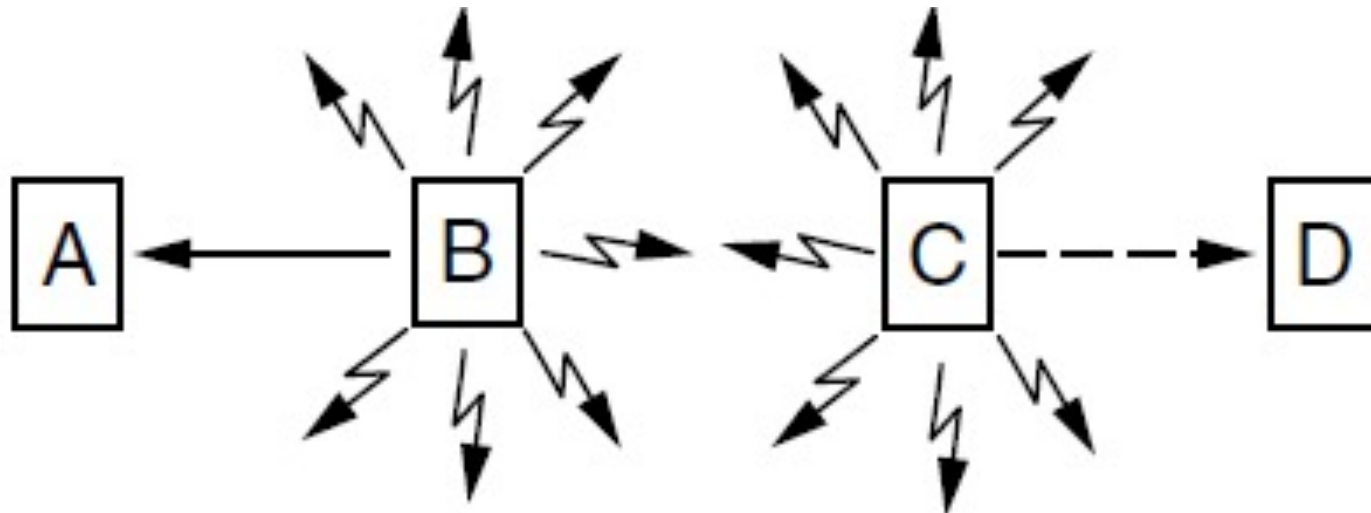
No CS: Hidden Terminals

- Node C is a hidden terminal when A sends to B
 - Similarly, A is a hidden terminal when C sends to B
 - A, C can't hear each other (to coordinate) yet collide at B
 - We want to avoid the inefficiency of collisions



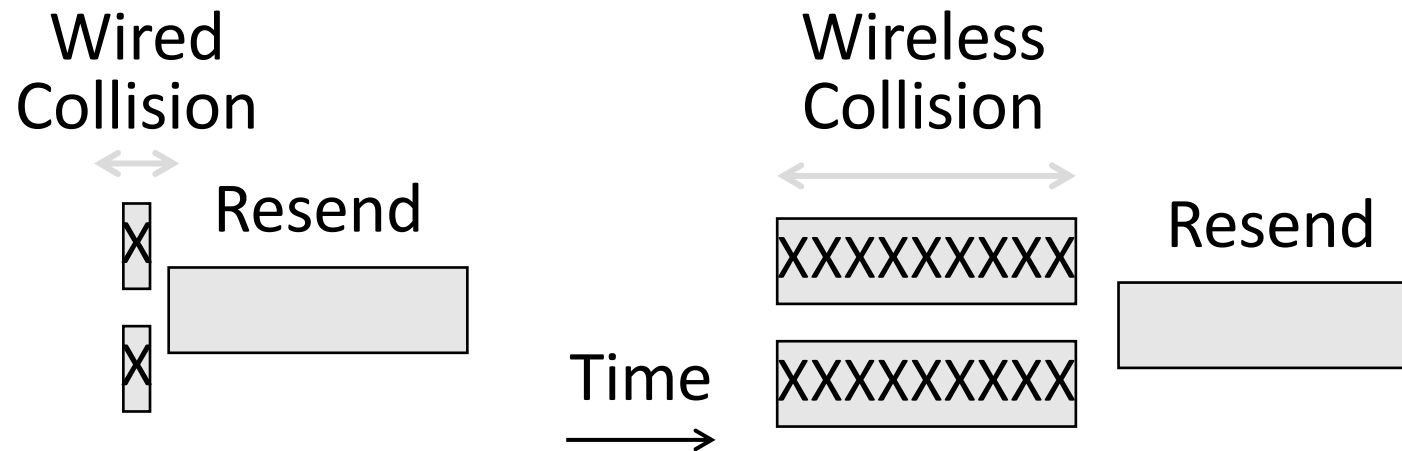
No CS: Exposed Terminals

- B, C are exposed terminals when sending to A, D
 - Can hear each other yet don't collide at receivers A and D
 - We want to send concurrently to increase performance



Nodes Can't Hear While Sending

- With wires, detecting collisions (and aborting) lowers their cost
- With wireless, more wasted time



Wireless Problems:

- Ideas?

MACA: Multiple Access w/ Collision Avoidance

- MACA uses a short handshake instead of CSMA (Karn, 1990)
 - 802.11 uses a refinement of MACA (later)
- Protocol rules:
 1. A sender node transmits a RTS (Request-To-Send, with frame length)
 2. The receiver replies with a CTS (Clear-To-Send, with frame length)
 3. Sender transmits the frame while nodes hearing the CTS stay silent
- Collisions on the RTS/CTS are still possible, but less likely

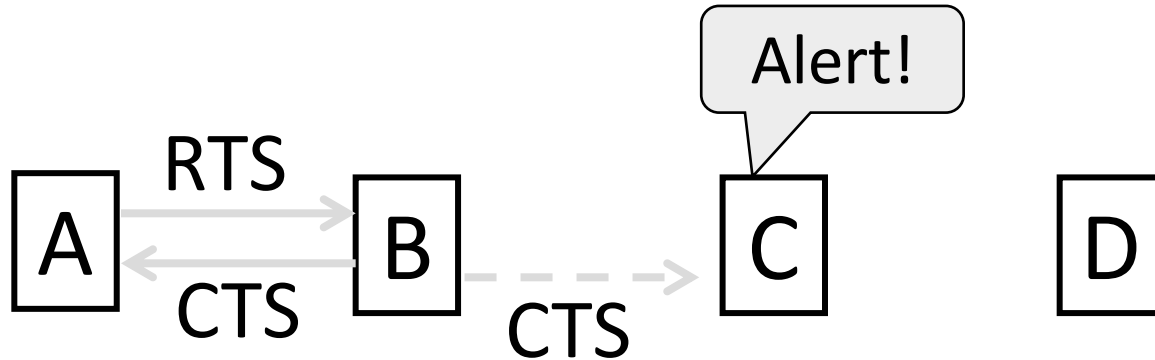
MACA – Hidden Terminals

- $A \rightarrow B$ with hidden terminal C
 1. A sends RTS, to B



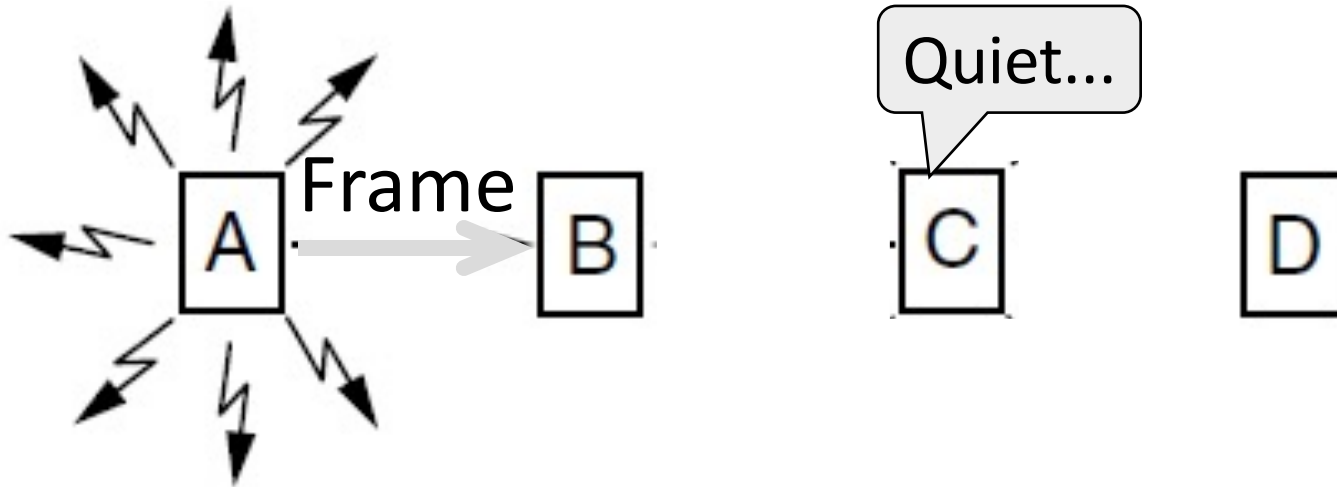
MACA – Hidden Terminals (2)

- A → B with hidden terminal C
 2. B sends CTS to A, and C overhears



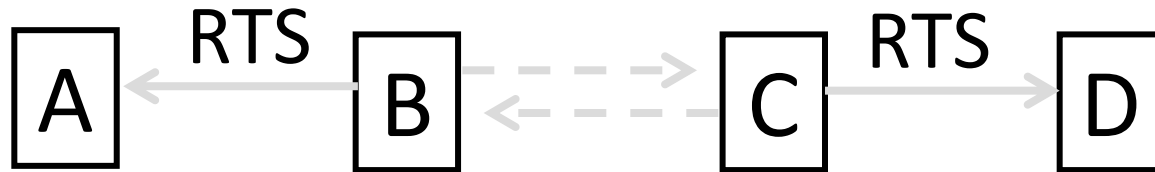
MACA – Hidden Terminals (3)

- $A \rightarrow B$ with hidden terminal C
 3. A sends frame while C defers



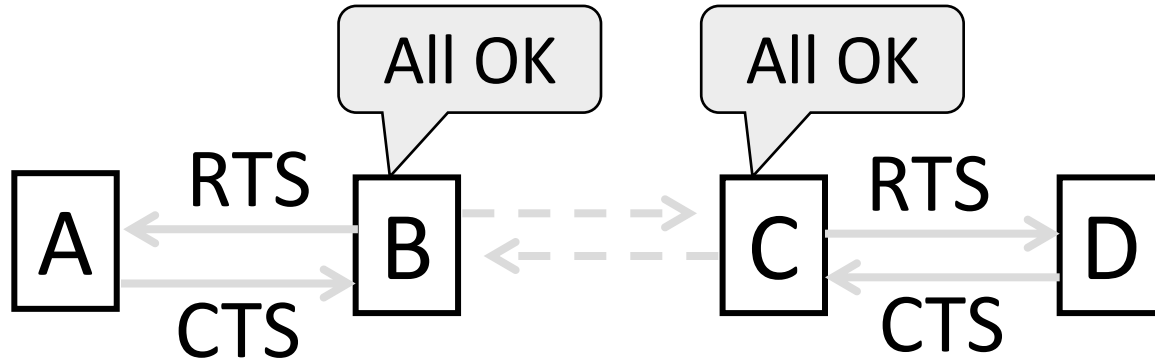
MACA – Exposed Terminals

- $B \rightarrow A$, $C \rightarrow D$ as exposed terminals
 - B and C send RTS to A and D



MACA – Exposed Terminals (2)

- $B \rightarrow A$, $C \rightarrow D$ as exposed terminals
 - A and D send CTS to B and C



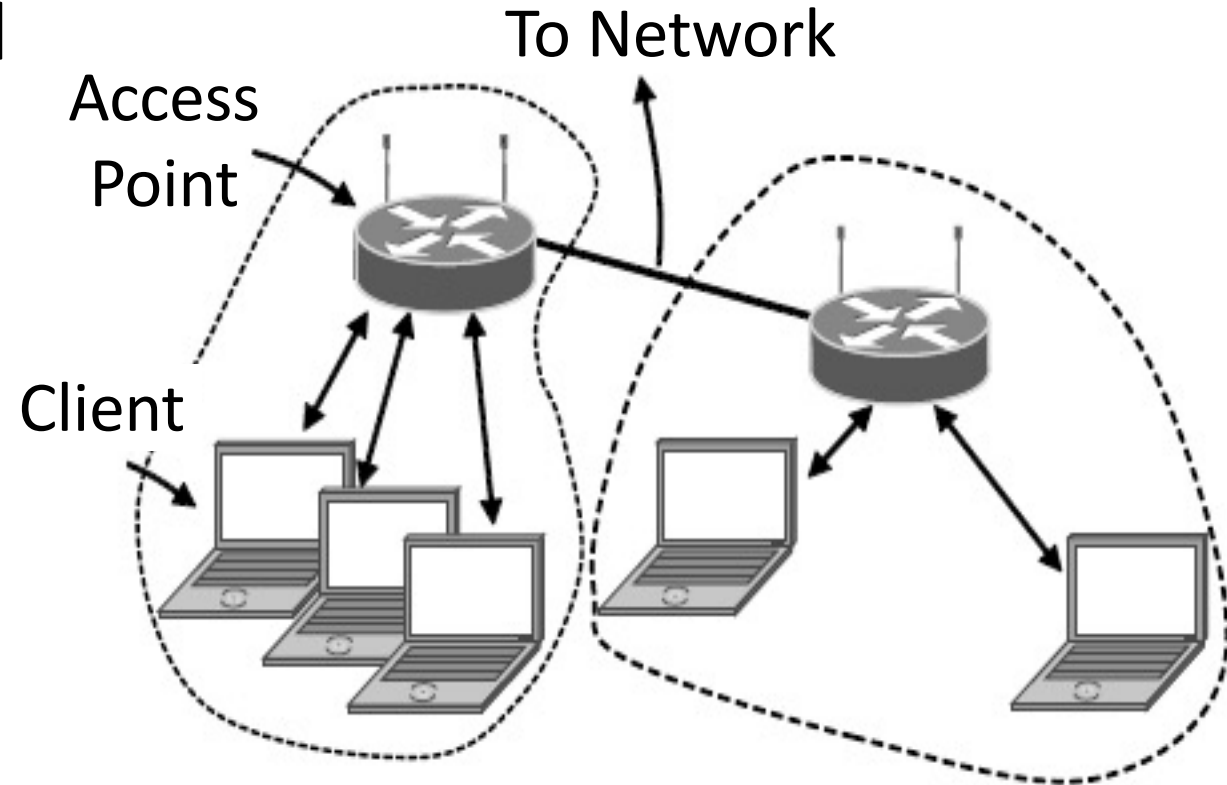
MACA – Exposed Terminals (3)

- $B \rightarrow A$, $C \rightarrow D$ as exposed terminals
 - A and D send CTS to B and C



802.11, or WiFi

- Very popular wireless LAN started in the 1990s
- Clients get connectivity from a (wired) AP (Access Point)
- It's a multi-access problem 😊
- Various flavors have been developed over time
 - Faster, more features

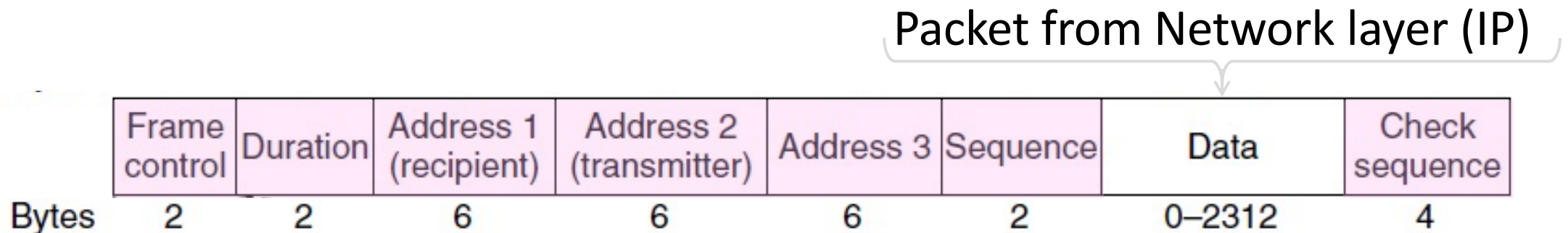


802.11 Physical Layer

- Uses 20/40 MHz channels on ISM (unlicensed) bands
 - 802.11b/g/n on 2.4 GHz
 - 802.11 a/n on 5 GHz
- OFDM modulation (except legacy 802.11b)
 - Different amplitudes/phases for varying SNRs
 - Rates from 6 to 54 Mbps plus error correction
 - 802.11n uses multiple antennas
 - Lots of fun tricks here

802.11 Link Layer

- Multiple access uses CSMA/CA (next); RTS/CTS optional
- Frames are ACKed and retransmitted with ARQ
- Funky addressing (three addresses!) due to AP
- Errors are detected with a 32-bit CRC
- Many, many features (e.g., encryption, power save)



Centralized MAC: Cellular

- Spectrum suddenly very scarce
 - We can't waste all of it sending JAMs
- We have QoS requirements
 - Can't be as loose with expectations
 - Can't have traffic fail
- We also have client/server
 - Centralized control
 - Not peer-to-peer/decentralized



GSM MAC

- FDMA/TDMA
- Use one channel for coordination – Random access w/BEB (no CSMA, can't detect)
- Use other channels for traffic
 - Dedicated channel for QoS

Nedlink (Basestasjon->Mobiltelefon)

0	1	2-5	6-9	10	11	12-19	20	21	22-29	30	31	32-39	40	41	42-49	50
FCH	SCH	BCCH	CCCH	FCH	SCH	CCCH	FCH	SCH	CCCH	FCH	SCH	CCCH	FCH	SCH	CCCH	IDLE

Opplink (Mobiltelefon->Basestasjon)

RACH 0	RACH 1	RACH	0-50	RACH	.	.	RACH	.	.	RACH	.	RACH	.	RACH	.	RACH 50
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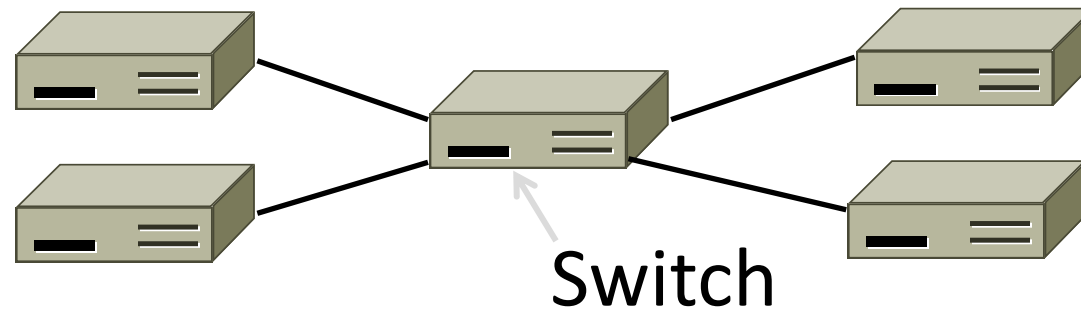
Recap: MAC layer ideas

- Random wait times upon collisions
- Carrier sense
 - Persistence
- Collision detection
- Binary exponential backoff
- RTS-CTS for hidden and exposed terminals

Link Layer: Switching

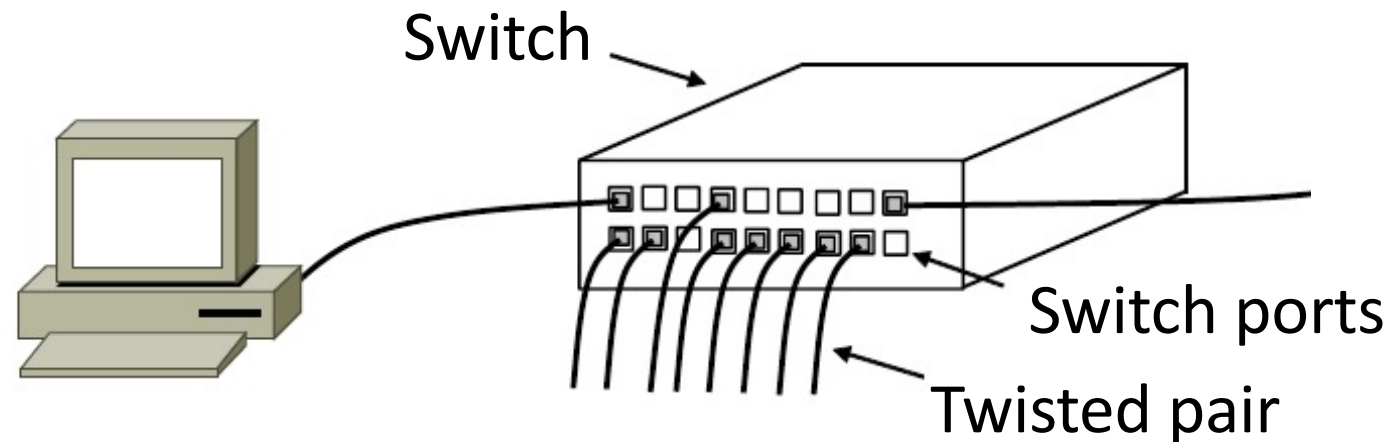
Switching

- How do we connect nodes with a switch instead of multiple access
 - Uses multiple links/wires
 - Basis of modern (switched) Ethernet



Switched Ethernet

- Hosts are wired to Ethernet switches with twisted pair
 - Switch serves to connect the hosts
 - Wires usually run to a closet



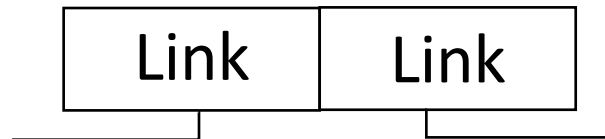
What's in the box?

- Remember from protocol layers:

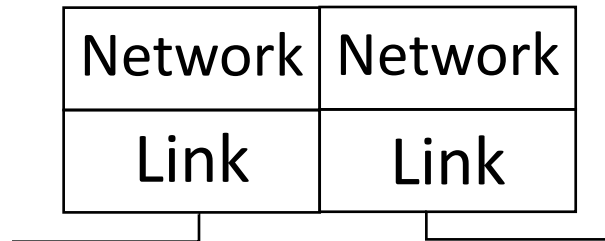
Hub, or
repeater



Switch



Router

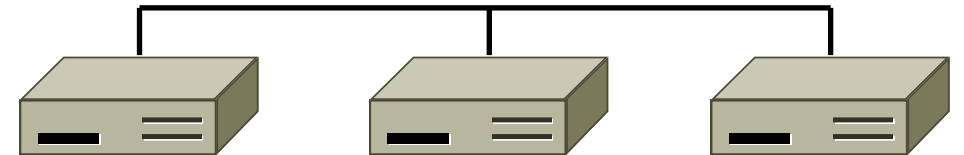
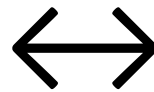
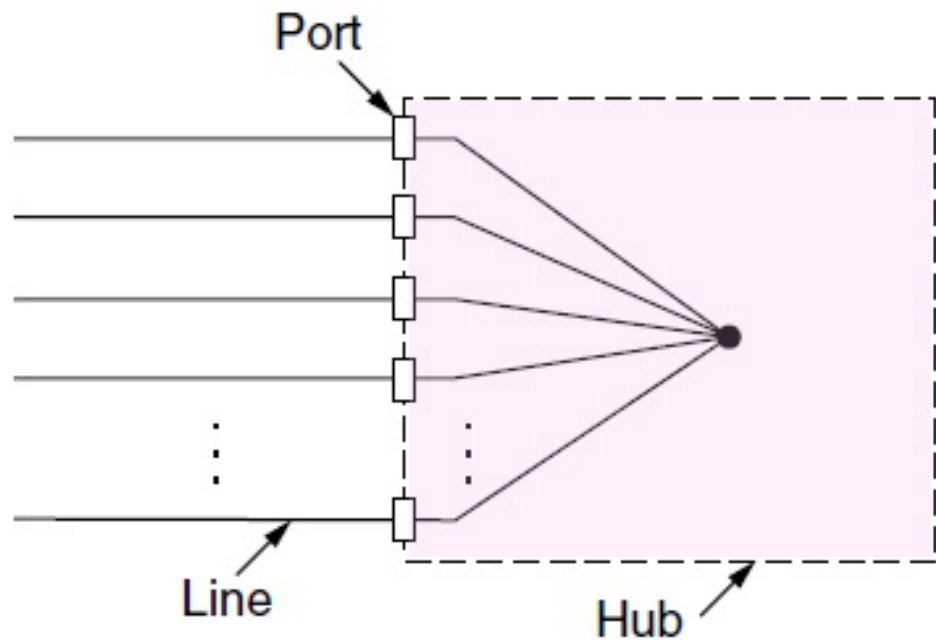


All look like this:



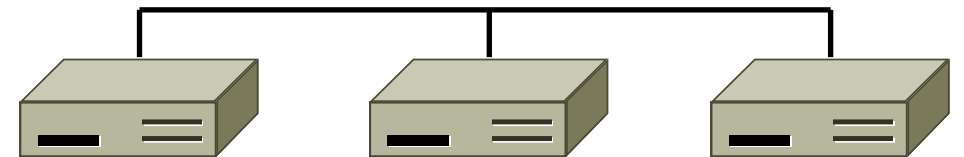
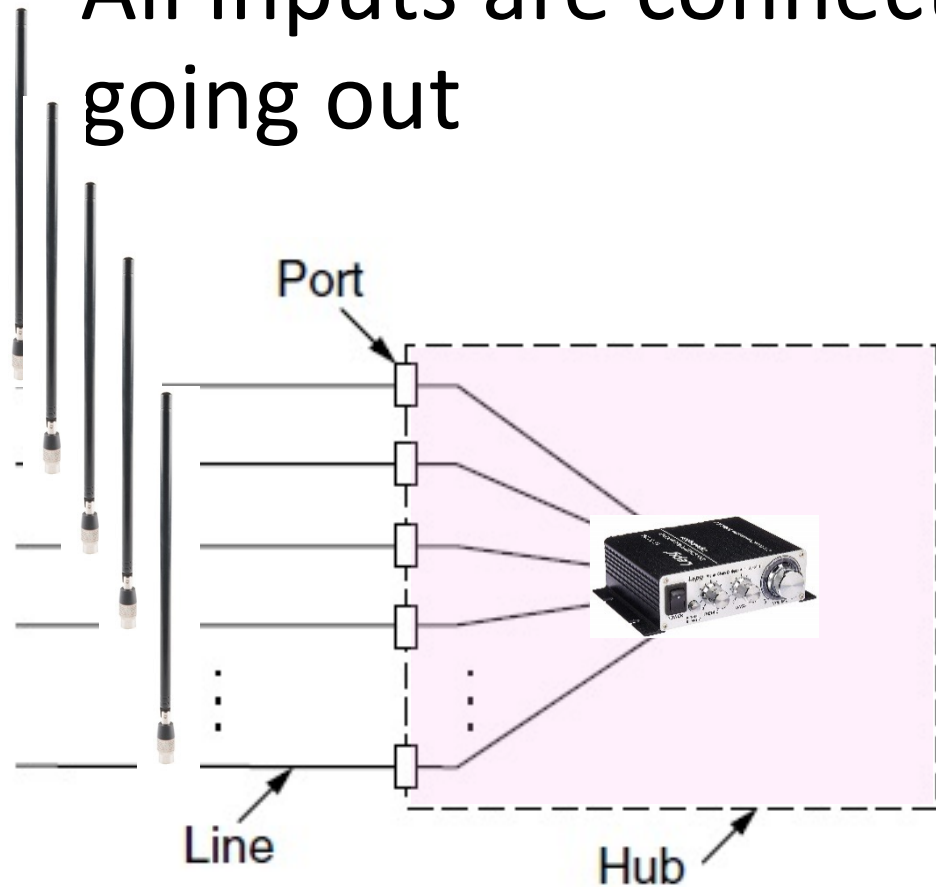
Inside a Hub

- All ports are wired together; more convenient and reliable than a single shared wire



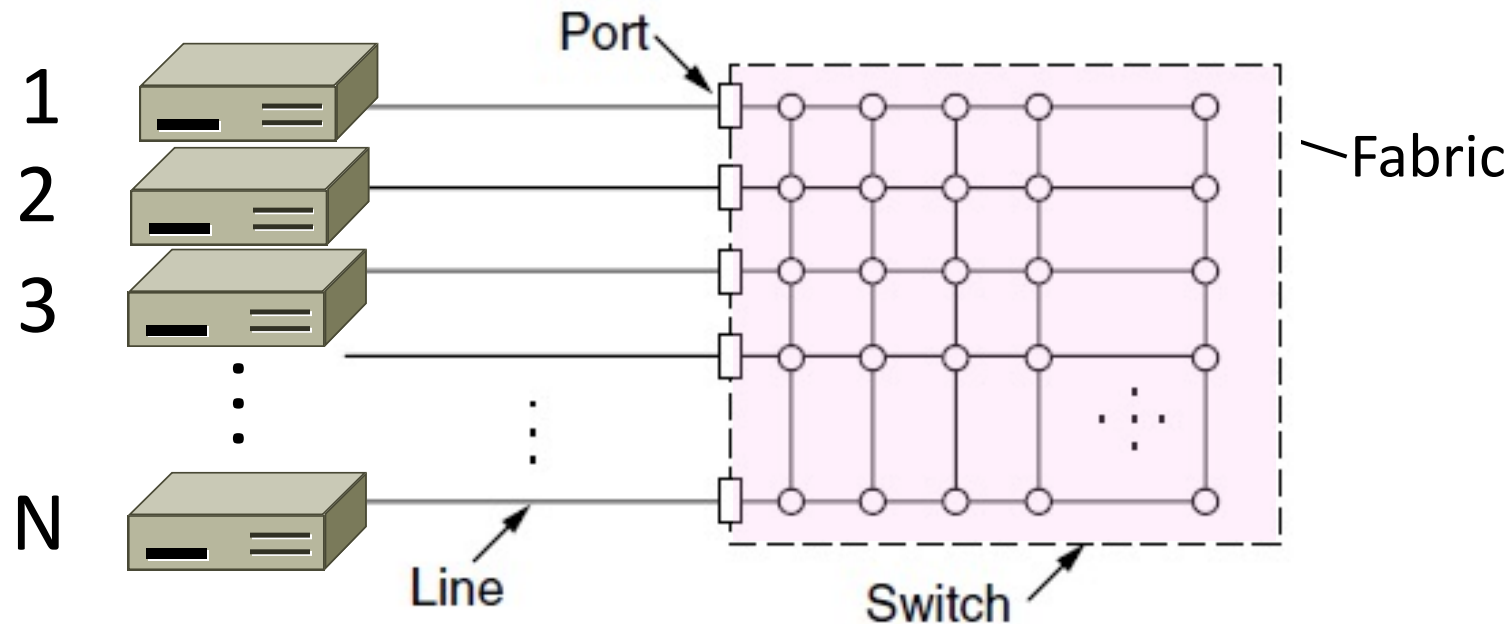
Inside a Repeater

- All inputs are connected; then amplified before going out



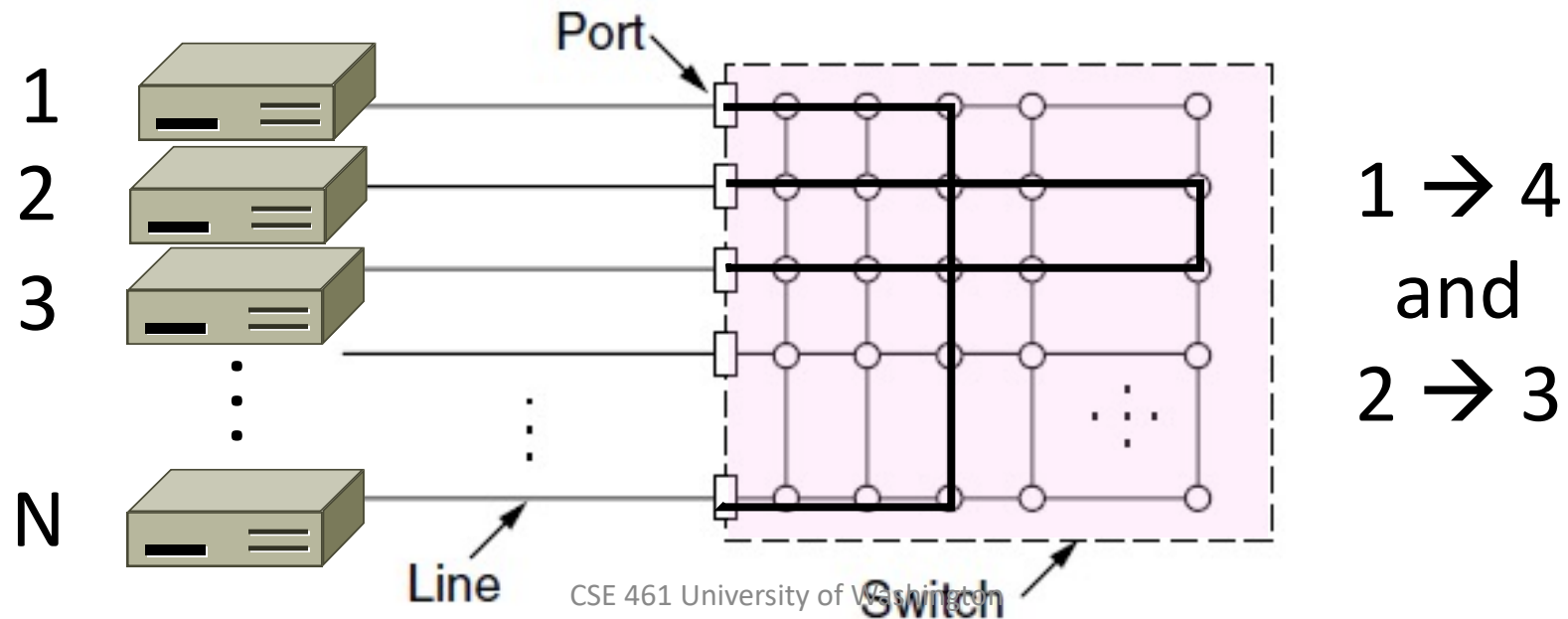
Inside a Switch

- Uses frame addresses (MAC addresses in Ethernet) to connect input port to the right output port; multiple frames may be switched in parallel



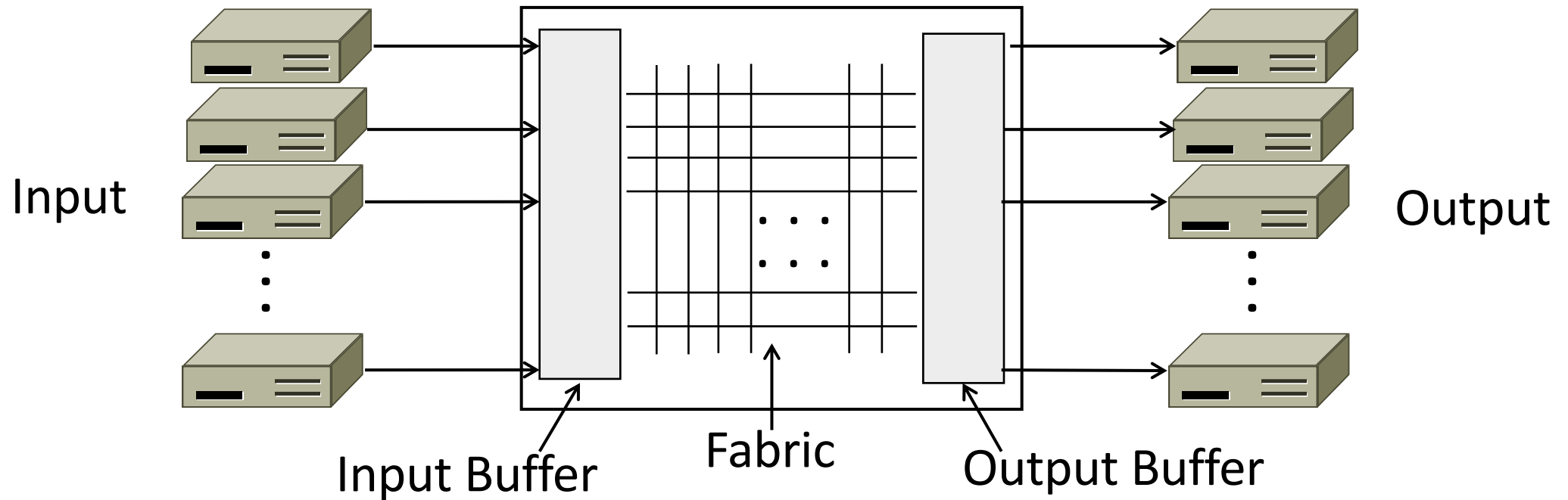
Inside a Switch (2)

- Port may be used for both input and output (full-duplex)
 - Just send, no multiple access protocol



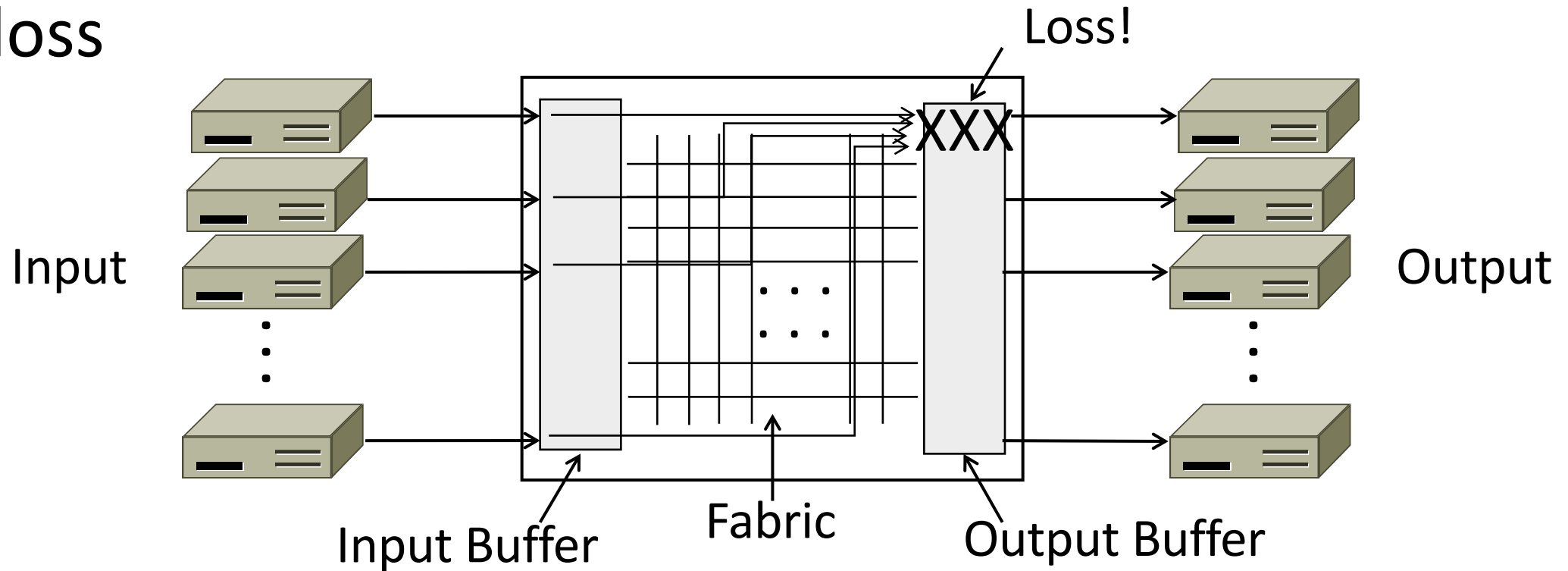
Inside a Switch (3)

- Need buffers for multiple inputs to send to one output



Inside a Switch (4)

- Sustained overload will fill buffer and lead to frame loss



Advantages of Switches

- Switches and hubs (mostly switches) have replaced the shared cable of classic Ethernet
 - Convenient to run wires to one location
 - More reliable; wire cut is not a single point of failure that is hard to find
- Switches offer scalable performance
 - E.g., 100 Mbps per port instead of 100 Mbps for all nodes of shared cable / hub