

Multiple access contd.

ALOHA recap

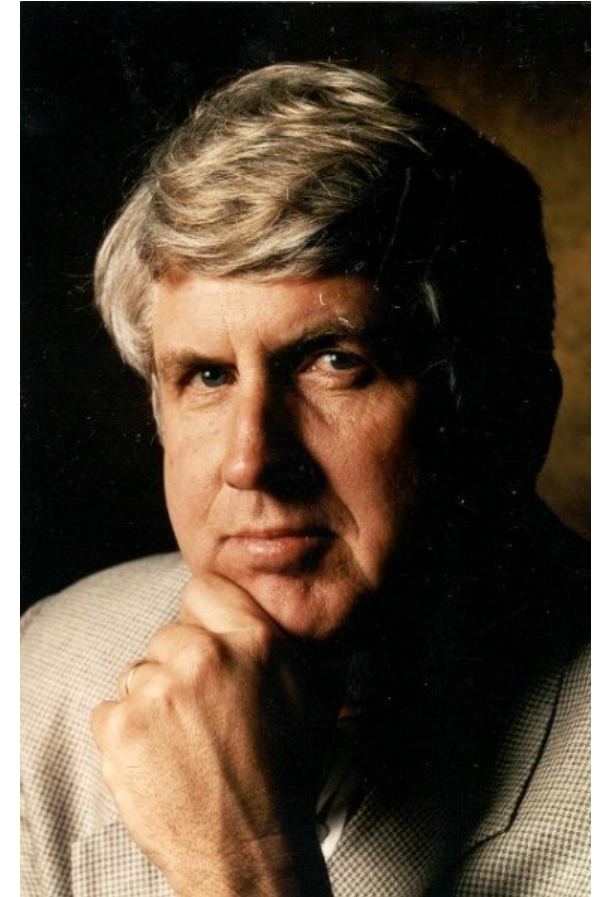
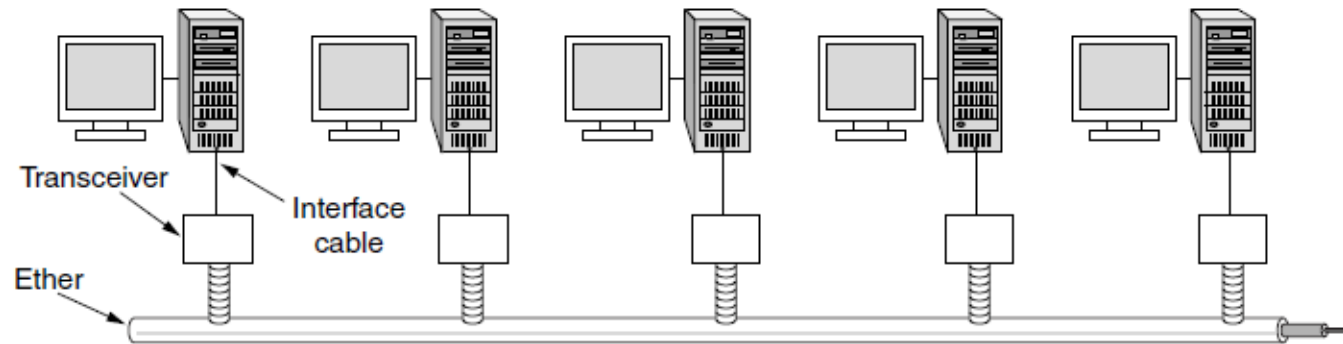
1. Send when you have to
2. If no ack, wait for a random time and send again

Super simple, but low efficiency

Classic Ethernet

ALOHA inspired Bob Metcalfe to invent Ethernet for LANs in 1973

- Nodes share 10 Mbps coaxial cable
- Hugely popular in 1980s, 1990s
- Turing award in 2023



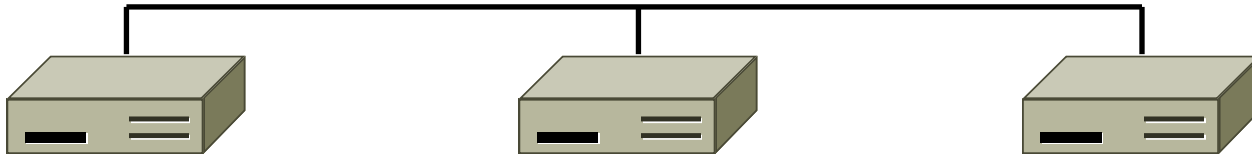
: © 2009 IEEE

CSMA (Carrier Sense Multiple Access)

- Improve ALOHA by listening for activity before we send (Doh!)
 - Easy with wires, recently made possible for wireless
- So does this eliminate collisions?
 - Why or why not?

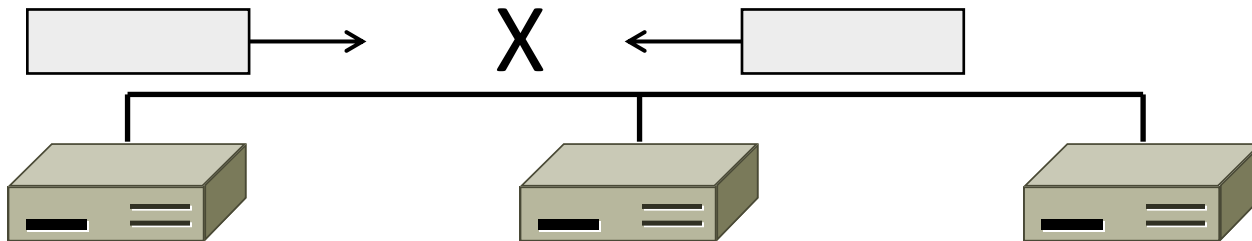
CSMA (2)

- Still possible to listen and hear nothing when another node is sending because of delay



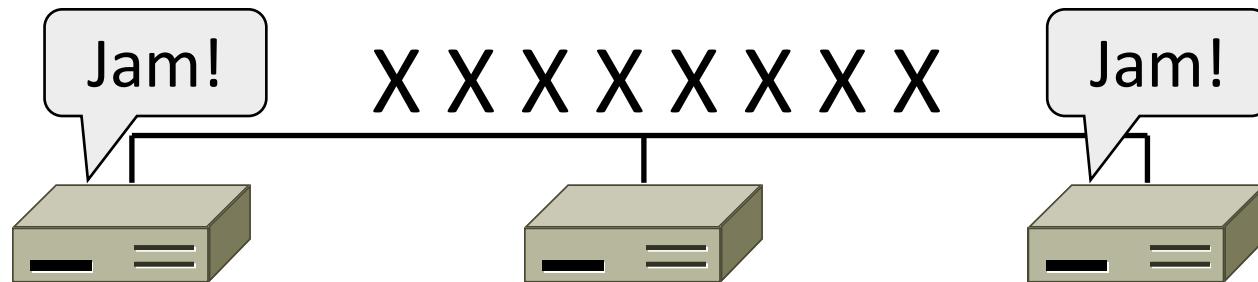
CSMA (3)

- CSMA is a good defense against collisions only when BD is small



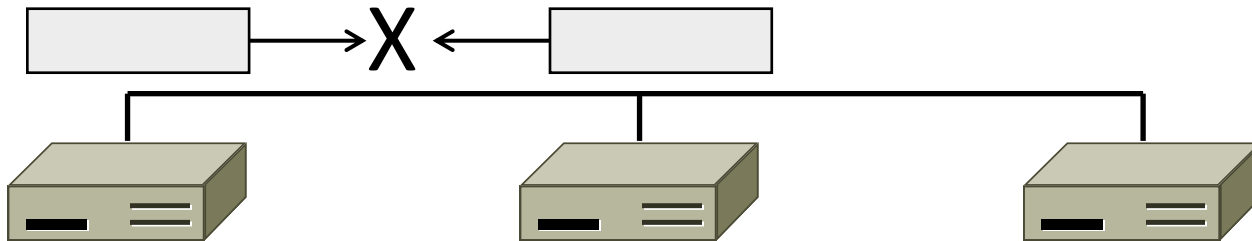
CSMA/CD (with Collision Detection)

- Can reduce the cost of collisions by detecting them and aborting (Jam) the rest of the frame time
 - Again, easy with wires, recently made possible for wireless



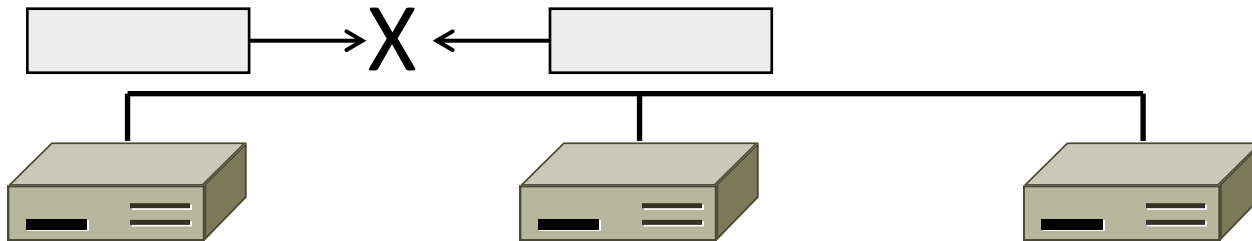
CSMA/CD Complications

- Everyone who collides needs to know it happened
 - How long do we need to wait to know there wasn't a JAM?



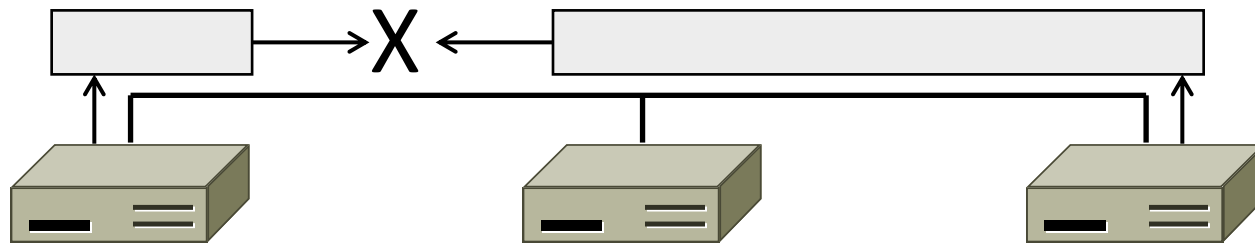
CSMA/CD Complications

- Everyone who collides needs to know it happened
 - How long do we need to wait to know there wasn't a JAM?
 - Time window in which a node may hear of a collision (transmission + jam) is $2D$ seconds



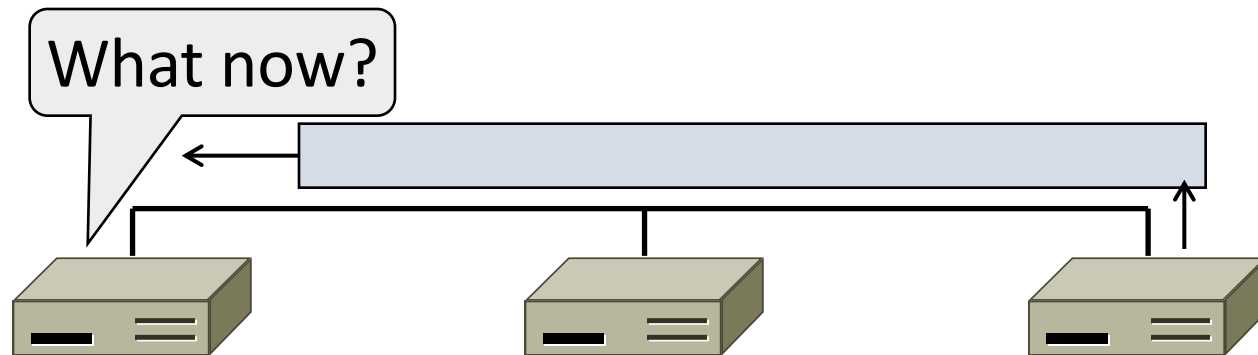
CSMA/CD Complications (2)

- Impose a minimum frame length of $2D$ seconds
 - So node can't finish before collision
 - Ethernet minimum frame is 64 bytes – Also sets maximum network length (500m w/ coax, 100m w/ Twisted Pair)



CSMA “Persistence”

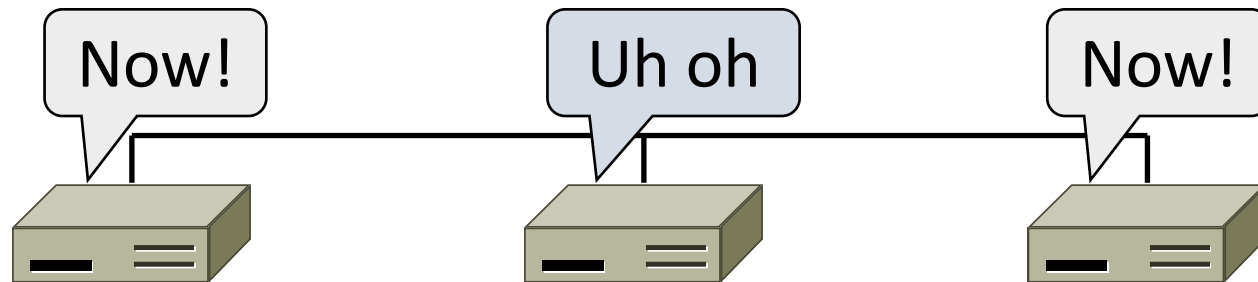
- What should a node do if another node is sending?



- Idea: Wait until it is done, and send

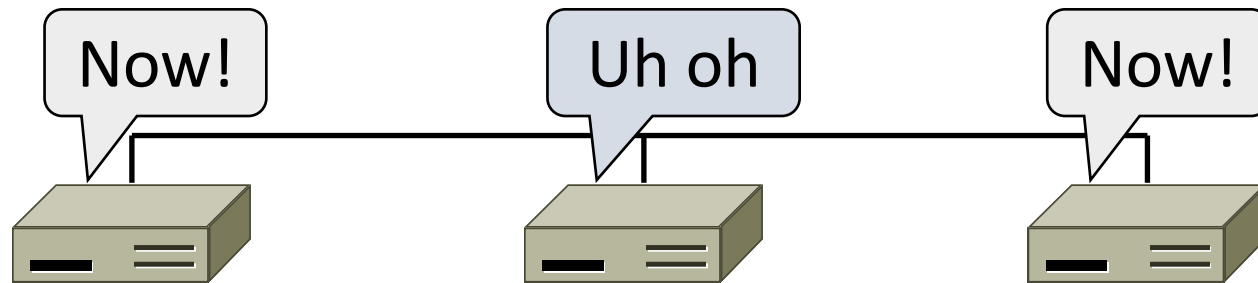
CSMA “Persistence” (2)

- Problem is that multiple waiting nodes will queue up then collide
 - More load, more of a problem



CSMA “Persistence” (2)

- Problem is that multiple waiting nodes will queue up then collide
 - Ideas?



CSMA “Persistence” (3)

- Intuition for a better solution
 - If there are N queued senders, we want each to send next with probability $1/N$

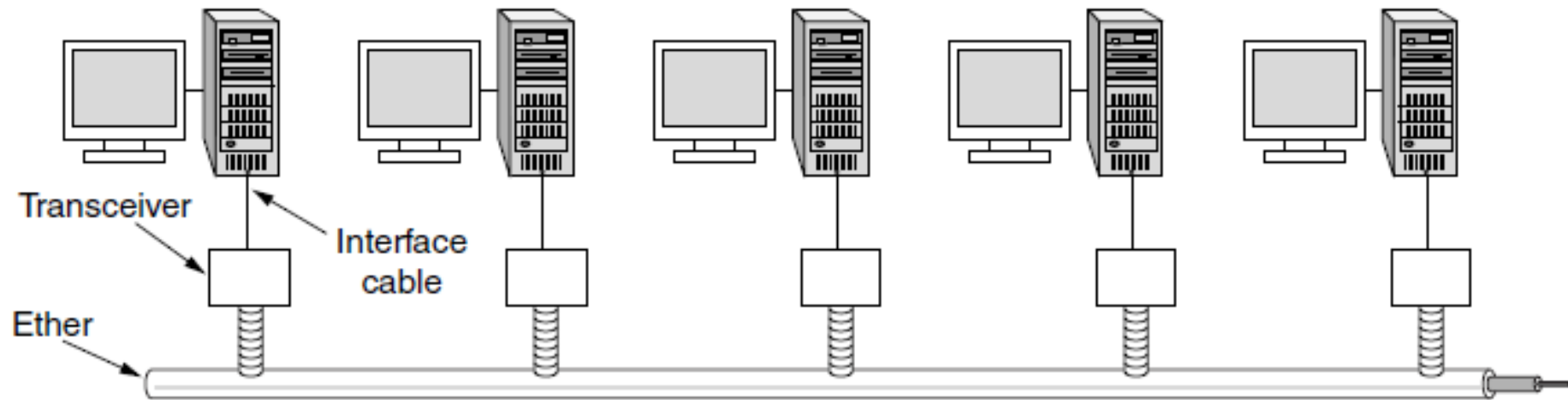


Binary Exponential Backoff (BEB)

- Cleverly estimates the probability
 - 1st collision, wait 0 or 1 frame times
 - 2nd collision, wait from 0 to 3 times
 - 3rd collision, wait from 0 to 7 times ...
- BEB doubles interval for each successive collision
 - Quickly gets large enough to work
 - Very efficient in practice

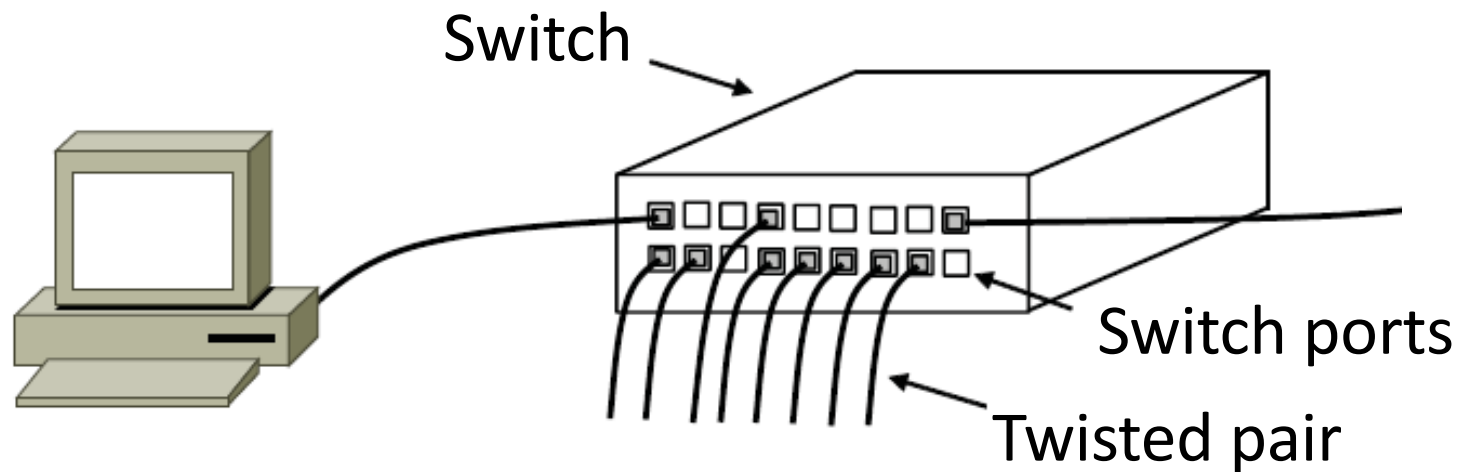
Classic Ethernet, or IEEE 802.3

- Most popular LAN of the 1980s, 1990s
 - 10 Mbps over shared coaxial cable
 - Multiple access with persistent CSMA/CD with BEB



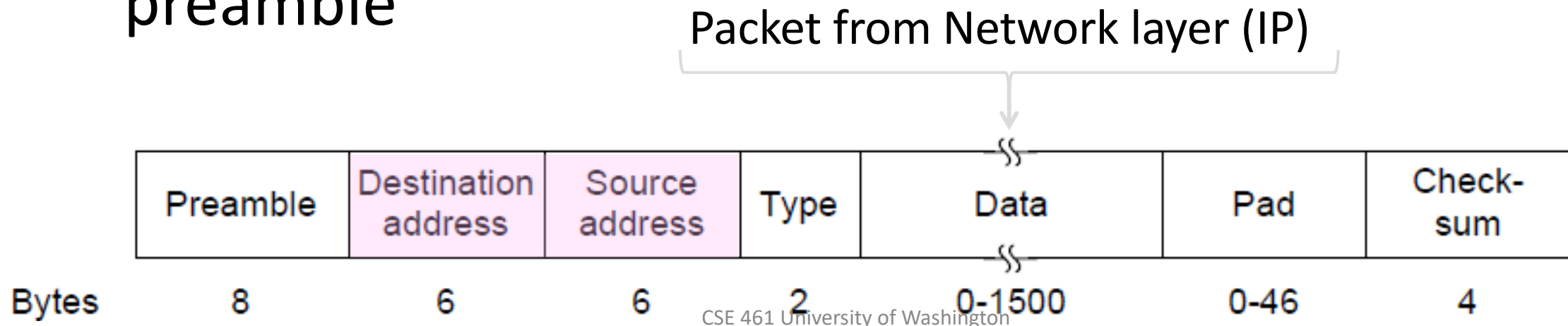
Modern Ethernet

- Based on switches, not multiple access, but still called Ethernet



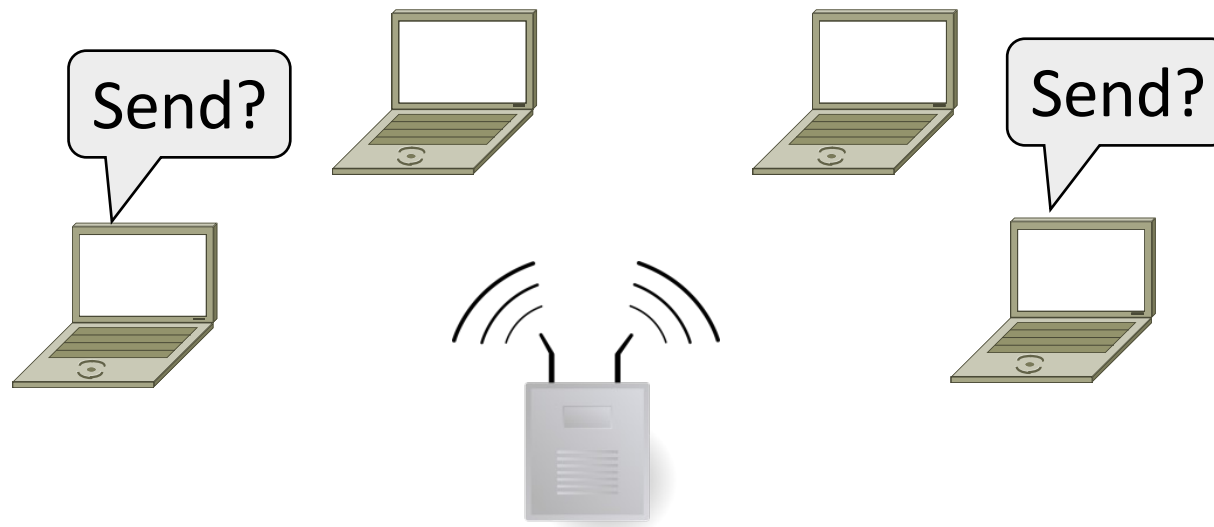
Ethernet Frame Format

- Has addresses to identify the sender and receiver
- CRC-32 for error detection; no ACKs or retransmission
- Start of frame identified with physical layer preamble



Wireless MACs

- How do wireless nodes share a single link? (Yes, this is WiFi!)
 - Build on our simple, wired model



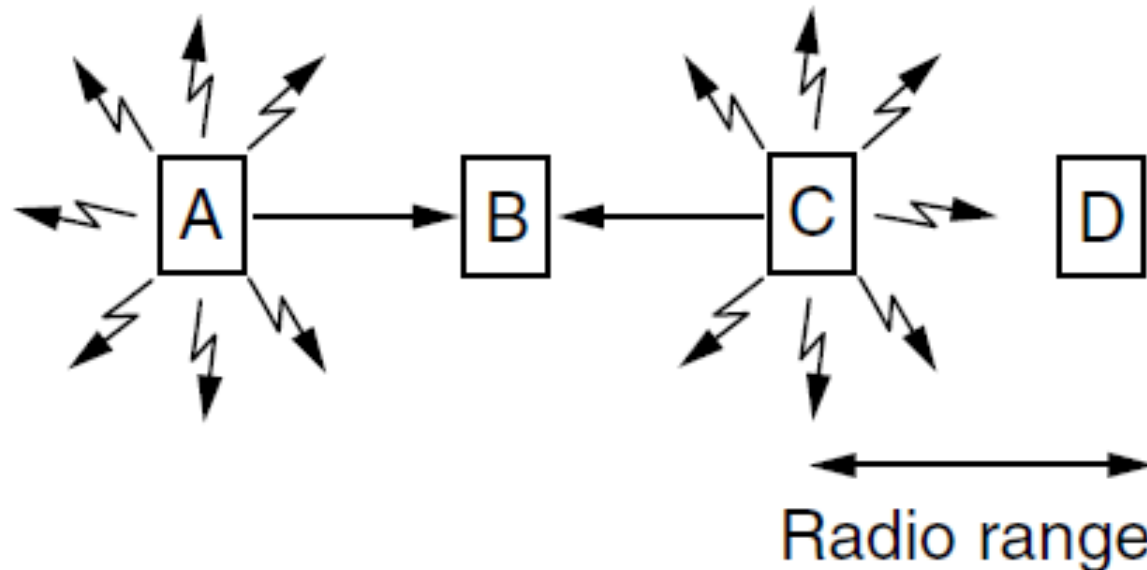
Wireless Complications

- Wireless is more complicated than wired (surprise!)
 1. Media is infinite – can't Carrier Sense
 2. Nodes usually can't hear while sending – can't Collision Detect



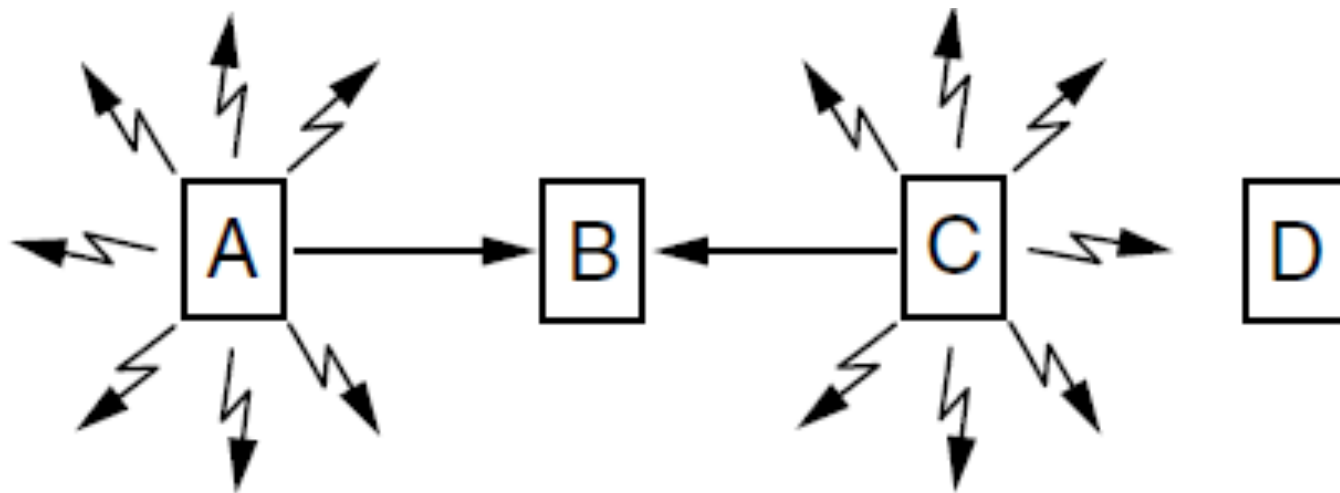
No CS: Different Coverage Areas

- Wireless signal is broadcast and received nearby, where there is sufficient SNR



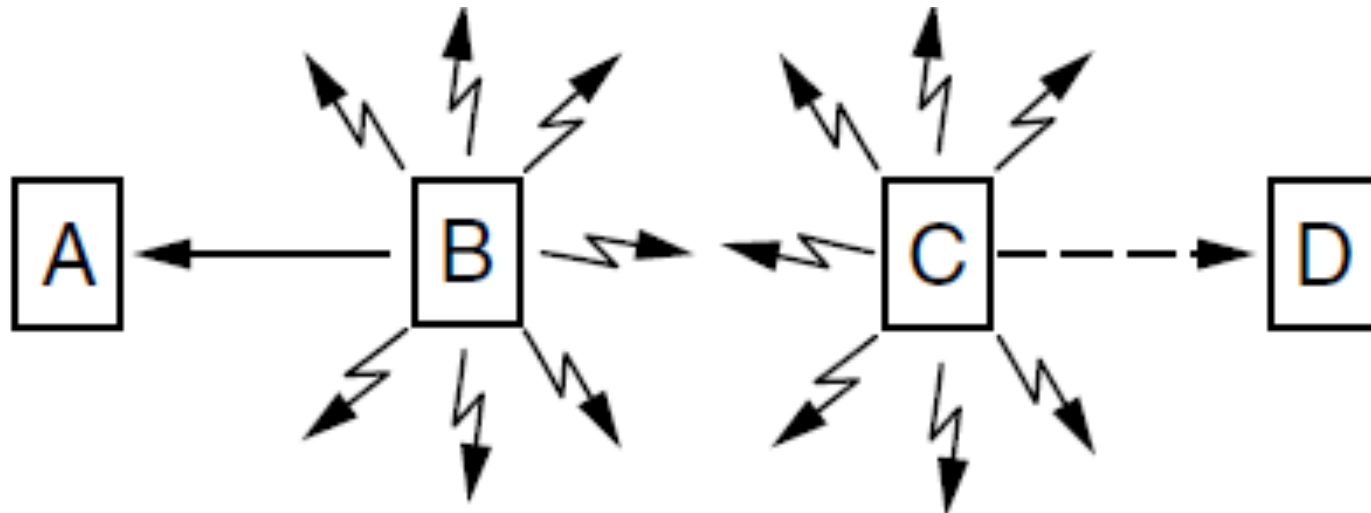
No CS: Hidden Terminals

- Node C is a hidden terminal when A sends to B
 - Similarly, A is a hidden terminal when C sends to B
 - A, C can't hear each other (to coordinate) yet collide at B
 - We want to avoid the inefficiency of collisions



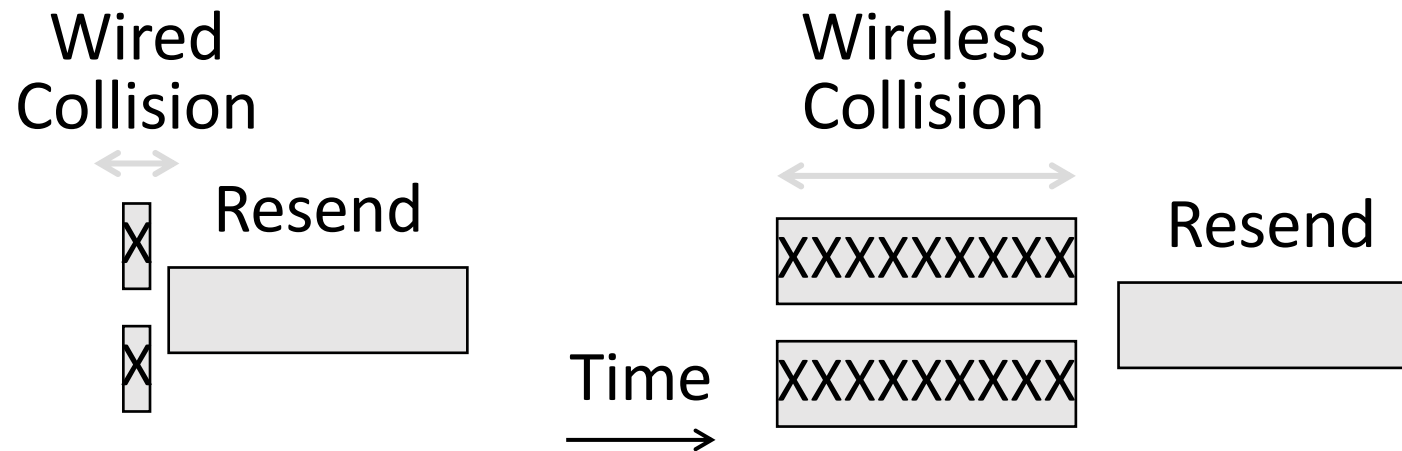
No CS: Exposed Terminals

- B, C are exposed terminals when sending to A, D
 - Can hear each other yet don't collide at receivers A and D
 - We want to send concurrently to increase performance



Nodes Can't Hear While Sending

- With wires, detecting collisions (and aborting) lowers their cost
- With wireless, more wasted time



Wireless Problems:

- Ideas?

MACA: Multiple Access w/ Collision Avoidance

- MACA uses a short handshake instead of CSMA (Karn, 1990)
 - 802.11 uses a refinement of MACA (later)
- Protocol rules:
 1. A sender node transmits a RTS (Request-To-Send, with frame length)
 2. The receiver replies with a CTS (Clear-To-Send, with frame length)
 3. Sender transmits the frame while nodes hearing the CTS stay silent
- Collisions on the RTS/CTS are still possible, but less likely

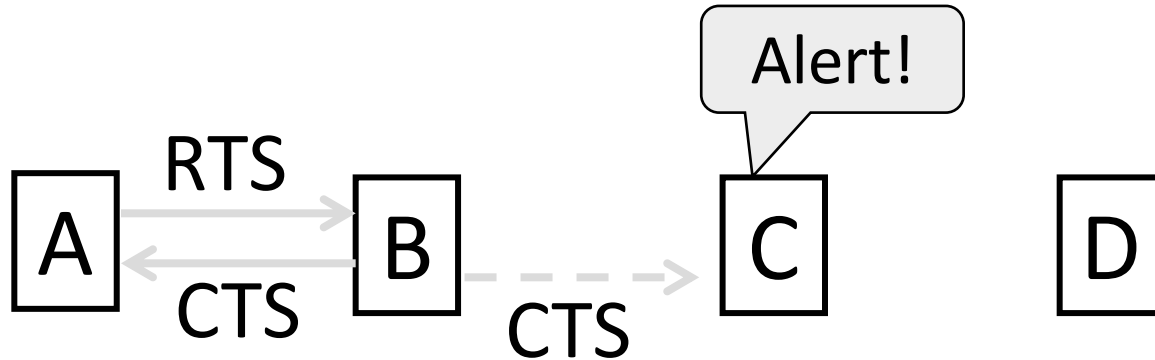
MACA – Hidden Terminals

- $A \rightarrow B$ with hidden terminal C
 1. A sends RTS, to B



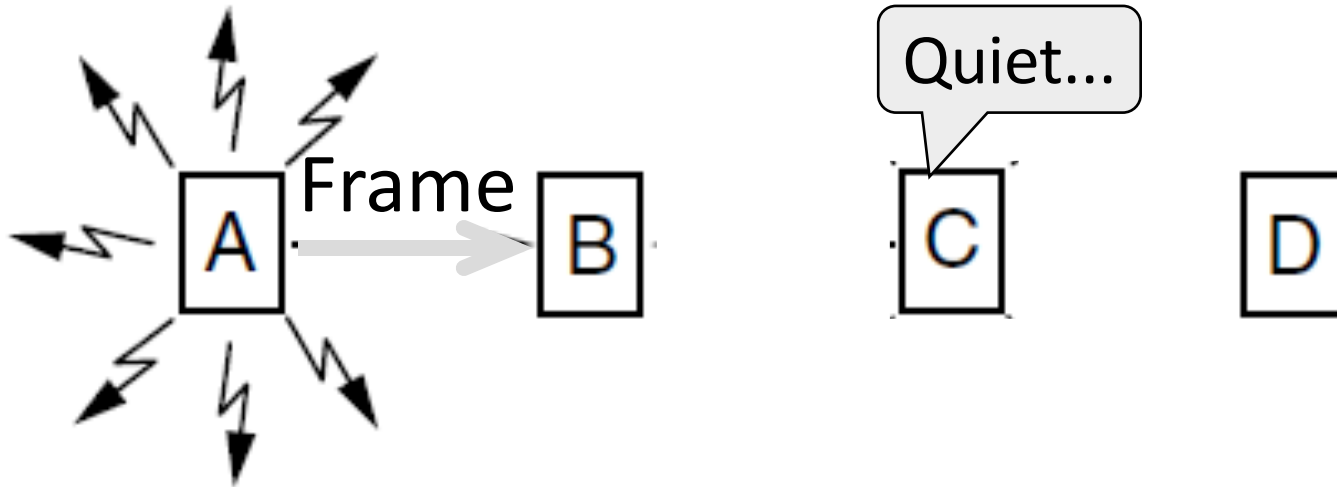
MACA – Hidden Terminals (2)

- A → B with hidden terminal C
 2. B sends CTS to A, and C overhears



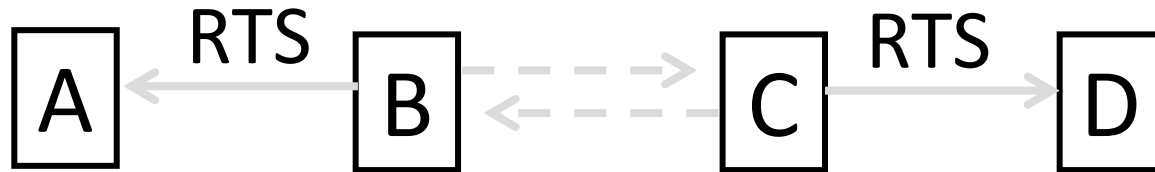
MACA – Hidden Terminals (3)

- $A \rightarrow B$ with hidden terminal C
 3. A sends frame while C defers



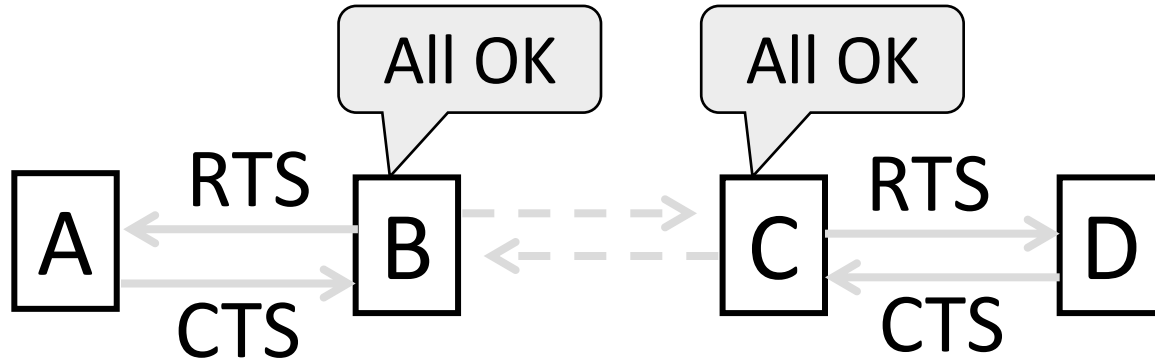
MACA – Exposed Terminals

- $B \rightarrow A$, $C \rightarrow D$ as exposed terminals
 - B and C send RTS to A and D



MACA – Exposed Terminals (2)

- $B \rightarrow A$, $C \rightarrow D$ as exposed terminals
 - A and D send CTS to B and C



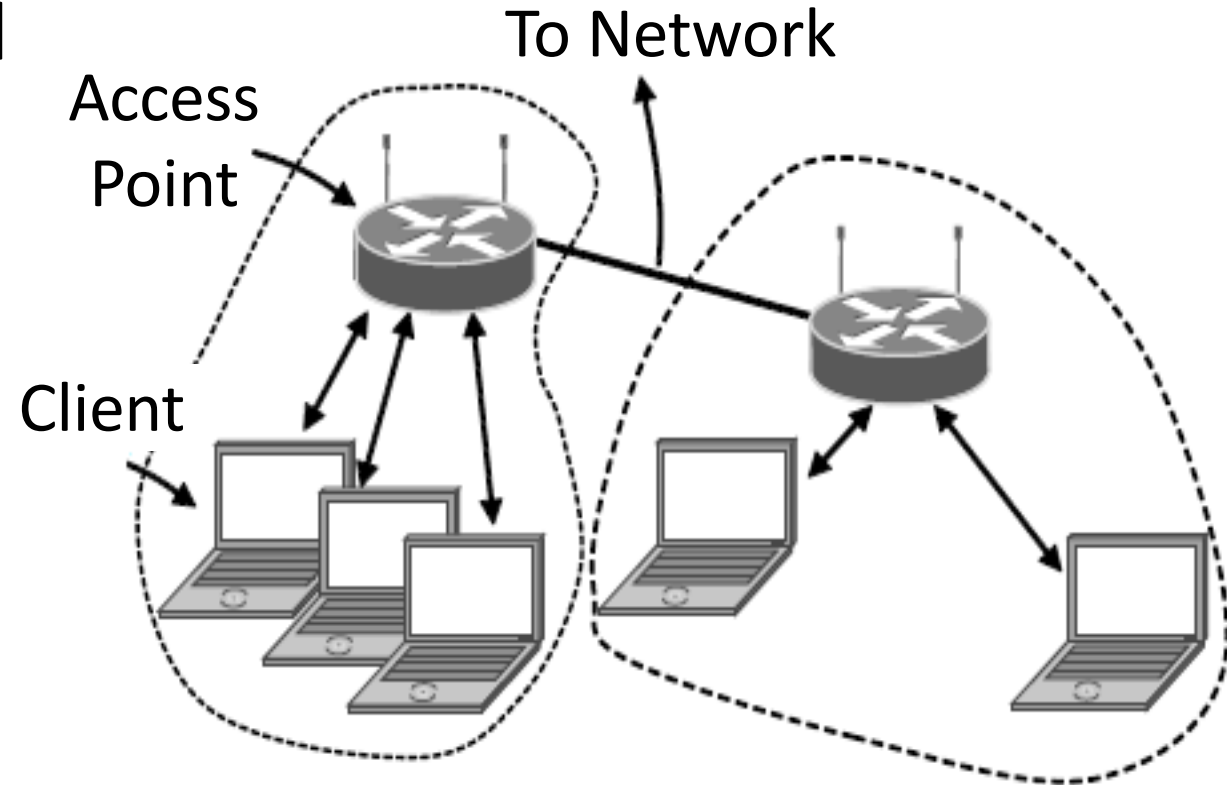
MACA – Exposed Terminals (3)

- $B \rightarrow A$, $C \rightarrow D$ as exposed terminals
 - A and D send CTS to B and C



802.11, or WiFi

- Very popular wireless LAN started in the 1990s
- Clients get connectivity from a (wired) AP (Access Point)
- It's a multi-access problem 😊
- Various flavors have been developed over time
 - Faster, more features

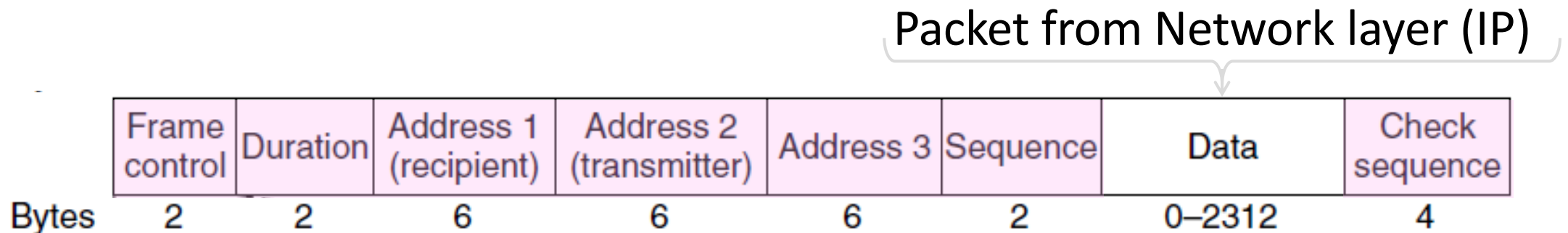


802.11 Physical Layer

- Uses 20/40 MHz channels on ISM (unlicensed) bands
 - 802.11b/g/n on 2.4 GHz
 - 802.11 a/n on 5 GHz
- OFDM modulation (except legacy 802.11b)
 - Different amplitudes/phases for varying SNRs
 - Rates from 6 to 54 Mbps plus error correction
 - 802.11n uses multiple antennas
 - Lots of fun tricks here

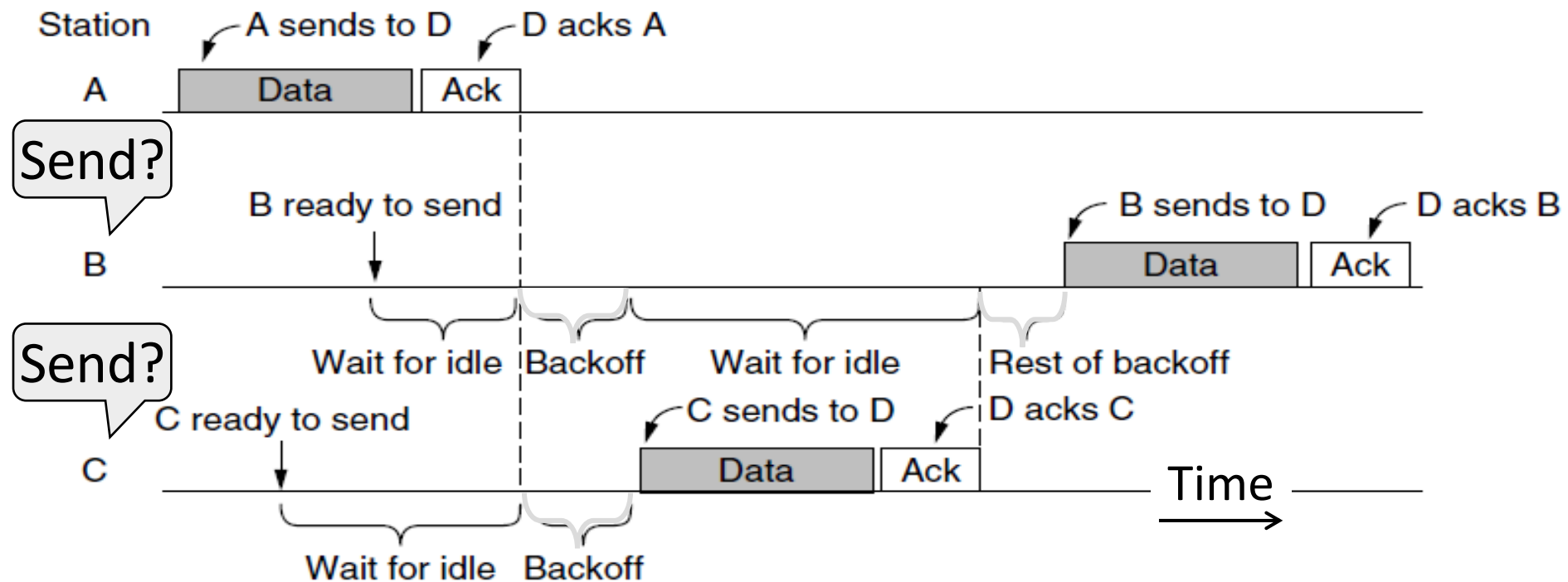
802.11 Link Layer

- Multiple access uses CSMA/CA (next); RTS/CTS optional
- Frames are ACKed and retransmitted with ARQ
- Funky addressing (three addresses!) due to AP
- Errors are detected with a 32-bit CRC
- Many, many features (e.g., encryption, power save)



802.11 CSMA/CA for Multiple Access

- Still using BEB!



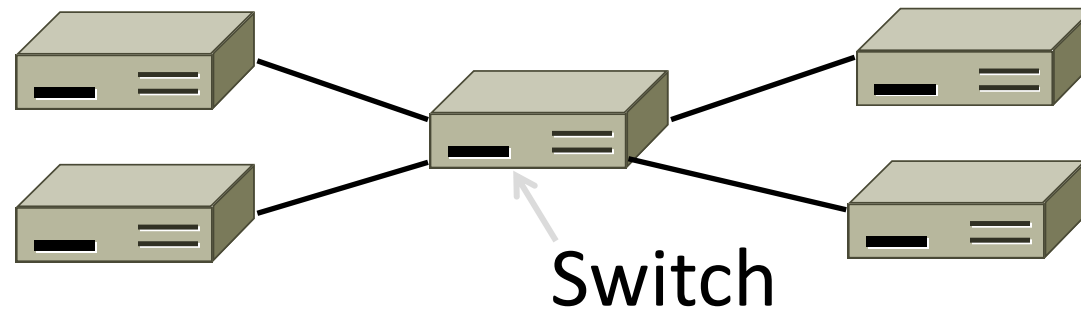
Recap: MAC layer ideas

- Random wait times upon collisions
- Carrier sense
 - Persistence
- Collision detection
- Binary exponential backoff
- RTS-CTS for hidden and exposed terminals

Link Layer: Switching

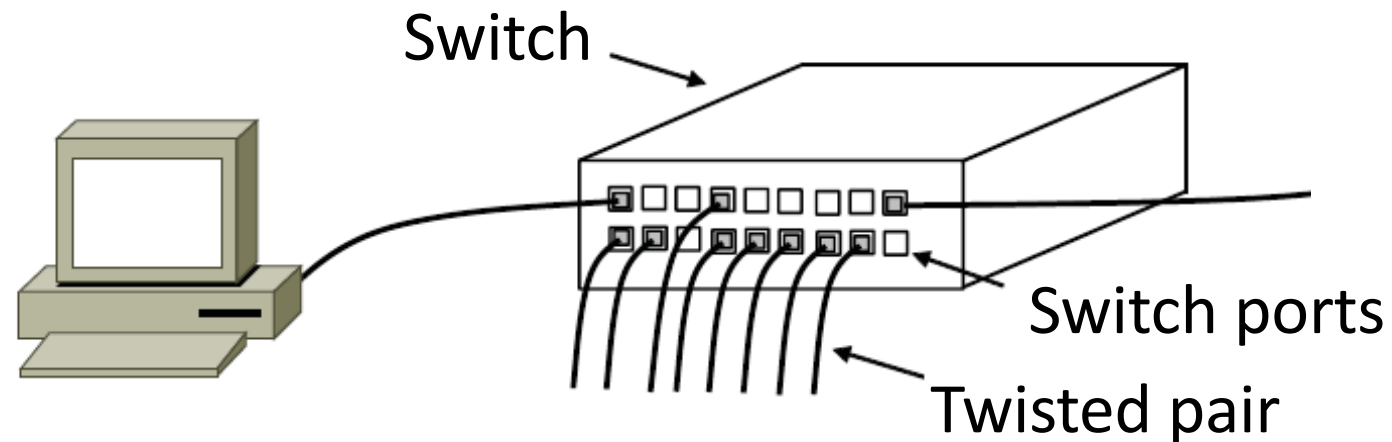
Switching

- How do we connect nodes with a switch instead of multiple access
 - Uses multiple links/wires
 - Basis of modern (switched) Ethernet



Switched Ethernet

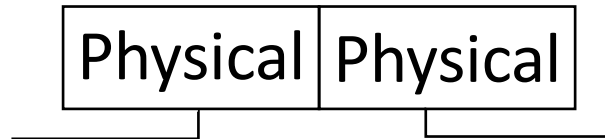
- Hosts are wired to Ethernet switches with twisted pair
 - Switch serves to connect the hosts
 - Wires usually run to a closet



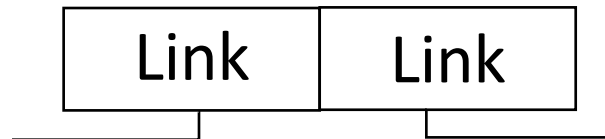
What's in the box?

- Remember from protocol layers:

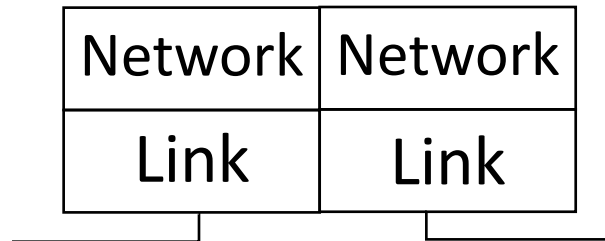
Hub, or
repeater



Switch



Router

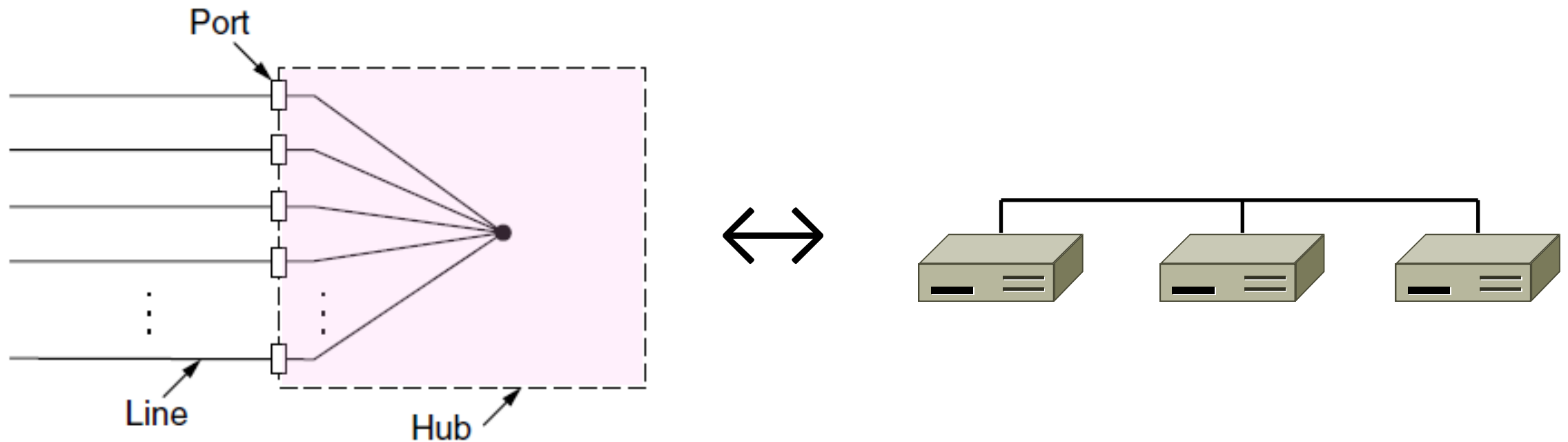


All look like this:



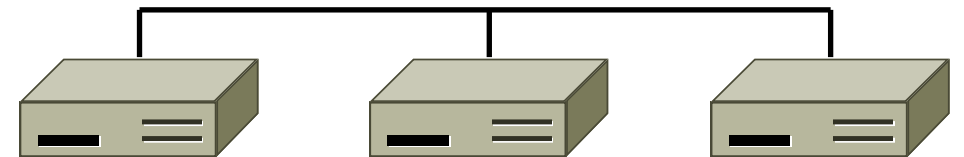
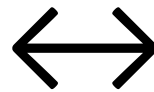
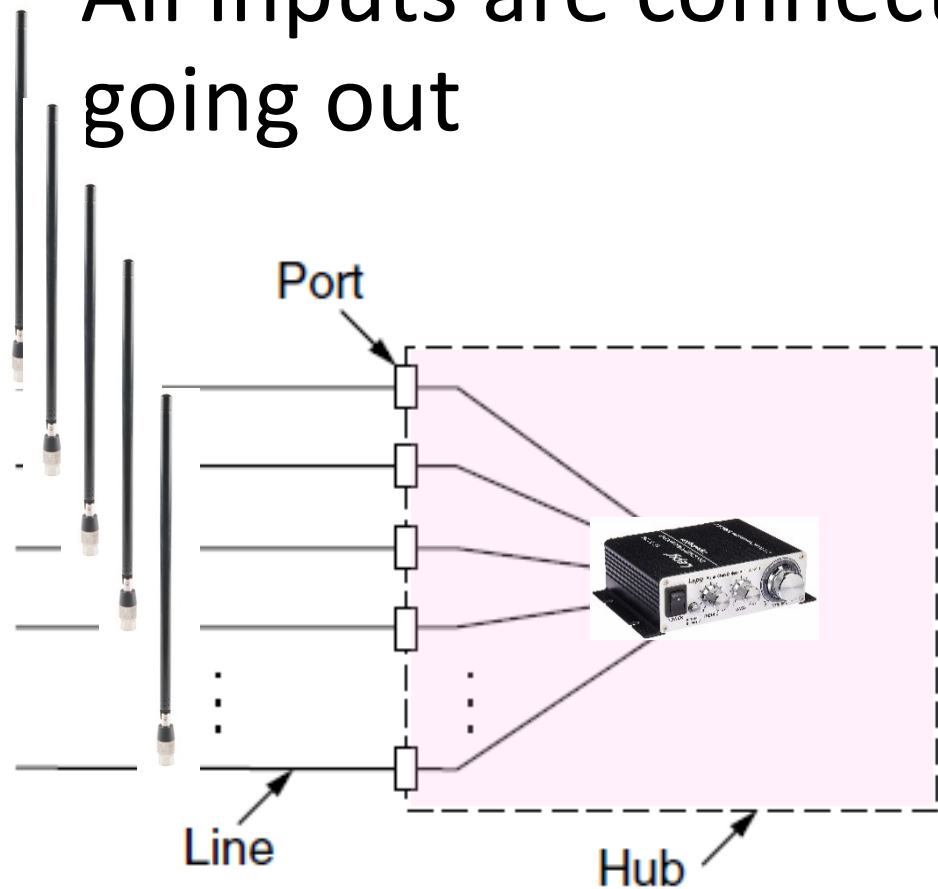
Inside a Hub

- All ports are wired together; more convenient and reliable than a single shared wire



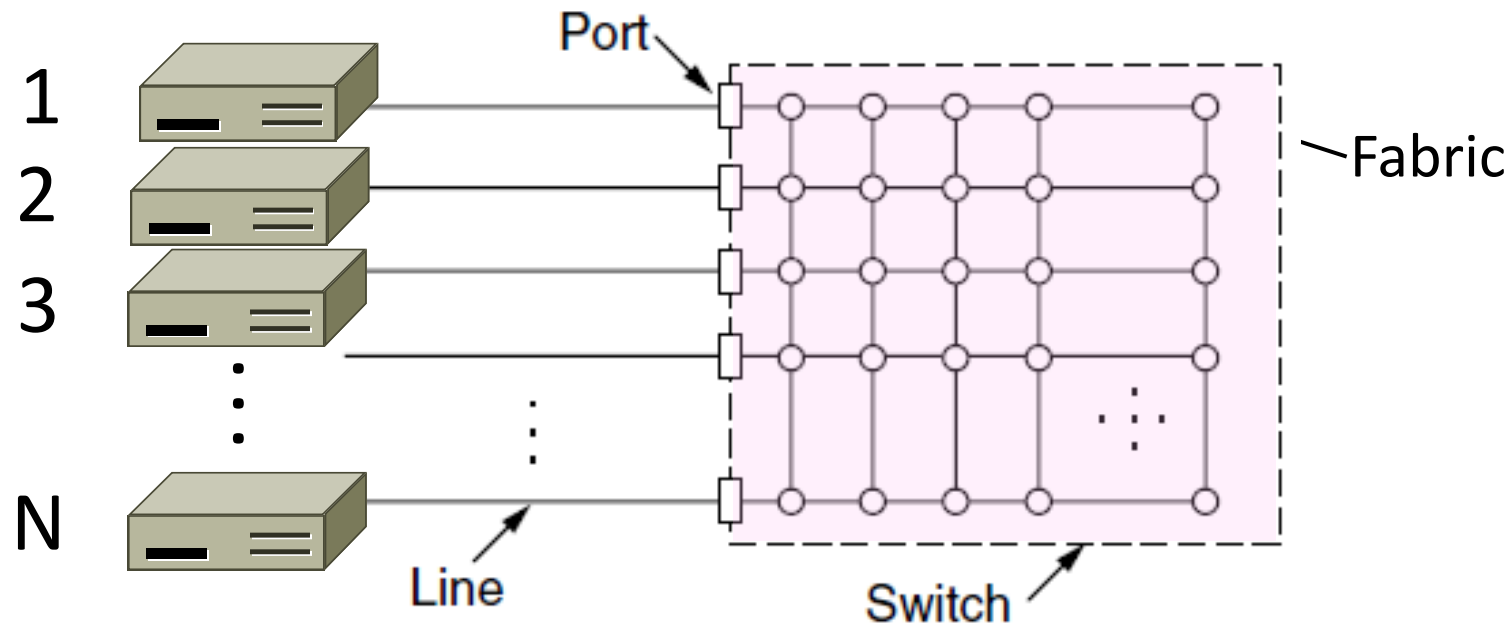
Inside a Repeater

- All inputs are connected; then amplified before going out



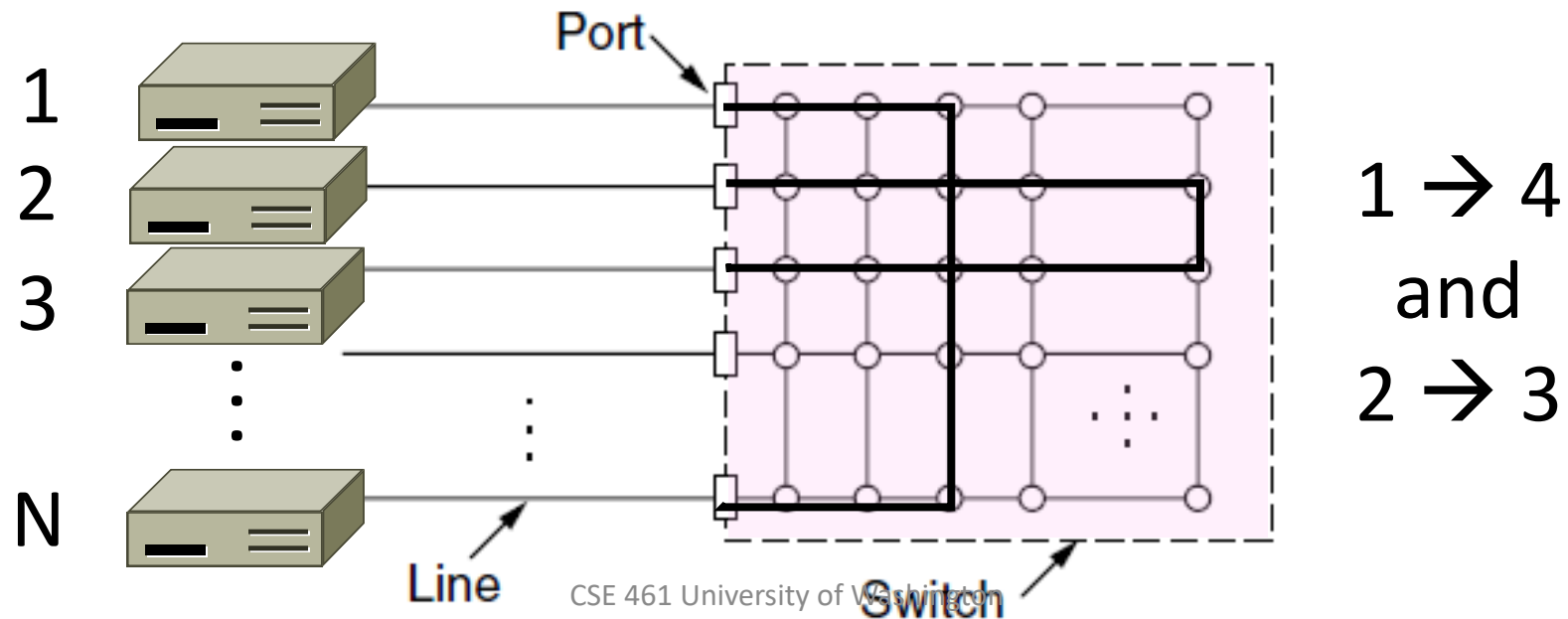
Inside a Switch

- Uses frame addresses (MAC addresses in Ethernet) to connect input port to the right output port; multiple frames may be switched in parallel



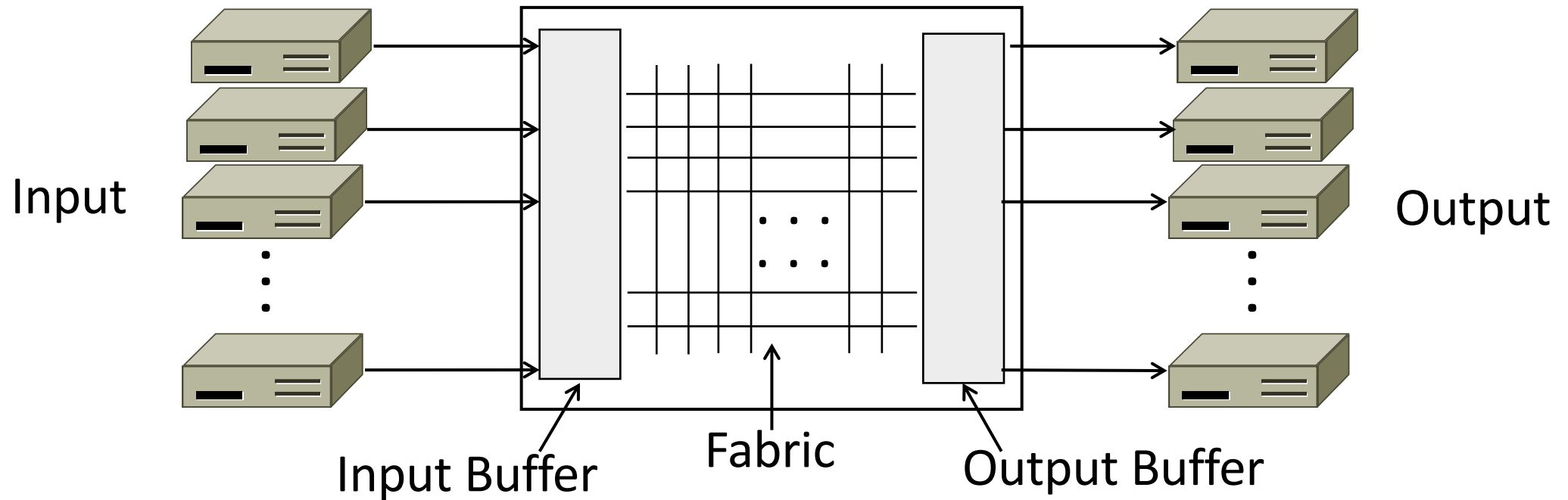
Inside a Switch (2)

- Port may be used for both input and output (full-duplex)
 - Just send, no multiple access protocol



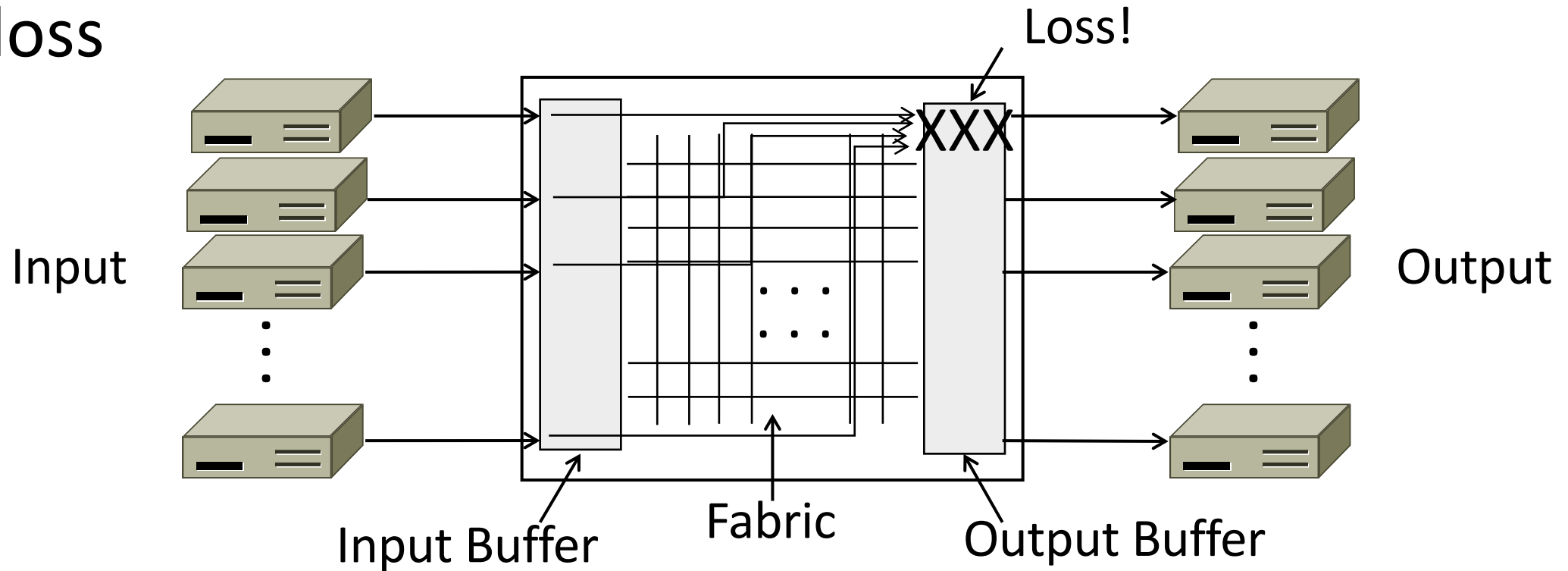
Inside a Switch (3)

- Need buffers for multiple inputs to send to one output



Inside a Switch (4)

- Sustained overload will fill buffer and lead to frame loss



Advantages of Switches

- Switches and hubs (mostly switches) have replaced the shared cable of classic Ethernet
 - Convenient to run wires to one location
 - More reliable; wire cut is not a single point of failure that is hard to find
- Switches offer scalable performance
 - E.g., 100 Mbps per port instead of 100 Mbps for all nodes of shared cable / hub