

Interrupts

- Fundamental concept in computation
- Interrupt execution of a program to “handle” an event
 - Don't have to rely on program relinquishing control
 - Can code program without worrying about others
- Issues
 - What can interrupt and when?
 - Where is the code that knows what to do?
 - How long does it take to handle interruption?
 - Can an interruption be, in turn, interrupted?
 - How does the interrupt handling code communicate its results?
 - How is data shared between interrupt handlers and programs?

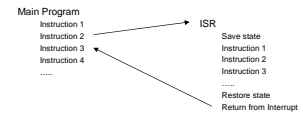
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What is an Interrupt?

- Reaction to something in I/O (human, comm link)
- Usually asynchronous to processor activities
- “interrupt handler” or “interrupt service routine” (ISR) invoked to take care of condition causing interrupt
 - Change value of internal variable (count)
 - Read a data value (sensor, receive)
 - Write a data value (actuator, send)



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Interrupts

- Code sample that does not interrupt

```
char SPI_SlaveReceive(void)
{
    /* Wait for reception complete */
    while(!(SPSR & (1<<SPIF)))
    ;
    /* Return data register */
    return SPDR;
}
```
- Instead of busy waiting until a byte is received the processor can generate an interrupt when it sets SPIF

```
SIGNAL(SIG_SPI) {
    RX_Byte = SPDR
}
```

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Saving and Restoring Context

- Processor and compiler dependent
- Where to find ISR code?
 - Different interrupts have separate ISRs
- Who does dispatching?
 - Direct
 - Different address for each interrupt type
 - Supported directly by processor architecture
 - Indirect
 - One top-level ISR
 - Switch statement on interrupt type
 - A mix of these two extremes?

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Saving and Restoring Context

- How much context to save?
 - Registers, flags, program counter, etc.
 - Save all or part?
 - Agreement needed between ISR and program
- Where should it be saved?
 - Stack, special memory locations, shadow registers, etc.
 - How much room will be needed on the stack?
 - Nested interrupts may make stack reach its limit – what then?
- Restore context when ISR completes

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Ignoring Interrupts

- Can interrupts be ignored?
 - It depends on the cause of the interrupt
 - No, for nuclear power plant temperature warning
 - Yes, for keypad on cell phone (human timescale is long)
- When servicing another interrupt
 - Ignore others until done
 - Can't take too long – keep ISRs as short as possible
 - Just do a quick count, or read, or write – not a long computation
- Interrupt disabling
 - Will ignored interrupt “stick”?
 - Rising edge sets a flip-flop
 - Or will it be gone when you get to it?
 - Level changes again and its as if it never happened
 - Don't forget to re-enable

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Prioritizing Interrupts

- When multiple interrupts happen simultaneously
 - Which is serviced first?
 - Fixed or flexible priority?
- Priority interrupts
 - Higher priority can interrupt
 - Lower priority can't
- Maskable interrupts
 - "don't bother me with that right now"
 - Not all interrupts are maskable, some are non-maskable

Interrupts in the ATmega16

- External interrupts
 - From I/O pins of microcontroller
- Internal interrupts
 - Timers
 - Output compare
 - Input capture
 - Overflow
 - Communication units
 - Receiving something
 - Done sending
 - ADC
 - Completed conversion

Interrupt Jump Vector Table

- Fixed location in memory to find first instruction for each type of interrupt
- Only room for one instruction
 - JMP to location of complete ISR

Address	Label	Code	Comments
0000	JMP	RESET	: RESET Handler
0004	JMP	INT0_vect	: INT0 Handler
0008	JMP	INT1_vect	: INT1 Handler
000C	JMP	TIMER0_COMP_vect	: Timer0 Compare Handler
0010	JMP	TIMER0_OVF_vect	: Timer0 Overflow Handler
0014	JMP	TIMER1_CAPT_vect	: Timer1 Capture Handler
0018	JMP	TIMER1_COMP_vect	: Timer1 Compare Handler
001C	JMP	TIMER1_OVF_vect	: Timer1 Overflow Handler
0020	JMP	TIMER2_OVF_vect	: Timer2 Overflow Handler
0024	JMP	SPI_STC_vect	: SPI Transfer Complete Handler
0028	JMP	USART_RX_vect	: USART RX Complete Handler
002C	JMP	USART_TX_vect	: USART TX Complete Handler
0030	JMP	ADC_COMP_vect	: ADC Conversion Complete Handler
0034	JMP	EE_READY_vect	: EEPROM Ready Handler
0038	JMP	ANALOG_COMP_vect	: Analog Comparator Handler
003C	JMP	TWI_vect	: Two-Wire Serial Interface Handler
0040	JMP	INT2_vect	: INT2 Handler
0044	JMP	TIMER2_COMP_vect	: Timer2 Compare Handler
0048	JMP	SLEEP_vect	: Sleep Program Memory Ready Handler
...
002A	RESET:	ldi r16, HIGH(RAMPD)	: Main program start
002B		out r16, r16	: Set stack pointer to top of RAM
002C		ldi r16, LOW(RAMPD)	
002D		out r16, r16	
002E		sei	: Enable interrupts
002F		-loop: rrr	
...

Chain of Events on Interrupt

- Finish executing current instruction
- Disable all interrupts **CLI** (Automatic)
- Push program counter on to stack
- Jump to interrupt vector table (Compiler)
- Jump to start of complete ISR
- Save any context that ISR may otherwise change
 - Registers and flags must be saved within ISR and restored before it returns – **this is very important!**
- Re-enable interrupts if nested interrupts are ok **SEI**
- Complete ISR's code
- Re-enable interrupts upon return
- Jump back to next instruction before interruption **RETI**

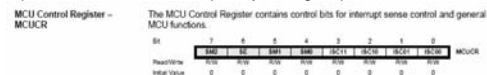
Shared Data Problem

- When you use interrupts you create the opportunity for multiple sections of code to update a variable.
- This might cause a problem in your logic if an interrupt updates a variable between two lines of code that are directly dependent on each other (e.g. if statement)
- One solution is to create critical sections where you disable the interrupts for a short period of time while you complete your logic on the shared variable

```
cli();
....critical section code goes here....
sei();
```

External Interrupts

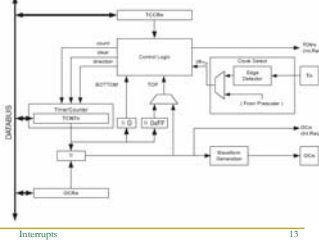
- Special pins: INT0, INT1, INT2
 - Can interrupt on edge or level
- Can interrupt even if set to be output pins
 - Implements "software interrupts" by setting output



ISC11	ISC10	Description
0	0	The low level of INT1 generates an interrupt request.
0	1	Any logical change on INT1 generates an interrupt request.
1	0	The falling edge of INT1 generates an interrupt request.
1	1	The rising edge of INT1 generates an interrupt request.

Closer Look at a Timer/Counter

- Timer0/Counter0
 - Clear timer on compare match (auto reload)
 - Prescaler (divide clock by up to 1024)
 - Overflow and compare match interrupts
 - Registers
 - Configuration
 - Count value
 - Output compare value



Timer/Counter Registers

- Timer/Counter Control Register TCCR0



• **Bit 7 – FOC0: Force Output Compare**
The FOC0 bit is only active when the WGM00 bit specifies a non-PWM mode. However, for ensuring compatibility with future devices, this bit must be set to zero when TCCR0 is written when operating in PWM mode. When writing a logical one to the FOC0 bit, an immediate compare match is forced on the Waveform Generation unit. The OC0 output is changed according to its COM01:0 bits setting. Note that the FOC0 bit is implemented as a sticky bit. Therefore it is the value present in the COM01:0 bits that determines the effect of the forced compare.
A FOC0 write will not generate any interrupt, nor will it clear the timer in CTC mode using OCR0 as TOP.
The FOC0 bit is always read as zero.

• **Bit 6:3 – WGM01:0: Waveform Generation Mode**
These bits control the counting sequence of the counter, the source for the maximum (TOP) counter value, and what type of Waveform Generation to be used. Modes of operation supported by the Timer/Counter unit are: Normal mode, Clear Timer on Compare Match (CTC) mode, and fast PWM (Fast PWM) modes. See Table 38 and 'Modes of Operation' on page 74.

Table 38. Waveform Generation Mode Bit Description¹⁾

WGM01 Mode (CTC)	WGM00 (PWM)	Waveform Generation Mode	TOP	Update of OCR0	OCV Flag
0	0	Normal	QVF	Immediate	MAX
1	0	PWM, Phase Correct	QVF	TOP	BOTTOM
2	1	CTC	OCR0	Immediate	MAX
3	1	Fast PWM	QVF	TOP	MAX

Note: 1) The CTC and PWM0 bit abbreviations are now obsolete. Use the WGM01:0 bit abbreviations. However, the functionality and location of these bits are compatible with previous versions of the timer.

Timer/Counter Registers (cont'd)

- Timer/Counter Control Register TCCR0

• **Bit 5:4 – COM01:0: Compare Match Output Mode**

These bits control the Output Compare pin (OC0) behavior. If one or both of the COM01:0 bits are set, the OC0 output overrides the normal port functionality of the I/O pin it is connected to. However, note that the Data Direction Register (DDR) bit corresponding to the OC0 pin must be set in order to enable the output driver.

When OC0 is connected to the pin, the function of the COM01:0 bits depends on the WGM01:0 bit setting. Table 39 shows the COM01:0 bit functionality when the WGM01:0 bits are set to a normal or CTC mode (non-PWM).

Table 39. Compare Output Mode, non-PWM Mode

COM01	COM00	Description
0	0	Normal port operation, OC0 disconnected
0	1	Toggle OC0 on compare match
1	0	Clear OC0 on compare match
1	1	Set OC0 on compare match

Timer/Counter Registers (cont'd)

- Timer/Counter Control Register TCCR0

• **Bit 2:0 – CS02:0: Clock Select**

The three Clock Select bits select the clock source to be used by the Timer/Counter.

Table 42. Clock Select Bit Description

CS02	CS01	CS00	Description
0	0	0	No clock source (Timer/Counter stopped)
0	0	1	Clk _{ICP} (No prescaling)
0	1	0	Clk _{ICP/8} (From prescaler)
0	1	1	Clk _{ICP/64} (From prescaler)
1	0	0	Clk _{ICP/256} (From prescaler)
1	0	1	Clk _{ICP/1024} (From prescaler)
1	1	0	External clock source on T0 pin. Clock on falling edge.
1	1	1	External clock source on T0 pin. Clock on rising edge.

If external pin modes are used for the Timer/Counter0, transitions on the T0 pin will clock the counter even if the pin is configured as an output. This feature allows software control of the counting.

Timer/Counter Registers (cont'd)

Timer/Counter Register – TCNT0



Output Compare Register – OCR0



Timer/Counter Interrupt Mask Register – TIMSK



Timer/Counter Interrupt Flag Register – TIFR



Setting Register Values

- Defined names for each register and bit
 - Set timer to clear on match
 - Set prescaler to 1024

$$TCCR0 = (1 \ll WGM01) | (1 \ll CS02) | (1 \ll CS00);$$

- Set count value to compare against

$$OCR0 = 150;$$

- Set timer to interrupt when it reaches count

$$TIMSK = (1 \ll OCIF0);$$

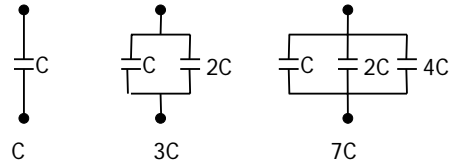
Writing an Interrupt Handler in C

- Set and clear interrupt enable
 - sei();
 - cli();
- Interrupt handler
 - SIGNAL(SIG_OUTPUT_COMPARE0)


```
{
                        i++;
                    }
```
- Setting I/O registers
 - TCCR0 = (1<<WGM01) | (1<<CS02) | (1<<CS00);
- Enabling specific interrupts
 - TIMSK = (1<<OCIE0);

Analog to digital conversion

- Use charge-redistribution technique
 - no sample and hold circuitry needed
 - even with perfect circuits quantization error occurs
- Basic capacitors
 - sum parallel capacitance



Analog to digital conversion (cont'd)

- Two reference voltage
 - mark bottom and top end of range of analog values that can be converted (V_L and V_H)
 - voltage to convert must be within these bounds (V_X)
- Successive approximation
 - most approaches to A/D conversion are based on this
 - 8 to 16 bits of accuracy
- Approach
 - sample value
 - hold it so it doesn't change
 - successively approximate
 - report closest match

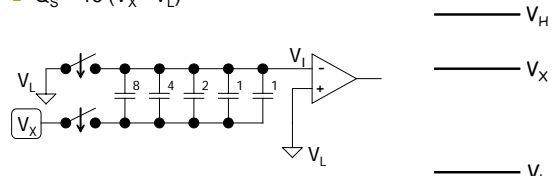
———— V_H

———— V_X

———— V_L

A-to-D – sample

- During the sample time the top plate of all capacitors is switched to reference low V_L
- Bottom plate is set to unknown analog input V_X
- $Q = CV$
- $Q_S = 16 (V_X - V_L)$



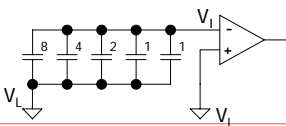
A-to-D – hold

- Hold state using logically controlled analog switches
 - Top plates disconnected from V_L
 - Bottom plates switched from V_X to V_L
- $Q_H = 16 (V_L - V_1)$
 - conservation of charge $Q_S = Q_H$
 - $16 (V_X - V_L) = 16 (V_L - V_1)$
 - $V_X - V_L = V_L - V_1$ (output of op-amp)

———— V_H

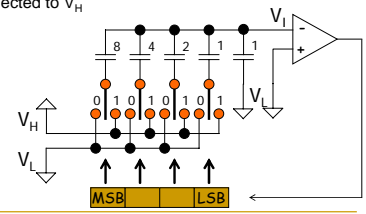
———— V_X

———— V_L



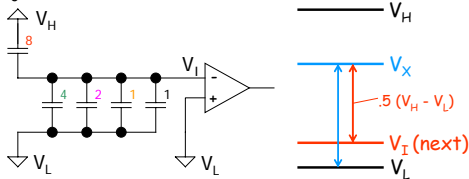
A-to-D – successive approximation

- Each capacitor successively switched from V_L to V_H
 - Largest capacitor corresponds to MSB
- Output of comparator determines bottom plate voltage of cap
 - > 0 : remain connected to V_H
 - < 0 : return to V_L



A-to-D example - MSB

- Suppose $V_X = 21/32 (V_H - V_L)$ and already sampled
- Compare after shifting half of capacitance to V_H
 - V_I goes up by $+ 8/16 (V_H - V_I) - 8/16 (V_L - V_I) = + 8/16 (V_H - V_L)$
 - original $V_L - V_I$ goes down and becomes
 - $V_L - (V_I + .5 (V_H - V_L)) = V_L - V_I - .5 (V_H - V_L)$
- Output > 0



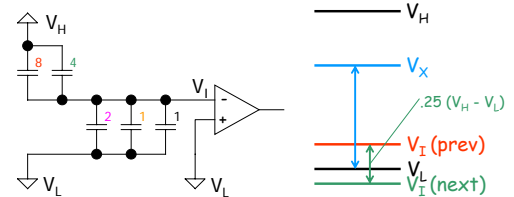
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A-to-D example - (MSB-1)

- Compare after shifting another part of cap. to V_H
 - V_I goes up by $+ 4/16 (V_H - V_I) - 4/16 (V_L - V_I) = + 4/16 (V_H - V_L)$
 - original $V_L - V_I$ goes down and becomes
 - $V_L - (V_I + .25 (V_H - V_L)) = V_L - V_I - .25 (V_H - V_L)$
- Output < 0 (went too far)



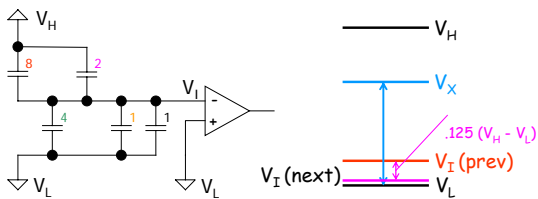
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A-to-D example - (MSB-2)

- Compare after shifting another part of cap. to V_H
 - V_I goes up by $+ 2/16 (V_H - V_I) - 2/16 (V_L - V_I) = + 2/16 (V_H - V_L)$
 - original $V_L - V_I$ goes down and becomes
 - $V_L - (V_I + .125 (V_H - V_L)) = V_L - V_I - .125 (V_H - V_L)$
- Output > 0



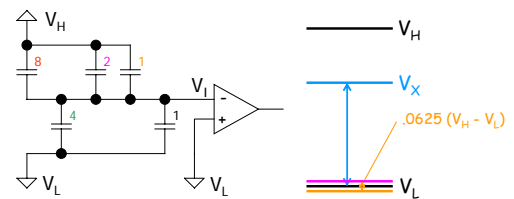
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A-to-D example - LSB

- Compare after shifting another part of cap. to V_H
 - V_I goes up by $+ 1/16 (V_H - V_I) - 1/16 (V_L - V_I) = + 1/16 (V_H - V_L)$
 - original $V_L - V_I$ goes down and becomes
 - $V_L - (V_I + .0625 (V_H - V_L)) = V_L - V_I - .0625 (V_H - V_L)$
- Output < 0 (went too far again)



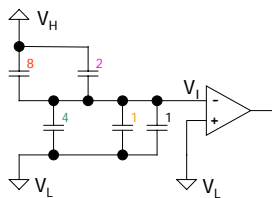
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A-to-D example final result

- Input sample of $21/32$
- Gives result of 1010 or $10/16 = 20/32$
- 3% error

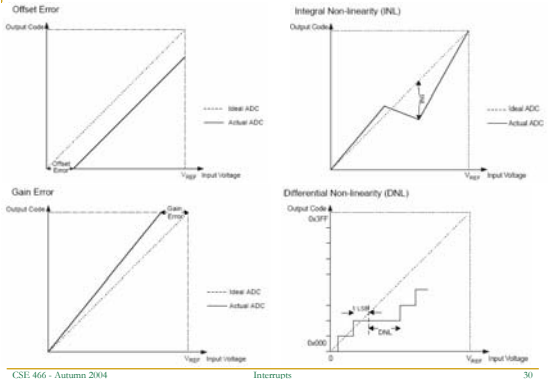


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A-to-D Conversion Errors



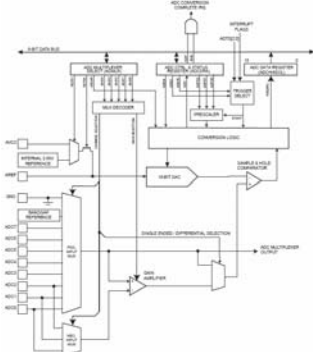
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Closer Look at A-to-D Conversion

- Needs a comparator and a D-to-A converter
- Takes time to do successive approximation
- Interrupt generated when conversion is completed



A-to-D Conversion on the ATmega16

- 10-bit resolution (adjusted to 8 bits as needed)
- 65-260 usec conversion time
- 8 multiplexed input channels
 - Difference of two pins
 - Optional gain on differential signal (amplifies difference)
- Interrupt on completion of A-to-D conversion
- 0-V_{CC} input range
 - Susceptible to noise – special analog supply pin (AVCC) and capacitor connection for reference voltage (AREF)

A-to-D Conversion (cont'd)

ADC Multiplexer Selection Register - ADMUX

Bit	7	6	5	4	3	2	1	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
Initial Value	0	0	0	0	0	0	0	0

- Bit 7-6 - REFSEL: Reference Selection Bits
- These bits select the voltage reference for the ADC, as shown in Table 83. If these bits are changed during a conversion, the change will not go in effect until this conversion is complete (ADIF in ADCSRA is set). The internal voltage reference options may not be used if an external reference voltage is being applied to the AREF pin.

Table 83. Voltage Reference Selections for ADC

REFSEL7	REFSEL6	Voltage Reference Selection
0	0	AREF: Internal V _{REF} turned off
0	1	AV _{CC} with external capacitor at AREF pin
1	0	Reserved
1	1	Internal 2.56V Voltage Reference with external capacitor at AREF pin

- Bit 5 - ADLAR: ADC Left Adjust Result
- The ADLAR bit affects the presentation of the ADC conversion result in the ADC Data Register. Write one to ADLAR to left adjust the result. Otherwise, the result is right adjusted. Changing the ADLAR bit will affect the ADC Data Register immediately, regardless of any ongoing conversions. For a complete description of this bit, see "The ADC Data Register - ADCL and ADCH" on page 218.

A-to-D Conversion (cont'd)

- Single-ended or differential
 - 1 of 8 single-ended
 - ADC_x - ADC₀ at 1x gain
 - ADC_{0,1} - ADC₀ at 200x
 - ADC_{2,3} - ADC₂ at 10x
 - ADC_{2,3} - ADC₃ at 200x
 - ADC_{0,1,2,3,4,5} - ADC₂ at 1x

Bits 4:0 - MUXA: Analog Channel and Gain Selection Bits
The value of these bits selects which combination of analog inputs are connected to the ADC. These bits also select the gain for the differential channels. See Table 84 for details. If these bits are changed during a conversion, the change will not go in effect until the conversion is complete (ADIF in ADCSRA is set).

Table 84. Input Channel and Gain Selections

MUXA:0	Single Ended Input	Positive Differential Input	Negative Differential Input	Gain
00000	ADC0			
00001	ADC1			
00010	ADC2			
00011	ADC3			
00100	ADC4			
00101	ADC5			
00110	ADC6			
00111	ADC7			
01000	ADC0	ADC0		10x
01001	ADC1	ADC0		10x
010010	ADC0	ADC0	ADC0	200x
010011	ADC1	ADC0	ADC0	200x
01100	ADC2	ADC2		10x
01101	ADC3	ADC2		10x
011010	ADC2	ADC2	ADC2	200x
011011	ADC3	ADC2	ADC2	200x
10000	ADC0	ADC1		1x
10001	ADC1	ADC1		1x
10010	ADC2	ADC1		1x
10011	ADC3	ADC1		1x
10100	ADC4	ADC1		1x
10101	ADC5	ADC1		1x
10110	ADC6	ADC1		1x
10111	ADC7	ADC1		1x
11000	ADC0	ADC2		1x
11001	ADC1	ADC2		1x
11010	ADC2	ADC2		1x
11011	ADC3	ADC2		1x
11100	AVCC	AVCC	AVCC	1x

A-to-D Conversion (cont'd)

The ADC Data Register - ADCL and ADCH

ADLAR = 0

Bit	15	14	13	12	11	10	9	8
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
Initial Value	0	0	0	0	0	0	0	0

ADLAR = 1

Bit	15	14	13	12	11	10	9	8
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
Initial Value	0	0	0	0	0	0	0	0

A-to-D Conversion (cont'd)

ADC Control and Status Register A - ADCSRA

Bit	7	6	5	4	3	2	1	0
Read/Write	R/W	R/W	R/W	R/W	R/W	R/W	R/W	R/W
Initial Value	0	0	0	0	0	0	0	0

- Bit 7 - ADEN: ADC Enable
- Bit 6 - ADSC: ADC Start Conversion
- Bit 5 - ADIF: ADC Auto Trigger Enable
- Bit 4 - ADIF: ADC Interrupt Flag
- Bit 3 - ADIF: ADC Interrupt Enable
- Bits 2:0 - ADPS2:0: ADC Prescaler Select Bits

ADPS2	ADPS1	ADPS0	Division Factor
0	0	0	2
0	0	1	2
0	1	0	4
0	1	1	8
1	0	0	16
1	0	1	32
1	1	0	64
1	1	1	128

A-to-D Conversion (cont'd)

Special Functions Register – SFOR



• **Bits 7:5 – ADTS2:0: ADC Auto Trigger Source**
 If ADATE in ADCSRA is written to one, the value of these bits selects which source will trigger an ADC conversion. If ADATE is cleared, the ADTS2:0 settings will have no effect. A conversion will be triggered by the rising edge of the selected interrupt flag. Note that switching from a trigger source that is cleared to a trigger source that is set, will generate a positive edge on the trigger signal. If ADON in ADCSRA is set, this will start a conversion. Switching to Free Running mode (ADTS2:0=0) will not cause a trigger event, even if the ADC interrupt flag is set.

Table 86. ADC Auto Trigger Source Selections

ADTS2	ADTS1	ADTS0	Trigger Source
0	0	0	Free Running mode
0	0	1	Analog Comparator
0	1	0	External Interrupt Request 0
0	1	1	Timer/Counter0 Compare Match
1	0	0	Timer/Counter0 Overflow
1	0	1	Timer/Counter0 Compare Match B
1	1	0	Timer/Counter1 Overflow
1	1	1	Timer/Counter1 Capture Event

• **Bit 4 – RES: Reserved Bit**
 This bit is reserved for future use. To ensure compatibility with future devices, this bit must be written to zero when SFOR is written.

A-to-D Conversion (cont'd)

Writing an Interrupt Handler in C (again)

- Ensure main program sets up all registers
- Enable interrupts as needed
- Enable global interrupts (SEI)
- Write handler routine for each enabled interrupt
 - What if an interrupt occurs and a handler isn't defined?
- Make sure routine does not disrupt others
 - Data sharing problem
 - Save any state that might be changed (done by compiler)
- Re-enable interrupts upon return
 - done by compiler with RETI

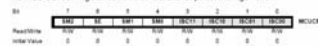
Power modes

- Processor can go to "sleep" and save power
- Different modes put different sets of modules to sleep
 - Which one to use depends on which modules are needed to wake up processor
 - Timers, external interrupts, ADC, serial communication lines, etc.
- set_sleep_mode (mode);
- sleep_mode ();

Power modes (cont'd)

MCU Control Register – MCUCR

The MCU Control Register contains control bits for power management.



• **Bits 7, 5, 4 – SM2:0: Sleep Mode Select Bits 2, 1, and 0**
 These bits select between the six available sleep modes as shown in Table 13.

Table 13. Sleep Mode Select

SM2	SM1	SM0	Sleep Mode
0	0	0	Idle
0	0	1	ADC Noise Reduction
0	1	0	Power-down
0	1	1	Power-save
1	0	0	Reserved
1	0	1	Reserved
1	1	0	Standby ¹⁾
1	1	1	Extended Standby ¹⁾

Note: 1. Standby mode and Extended Standby mode are only available with external crystals or resonators.

• **Bit 6 – SE: Sleep Enable**

The SE bit must be written to logic one to make the MCU enter the sleep mode when the SLEEP instruction is executed. To avoid the MCU entering the sleep mode unless it is the programmer's purpose, it is recommended to write the Sleep Enable (SE) bit to one just before the execution of the SLEEP instruction and to clear it immediately after waking up.

Power modes (cont'd)

- Wake up sources and active clocks

Sleep Mode	Active Clock domains				Oscillators		Wake-up Sources					
	clk _{cpu}	clk _{cpu}	clk _{cpu}	clk _{cpu}	Main Clock Source Enabled	Timer One Enabled	INT2/INT1	TWB Address Match	Timer 2	SPM EEPROM Ready	ADC	Other I/O
Idle					X	X ⁽¹⁾	X	X	X	X	X	X
ADC Noise Reduction			X	X	X	X ⁽¹⁾	X ⁽²⁾	X	X	X	X	X
Power-down							X ⁽²⁾	X				
Power-save					X ⁽³⁾		X ⁽²⁾	X				
Standby ⁽¹⁾					X		X ⁽²⁾	X				
Extended Standby ⁽¹⁾					X ⁽³⁾	X	X ⁽²⁾	X ⁽³⁾	X	X ⁽³⁾		

Notes: 1. External Crystal or resonator selected as clock source.
 2. If AS2 or AS1R is set.
 3. Only INT2 or level interrupt INT1 and INT0.