

CSE 484 / CSE M 584 (Autumn 2011)

Security and Networks

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Thanks to Dan Boneh, Dieter Gollmann, John Manferdelli, John Mitchell, Vitaly Shmatikov, Bennet Yee, and many others for sample slides and materials ...

Class updates

- (Short) Homework 3
 - Due Wednesday
 - Individual assignment
- My office hours this week:
 - **CSE 210:** M,W,F after class. T,Th afternoons
 - others by appointment
 - come pick up graded Homework #2

Lab 3

- **Posted on website and on Catalyst.**

- <https://catalyst.uw.edu/collectit/assignment/dhalperi/17513/72548>
- Hack my privacy!

Lab 3

- **Posted on website and on Catalyst.**

- <https://catalyst.uw.edu/collectit/assignment/dhalperi/17513/72548>

- Hack my privacy!

- ***This lab is optional***

- Can only help your grade.

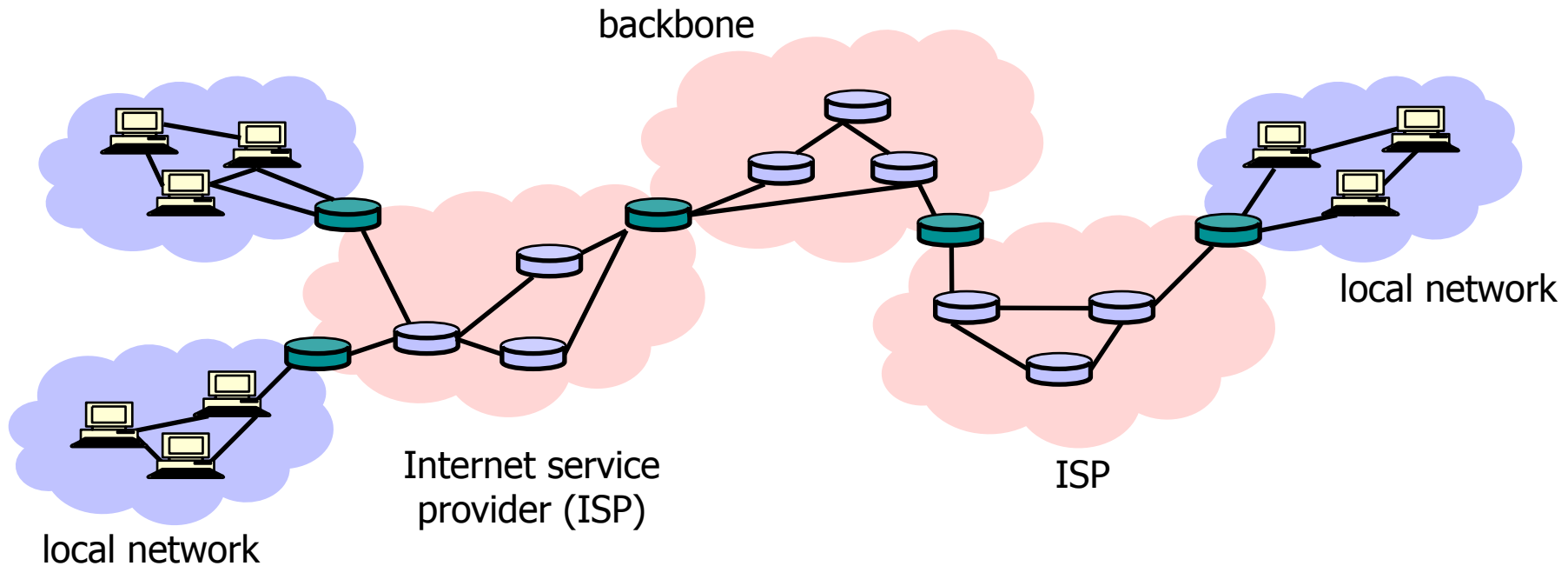
- Lots of opportunity for extra credit.

- I really think this lab is fun, and encourage you to do it, but we're not going to require it.

This week

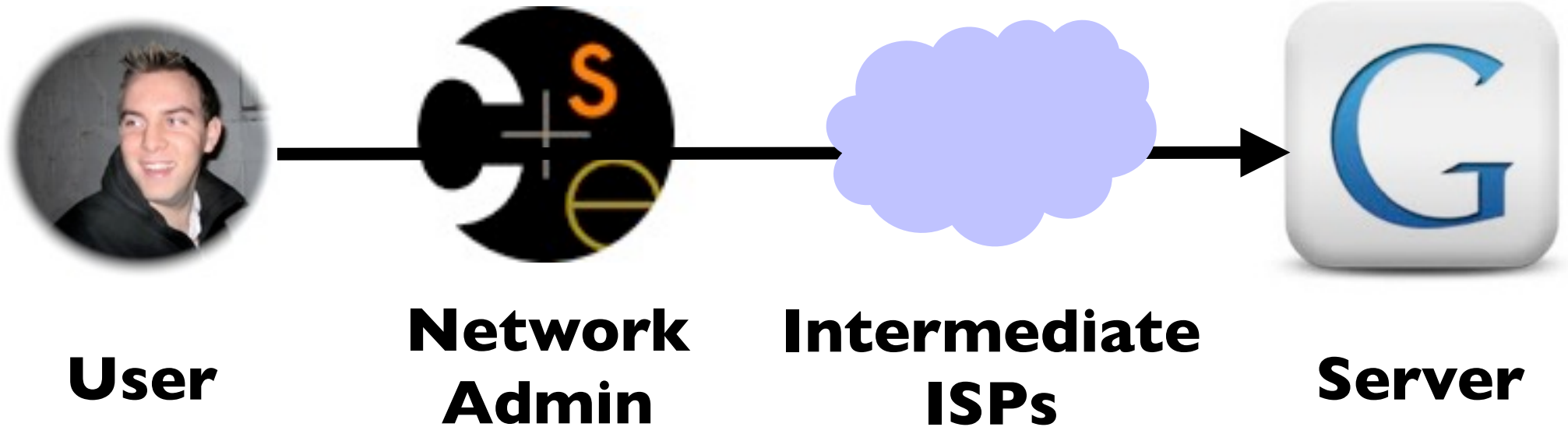
- **Today:** Network security
- **Wednesday:** Potpourri
- **Friday:** Any questions you have
 - Submit to my email, cse484-tas
 - Submit anonymously via the feedback form on the website

Internet Infrastructure



- ◆ TCP/IP for packet routing and connections
- ◆ Border Gateway Protocol (BGP) for route discovery
- ◆ Domain Name System (DNS) for IP address discovery

(Some) Entities



(Some) Goals



User

(Some) Goals



User

- Service (can get to Internet)
- Privacy (middle-entities shouldn't know what communicating or with whom)
- Fairness (e.g., get service I paid for)
- Integrity (can't impersonate me)
- Safety (network shouldn't attack me)

(Some) Goals



**Network
Admin**

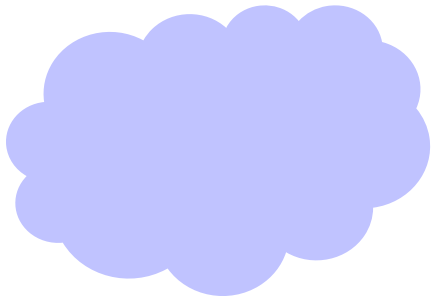
(Some) Goals



Network Admin

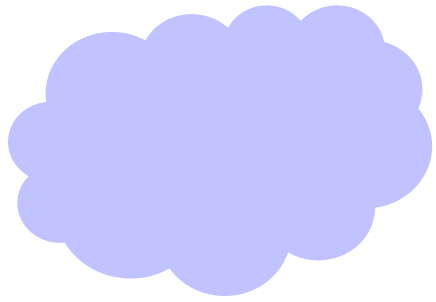
- Service (clients can get to Internet)
- Performance (network works well)
- Identity (know what's on network)
- Safety (no one launching attacks)
- Accountability (can find bad users)

(Some) Goals



**Intermediate
ISPs**

(Some) Goals



- Service (deliver traffic -> earn \$\$)
- Reliability & Performance (network works well)
- Integrity of delivered traffic (can bill customers properly, you're not over-charged by providers)

**Intermediate
ISPs**

(Some) Goals



Server

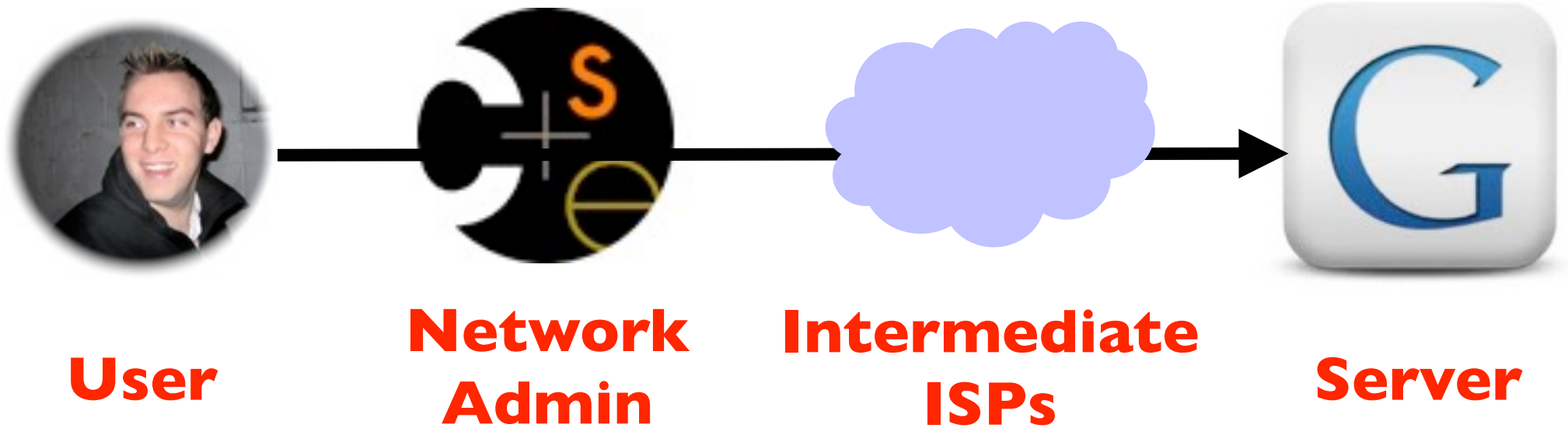
(Some) Goals



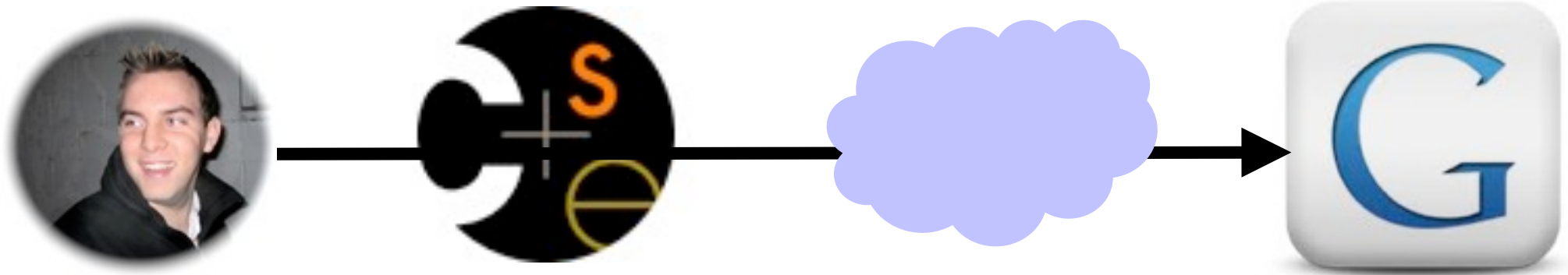
Server

- Service (deliver traffic -> earn \$\$)
- Reliability & Performance (network works well)
- Analytics (better delivery)
- Accounting (can bill customers properly)
- Safety (not being attacked)

(Some) Malicious Goals



(Some) Malicious Goals



User

**Network
Admin**

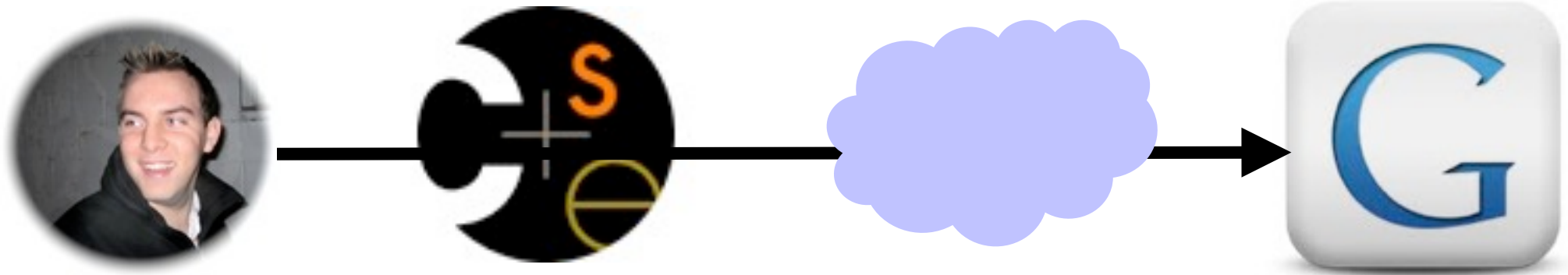
**Intermediate
ISPs**

Server

Launch
undetectable
attacks

Probe for
vulnerabilities

(Some) Malicious Goals



User

**Network
Admin**

**Intermediate
ISPs**

Server

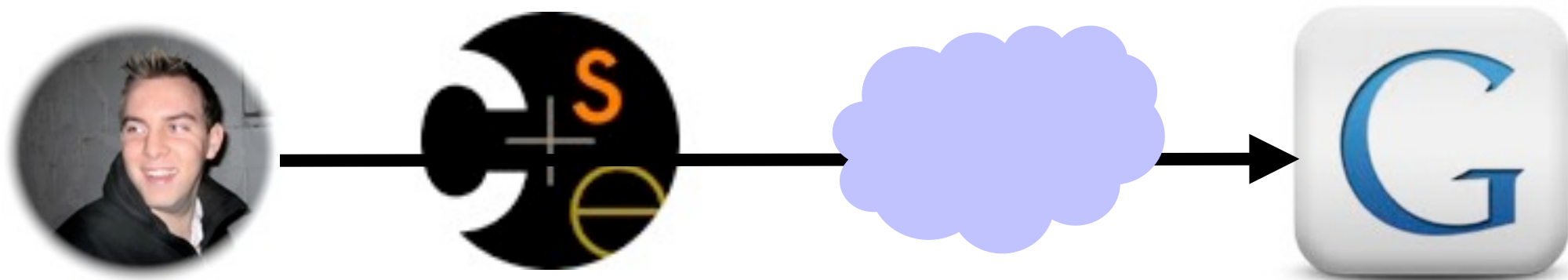
Launch
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Spy on/tamper with traffic

Impersonate servers

Probe for
vulnerabilities

(Some) Malicious Goals



User

Launch undetectable attacks

Probe for vulnerabilities

Network Admin

Spy on/tamper with traffic

Impersonate servers

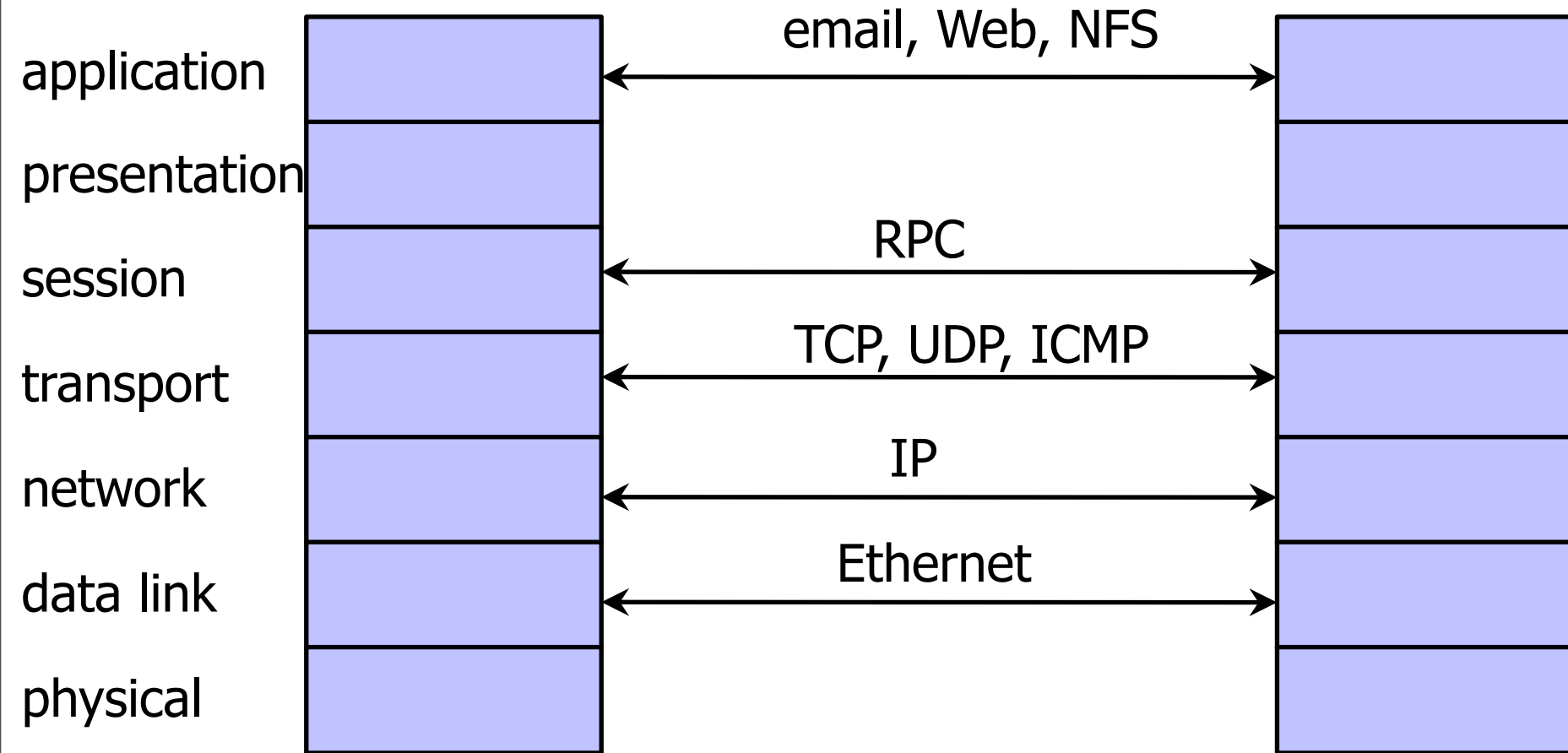
Intermediate ISPs

Server

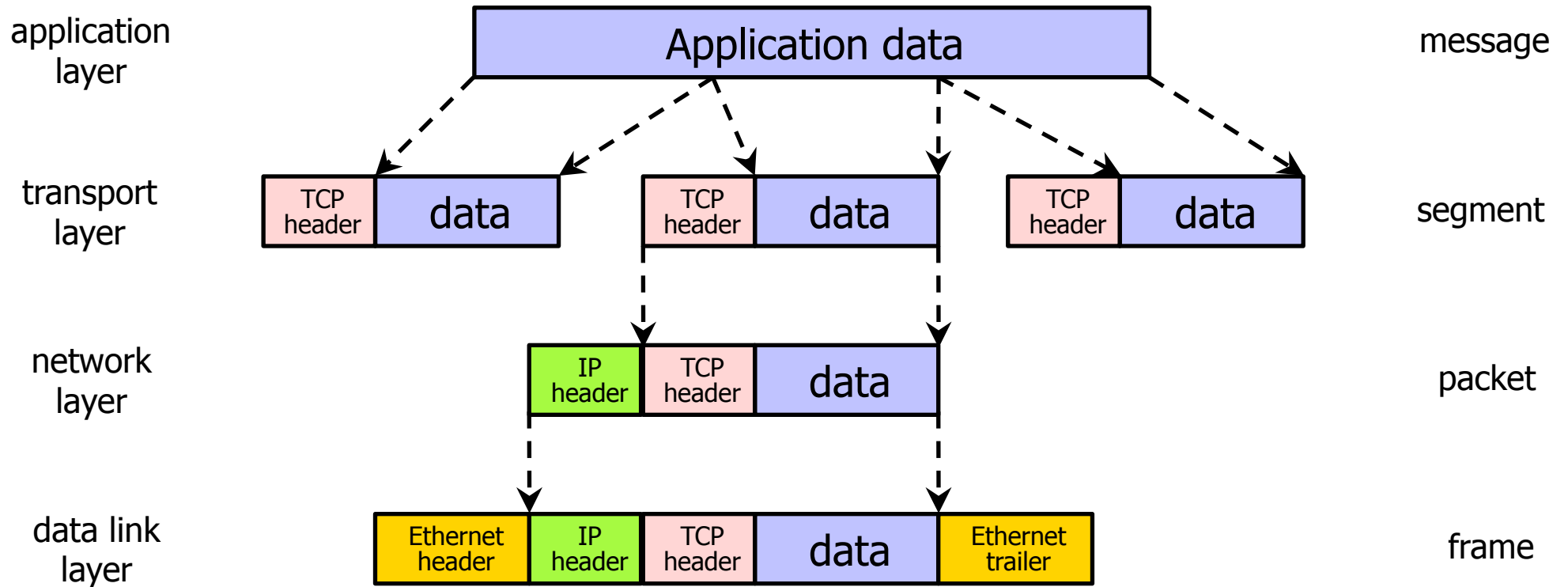
Spy on users

Identify anonymous users

OSI Protocol Stack

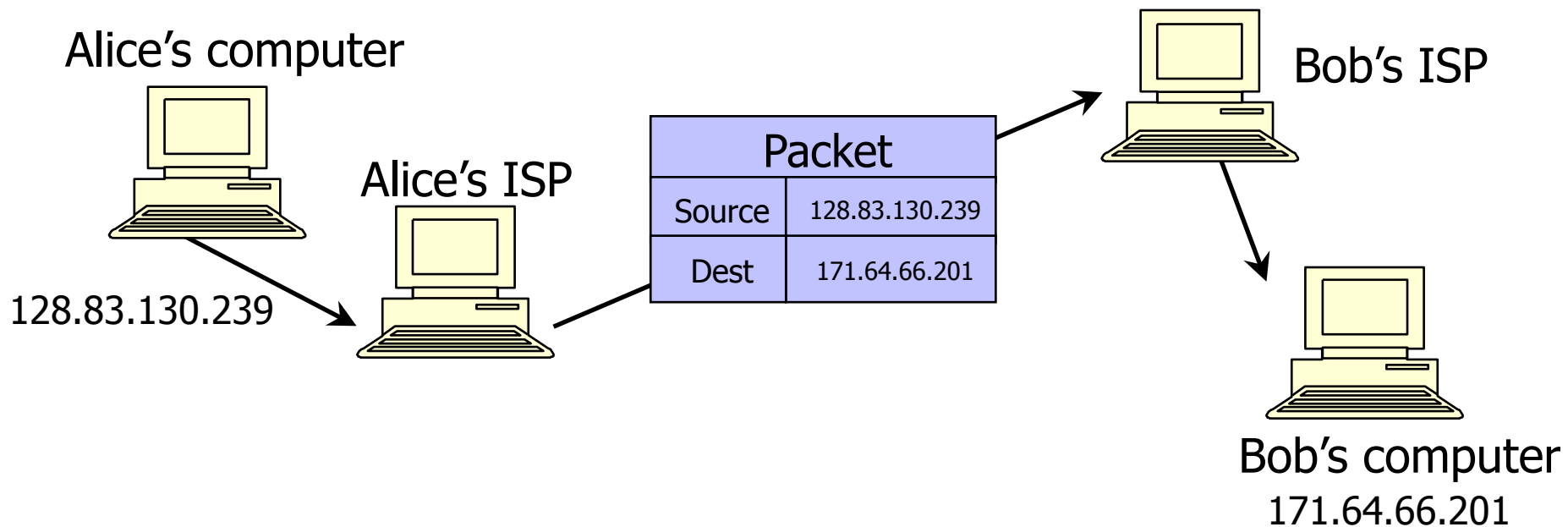


Data Formats



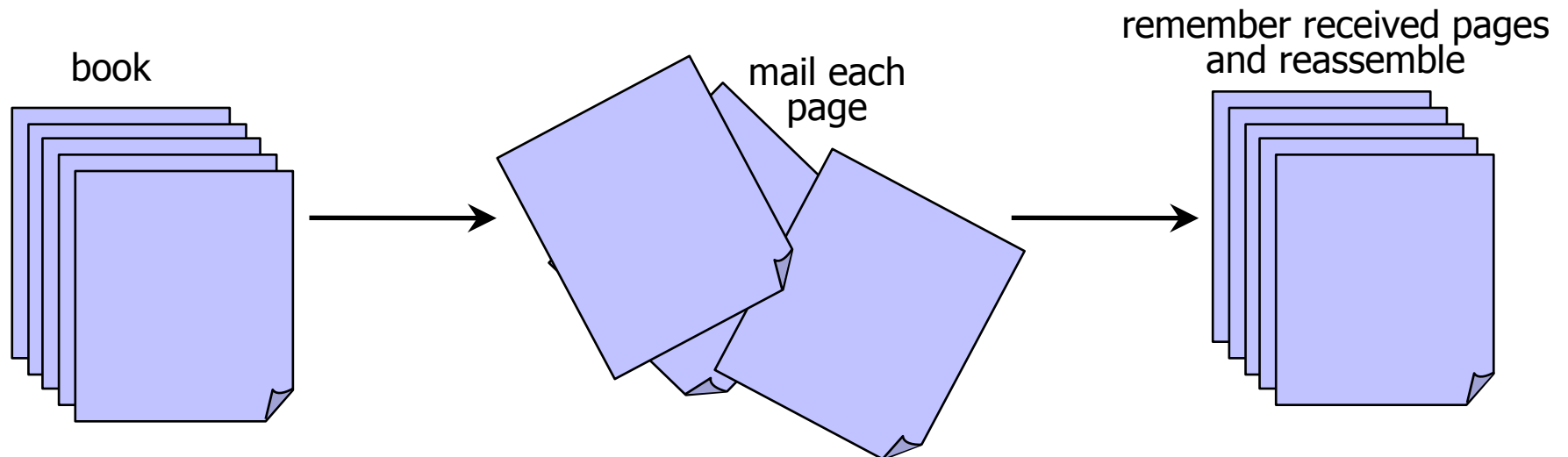
IP (Internet Protocol)

- ◆ Connectionless
 - Unreliable, “best-effort” protocol
- ◆ Uses numeric addresses for routing
 - Typically several hops in the route



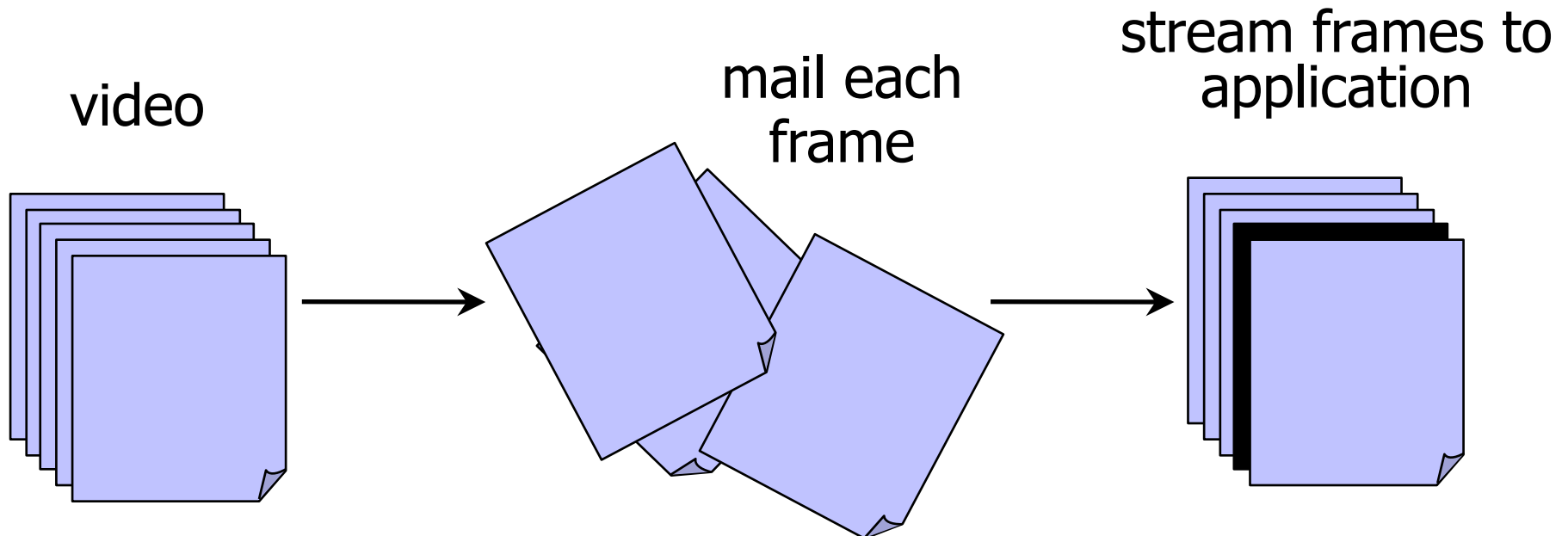
TCP (Transmission Control Protocol)

- ◆ Sender: break data into packets
 - Sequence number is attached to every packet
- ◆ Receiver: reassemble packets in correct order
 - Acknowledge receipt; lost packets are re-sent
- ◆ Connection state maintained on both sides



UDP (User Datagram Protocol)

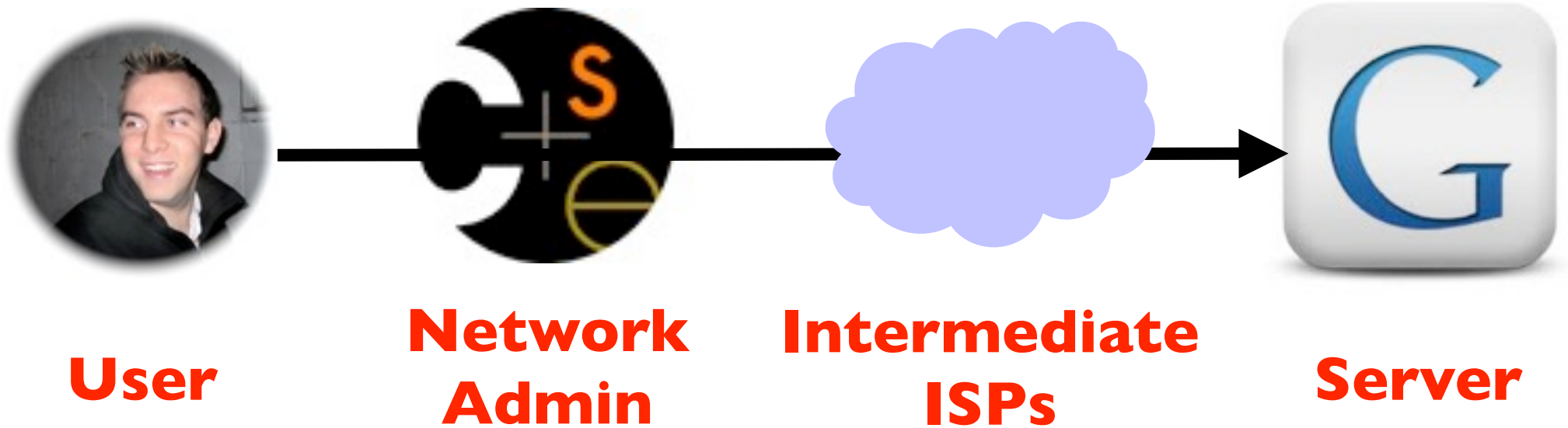
- ◆ Sender: break data into packets
 - Sequence number - maybe? If Application wants them
- ◆ Receiver: receive packets
 - No acknowledgement
 - Dropped packets are skipped - no retransmission



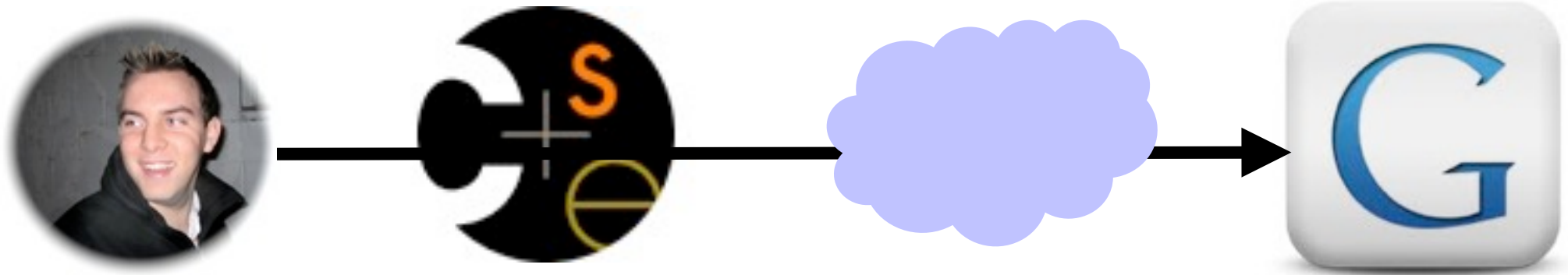
ICMP (Control Message Protocol)

- ◆ Provides feedback about network operation
 - “Out-of-band” messages carried in IP packets
 - Error reporting, congestion control, reachability, etc.
- ◆ Example messages:
 - Destination unreachable
 - Time exceeded
 - Parameter problem
 - Redirect to better gateway
 - Reachability test (echo / echo reply)
 - Message transit delay (timestamp request / reply)

(Some) Malicious Goals



(Some) Malicious Goals



User

**Network
Admin**

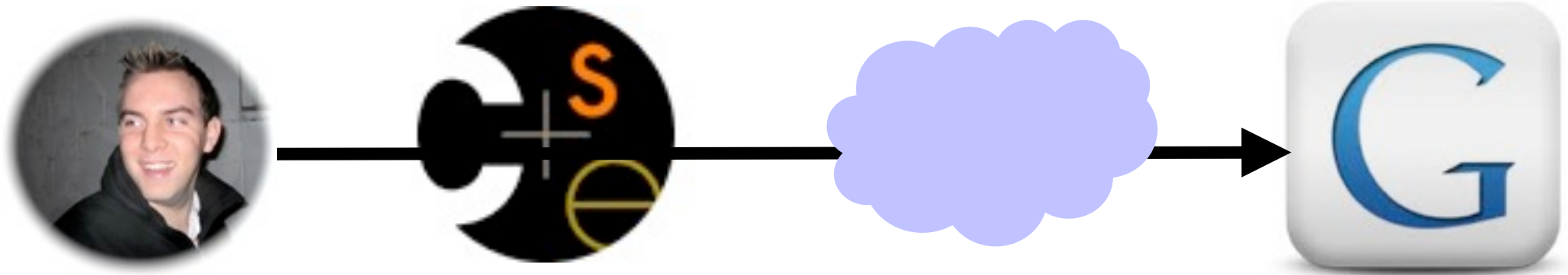
**Intermediate
ISPs**

Server

Launch
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(Some) Malicious Goals



User

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**Intermediate
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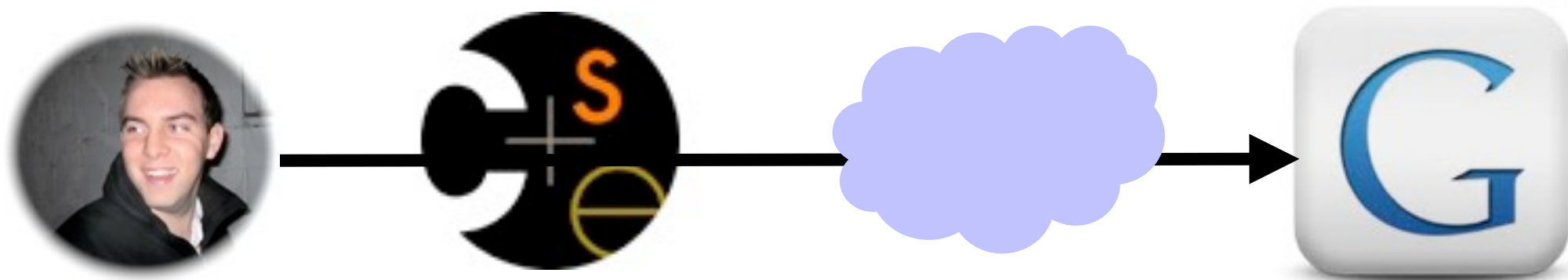
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Spy on/tamper with traffic

Impersonate servers

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(Some) Malicious Goals



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**Intermediate
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Launch
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Spy on/tamper with traffic

Identify
anonymous
users

Impersonate servers

Probe for
vulnerabilities

Detecting attacks



User

Launch
undetectable
attacks

Detecting attacks



User

Launch
undetectable
attacks

- **Problem:** IP packets contain source IP address

Detecting attacks



User

Launch
undetectable
attacks

- **Problem:** IP packets contain source IP address

Detecting attacks

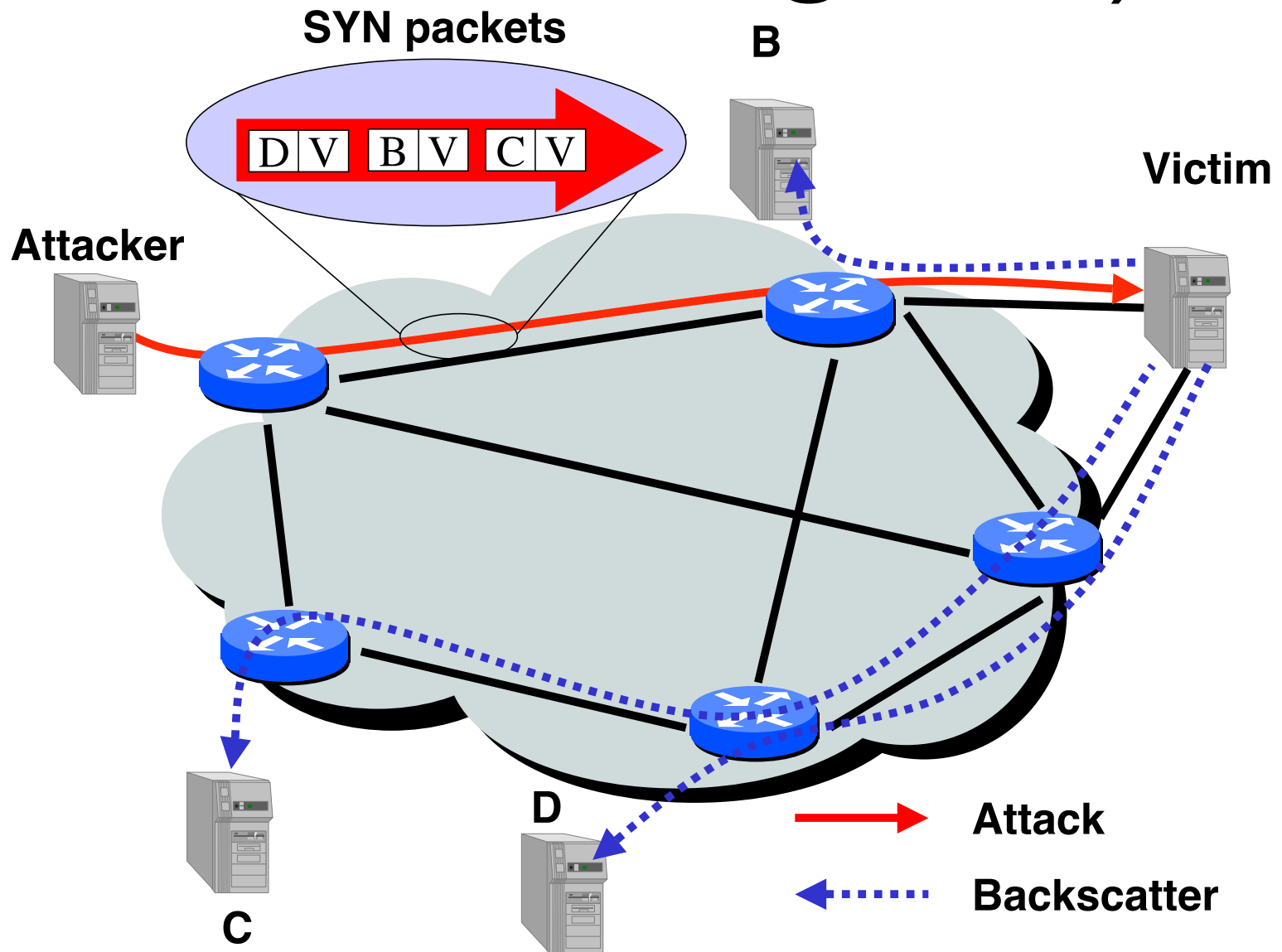


User

Launch
undetectable
attacks

- **Problem:** IP packets contain source IP address
- **Solution:** Spoof IP address

Inferring DDOS (Moore, Voelker, Savage '01)



Finding vulnerabilities



User

Probe for
vulnerabilities

Finding vulnerabilities



- **Many, many tools**

User

Probe for
vulnerabilities

Finding vulnerabilities



User

Probe for
vulnerabilities

- **Many, many tools**
- One example: **Nmap**
 - Many services have known TCP/UDP ports
 - These give away what services you're running

Nmap example (me)

```
dhalperi@dhm cse484 % nmap dsp.cs.washington.edu
```

```
Starting Nmap 5.51 ( http://nmap.org ) at 2011-12-05 14:05 PST
```

```
Nmap scan report for dsp.cs.washington.edu (128.208.4.246)
```

```
Host is up (0.0062s latency).
```

```
Not shown: 996 closed ports
```

```
PORT      STATE SERVICE
```

```
22/tcp    open  ssh
```

```
139/tcp   open  netbios-ssn
```

```
443/tcp   open  https
```

```
445/tcp   open  microsoft-ds
```

```
Nmap done: 1 IP address (1 host up) scanned in 1.36 seconds
```

Nmap example (aqua)

```
dhalperi@dhm cse484 % nmap aqua.cs.washington.edu
```

```
Starting Nmap 5.51 ( http://nmap.org ) at 2011-12-05 14:06 PST
```

```
Nmap scan report for aqua.cs.washington.edu (128.208.4.187)
```

```
Host is up (0.0022s latency).
```

```
Not shown: 990 filtered ports
```

```
PORT      STATE SERVICE
```

```
80/tcp    open  http
```

```
135/tcp   open  msrpc
```

```
139/tcp   open  netbios-ssn
```

```
445/tcp   open  microsoft-ds
```

```
1025/tcp  open  NFS-or-IIS
```

```
1026/tcp  open  LSA-or-nterm
```

```
1027/tcp  open  IIS
```

```
1028/tcp  open  unknown
```

```
1048/tcp  open  neod2
```

```
3389/tcp  open  ms-term-serv
```

```
Nmap done: 1 IP address (1 host up) scanned in 5.29 seconds
```

telnet example

Fingerprinting users



Server

Identify
anonymous
users

Fingerprinting users



Server

Identify
anonymous
users

- **Browser**

Fingerprinting users



Server

Identify
anonymous
users

- **Browser**
- **Clocks**

Fingerprinting users



Server

Identify
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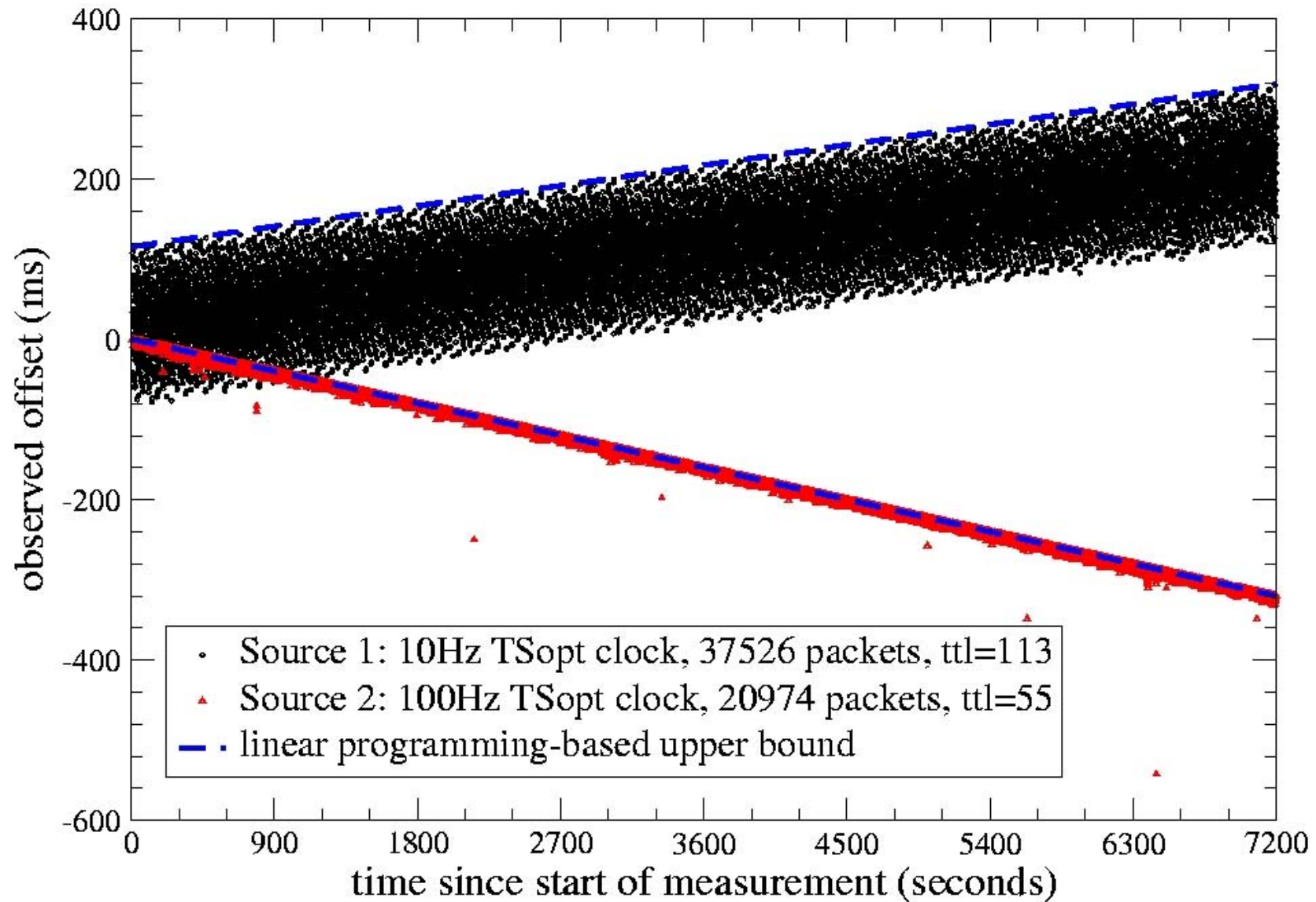
- **Browser**
- **Clocks**
- More

Browser example

<http://panopticklick.eff.org/>

Clocks

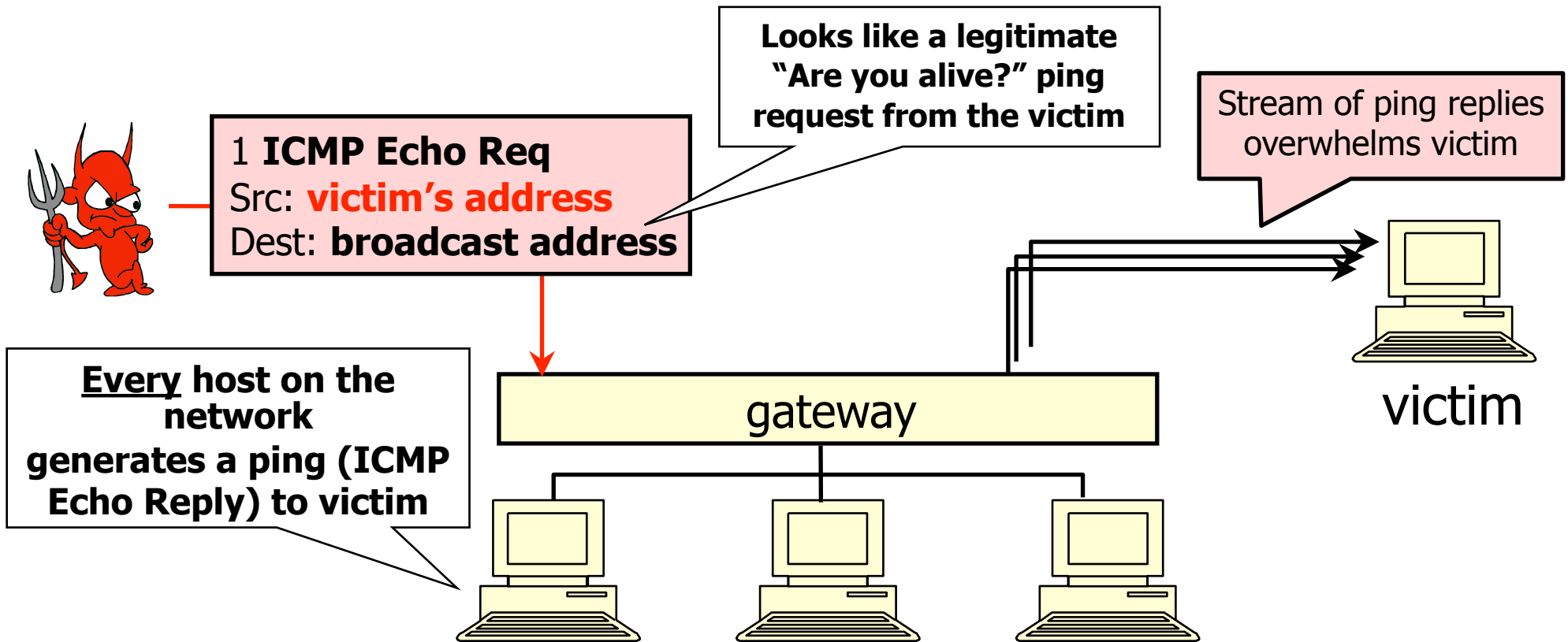
Clocks



Security Issues in TCP/UDP

- ◆ Network packets pass through/by untrusted hosts
 - Eavesdropping (packet sniffing)
 - Modifications
- ◆ IP addresses are public
 - Smurf attacks
 - Anonymity?
- ◆ TCP connection requires state
 - SYN flooding
- ◆ TCP state is easy to guess
 - TCP spoofing and connection hijacking

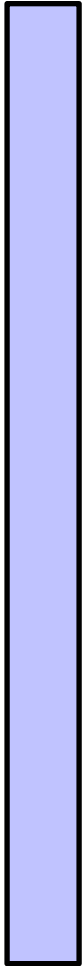
Smurf Attack



Solution: reject external packets to broadcast addresses

TCP Handshake

C

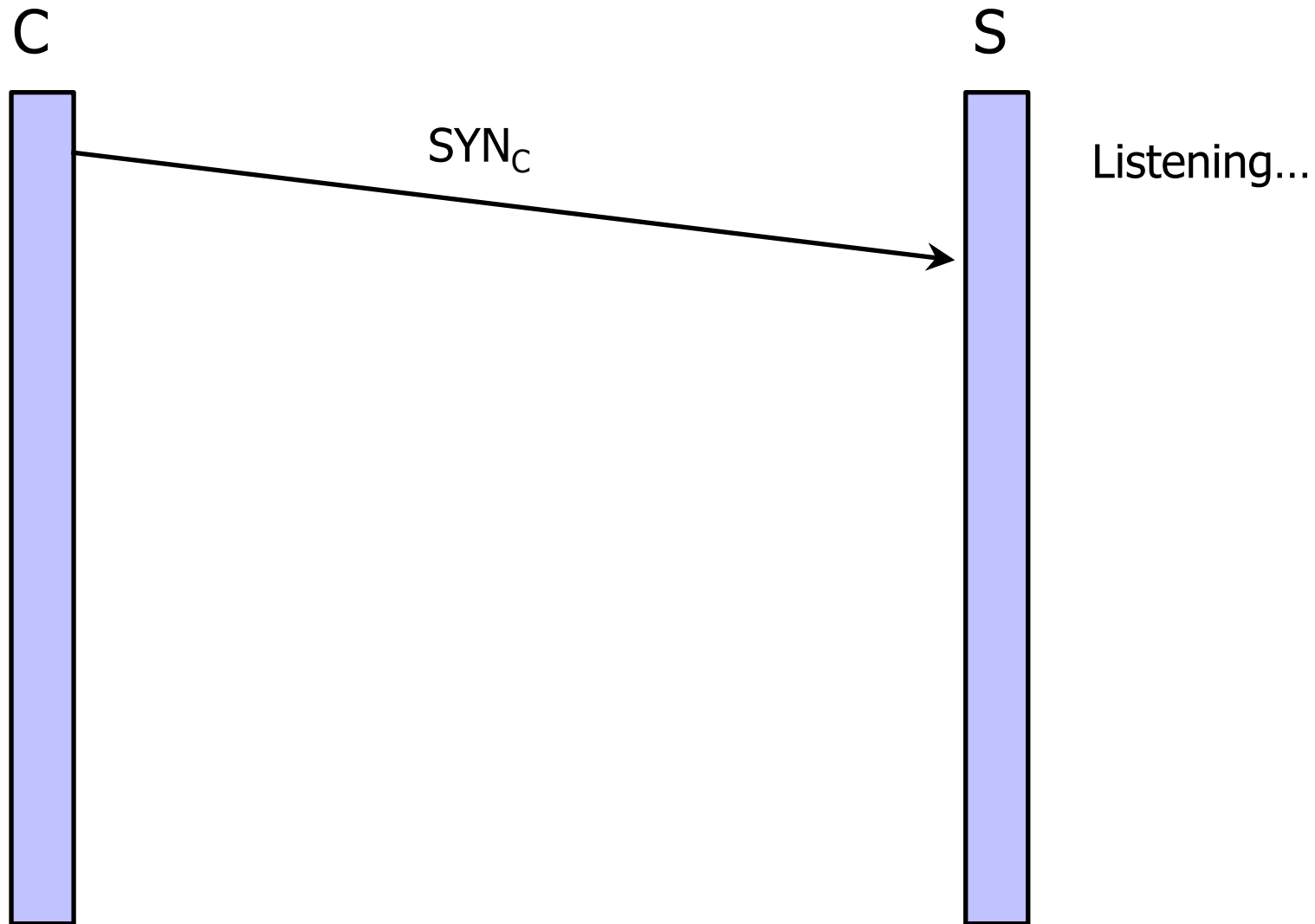


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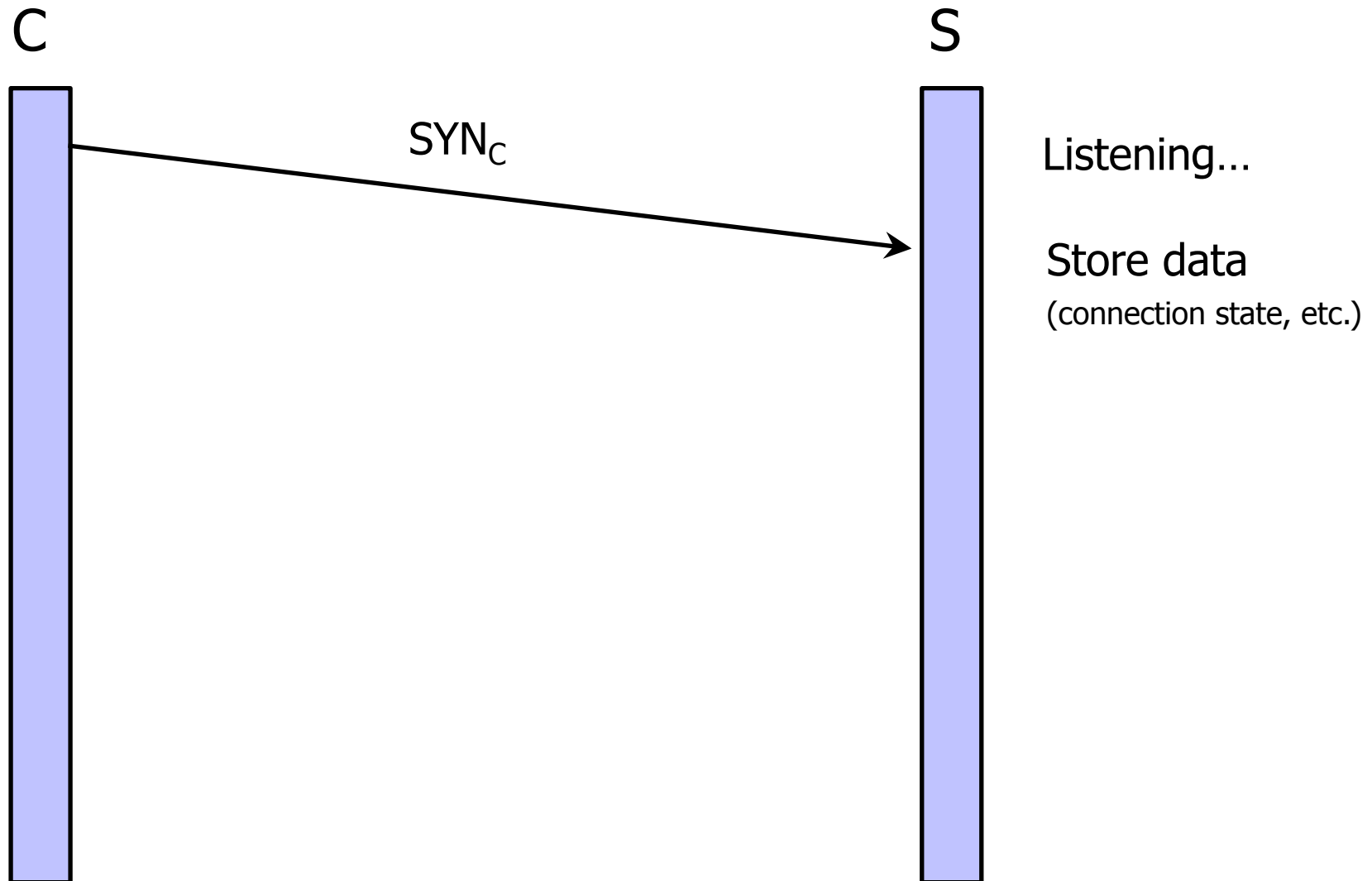


Listening...

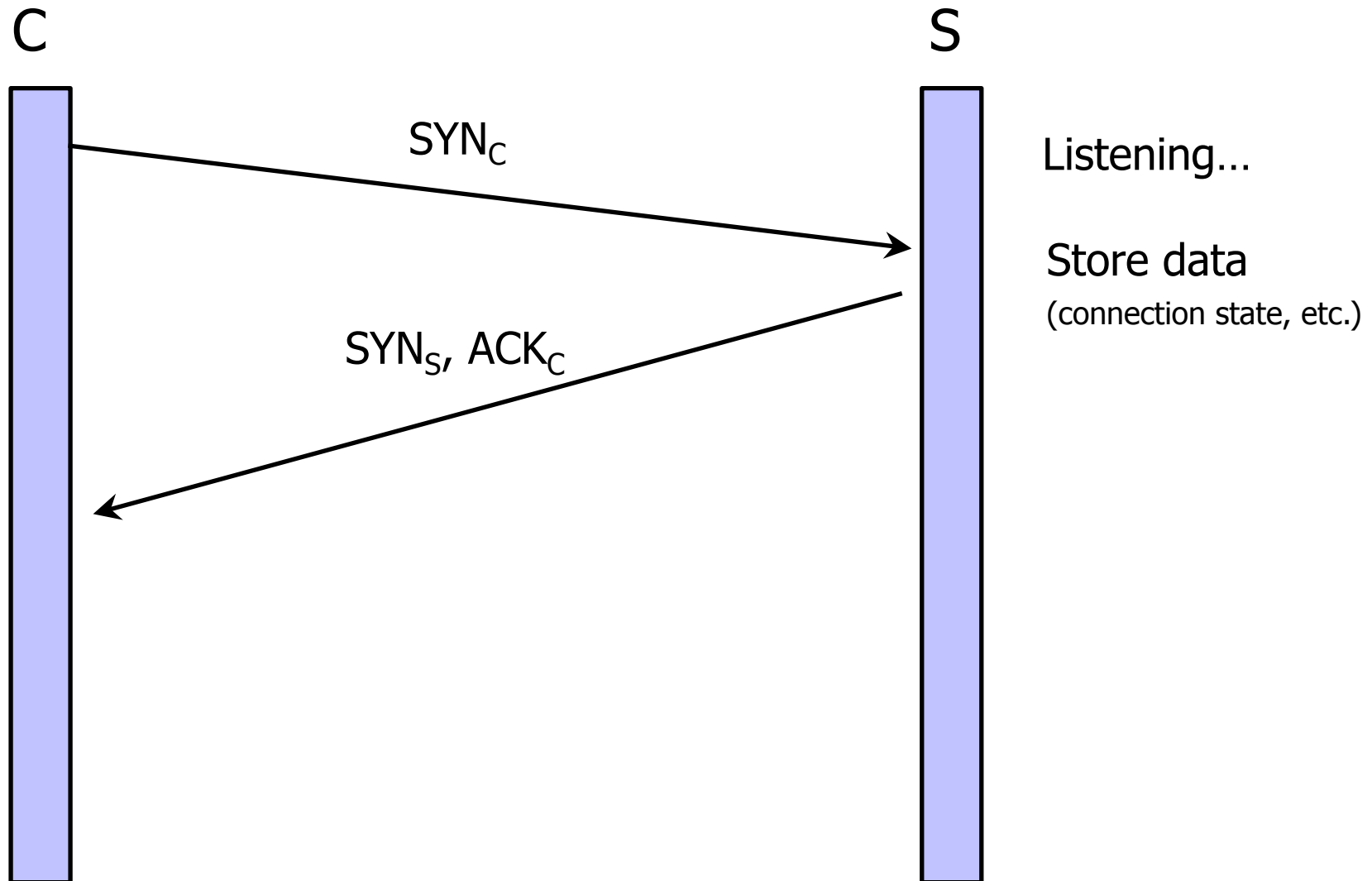
TCP Handshake



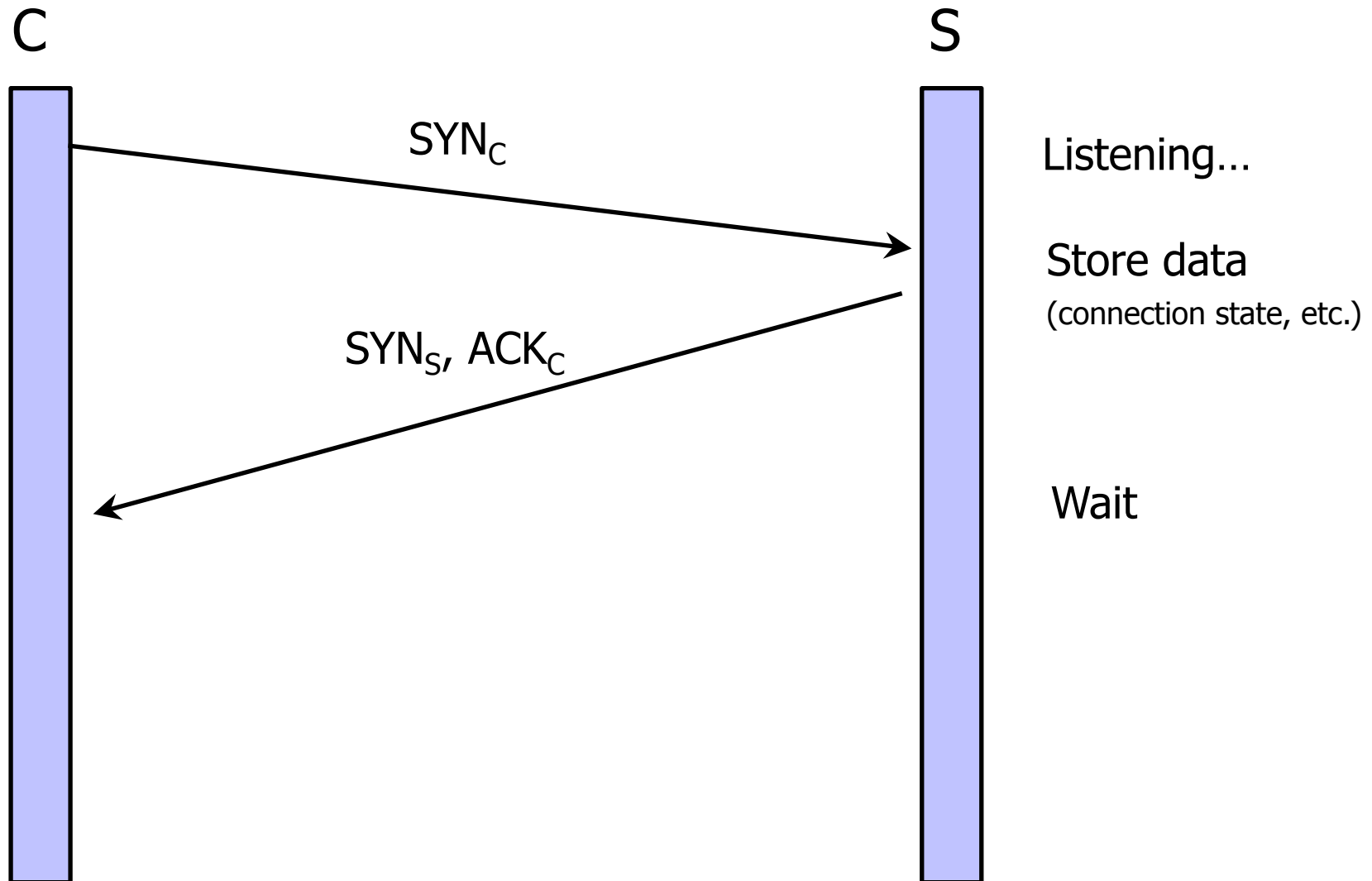
TCP Handshake



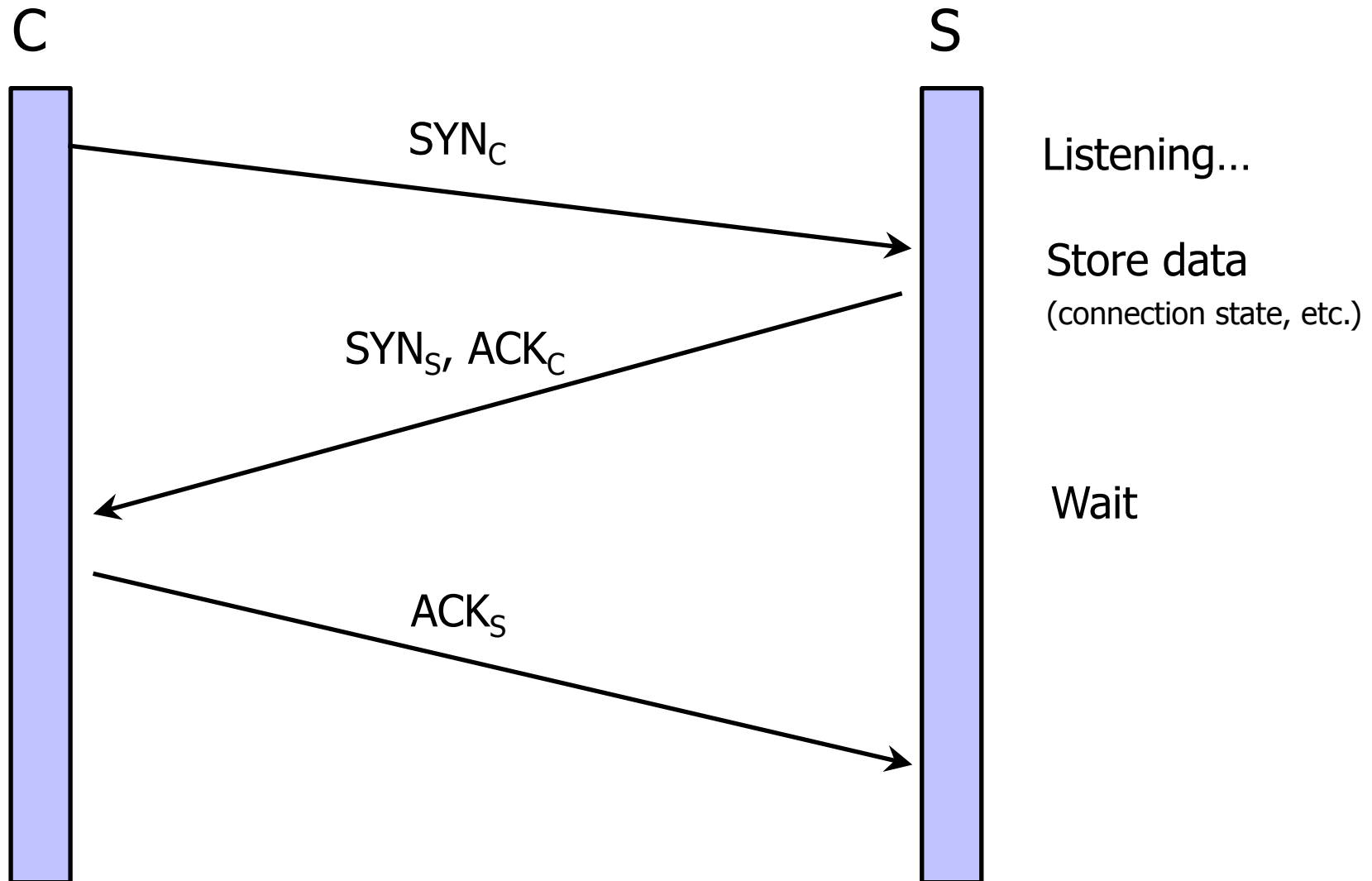
TCP Handshake



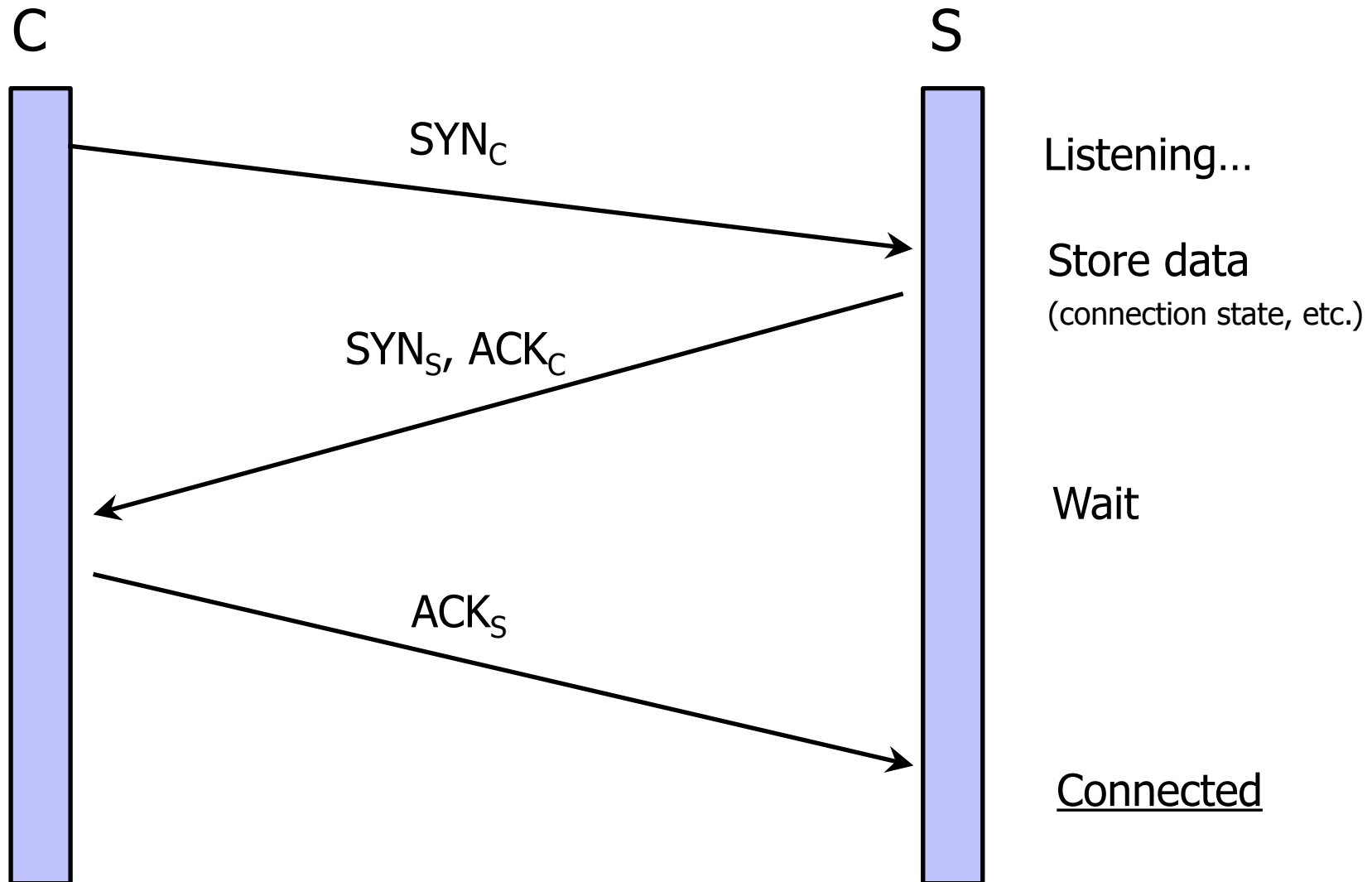
TCP Handshake



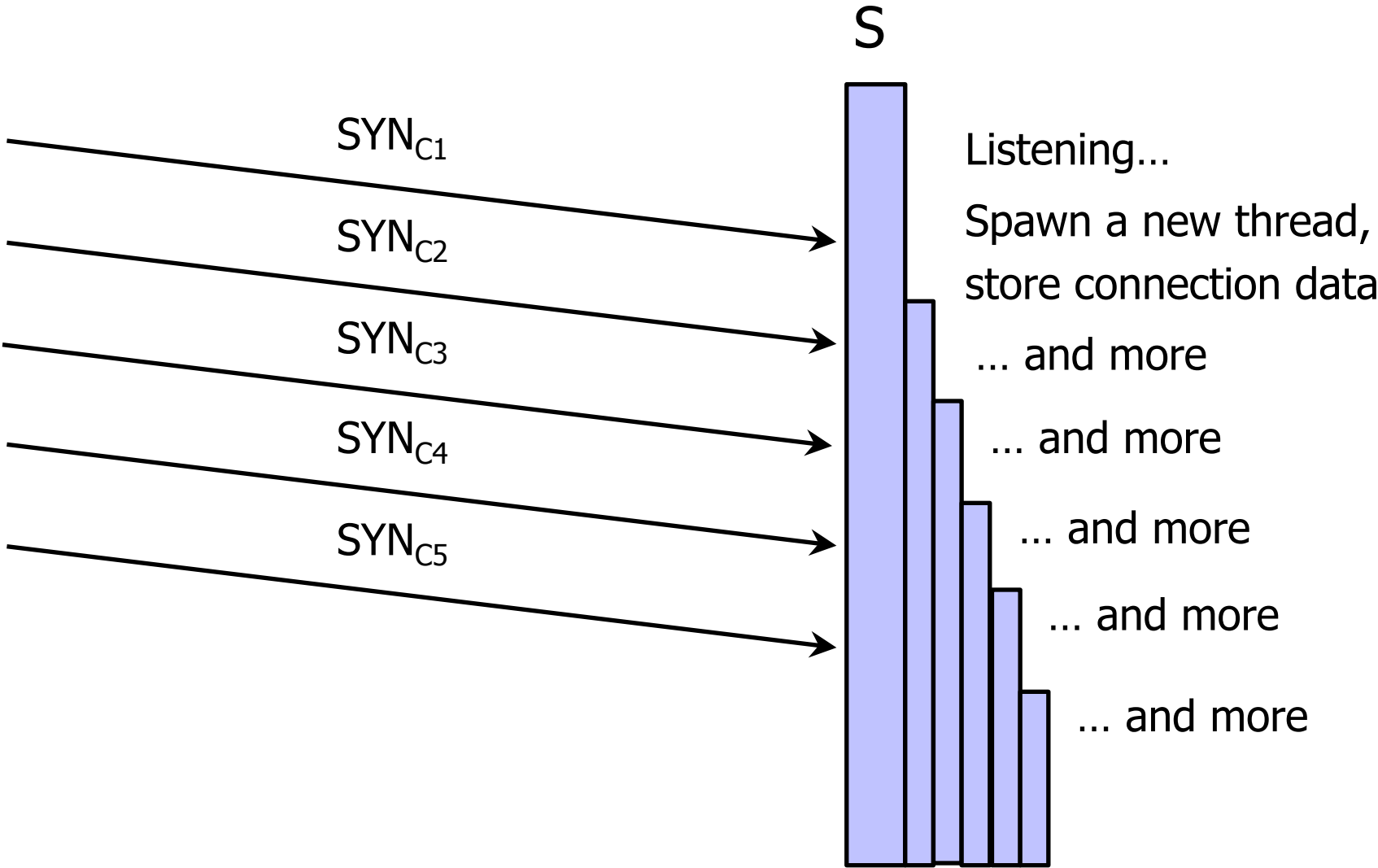
TCP Handshake



TCP Handshake



SYN Flooding Attack



SYN Flooding Explained

- ◆ Attacker sends many connection requests with spoofed source addresses
- ◆ Victim allocates resources for each request
 - Connection state maintained until timeout
 - Fixed bound on half-open connections
- ◆ Once resources exhausted, requests from legitimate clients are denied
- ◆ This is a classic **denial of service (DoS)** attack
 - Common pattern: it costs nothing to TCP initiator to send a connection request, but TCP responder must allocate state for each request (**asymmetry!**)