

CSE 490 GZ
Assignment 6
Due February 20, 2004

1. In this problem you will investigate how the GLA algorithm works in 2-dimensions. Consider 2-dimensional vectors with each coordinate having 15 values, 0 to 15. Let our training set be $X = \{(0, 0), (1, 1), (2, 2), \dots, (15, 15)\}$. (Note: let's assume that the rounding function rounds down on .5, for example round of 6.5 is 6.)
 - (a) Starting with vectors $c(0) = (0, 7)$ and $c(1) = (15, 7)$ run the GLA algorithm until there is no decrease in distortion.
 - (b) What happens to the GLA if the starting vectors are $c(0) = (0, 15)$ and $c(1) = (14, 1)$?
 - (c) Run the GLA algorithm with the splitting strategy. When splitting a codeword c , create a new codeword $c' = c + (1, 1)$.

2. In this problem you will explore Orchard's algorithm, an alternative to the k-d tree for fast full search. The algorithm is explained in the paper on fast full search in the reading on the web. Suppose there are n codewords $c[i]$, $0 \leq i \leq n - 1$. We construct a $n \times n$ array A where $A[i, j]$ where $A[i, 0] = i$, $A[i, 1]$ is the index of the closest codeword to codeword i , $A[i, 2]$ is the index of the second closest codeword to codeword i , and so on. Generally, $A[i, j]$ is the index of the j -th closest codeword to codeword i . In a simple version of Orchard's algorithm, suppose we are searching for the closest codeword to vector c . Every search starts with codeword indexed $k = 1$, $c[1]$. Compute $r = \|c - c[1]\|$. Generally, r is the distance from c to $c[k]$, where k is index of the closest codeword to c found so far. In each iteration of the algorithm we search the subarray $A[k, j]$, $j = 1, 2, \dots, n$, until either (i) we find the first k' such that $\|c - c[k']\| < r$ or (ii) we find the first k' such that $\|c[k] - c[k']\| \geq 2r$ or (iii) the subarray is searched and neither (i) nor (ii) occurs. In case (i) we set k to k' and recompute $r = \|c - c[k]\|$ and do another iteration. In cases (ii) and (iii) we are done and k is the index of the closest codeword to c .
 - (a) Consider the codewords $(4, 4), (10, 4), (7, 7), (4, 10), (10, 10)$ indexed 0 to 4, respectively. Construct the array A needed to run Orchard's algorithm.
 - (b) Run Orchard's algorithm on the input $(9, 7)$.
 - (c) Generally, Orchard's algorithm will be run for many inputs in the encoding process.

- i. Explain in detail how taking square roots can be avoided.
- ii. Explain how an additional data structure (array) can be constructed during the sorting phase so that no distance computation has to be made in the test used in (ii).
- iii. Explain how an additional data structure (array) can be used during the algorithm to avoid computing the distance twice between c and a codeword in the test used in (i).