#### CSE 521: Design and Analysis of Algorithms I

Fall 2020

# Background / Cheat Sheet

In this note I will discuss several background materials that we will discuss and exploit many times throughout this course.

## 1 Randomized Algorithm

**Expectation:** For a random variable X with domain, the discrete set S,

$$\mathbb{E}\left[X\right] = \sum_{s \in S} \mathbb{P}\left[X = s\right] s.$$

**Linearity of Expectation:** For any two Random variables X, Y,

$$\mathbb{E}\left[X+Y\right] = \mathbb{E}\left[X\right] + \mathbb{E}\left[Y\right].$$

**Variance:** The variance of a random variable X is defined as  $Var(X) = \mathbb{E}[(X - \mathbb{E}[X])^2]$ . The following identity always holds,

$$\operatorname{Var}(X) = \mathbb{E}\left[X^2\right] - (\mathbb{E}\left[X\right])^2.$$

The standard deviation of X,  $\sigma(X) = \sqrt{\operatorname{Var}(X)}$ .

**Mutual Independence** A set of random variables  $X_1, \ldots, X_n$  are mutually independent if for any  $S \subseteq \{1, \ldots, n\}$ ,

$$\mathbb{E}\left[\prod_{i\in S} X_i\right] = \prod_{i\in S} \mathbb{E}\left[X_i\right].$$

**k-wise Independence** For an integer  $k \geq 2$ , a set of random variables  $X_1, \ldots, X_n$  is set to be k-wise independent if for any set  $S \subseteq \{1, \ldots, n\}$  of size k,

$$\mathbb{E}\left[\prod_{i\in S} X_i\right] = \prod_{i\in S} \mathbb{E}\left[X_i\right].$$

Sum of Variance: Let  $X_1, \ldots, X_n$  be pairwise independent random variables, then

$$\operatorname{Var}(X_1 + \dots + X_n) = \operatorname{Var}(X_1) + \dots + \operatorname{Var}(X_n).$$

Markov's Inequality Let X be a nonnegative random variable, then for any  $k \geq 0$ ,

$$\mathbb{P}\left[X \ge k\right] \le \frac{\mathbb{E}\left[X\right]}{k}.$$

Chebyshev's Inequality For any random variable X and any  $\epsilon > 0$ ,

$$\mathbb{P}\left[\left|X - \mathbb{E}\left[X\right]\right| > \epsilon\right] \le \frac{\operatorname{Var}(X)}{\epsilon^2}.$$

So, equivalently,

$$\mathbb{P}\left[|X - \mathbb{E}\left[X\right]| > k\sigma(X)\right] \le \frac{1}{k^2}.$$

**Hoeffding's Inequality** Let  $X_1, ..., X_n$  be independent random variables where for all  $i, X_i \in [a_i, b_i]$ . Then, for any  $\epsilon > 0$ ,

$$\mathbb{P}\left[\left|\sum_{i=1}^{n} X_i - \mathbb{E}\sum_{i=1}^{n} X_i\right| > \epsilon\right] \le 2\exp\left(\frac{-2\epsilon^2}{\sum_{i=1}^{n} (a_i - b_i)^2}\right)$$

**Multiplicative Chernoff Bound** Let  $X_1, \ldots, X_n$  be independent Bernoulli random variables, i.e., for all  $i, X_i \in \{0, 1\}$ , and let  $X = X_1 + \cdots + X_n$  and  $\mu = \mathbb{E}[X]$ . Then, for any  $\epsilon > 0$ ,

$$\mathbb{P}\left[X>(1+\epsilon)\mu\right] \leq \left(\frac{e^{\epsilon}}{(1+\epsilon)^{1+\epsilon}}\right)^{\mu} \leq e^{-\frac{\epsilon^2\mu}{2+\epsilon}},$$

and

$$\mathbb{P}\left[X < (1 - \epsilon)\mu\right] \le e^{-\epsilon^2 \mu/2}$$

**McDiarmid's Inequality** Let  $X_1, \ldots, X_n \in \mathcal{X}$  be independent random variables. Let  $f: \mathcal{X}^n \to \mathbb{R}$ . If for all  $1 \le i \le n$  and for all  $x_1, \ldots, x_n$  and  $\tilde{x}_i$ ,

$$|f(x_1,\ldots,x_n)-f(x_1,\ldots,x_{i-1},\tilde{x}_i,x_{i+1},\ldots,x_n)| \le c_i,$$

then,

$$\mathbb{P}\left[|f(X_1,\ldots,X_n) - \mathbb{E}f(X_1,\ldots,X_n)| > \epsilon\right] \le 2\exp\left(-\frac{-2\epsilon^2}{\sum_i c_i^2}\right).$$

Concentration of Gaussians Let  $X_1, \ldots, X_n$  be independent standard normal random variables i.e., for all  $i, X_i \sim \mathcal{N}(0, 1)$ . Then, for any  $\epsilon > 0$ ,

$$\mathbb{P}\left[\left|\sum_{i=1}^{n} X_i^2 - n\right| > \epsilon\right] \le 2\exp\left(\frac{\epsilon^2}{8}\right)$$

Gaussian Density Function The density function of a 1-dimensional normal random variable  $X \sim \mathcal{N}(\mu, \sigma^2)$  is as follows:

$$\frac{1}{\sqrt{2\pi\sigma^2}}e^{-(x-\mu)^2/2\sigma^2}.$$

More generally, we say  $X_1, \ldots, X_n$  form a multivariate normal random variable when they have following density function:

$$\det(2\pi\Sigma)^{-1/2}e^{-(x-\mu)^{\mathsf{T}}\Sigma^{-1}(x-\mu)/2}$$

where  $\Sigma$  is the covariance matrix of  $X_1, \ldots, X_n$ . In particular, for all i, j,

$$\Sigma_{i,j} = \operatorname{Cov}(X_i, X_j) = \mathbb{E}\left[X_i - \mathbb{E}\left[X_i\right]\right] \mathbb{E}\left[X_j - \mathbb{E}\left[X_j\right]\right] = \mathbb{E}\left[X_i X_j\right] - \mathbb{E}\left[X_i\right] \mathbb{E}\left[X_j\right].$$

As a special case, if  $X_1, \ldots, X_n$  are standard normals chosen independently then  $\Sigma$  is just the identity matrix.

### 2 Spectral Algorithms

**Determinant** Let  $A \in \mathbb{R}^{n \times n}$ , the determinant of A can be written as follows:

$$\det(A) = \sum_{\sigma} \prod_{i=1}^{n} A_{i,\sigma(i)} \operatorname{sgn}(\sigma).$$

where the sum is over all permutations  $\sigma$  of the numbers  $1, \ldots, n$ , and  $\operatorname{sgn}(\sigma) \in \{+1, -1\}$ . For a permutation  $\sigma$ ,  $\operatorname{sgn}(\sigma)$  is the parity of the number of swaps one needs to transform  $\sigma$  into the identity permutations. For example, for n = 4,  $\operatorname{sgn}(1, 2, 3, 4) = +1$  because we need no swaps,  $\operatorname{sgn}(2, 1, 3, 4) = -1$  because we can transform it to the identity just by swapping 1, 2 and  $\operatorname{sgn}(3, 1, 2, 4) = +1$ .

#### **Properties of Determinant**

• For a matrix  $A \in \mathbb{R}^{n \times n}$ ,  $\det(A) \neq 0$  if and only if the columns of A are linearly independent. Recall that for a set of vectors  $v_1, \ldots, v_n \in \mathbb{R}^n$ , we say they are linearly independent if for any set of coefficients  $c_1, \ldots, c_n$ 

$$c_1 v_1 + c_2 v_2 + \dots + c_n v_n = 0$$

only when  $c_1 = c_2 = \cdots = c_n = 0$ . In other words,  $v_1, \ldots, v_n$  are linearly independent if no  $v_i$  can be written as a linear combination of the rest of the vectors.

• For any matrix  $A \in \mathbb{R}^{n \times n}$ , with eigenvalues  $\lambda_1, \ldots, \lambda_n$ ,

$$\det(A) = \prod_{i=1}^{n} \lambda_i$$

So, det(A) = 0 iff A has at least one zero eigenvalue. So, it follows from the previous fact that A has a zero eigenvalue iff columns of A are linearly independent.

• For any two square matrices  $A, B \in \mathbb{R}^{n \times n}$ ,

$$\det(AB) = \det(A)\det(B).$$

Characteristic Polynomial For a matrix  $A \in \mathbb{R}^{n \times n}$  we write  $\det(xI - A)$  for an indeterminant (variable) x is called the characteristic polynomial of A. The roots of this polynomial are the eigenvalues of A. In particular,

$$\det(xI - A) = (x - \lambda_1)(x - \lambda_2)\dots(x - \lambda_n),$$

where  $\lambda_1, \ldots, \lambda_n$  are the eigenvalues of A. It follows from the above identity that for x = 0,  $\det(-A) = \prod_{i=1}^{n} \lambda_i$  or equivalently,  $\det(A) = \prod_{i=1}^{n} \lambda_i$ .

**Rank** The rank of a matrix  $A \in \mathbb{R}^{n \times n}$  is the number of nonzero eigenvalues of A. More generally, the rank of a matrix  $A \in \mathbb{R}^{m \times n}$  is the number of nonzero eigenvalues of A. Or in other words, it is the number of nonzero eigenvalues of  $AA^{\mathsf{T}}$ .

**PSD matrices** We discuss several equivalent definitions of PSD matrices. A symmetric matrix  $A \in \mathbb{R}^{n \times n}$  is positive semidefinite (PSD) iff

- All eigenvalues of A are nonnegative
- A can be written as  $BB^{\intercal}$  for some matrix  $B \in \mathbb{R}^{n \times m}$ .
- $x^{\intercal}Ax > 0$  for all vectors  $x \in \mathbb{R}^n$ .
- $\det(A_{S,S}) \geq 0$  for all  $S \subseteq \{1,\ldots,n\}$  where  $A_{S,S}$  denotes the square submatrix of A with rows and columns indexed by S.

The following fact about PSD matrices is immediate. If  $A \succeq 0$  is an  $n \times n$  matrix, then for any matrix  $C \in \mathbb{R}^{k \times n}$ ,

$$CAC^T \succ 0.$$

This is because for any vector  $x \in \mathbb{R}^k$ ,

$$x^T C A C^T x = (C^T x)^T A (C^T x) = y^T A y \ge 0,$$

where  $y = C^T x$ .

For two symmetric  $A, B \in \mathbb{R}^n$  we write  $A \leq B$  if and only if  $B - A \succeq 0$ . In other words,  $A \leq B$  if and only if for any vector  $x \in \mathbb{R}^n$ ,

$$x^T A x \le x^T B x$$
.

Let  $\lambda_1, \ldots, \lambda_n$  be the eigenvalues of A, and  $\tilde{\lambda}_1, \ldots, \tilde{\lambda}_n$  be the eigenvalues of B. If  $A \leq B$ , then for all i,  $\lambda_i \leq \tilde{\lambda}_i$ .

**Nonsymmetric Matrices** Any matrix  $A \in \mathbb{R}^{m \times n}$  (for  $m \leq n$ ) can be written as

$$A = \sum_{i=1}^{m} \sigma_i u_i v_i^{\mathsf{T}}$$

where

- $u_1, \ldots, u_m \in \mathbb{R}^m$  form an orthonormal set of vectors. These are called left singular vectors of A and they have the property,  $u_i A = \sigma_i v_i$ . These vectors are the eigenvectors of the matrix  $AA^{\mathsf{T}}$ .
- $v_1, \ldots, v_m \in \mathbb{R}^n$  form an orthonormal set of vectors. Note that these vectors do not necessarily span the space. These vectors are eigenvectors of the matrix  $A^{\mathsf{T}}A$ .
- $\sigma_1, \ldots, \sigma_m$  are called the singular values of A. They are always real an nonnegative. In fact they are eigenvalues of the PSD matrix  $AA^{\mathsf{T}}$ .

**Rotation Matrix** A matrix  $R^{n \times n}$  is a rotation matrix iff  $||Rx||_2 = ||x||_2$  for all vectors  $x \in \mathbb{R}^n$ . In other words, R as an *operator* preserves the norm of all vectors. Next, we discuss equivalent definitions of R being a rotation matrix. R is a rotation matrix iff

- $RR^{\intercal} = I$ .
- All singular values of R are 1.
- Columns of R form an orthonormal set of vectors in  $\mathbb{R}^n$ .

**Projection Matrix** A symmetric matrix  $P \in \mathbb{R}^{n \times n}$  is a projection matrix iff

- It can be written as  $P = \sum_{i=1}^{k} v_i v_i^{\mathsf{T}}$  for some  $1 \leq k \leq n$ .
- All eigenvalues of P are 0 or 1.
- PP = P.

It follows from the spectral theorem that there is a unique projection matrix of rank n and that is the identity matrix. In general a projection matrix projects any given vector x to the linear subspace corresponding to span of the vectors  $v_1, \ldots, v_k$ .

**Trace** For a square matrix  $A \in \mathbb{R}^{n \times n}$  we write

$$Tr(A) = \sum_{i=1}^{n} A_{i,i}$$

to denote the sum of entries on the diagonal of A. Next, we discuss several properties of the trace.

- Trace of A is equal to the sum of all eigenvalues of A.
- Trace is a linear operator, for any two square matrices  $A, B \in \mathbb{R}^{n \times n}$

$$\operatorname{Tr}(A+B) = \operatorname{Tr}(A) + \operatorname{Tr}(B)$$
  
 $\operatorname{Tr}(tA) = t\operatorname{Tr}(A), \forall t \in \mathbb{R}.$ 

- It follows by the previous fact that for a random matrix X,  $\mathbb{E}[\operatorname{Tr}(X)] = \operatorname{Tr}(\mathbb{E}[X])$ .
- For any pair of matrices  $A \in \mathbb{R}^{n \times k}$  and  $B \in \mathbb{R}^{k \times n}$  such that AB is a square matrix we have

$$Tr(AB) = Tr(BA).$$

So, in particular, for any vector  $v \in \mathbb{R}^n$ ,

$$\operatorname{Tr}(vv^{\mathsf{T}}) = \operatorname{Tr}(v^{\mathsf{T}}v) = ||v||^2.$$

• For any matrix  $A \in \mathbb{R}^{m \times n}$ 

$$||A||_F^2 = \sum_{i=1}^m \sum_{j=1}^n A_{i,j}^2 = \text{Tr}(AA^{\mathsf{T}}).$$

**Matrix Chernoff Bound** Let X be a random  $n \times n$  PSD matrix. Suppose that  $X \leq \alpha \mathbb{E}[X]$  with probability 1 for some  $\alpha \geq 0$ . Let  $X_1, \ldots, X_k$  be independent copies of X. Then, for any  $0 < \epsilon < 1$ ,

$$\mathbb{P}\left[(1-\epsilon)\mathbb{E}\left[X\right] \leq \frac{1}{k}(X_1+\cdots+X_k) \leq (1+\epsilon)\mathbb{E}\left[X\right]\right] \geq 1-2ne^{-\epsilon^2k/4\alpha}.$$

### 3 Optimization

**Convex Functions** A function  $f: \mathbb{R}^n \to \mathbb{R}$  is convex on a set  $S \subseteq \mathbb{R}^n$  if for any two points  $x, y \in S$ , we have

$$f\left(\frac{x+y}{2}\right) \le \frac{1}{2}(f(x) + f(y)).$$

We say f is concave if for any such  $x, y \in S$ , we have

$$f\left(f(\alpha x + (1 - \alpha)y\right) \ge \alpha f(x) + (1 - \alpha)f(y),$$

for any  $0 \le \alpha \le 1$ . There is an equivalent definition of convexity: For a function  $f : \mathbb{R}^n \to \mathbb{R}$ , the Hessian of  $f, \nabla^2 f$  is a  $n \times n$  matrix defined as follows:

$$(\nabla^2 f)_{i,j} = \partial_{x_i} \partial_{x_j} f$$

for all  $1 \le i, j \le n$ . We can show that f is convex over S if and only if for all  $a \in S$ ,

$$\nabla^2 f \Big|_{x=a} \succeq 0.$$

For example, consider the function  $f(x) = x^T A x$  for  $x \in \mathbb{R}^n$  and  $A \in \mathbb{R}^{n \times n}$ . Then,  $\nabla^2 f = A$ . So, f is convex (over  $\mathbb{R}^n$ ) if and only if  $A \succeq 0$ .

For another example, let  $f: \mathbb{R} \to \mathbb{R}$  be  $f(x) = x^k$  for some integer  $k \geq 2$ . Then,  $f''(x) = k(k-1)x^{k-2}$ . If k is an even integer,  $f''(x) \geq 0$  over all  $x \in \mathbb{R}$ , so f is convex over all real numbers. On the other hand, if k is an odd integer then  $f''(x) \geq 0$  if and only if  $x \geq 0$ . So, in this f is convex only over non-negative reals.

Similarly, f is concave over S, if  $\nabla^2 f\Big|_{x=a} \leq 0$  for all  $a \in S$ . For example,  $x \mapsto \log x$  is concave over all positive reals.

Convex set We say a set  $S \subseteq \mathbb{R}^n$  is convex if for any pair of points  $x, y \in S$ , the line segment connecting x to y is in S.

For example, let  $f: \mathbb{R}^n \to \mathbb{R}$  be a convex function over a set  $S \subseteq \mathbb{R}^n$ . Let  $t \in \mathbb{R}$ , and define

$$T = \{x \in \mathbb{R}^n : f(x) \le t\}.$$

Then, T is convex. This is because if  $x, y \in T$ , then for any  $0 \le \alpha \le 1$ ,

$$f(\alpha x + (1 - \alpha)y) \le \alpha f(x) + (1 - \alpha)f(y) \le \alpha t + (1 - \alpha)t = t$$

where the first inequality follows by convexity of f. So,  $\alpha x + (1 - \alpha) \in T$  and T is convex.

**Norms are Convex functions** A norm  $\|\cdot\|$  is defined as a function that maps  $\mathbb{R}^n$  to  $\mathbb{R}$  and satisfies the following three properties,

- i) ||x|| > 0 for all  $x \in \mathbb{R}^n$ ,
- ii)  $\|\alpha x\| = \alpha \|x\|$  for all  $\alpha \ge 0$  and  $x \in \mathbb{R}^n$ ,
- iii) Triangle inequality:  $||x+y|| \le ||x|| + ||y||$  for all  $x, y \in \mathbb{R}^n$ .

It is easy to see that any norm function is a convex function: This is because for any  $x, y \in \mathbb{R}^n$ , and  $0 \le \alpha \le 1$ ,

$$\|\alpha x + (1 - \alpha)y\| \le \|\alpha x\| + \|(1 - \alpha)y\| = \alpha \|x\| + (1 - \alpha)\|y\|.$$

# 4 Useful Inequalities

• For real numbers,  $a_1, \ldots, a_n$  and nonnegative reals  $b_1, \ldots, b_n$ ,

$$\min_{i} \frac{a_i}{b_i} \le \frac{a_1 + \dots + a_n}{b_1 + \dots + b_n} \le \max_{i} \frac{a_i}{b_i}$$

• Cauchy-Schwartz inequality: For real numbers  $a_1, \ldots, a_n, b_1, \ldots, b_n$ ,

$$\sum_{i=1}^{n} a_i \cdot b_i \le \sqrt{\sum_i a_i^2} \cdot \sqrt{\sum_i b_i^2}$$

There is an equivalent vector-version of the above inequality. For any two vectors  $u, v \in \mathbb{R}^n$ ,

$$\sum_{i=1}^{n} u_i \cdot v_i = \langle u, v \rangle \le ||u|| \cdot ||v||$$

The equality in the above holds only when u, v are parallel.

• AM-GM inequality: For any n nonnegative real numbers  $a_1, \ldots, a_n$ ,

$$\frac{a_1 + \dots + a_n}{n} \ge (a_1 \cdot a_2 \cdot \dots a_n)^{1/n}.$$

• Relation between norms: For any vector  $a \in \mathbb{R}^n$ ,

$$||a||_2 \le ||a||_1 \le \sqrt{n} \cdot ||a||_2$$

The right inequality is just Cauchy-Schwartz inequality.

• For any real numbers  $a_1, \ldots, a_n$ ,

$$(|a_1| + \dots + |a_n|)^2 \le n(a_1^2 + \dots + a_n^2).$$

This is indeed a special case of Cauchy-Schwartz inequality.

• For any real number  $x, 1-x \le e^{-x}$ . In this course we use  $1-x \approx e^{-x}$  to simplify calculations.