## Lecture 7

## The Brunn-Minkowski Theorem and Influences of Boolean Variables

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**Theorem 7.1 (Brunn-Minkowski).** If  $A, B \subseteq \mathbb{R}^n$  satisfy some mild assumptions (in particular, convexity suffices), then

$$[\operatorname{vol}(A+B)]^{\frac{1}{n}} \ge [\operatorname{vol}(A)]^{\frac{1}{n}} + [\operatorname{vol}(B)]^{\frac{1}{n}}$$

where  $A + B = \{a + b : a \in A \text{ and } b \in B\}.$ 

*Proof.* First, suppose that A and B are axis aligned boxes, say  $A = \prod_{j=1}^{n} I_j$  and  $B = \prod_{i=1}^{n} J_i$ , where each  $I_j$  and  $J_i$  is an interval with  $|I_j| = x_j$  and  $|J_i| = y_i$ . We may assume WLOG that  $I_j = [0, x_j]$  and  $J_i = [0, y_i]$  and hence  $A + B = \prod_{i=1}^{n} [0, x_i + y_i]$ . For this case, the BM inequality asserts that

$$\prod_{i=1}^{n} (x_i + y_i)^{\frac{1}{n}} \ge \prod_{i=1}^{n} x_i^{\frac{1}{n}} \cdot \prod_{i=1}^{n} y_i^{\frac{1}{n}}$$

$$\Rightarrow \qquad 1 \ge \left[ \prod \left( \frac{x_i}{x_i + y_i} \right) \right]^{\frac{1}{n}} \cdot \left[ \prod \left( \frac{y_i}{x_i + y_i} \right) \right]^{\frac{1}{n}}$$

Now, since the arithmetic mean of n numbers is bounded above by their harmonic mean, we have  $(\prod \alpha_i)^{\frac{1}{n}} \leq \frac{\sum \alpha_i}{n}$  and  $(\prod (1 - \alpha_i))^{\frac{1}{n}} \leq \frac{\sum (1 - \alpha_i)}{n}$ . Taking  $\alpha_i = \frac{x_i}{x_i + y_i}$  and hence  $1 - \alpha_i = \frac{y_i}{x_i + y_i}$ , we see that the above inequality always holds. Hence the BM inequality holds whenever A and B are axis aligned boxes.

Now, suppose that A and B are the disjoint union of axis aligned boxes. Suppose that  $A = \bigcup_{\alpha \in \mathcal{A}} A_{\alpha}$ and  $B = \bigcup_{\beta \in \mathcal{B}} B_{\beta}$ . We proceed by induction on  $|\mathcal{A}| + |\mathcal{B}|$ . We may assume WLOG that  $|\mathcal{A}| > 1$ . Since the boxes are disjoint, there is a hyperplane separating two boxes in  $\mathcal{A}$ . We may assume WLOG that this hyperplane is  $x_1 = 0$ .



Let  $A^+ = \{x \in A : x_1 \ge 0\}$  and  $A^- = \{x \in A : x_1 \le 0\}$  as shown in the figure above. It is clear that both  $A^+$  and  $A^-$  are the disjoint union of axis aligned boxes. In fact, we may let  $A^+ = \bigcup_{\alpha \in \mathcal{A}^+} A_{\alpha}$  and  $A^- = \bigcup_{\alpha \in \mathcal{A}^-} A_{\alpha}$  where  $|\mathcal{A}^+| < |\mathcal{A}|$  and  $|\mathcal{A}^-| < |\mathcal{A}|$ . Suppose that  $\frac{\operatorname{vol}(A^+)}{\operatorname{vol}(A)} = \alpha$ . Pick a  $\lambda$  so that

$$\frac{\operatorname{vol}\left(\left\{x \in B : x_1 \ge \lambda\right\}\right)}{\operatorname{vol}\left(B\right)} = \alpha$$

We can always do this by the mean value theorem because the function  $f(\lambda) = \frac{\operatorname{vol}(\{x \in B : x_1 \ge \lambda\})}{\operatorname{vol}(B)}$  is continuous, and  $f(\lambda) \to 0$  as  $\lambda \to \infty$  and and  $f(\lambda) \to 1$  as  $\lambda \to -\infty$ .

Let  $B^+ = \{x \in B : x_1 \ge \lambda\}$  and  $B^- = \{x \in B : x_1 \le \lambda\}$ . By induction, we may apply BM to both  $(A^+, B^+)$  and  $(A^-, B^-)$ , obtaining

$$\left[ \operatorname{vol} \left( A^{+} + B^{+} \right) \right]^{\frac{1}{n}} \ge \left[ \operatorname{vol} \left( A^{+} \right) \right]^{\frac{1}{n}} + \left[ \operatorname{vol} \left( B^{+} \right) \right]^{\frac{1}{n}}$$
$$\left[ \operatorname{vol} \left( A^{-} + B^{-} \right) \right]^{\frac{1}{n}} \ge \left[ \operatorname{vol} \left( A^{-} \right) \right]^{\frac{1}{n}} + \left[ \operatorname{vol} \left( B^{-} \right) \right]^{\frac{1}{n}}$$

Now,

$$\left[ \operatorname{vol} \left( A^+ \right) \right]^{\frac{1}{n}} + \left[ \operatorname{vol} \left( B^+ \right) \right]^{\frac{1}{n}} = \alpha^{\frac{1}{n}} \left[ \left[ \operatorname{vol} \left( A \right) \right]^{\frac{1}{n}} + \left[ \operatorname{vol} \left( B \right) \right]^{\frac{1}{n}} \right]$$
$$\left[ \operatorname{vol} \left( A^- \right) \right]^{\frac{1}{n}} + \left[ \operatorname{vol} \left( B^- \right) \right]^{\frac{1}{n}} = (1 - \alpha)^{\frac{1}{n}} \left[ \left[ \operatorname{vol} \left( A \right) \right]^{\frac{1}{n}} + \left[ \operatorname{vol} \left( B \right) \right]^{\frac{1}{n}} \right]$$

Hence

$$\left[\operatorname{vol}\left(A^{+}+B^{+}\right)\right]^{\frac{1}{n}}+\left[\operatorname{vol}\left(A^{-}+B^{-}\right)\right]^{\frac{1}{n}} \ge \left[\left[\operatorname{vol}\left(A\right)\right]^{\frac{1}{n}}+\left[\operatorname{vol}\left(B\right)\right]^{\frac{1}{n}}\right]$$

The general case follows by a limiting argument (without the analysis for the case where equality holds).  $\Box$ 

Suppose that  $f : \mathbb{S}^1 \to \mathbb{R}$  is a mapping having a Lipshitz constant 1. Hence

$$||f(x) - f(y)|| \le ||x - y||_2$$

Let  $\mu$  be the median of f, so

$$\mu = \text{prob} \left[ \{ x \in \mathbb{S}^n : f(x) < \mu \} \right] = \frac{1}{2}$$

We assume that the probability distribution always admits such a  $\mu$  (at least approximately). The following inequality holds for every  $\epsilon > 0$  as a simple consequence of the isoperimetric inequality on the sphere.

$$\{\mathbf{x}\in\mathbb{S}^n: |\mathbf{f}-\boldsymbol{\mu}|>\epsilon\}<2e^{-\epsilon n/2}$$

For  $A \subseteq \mathbb{S}^n$  and for  $\epsilon > 0$ , let

$$A_{\epsilon} = \{ x \in \mathbb{S}^n : \text{dist} \, x, A < \epsilon \}$$

**Question 7.1.** Find a set  $A \subseteq \mathbb{S}^n$  with A = a for which  $A_{\epsilon}$  is the smallest.

The probability used here is the (normalized) Haar measure. The answer is always a spherical cap, and in particular if  $a = \frac{1}{2}$ , then the best A is the hemisphere (and so  $A_{\epsilon} = \{x \in \mathbb{S}^n : x_1 < \epsilon\}$ ). We will show that for  $A \subseteq \mathbb{S}^n$  with  $A = \frac{1}{2}$ ,  $A_{\epsilon} \ge 1 - 2e^{-\epsilon^2 n/4}$ . If A is the hemisphere, then  $A_{\epsilon} = 1 - \Theta(e^{-\epsilon^2 n/2})$ , and so the hemisphere is the best possible set.

But first, a small variation on BM :

$$\operatorname{vol}\left(\frac{A+B}{2}\right) \ge \sqrt{\operatorname{vol}\left(A\right) \cdot \operatorname{vol}\left(B\right)}$$

This follows from BM because

$$\operatorname{vol}\left(\frac{A+B}{2}\right)^{\frac{1}{n}} \ge \operatorname{vol}\left(\frac{A}{2}\right)^{\frac{1}{n}} + \operatorname{vol}\left(\frac{B}{2}\right)^{\frac{1}{n}}$$
$$= \frac{1}{2}\left[\operatorname{vol}\left(A\right)^{\frac{1}{n}} + \operatorname{vol}\left(B\right)^{\frac{1}{n}}\right]$$
$$\ge \sqrt{\operatorname{vol}\left(A\right)^{\frac{1}{n}} + \operatorname{vol}\left(B\right)^{\frac{1}{n}}}$$

For  $A \subseteq \mathbb{S}^n$ , let  $\tilde{A} = \{\lambda a : a \in A, 1 \ge \lambda \ge 0\}$ . Then  $A = \mu_{n+1}(\tilde{A})$ . Let  $B = \mathbb{S}^n \setminus A_{\epsilon}$ . Lemma 7.2. If  $\tilde{x} \in \tilde{A}$  and  $\tilde{y} \in \tilde{B}$ , then

$$\left|\frac{\tilde{x}+\tilde{y}}{2}\right| \le 1 - \frac{\epsilon^2}{8}$$

It follows that  $\frac{\tilde{A}+\tilde{B}}{2}$  is contained in a ball of radius at most  $1-\frac{\epsilon^2}{8}$ . Hence

$$\left(1 - \frac{\epsilon^2}{8}\right)^{n+1} \ge \operatorname{vol}\left(\frac{\tilde{A} + \tilde{B}}{2}\right)$$
$$\ge \sqrt{\operatorname{vol}\left(\tilde{A}\right) \cdot \operatorname{vol}\left(\tilde{B}\right)}$$
$$\ge \sqrt{\frac{\operatorname{vol}\left(\tilde{B}\right)}{2}}$$

Therefore,  $2e^{-\epsilon^2 n/4} \ge \operatorname{vol}\left(\tilde{B}\right)$ .

## 7.1 Boolean Influences

Let  $f : \{0,1\}^n \to \{0,1\}$  be a boolean function. For a set  $S \subseteq [n]$ , the influence of S on f,  $I_f(S)$  is defined as follows. When we pick  $\{x_i\}_{i \notin S}$  uniformly at random, three things can happen.

1. f = 0 regardless of  $\{x_i\}_{i \in S}$  (suppose that this happens with probability  $q_0$ ).

2. f = 1 regardless of  $\{x_i\}_{i \in S}$  (suppose that this happens with probability  $q_1$ ).

3. With probability  $\text{Inf}_{f}(S) := 1 - q_0 - q_1$ , f is still undetermined.

Some examples:

• (Dictatorship)  $f(x_1, x_2, \dots, x_n) = x_1$ . In this case

$$\operatorname{Inf}_{\mathsf{dictatorship}}\left(S\right) = \begin{cases} 1 & \text{ if } i \in S\\ 0 & \text{ if } i \notin S \end{cases}$$

• (Majority) For n = 2k + 1,  $f(x_1, x_2, ..., x_n)$  is 1 if and only if a majority of the  $x_i$  are 1. For example, if  $S = \{1\}$ ,

$$\operatorname{Inf}_{\mathsf{majority}}(\{1\}) = \operatorname{prob}(x_1 \text{ is the tie breaker })$$

$$=\frac{\binom{2k}{k}}{2^{2k}}=\Theta\left(\frac{1}{\sqrt{k}}\right)$$

For fairly small sets S,

$$\operatorname{Inf}_{\mathsf{majority}}(S) = \Theta\left(\frac{|S|}{\sqrt{n}}\right)$$

• (Parity)  $f(x_1, x_2, ..., x_n) = 1$  if and only if an even number of the  $x_i$ 's are 1. In this case

$$\operatorname{Inf}_{\mathsf{parity}}\left(\{x_i\}\right) = 1$$

for every  $1 \le i \le n$ .

Question 7.2. What is the smallest  $\delta = \delta(n)$  such that there exists a function  $f : \{0,1\}^n \to \{0,1\}$  which is balanced (i.e.,  $Ef = \frac{1}{2}$ ) for which  $\operatorname{Inf}_f(\{x_i\}) < \delta$  for all  $x_i$ ?

Consider the following example, called tribes. The set of inputs  $\{x_1, x_2, \ldots, x_n\}$  is partitioned into tribes of size *b* each. Here,  $f(x_1, x_2, \ldots, x_n) = 1$  if and only if there is a tribe that unanimously 1.



Since we want  $Ef = \frac{1}{2}$ , we must have  $\operatorname{prob}(f = 0) = \left(1 - \frac{1}{2^b}\right)^{\frac{n}{b}} = \frac{1}{2}$ . Therefore,  $\frac{n}{b}\ln\left(1 - \frac{1}{2^b}\right) = -\ln 2$ . We use the Taylor series expansion for  $\ln(1 - \epsilon) = -\epsilon - \epsilon^2/2 - \cdots = -\epsilon - O(\epsilon^2)$  to get  $\frac{n}{b}\left(\frac{1}{2^b} + O\left(\frac{1}{4^b}\right)\right) = -\ln 2$ . This yields  $n = b \cdot 2^b \ln 2 (1 + O(1))$ . Hence  $b = \log_2 n - \log_2 \ln n + \Theta(1)$ . Hence,

$$\begin{aligned} \text{Inf}_{\text{tribes}} \left( x \right) &= \left( 1 - \frac{1}{2^b} \right)^{\frac{n/b}{-}1} \cdot \left( \frac{1}{2} \right)^{b-1} \\ &= \frac{\left( 1 - \frac{1}{2^b} \right)^{\frac{n}{b}}}{1 - \frac{1}{2^b}} \cdot \frac{1}{2^{b-1}} \\ &= \frac{1}{1 - \frac{1}{2^b}} \cdot \frac{1}{2^b} \\ &= \frac{1}{2^{b-1}} = \Theta \left( \frac{\log b}{n} \right) \end{aligned}$$

In this example, each individual variable has influence  $\Theta(\log n/n)$ . It was later shown that this is lowest possible influence.

**Proposition 7.3.** If  $Ef = \frac{1}{2}$ , then  $\sum_{x} \text{Inf}_{f}(x) \ge 1$ .

This is a special case of the edge isoperimetric inequality for the cube, and the inequality is tight if f is dictatorship.



The variable x is influential in the cases indicated by the solid lines, and hence

$$\operatorname{Inf}_{f}(x) = \frac{\# \text{ of mixed edges}}{2^{n-1}}$$

Let  $S = f^{-1}(0)$ . Then  $\sum Inf_f(x) = \frac{1}{2^{n-1}}e(S, S^c)$ .

One can use  $\hat{f}$  to compute influences. For example, if f is monotone (so  $x \prec y \Rightarrow f(x) \leq f(y)$ ), then

$$\hat{f}(S) = \sum_{T} \frac{(-1)^{|S \cap T|}}{2^n}$$

Therefore,

$$\begin{split} \hat{f}(\{i\}) &= \frac{1}{2^n} \sum_{i \notin T} f(T) - \frac{1}{2^n} \sum_{i \in T} f(T) \\ &= \frac{1}{2^n} \sum_{i \notin T} \left( f(T) - f(T \cup \{i\}) \right) \\ &= \frac{-1}{2^n} \cdot \# \text{ mixed edges in the direction of } i \\ &= -\frac{1}{2} \text{Inf}_f(x_i) \end{split}$$

Hence  $\text{Inf}_{f}(x_{i}) = -2\hat{f}(\{i\})$ . What can be done to express  $\text{Inf}_{f}(x)$  for a general f? Define

$$f^{(i)}(z) = f(z) - f(z \oplus e_i)$$



The last term will be evaluated using Parseval. For this, we need to compute the Fourier expression of  $f^{(i)}$  (expressed in terms of  $\hat{f}$ ).

$$\begin{split} \widehat{f^{(i)}}(S) &= \frac{1}{2^n} \sum_T f^{(i)}(T)(-1)^{|S \cap T|} \\ &= \frac{1}{2^n} \sum_T \left[ f(T) - f(T \oplus \{i\}) \right] (-1)^{|S \cap T|} \\ &= \frac{1}{2^n} \sum_{i \notin T} \left( \left[ f(T) - f(T \cup \{i\}) \right] (-1)^{|S \cap T|} + \left[ f(T \cup \{i\}) - f(T) \right] (-1)^{|S \cap (T \cup \{i\})|} \right) \\ &= \frac{1}{2^n} \sum_{i \notin T} \left[ f(T) - f(T \cup \{i\}) \right] \left( (-1)^{|S \cap T|} - (-1)^{|S \cap (T \cup \{i\})|} \right) \\ &= \begin{cases} 0 & \text{if } i \notin S \\ 2\hat{f}(S) & \text{if } i \in S \end{cases} \end{split}$$

Using Parseval on  $\widehat{f^{(i)}}$  along with the fact that  $\widehat{f^{(i)}}$  takes on only values  $\{0, \pm 1\}$ , we conclude that

$$\operatorname{Inf}_{f}(x_{i}) = 4 \sum_{i \in S} |hatf(S)|^{2}$$

Next time, we will show that if  $Ef = \frac{1}{2}$ , then there exists a *i* such that  $\sum_{i \in S} \left( \hat{f}(S) \right)^2 > \Omega(\ln n/n)$ . Lemma 7.4. For every  $f : \{0,1\}^n \to \{0,1\}$ , there is a monotone  $g : \{0,1\}^n \to \{0,1\}$  such that

- Eg = Ef.
- For every  $s \subseteq [n]$ ,  $\operatorname{Inf}_{g}(S) \leq \operatorname{Inf}_{f}(S)$ .

Proof. We use a shifting argument.



Clearly  $E\tilde{f} = Ef$ . We will show that for all S,  $\mathrm{Inf}_{\tilde{f}}(S) \leq \mathrm{Inf}_{f}(S)$ . We may keep repeating the shifting step until we obtain a monotone function g. It is clear that the process will terminate by considering the progress measure  $\sum f(x) |x|$  which is strictly increasing. Therefore, we only need show that  $\mathrm{Inf}_{\tilde{f}}(()S) \leq \mathrm{Inf}_{f}(S)$ .