

# Finite Model Theory – Homework 3

April 27, 2018

## 1 The Satisfiability Problem

1. (0 points)

- (a) Let  $\sigma = \{U_1, \dots, U_m\}$  be a relational vocabulary with  $k$  unary predicate symbols. Prove that sentences over this vocabulary satisfy the small model property: if  $\varphi$  has a model, then it has a finite model of size  $\leq f(|\varphi|)$ , for some computable function  $f$ . (Hint: given any infinite structure  $\mathbf{A}$  and a number  $k$  describe a finite model  $\mathbf{B}$  s.t.  $\mathbf{A} \sim_k \mathbf{B}$ .)
- (b) Prove that the satisfiability problem for a relational vocabulary consisting of only unary predicate symbols is decidable.
- (c) Let  $\varphi(x)$  be formula with a free variable  $x$ , and  $R$  be a unary relational symbol. We say that  $\varphi$  is monotone in a relational symbol  $R$  if for any two structures  $\mathbf{A}, \mathbf{B}$  with the same domain and satisfying  $R^{\mathbf{A}} \subseteq R^{\mathbf{B}}$ , and  $S^{\mathbf{A}} = S^{\mathbf{B}}$  for every other relational symbol  $S$ , we have  $\{a \in A \mid \mathbf{A} \models \varphi(a)\} \subseteq \{b \in B \mid \mathbf{B} \models \varphi(b)\}$ . (Note: this is the semantic property needed for the least fixpoint,  $[\text{lfp}_{R,x}\varphi]$ .) Prove that, if the vocabulary includes at least one binary relational symbol other than  $R$ , then the problem “given  $\varphi$  check if it is monotone in  $R$  over all finite structures” is undecidable.

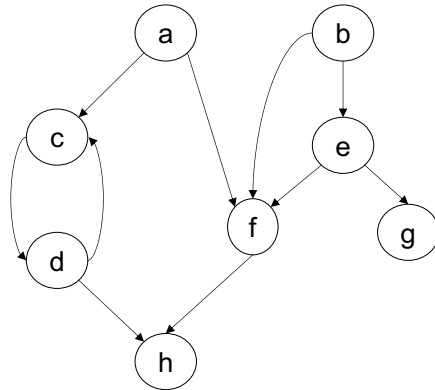
## 2 Descriptive Complexity

2. (0 points)

(a) Let  $G = (V, E)$  be a finite graph, and consider the following query:

$$q(x) = [\text{lfp}_{S,x}(\forall y(E(x, y) \rightarrow S(y)))](x)$$

i. Which nodes  $x$  does the query return on the graph below?



ii. Write an FO sentence (without fixpoints!) that is equivalent to  $\forall x \neg q(x)$ .

iii. Consider these complexity classes:  $AC^0$ ,  $P$ TIME,  $NP$ ,  $PSPACE$ . Indicate the lowest complexity class to which  $q$  belongs. You can just indicate the lowest complexity class, no need to prove that it's not lower than that (but you are welcome to do so).

(b) Consider the vocabulary  $(\prec, P_a, P_b, P_c)$  of strings over the alphabet  $\Sigma = \{a, b, c\}$ .

i. Write each of the regular expressions below in FO or in MSO. Use succ,  $\leq$ , min, max when needed, since these are expressible using  $\prec$ .

$$E_1 = (a|b)^*.c^*$$

$$E_2 = (a.b)^*$$

$$E_3 = (a.a.a)^*$$

ii. Write a regular expression describing the following language:

$$\forall S(\exists x(S(x) \wedge P_a(x))) \rightarrow (\exists y(S(y) \wedge P_b(y)))$$