## Digital Representation

INFO/CSE 100, Autumn 2004
Fluency in Information Technology
http://www.cs.washington.edu/100

## Readings and References

- Reading
» Fluency with Information Technology
- Chapter 8, Bits and the "Why" of Bytes


## Are two symbols enough?



We can represent each player's move this way, but what about representing the whole game?

Empty position: $\otimes$
use this set of symbols

- empty cell: $\otimes$
- player 1: $\times$
- player 2: O

- Now we can represent this game as one 9-digit string: $\mathrm{O} \otimes \otimes \boldsymbol{X} \mathbf{X} \mathrm{O}_{\otimes \otimes \otimes}$
- How many possible game states are there?
» $3 \times 3 \times 3 \times 3 \times 3 \times 3 \times 3 \times 3 \times 3=3^{9}=19683$
- Telephone Tone dialing
» decide how many different items of information you want to represent
- 16 keypad buttons (including rarely used A, B, C, D) » decide how many "digits" or positions you want to use
- 2 simultaneous tones
» decide on a set of symbols
- 8 different tones

|  | 1209hz | 1336 hz | 1477 hz | 1633 hz |
| :---: | :---: | :---: | :---: | :---: |
| 697 hz | 1 | 2 | 3 | A |
| 770 hz | 4 | 5 | 6 | B |
| 852 hz | 7 | 8 | 9 | C |
| 941 hz | $*$ | 0 | $\#$ | $D$ |

## Another encoding

use a different set of symbols

- empty cell: 0
- player 1: $\mathbf{1}$
- player 2: 2

| $\mathbf{2}$ | 0 | 0 |
| :--- | :--- | :--- |
| $\mathbf{1}$ | $\mathbf{1}$ | $\mathbf{2}$ |
| 0 | 0 | 0 |

- Now we can represent this game as one 9-digit number: 200112000
- How many possible game states are there?
» $3 \times 3 \times 3 \times 3 \times 3 \times 3 \times 3 \times 3 \times 3=39=19683$

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## Telephone Tones

Use this set of symbols

- tone 1 :
[697 hz], [770 hz], [850 hz], or [941 hz]
- tone 2 :
[1209 hz], [1336 hz], [1477 hz], or [1633 hz]

|  | 1209hz | 1336hz | 1477 hz | 1633 hz |
| :---: | :---: | :---: | :---: | :---: |
| 697hz | 1 | 2 | 3 | A |
| 770hz | 4 | 5 | 6 | B |
| 852hz | 7 | 8 | 9 | C |
| 944hz | $*$ | 0 | $\#$ | $D$ |

- Now we can represent each button as a 2-tone sound
- How many possible combinations of tones are there?
» $4 \times 4=4^{2}=16$



## Info in the Physical World

- Physical world:
» The most fundamental representation of information is presence/absence of a phenomenon
- matter, light, magnetism, flow, charge, ...

The PandA representation

- detect: "Is the phenomenon present?"
- set: make phenomenon present or absent

Any controllable phenomenon works: define it right


## Connect Physical/Logical

- The power of IT comes from the fact that physical and logical worlds can be connected
Present represents true / Absent represents false
- or maybe vice versa --


## Pavement Memory

false true false false false true true false true false true false false false

| 0 | 1 | 0 | 0 | 0 | 1 | 1 | 0 | 1 | 0 | 1 | 0 | 0 | 0 |
| :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- | :--- |

## Info in the Logical World

- Logical World:
» Information, reasoning, computation are formulated by true/false and logic
- All men are mortal
- Aristotle is a man
- Aristotle is mortal
- True and false can be the patterns for encoding information

$0 \quad 0 \quad 1 \quad 0$

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## Bits

- PandA is a binary representation because it uses 2 patterns
- The word "bit"
» is a contraction for "binary digit"
" represents a position in space/time capable of being set and detected in 2 patterns

Sherlock Holmes's Mystery of Silver Blaze -a popular example where "absent" gives information ... the dog didn't bark, that is the phenomenon wasn't detected

## Bytes

- A byte is eight bits treated as a unit
» Adopted by IBM in 1960s
» A standard measure ever since
» Bytes encode the Latin alphabet using ASCII -the American Standard Code for Information Interchange

01000110
01001001
01010100

| ASCII | $\begin{aligned} & 0 \\ & 0 \\ & 0 \\ & 0 \\ & \hline \end{aligned}$ | $\begin{array}{\|l\|} \hline 0 \\ 0 \\ 0 \\ 1 \\ \hline \end{array}$ | $\begin{aligned} & 0 \\ & 1 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 1 \end{aligned}$ | $\begin{aligned} & 1 \\ & 0 \\ & 0 \end{aligned}$ |  |  | $\begin{aligned} & 0 \\ & 1 \\ & 1 \\ & 0 \end{aligned}$ | $\begin{aligned} & 0 \\ & 1 \\ & 1 \\ & 1 \\ & \hline \end{aligned}$ |  |  | $\begin{aligned} & 1 \\ & 0 \\ & 0 \\ & 1 \end{aligned}$ | $\begin{array}{\|l\|} \hline 1 \\ 0 \\ 1 \\ 1 \\ 0 \end{array}$ | $\begin{aligned} & 1 \\ & 0 \\ & 1 \\ & 1 \end{aligned}$ |  | $\begin{aligned} & 1 \\ & 1 \\ & 0 \\ & 0 \end{aligned}$ | $\begin{array}{\|l\|} \hline 1 \\ 1 \\ 0 \\ 1 \\ \hline \end{array}$ | $\begin{aligned} & 1 \\ & 1 \\ & 1 \\ & 1 \\ & 0 \end{aligned}$ | 1 <br> 1 <br> 1 <br> 1 | 1 1 1 1 |
| :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: | :---: |
| 0000 | " | $\mathrm{s}_{4}$ | ${ }^{1}$ | $5_{5}$ | \& |  |  | \% | $\mathrm{E}_{1}$ | ${ }_{5}$ | s | ${ }_{T}$ | ${ }^{\circ}$ | ${ }_{T}$ |  | ${ }^{\text {F }}$ | ${ }^{\circ}$ | ${ }_{5}$ |  | ${ }_{1}$ |
| 0001 | ${ }^{\circ} \mathrm{L}$ | $\mathrm{O}_{1}$ | $\mathrm{O}_{2}$ | $0_{3}$ | $0_{4}$ | " |  | \% | $\mathrm{E}_{8}$ | - | ${ }^{\prime}$ | $\mathrm{in}^{\text {n }}$ | $5_{8}$ | $\varepsilon_{0}$ |  | ${ }_{5}$ | ${ }_{5}$ | $\mathrm{F}_{5}$ |  | $\mathrm{v}_{5}$ |
| 0010 |  | $!$ | " | \# | 8 | \% |  | ${ }_{8}$ |  |  | ( | ) | * | + |  |  |  |  |  | ' |
| 0011 | 0 | 1 | 2 | 3 | 4 | 5 |  | 6 | 7 | 8 | 8 | 9 |  | , |  | < | $=$ | $>$ |  | ? |
| 0100 | @ | A | B | C | D |  |  | F | G |  | H | I | J | K |  | L | M | N |  | 0 |
| 0101 | P | Q | R | S | T |  |  | V | W | X | X | Y | z | [ |  | 1 | ] | ${ }^{\circ}$ |  |  |
| 0110 |  | a | b | c | d | e |  | f | g |  | h | i | j | k |  | 1 | m | n |  | 0 |
| 0111 | p | q | r | $s$ | t | u |  | v | w |  | 8 | y | $z$ |  |  | 1 | \} | $\sim$ |  | ${ }^{\circ}$ |
| 1000 | \% | $E_{1}$ | : | $\mathrm{B}_{3}$ | ${ }_{5}$ | " |  | 5 | $\mathrm{E}_{5}$ | - | s | ${ }^{\prime}$ | ${ }^{5}$ | \% |  | ${ }^{\circ}$ | ${ }^{1}$ | 5 |  | 3 |
| 1001 | \% | $\mathrm{P}_{1}$ | $\stackrel{\circ}{\text { a }}$ | $s_{8}$ | $\circ$ | " |  | ${ }^{\text {sp }}$ | ${ }_{\text {F }}$ |  | = | $\because$ | ${ }^{\circ}{ }^{*}$ | ${ }^{\circ}$ |  | ${ }_{5}$ | ${ }^{\circ} \mathrm{s}$ | ${ }^{\circ}$ |  | \% |
| 1010 | \% | 1 | ¢ | E | 0 | \# |  | + | § |  |  | © | \% | " |  | $\checkmark$ | - | (1) |  |  |
| 1011 | - | $\pm$ | ${ }^{2}$ | = |  | म |  | $\pi$ | $\cdot$ |  |  | 1 | ${ }^{\circ}$ | B |  | 1/4 | 1/2 | \% |  | $\iota$ |
| 1100 | A | A | A | A | A |  |  | E | C |  |  | Ė | Ê | E |  | I | f | I |  | İ |
| 1101 | Đ | N | - | ó | ó |  |  | 0 | $\times$ |  |  | U̇ | Ú | 0 |  | U | Y | p |  | B |
| 1110 | à | á | à | å | à |  |  | ※ | 9 |  |  | é | ê |  |  | i | i | í |  | 1 |
| 1111 | ŏ | ni | ò | ó | ô |  |  | 0 | $\div$ |  |  | ù | ú | û |  | ü | ¢́ | p |  | y |

## Encoding Information

- Bits and bytes encode the information, but that's not all
» Tags encode format and some structure in word processors
» Tags encode format and some structure in HTML
» In the Oxford English Dictionary tags encode structure and some formatting


## OED Entry For Byte

[^0]- IT joins physical \& logical domains so physical devices do our logical work
» Symbols represent things 1-to-1
» Create symbols by grouping patterns
» PandA representation is fundamental
- presence and absence
» Bit, a place where 2 patterns set/detect
" ASCII is a byte encoding of Latin alphabet
» In addition to content, encode structure


[^0]:    byte (balt). Computers. [Arbitrary, prob. influenced by bit $s b .{ }^{4}$ and bite $s b$.] A group of eight consecutive bits operated on as a unit in a computer. 1964 Blaauw \& Brooks in IBM Systems Jrnl. III. 122 An 8-bit unit of information is fundamental to most of the formats [of the System/360]. A consecutive group of $n$ such units constitutes a field of length $n$. Fixed-length fields of length one, two, four, and eight are ermed bytes, halfwords, words, and double words respectively. 1964 IBM Jrnl. Res. \& Developm. VIII. $97 / 1$ When a byte of data appears from an I/O device, the CPU is seized, dumped, used and restored. 1967 P. A. Stark Digital Computer Programming xix. 351 The normal operations in fixed point are done on four bytes at a time. 1968 Dataweek 24 Jan. 1/1 Tape reading and writing is at from 34,160 to 192,000 bytes per second.
    <e><hg><hw>byte</hw> <pr><ph>balt</ph></pr></hg>. <la>Computers</la>. <etym>Arbitrary, prob influenced by $<\mathrm{xr}><\mathrm{x}>\mathrm{bit}</ \mathrm{x}></ \mathrm{xr}><\mathrm{ps}>\mathrm{n}$. $<\mathrm{hm}>4</ \mathrm{hm}></ \mathrm{ps}>$ and $<\mathrm{xr}><\mathrm{x}>\mathrm{bite}</ \mathrm{x}><\mathrm{ps}>\mathrm{n}$. $</ \mathrm{ps}>$ $</ \mathrm{xr}></$ etym $><$ s $4>$ A group of eight consecutive bits operated on as a unit in a computer.</s4> $<$ qp $><$ q $><$ qd $>1964<$ qd $><$ a $>$ Blaauw $<$ a $>$ \&amp. $<$ a $>$ Brooks $<$ a $><$ bib $>$ in $</$ bib $><w>$ IBM Systems Jrnl. $</ \mathrm{w}><$ lc $>$ III. $122</ \mathrm{lc}><\mathrm{ql}>$ An 8 -bit unit of information is fundamental to most of the formats ed $>$ of the System $/ 360</$ ed $>$. \&es.A consecutive group of $<\mathrm{i}>\mathrm{n}<\mathrm{i}\rangle$ such units constitutes a field of
    length $\langle\mathrm{i}>\mathrm{n}<\mathrm{i}\rangle$. \&es. Fixed-length fields of length one, two, four, and eight are termed bytes, halfwords, length $<\mathrm{i}>\mathrm{n}<\mathrm{i}>$. \&es. Fixed-length fields of length one, two, four, and eight are termed bytes, halfwords,
    words, and double words respectively. $<$ q $\downarrow><\mathrm{q}><\mathrm{q}><\mathrm{qd}>1964</ \mathrm{qd}><\mathrm{w}>$ IBM Jrnl. Res. \&amp. Developm. $</ \mathrm{w}><\mathrm{lc}>$ VIII. $97 / 1</ \mathrm{lc}><\mathrm{q} \mid>$ When a byte of data appears from an I/O device, the CPU seized, dumped, used and restored. $</ \mathrm{q} \mid></ \mathrm{q}><\mathrm{q}><\mathrm{qd}>1967</ \mathrm{qd}><\mathrm{a}>$ P. A. Stark $</ \mathrm{a}><\mathrm{w}>$ Digital Computer Programming $</ \mathrm{w}><\mathrm{lc}>x$ xix. $351</ \mathrm{lc}><\mathrm{q} \mid>$ The normal operations in fixed point are done on four bytes at a time. $</ \mathrm{q}\rangle></ \mathrm{q}\rangle<\mathrm{q}\rangle<\mathrm{qd}>1968</ \mathrm{qd}\rangle<\mathrm{w}>$ Dataweek $</ \mathrm{w}><\mathrm{lc}>24$ Jan. $1 / 1</ \mathrm{lc}><\mathrm{q}\rangle$ Tape reading and writing is at from 34,160 to 192,000 bytes per second. $</ \mathrm{q} \mid></ \mathrm{q}></ \mathrm{qp}></ \mathrm{e}>$

