

Algorithmic Complexity I

CSE 120 Spring 2017

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Administrivia

- ❖ Assignments:
 - Binary Practice (4/21)
 - Creativity Assignment (4/24)

- ❖ Midterm in class on Wednesday, 4/26
 - 1 sheet of notes (2-sided, letter, handwritten)
 - Fill-in-the-blank(s), short answer questions, maybe simple drawing
 - Questions will cover lectures, assignments, and readings
 - Midterm Review sheet will be released tonight (4/19), will be covered in lab next week (4/25)

Outline

- ❖ **Algorithm Analysis: The Basics**
- ❖ Comparing Algorithms
- ❖ Orders of Growth

Algorithm Correctness

- ❖ An algorithm is considered **correct** if for every input, it reports the correct output and doesn't run forever or cause an error
- ❖ Incorrect algorithms may run forever, crash, or not return the correct answer
 - But they could still be useful!
 - *e.g.* an approximation algorithm
- ❖ Showing correctness
 - Mathematical proofs for algorithms
 - Empirical verification of implementations

Algorithm Analysis

- ❖ One commonly used criterion for analyzing algorithms is **computation time**
 - How long does the algorithm take to run and finish its task?
 - Can be used to compare different algorithms for the same computational problem
- ❖ How to measure this time?
 - Counting in my head
 - Stopwatch
 - Within your program

Aside: Computation Time

- ❖ Computers take time to complete their tasks
 - Under the hood, it's sort of like a bunch of buckets of water filling up – you have to wait for water to reach the top of a bucket for a single computation to complete
 - Buckets take about a billionth of a second to fill (~ 1 nanosecond)
 - There are billions of them on a single chip!
- ❖ A CPU can generally only execute one instruction at a time

Timing in Processing

- ❖ The function `millis()` returns the number of milliseconds since starting your program (as an `int`)
 - To start timing, call and store the value in a variable
 - Call again after your computation and subtract the values

```
13 void draw() {  
14     int time = millis();  
15     someComputation();  
16     println("Took " + (millis()-time) + " milliseconds to compute.");  
17     noLoop();  
18 }
```

Outline

- ❖ Algorithm Analysis: The Basics
- ❖ **Comparing Algorithms**
- ❖ Orders of Growth

Algorithm: Searching A Sorted Array

- ❖ Input: Numbers in a sorted array, desired number
- ❖ Output: If desired number is in the array (**true/false**)

- ❖ Algorithm 1:
 - Check each index starting from 0 for desired number
 - If equal, then report **true**
 - If not equal, then move to next index
 - If at end of array, then report **false**
 - Called **Linear Search** (also works for unsorted array)

Algorithm: Searching A Sorted Array

- ❖ Input: Numbers in a sorted array, desired number
- ❖ Output: If desired number is in the array (**true/false**)

- ❖ Algorithm 2:
 - Check “middle” index for desired number
 - If equal, then report **true**
 - If less than desired number, check *halfway* forwards next
 - If greater than desired number, check halfway backwards next
 - If no halfway point left, then report **false**
 - Called **Binary Search**
 - <http://www.cs.armstrong.edu/liang/animation/web/BinarySearch.html>

Peer Instruction Question

- ❖ On average, which algorithm would take less time to complete a search?
 - Vote at <http://PollEv.com/justinh>

- A. **Algorithm 1 (Linear Search)**
- B. **Algorithm 2 (Binary Search)**
- C. **They'd take about the same amount of time**

Measuring Linear Search

- ❖ Let's time Linear Search:

```
16 void draw() {  
17     int n = 3;  
18     println("Is " + n + " in intArr?");  
19     int time = millis();  
20     println(linearSearch(n));  
21     println("Took " + (millis()-time) + " milliseconds to compute.");  
22     noLoop();  
23 }
```

- ❖ One issue: our algorithm seems to be too fast to measure!
 - How can we fix this?

Best Case vs. Worst Case vs. Average Case

- ❖ We were measuring close to the best case!
 - Didn't matter how long our array was
- ❖ Could measure average case instead
 - Run many times on random numbers and average results
- ❖ Instead, we'll do worst case analysis. Why?
 - Nice to know the most time we'd ever spend
 - Worst case happens often
 - Average case is often similar to worst case



Example of Worst Case in Action

- ❖ Many web servers out there run something called “The Apache HTTP Server” (or just Apache for short)
 - When a user enters a particular URL, Apache delivers the correct files from the server to the person on the internet
 - An old version of Apache had a bug where if you entered a URL with tons of consecutive slashes, it could take *hours* to complete the request
 - e.g. <http://someurl.com//////////>
- ❖ Bottom line: an algorithm is often judged by its worst case behavior

What is the Worst Case?

- ❖ Discuss with your neighbor (no voting):
 - Assume `intArr.length` is 1000000 and `intArr[i] = i`;
 - What is a worst case argument for `num` for Linear Search?
 - What is a worst case argument for `num` for Binary Search?

- A. 1
- B. 500000
- C. 1000000
- D. 1000001
- E. Something else

```
25 boolean linearSearch(int num) {
26     for(int i = 0; i < intArr.length; i = i + 1) {
27         if(intArr[i] == num) {
28             return true;
29         }
30     }
31     return false;
32 }
```

Timing Experiments

- ❖ Let's try running Linear Search on a worst case argument value
 - Results:

- ❖ Now let's run Binary Search on a worst case argument value
 - Results:

Runtime Intuition

- ❖ Does it seem reasonable that the runtimes were inconsistent?

- ❖ Some reasons:
 - Your computer isn't just running Processing – there's a lot of other stuff running (*e.g.* operating system, web browser)
 - The computer hardware does lots of fancy stuff to avoid slowdown due to physical limitations
 - These may not work as well each execution based on other stuff going on in your computer at the time

Empirical Analysis Conclusion

- ❖ We've shown that Binary Search is seemingly much faster than Linear Search
 - Similar to having two sprinters race each other
- ❖ Limitations:
 - Different computers may have different runtimes
 - Same computer may have different runtime on same input
 - Need to implement the algorithm in order to run it
- ❖ Goal: come up with a “universal algorithmic classifier”
 - Analogous to coming up with a metric to compare all athletes (or fighters)

Outline

- ❖ Algorithm Analysis: The Basics
- ❖ Comparing Algorithms
- ❖ **Orders of Growth**

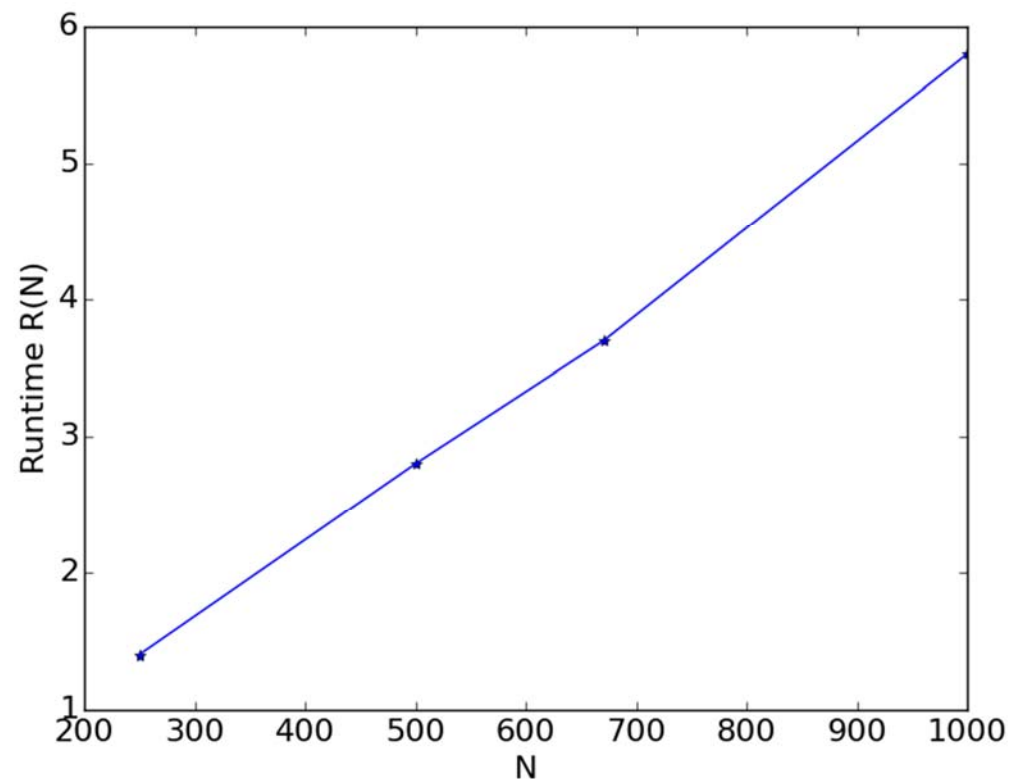
Characterizing Algorithms

- ❖ The method computer scientists use is roughly:
 - 1) Measure the algorithm's runtime on many different input sizes N (e.g. arrays of length 100, 200, 400, 800, ...)
 - To avoid runtime issues, can also count number of “steps” involved
 - 2) Make a plot of the runtime as a function of N , which we'll call $R(N)$
 - 3) Determine the general *shape* of $R(N)$
 - Does $R(N)$ look like N (linear), N^2 (quadratic), N^3 (cubic), $\log N$ (logarithmic), etc.

Linear Search

- ❖ As the name implies, Linear Search is linear
 - If you double N , then $R(N)$ should roughly double

N (input size)	$R(N)$ (time)
250 items	1.4 sec
500 items	2.8 sec
671 items	3.8 sec
1000 items	5.7 sec



Peer Instruction Question

- ❖ Algorithm for: do any pairs in array sum to zero?
- ❖ Which function does $R(N)$ look like?
 - Vote at <http://PollEv.com/justinh>

A. $\text{sqrt}(N)$

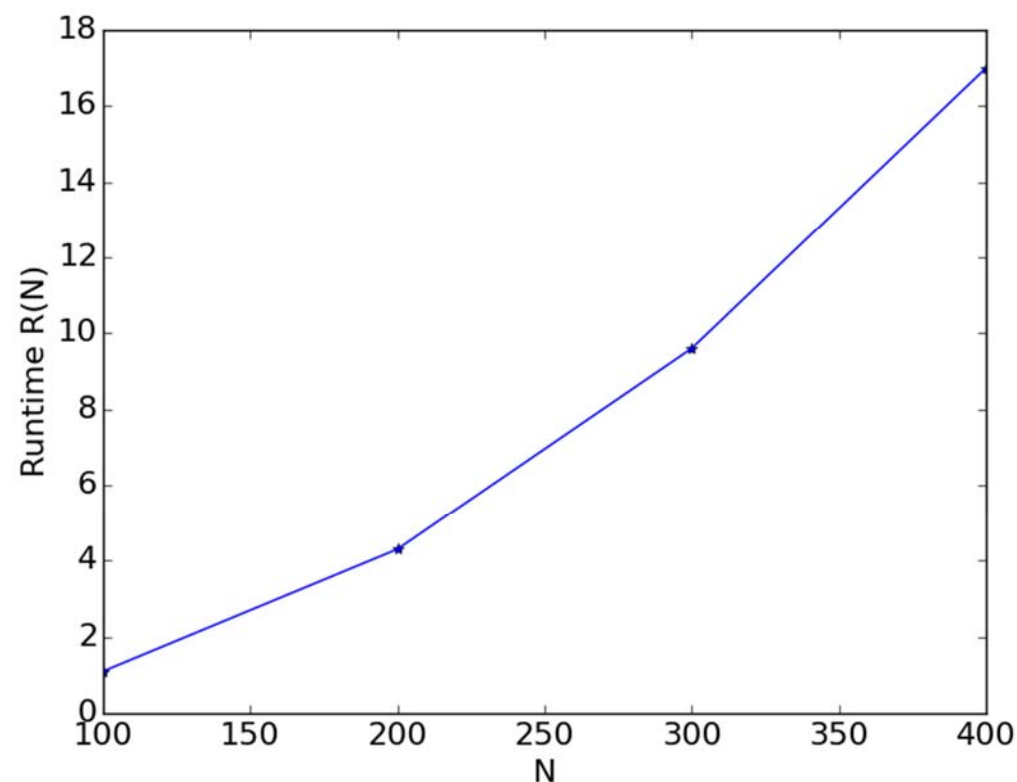
B. $\log(N)$

C. N

D. N^2

E. 2^N

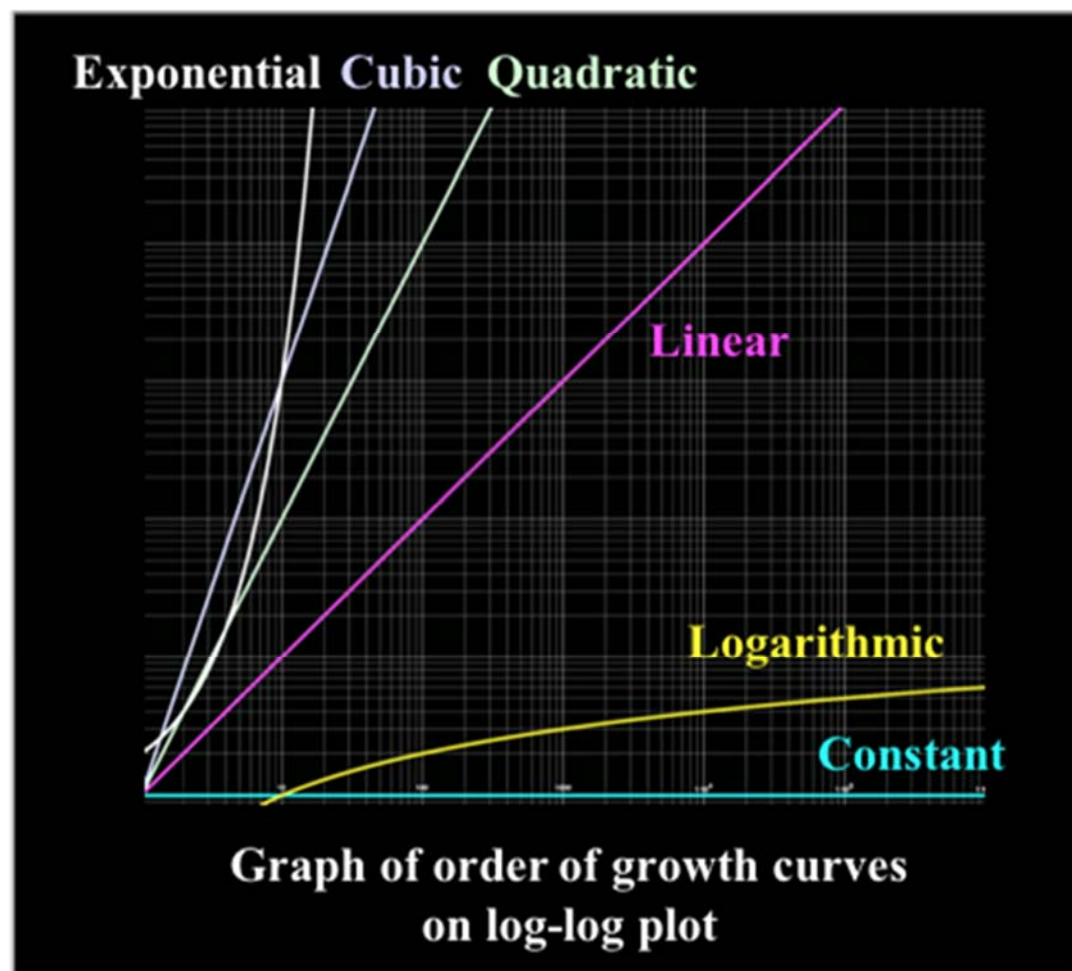
N (input size)	R(N) (time)
100 items	1.1 seconds
200 items	4.3 seconds
300 items	9.6 seconds
400 items	17.0 seconds



Orders of Growth

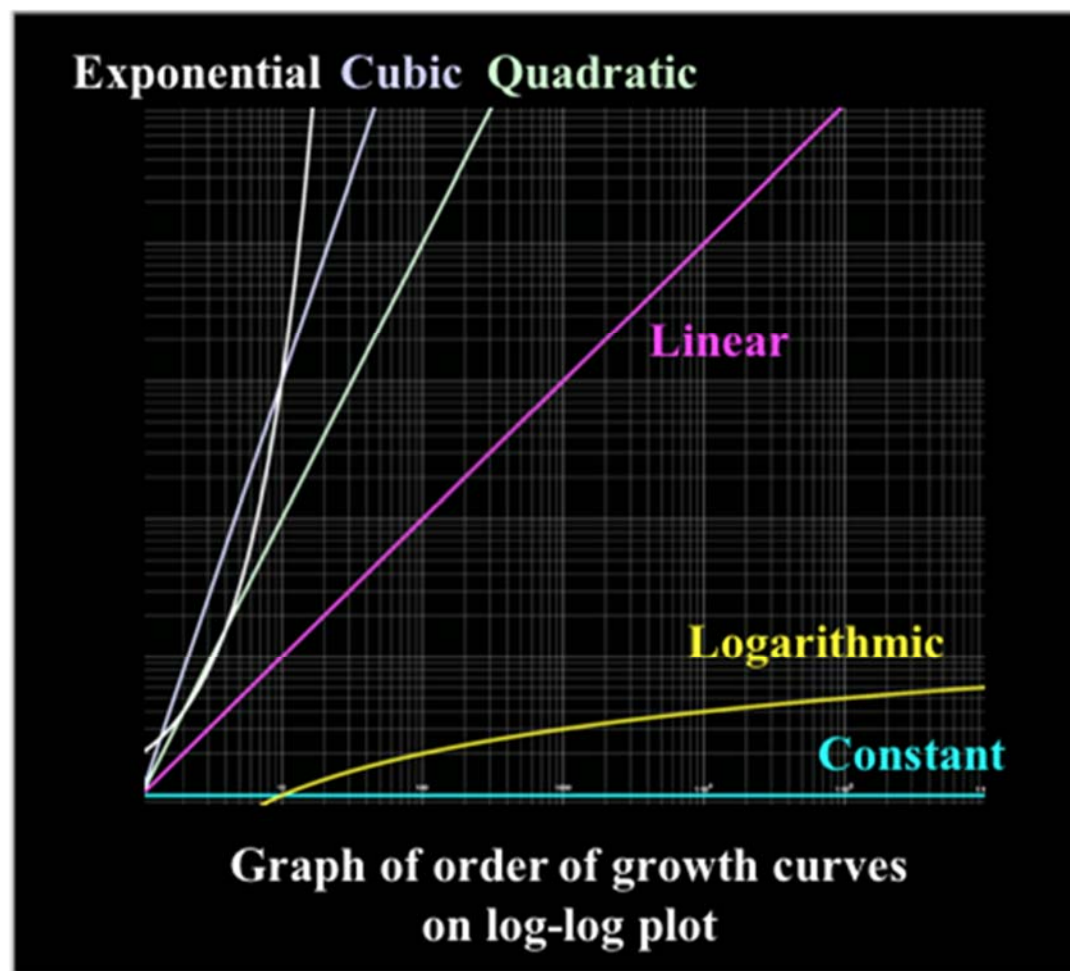
❖ The **order of growth** of $R(N)$ is its general shape:

- Constant 1
- Logarithmic $\log N$
- Linear N
- Quadratic N^2
- Cubic N^3
- Exponential 2^N
- Factorial $N!$



Orders of Growth

- ❖ The **order of growth** of $R(N)$ is its general shape:
 - Use *dominant* term
 - e.g. $10N^2 + 4 \log N$ is quadratic



Binary Search

- ❖ What order of growth is Binary Search?
 - Analyze using number of “steps” in worst case

N (input size)	Indices to Check
1 items	
2 items	
4 items	
8 items	
16 items	

Which is Faster?

- ❖ Suppose we have two algorithms: one is linear in N and the other is quadratic in N
 - No voting

- ❖ Which is faster?
 - A. **Linear Algorithm**
 - B. **Quadratic Algorithm**
 - C. **It depends**

The Reason Order of Growth Matters

- ❖ Roughly speaking, we care about really big N in real world applications
 - *e.g.* For Facebook, N (users) is ~ 1 billion
 - Want to generate list of suggested friends? Better be a fast algorithm as a function of N
- ❖ Order of growth is just a rough rule of thumb
 - There are limited cases where an algorithm with a worse order of growth can actually be faster
 - In almost all cases, order of growth works very well as a representation of an algorithm's speed

Orders of Growth Comparison

- ❖ The numbers below are rough estimates for a “typical” algorithm on a “typical” computer – provides a qualitative difference between the orders of growth

	Linearithmic						
	Linear		Quadratic	Cubic	Exponential	Exponential	Factorial
	n	$n \log_2 n$	n^2	n^3	1.5^n	2^n	$n!$
$n = 10$	< 1 sec	< 1 sec	< 1 sec	< 1 sec	< 1 sec	< 1 sec	4 sec
$n = 30$	< 1 sec	< 1 sec	< 1 sec	< 1 sec	< 1 sec	18 min	10^{25} years
$n = 50$	< 1 sec	< 1 sec	< 1 sec	< 1 sec	11 min	36 years	very long
$n = 100$	< 1 sec	< 1 sec	< 1 sec	1 sec	12,892 years	10^{17} years	very long
$n = 1,000$	< 1 sec	< 1 sec	1 sec	18 min	very long	very long	very long
$n = 10,000$	< 1 sec	< 1 sec	2 min	12 days	very long	very long	very long
$n = 100,000$	< 1 sec	2 sec	3 hours	32 years	very long	very long	very long
$n = 1,000,000$	1 sec	20 sec	12 days	31,710 years	very long	very long	very long

Table 2.1 The running times (rounded up) of different algorithms on inputs of increasing size, for a processor performing a million high-level instructions per second. In cases where the running time exceeds 10^{25} years, we simply record the algorithm as taking a very long time.