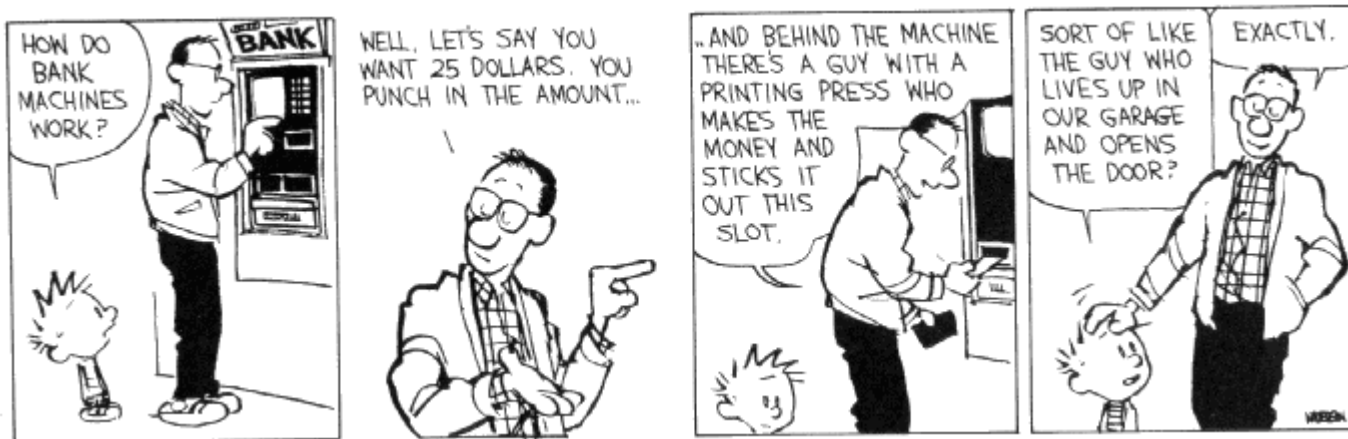


# CSE 311: Foundations of Computing

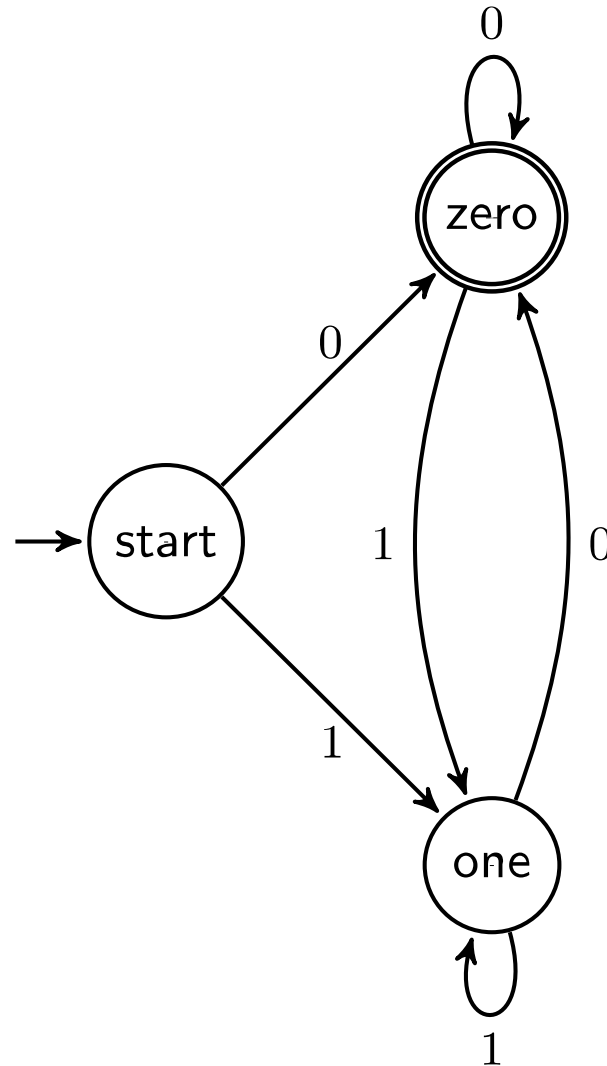
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## Lecture 22: Finite State Machines



# Last class: Strings this machine says are OK?

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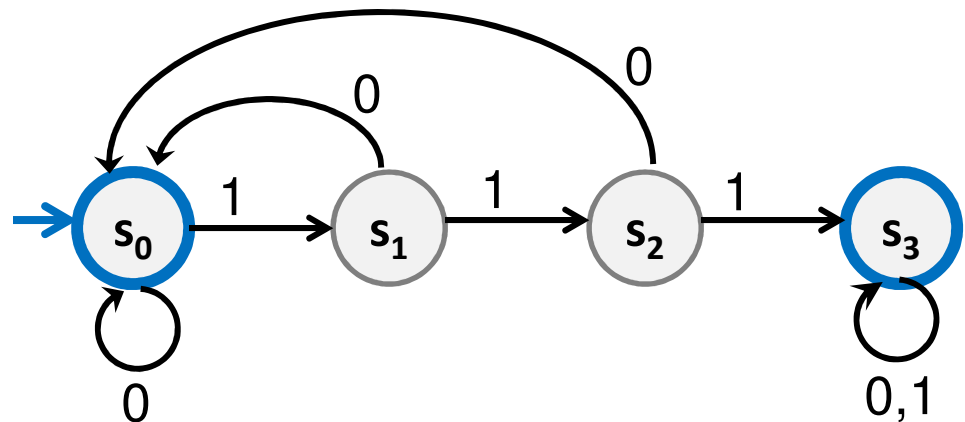
The set of all binary strings that end in 0

# Finite State Machines

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- States
- Transitions on input symbols
- Start state and final states
- The “language recognized” by the machine is the set of strings that reach a final state from the start

Old State	0	1
$s_0$	$s_0$	$s_1$
$s_1$	$s_0$	$s_2$
$s_2$	$s_0$	$s_3$
$s_3$	$s_3$	$s_3$

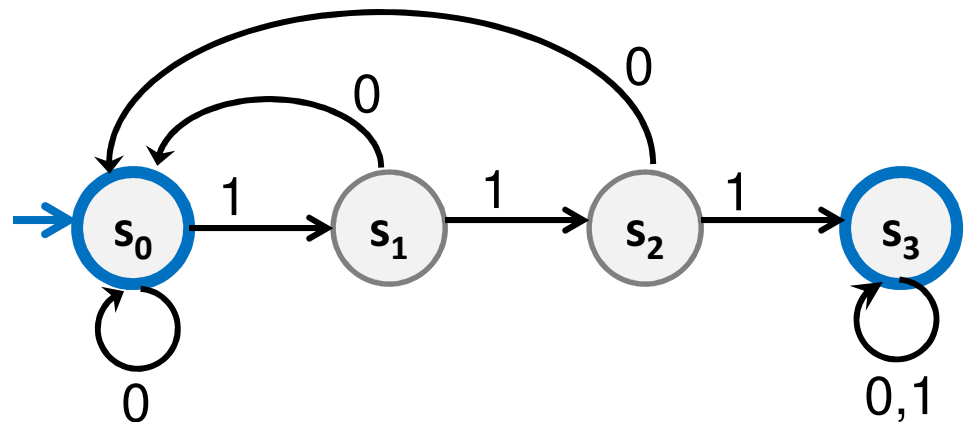


# Finite State Machines

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- Each machine designed for strings over some fixed alphabet  $\Sigma$ .
- Must have a transition defined from each state for every symbol in  $\Sigma$ .

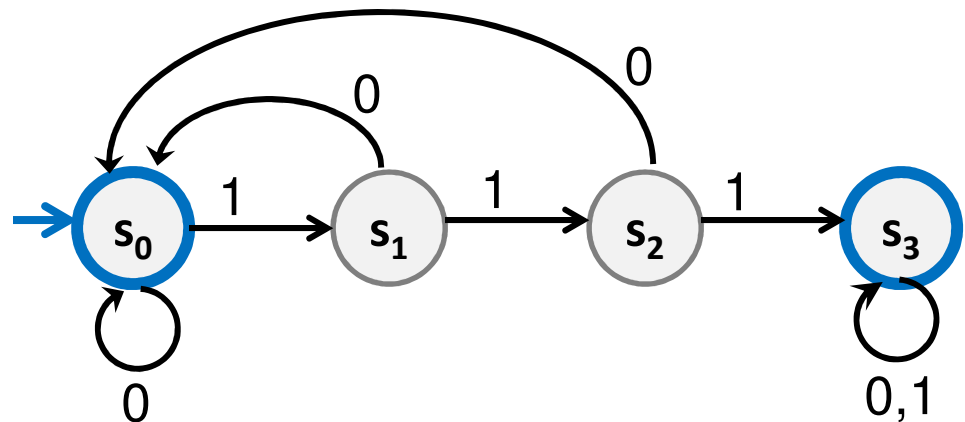
Old State	0	1
$s_0$	$s_0$	$s_1$
$s_1$	$s_0$	$s_2$
$s_2$	$s_0$	$s_3$
$s_3$	$s_3$	$s_3$



# What language does this machine recognize?

---

Old State	0	1
$s_0$	$s_0$	$s_1$
$s_1$	$s_0$	$s_2$
$s_2$	$s_0$	$s_3$
$s_3$	$s_3$	$s_3$

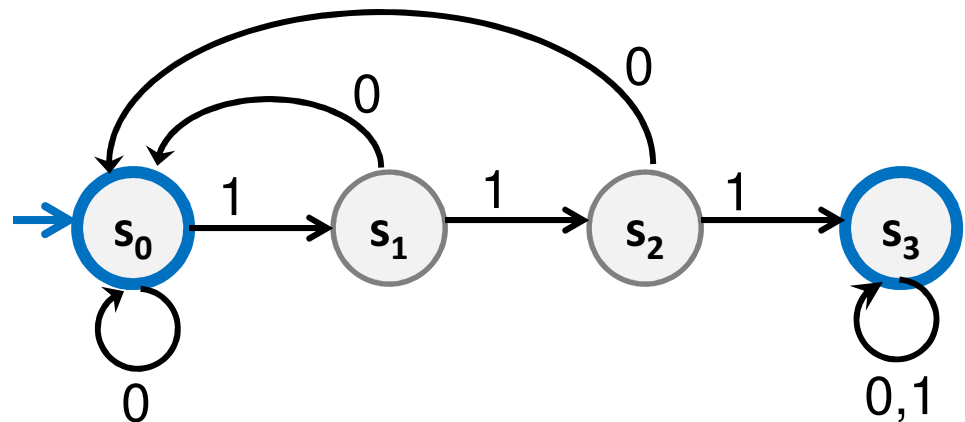


# What language does this machine recognize?

---

The set of all binary strings that contain **111** or don't end in **1**

Old State	0	1
$s_0$	$s_0$	$s_1$
$s_1$	$s_0$	$s_2$
$s_2$	$s_0$	$s_3$
$s_3$	$s_3$	$s_3$



# Strings over $\{0, 1, 2\}$

---

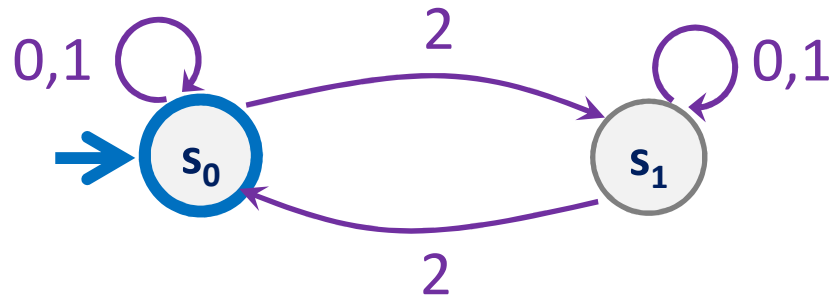
$M_1$ : Strings with an even number of 2's



# Strings over $\{0, 1, 2\}$

---

$M_1$ : Strings with an even number of 2's





# State Machine Design Recipe

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Given a language, how do you design a state machine for it?

**Create states to remember enough**

(about the portion of the input string that it has already seen)

**to correctly answer “accept/reject”** on the whole string after seeing the rest.

**Add labeled edges to show how the memory (state) should be updated for each new symbol.**

## Strings over $\{0, 1, 2\}$

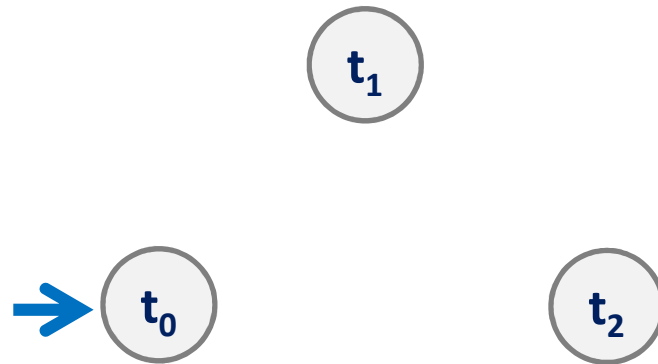
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$M_2$ : Strings where the sum of digits mod 3 is 0

# Strings over $\{0, 1, 2\}$

---

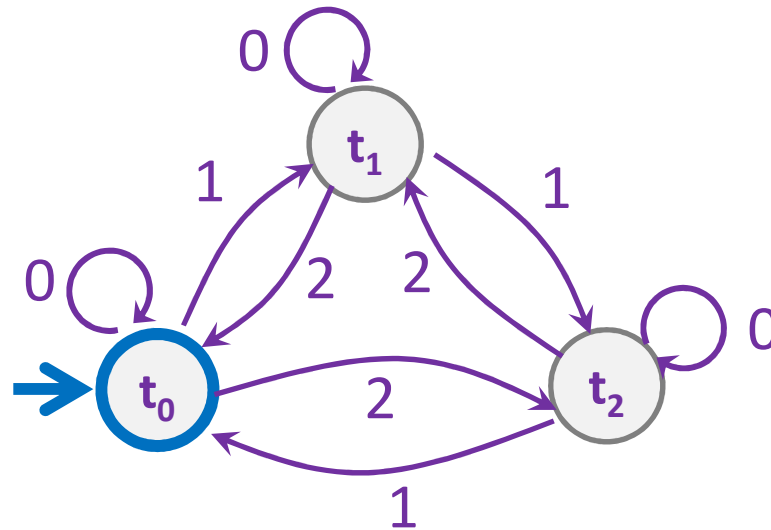
$M_2$ : Strings where the sum of digits mod 3 is 0



# Strings over $\{0, 1, 2\}$

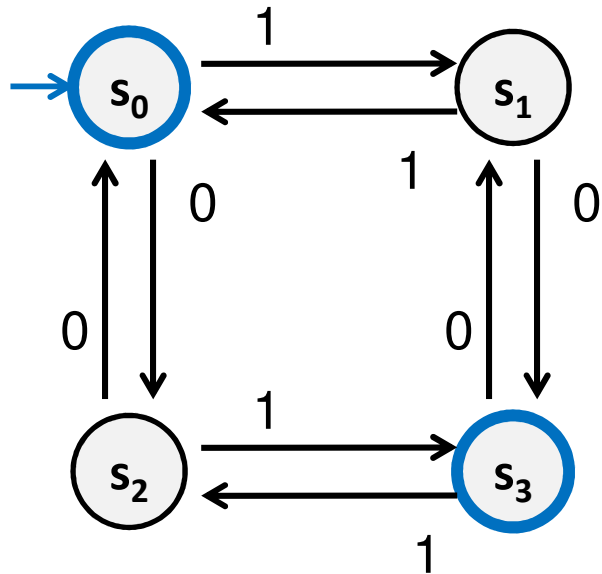
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$M_2$ : Strings where the sum of digits mod 3 is 0



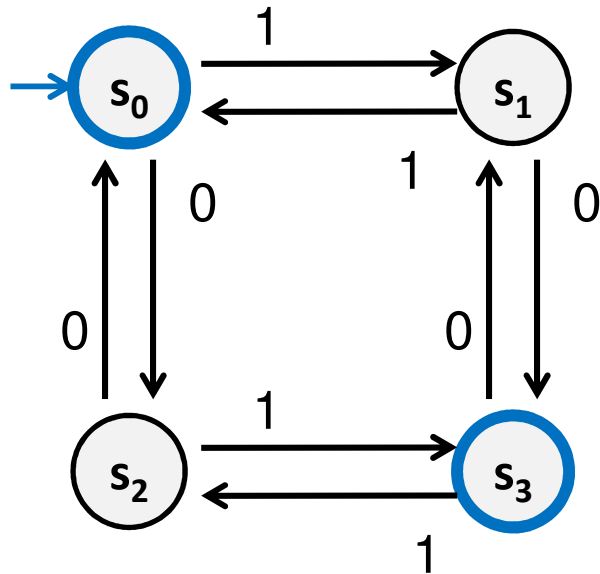
# What language does this machine recognize?

---



# What language does this machine recognize?

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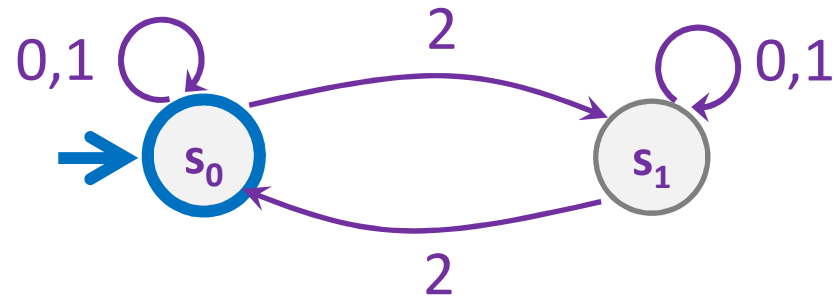
The set of all binary strings with # of 1's  $\equiv$  # of 0's (mod 2)  
(both are even or both are odd).

Can you think of a simpler description?

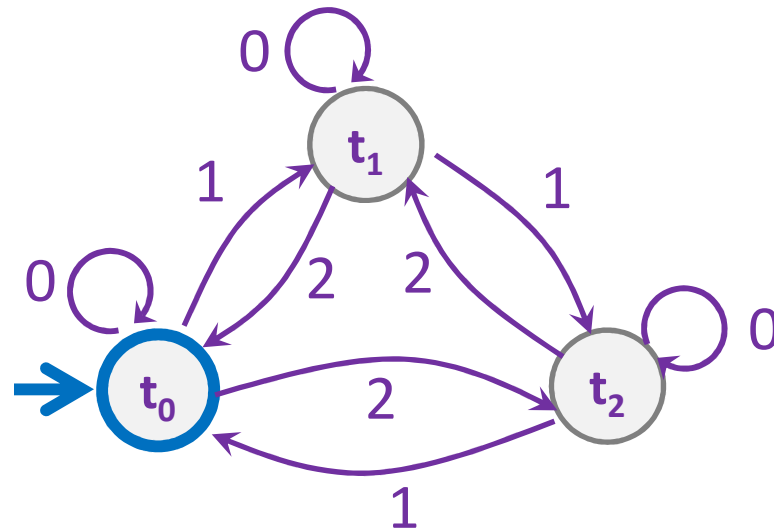
# Strings over $\{0, 1, 2\}$

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$M_1$ : Strings with an even number of 2's

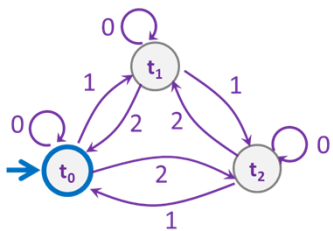
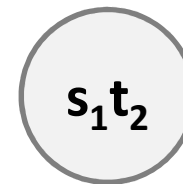
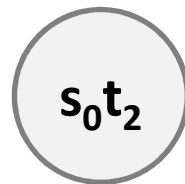
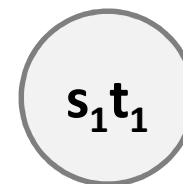
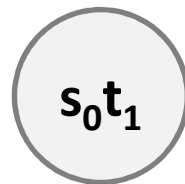
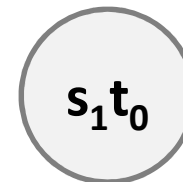
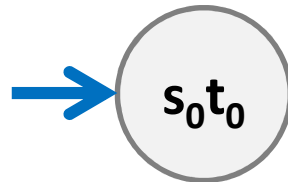
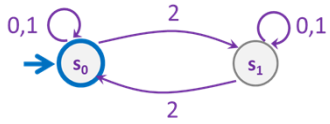


$M_2$ : Strings where the sum of digits mod 3 is 0



# Strings over $\{0,1,2\}$ w/ even number of 2's and mod 3 sum 0

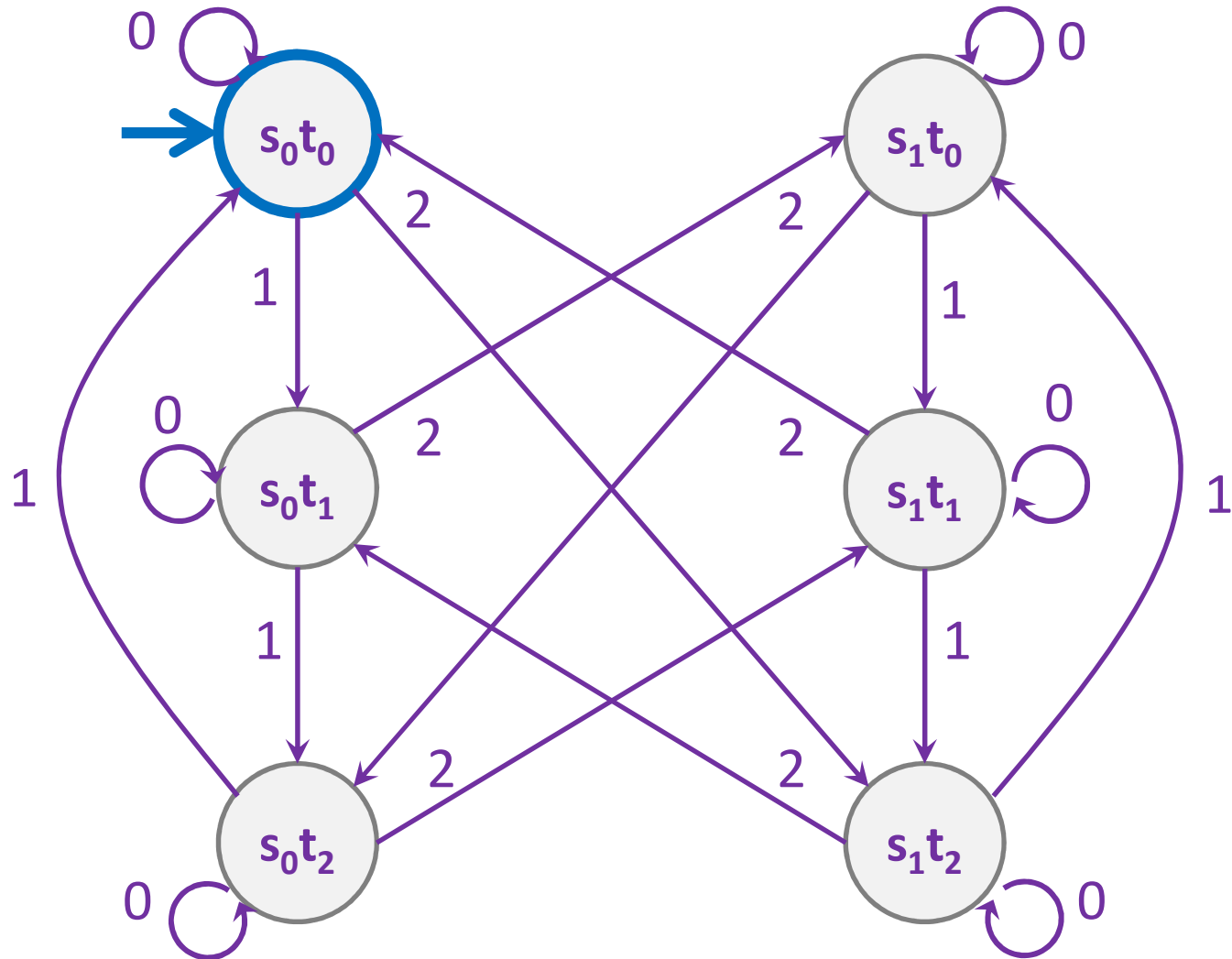
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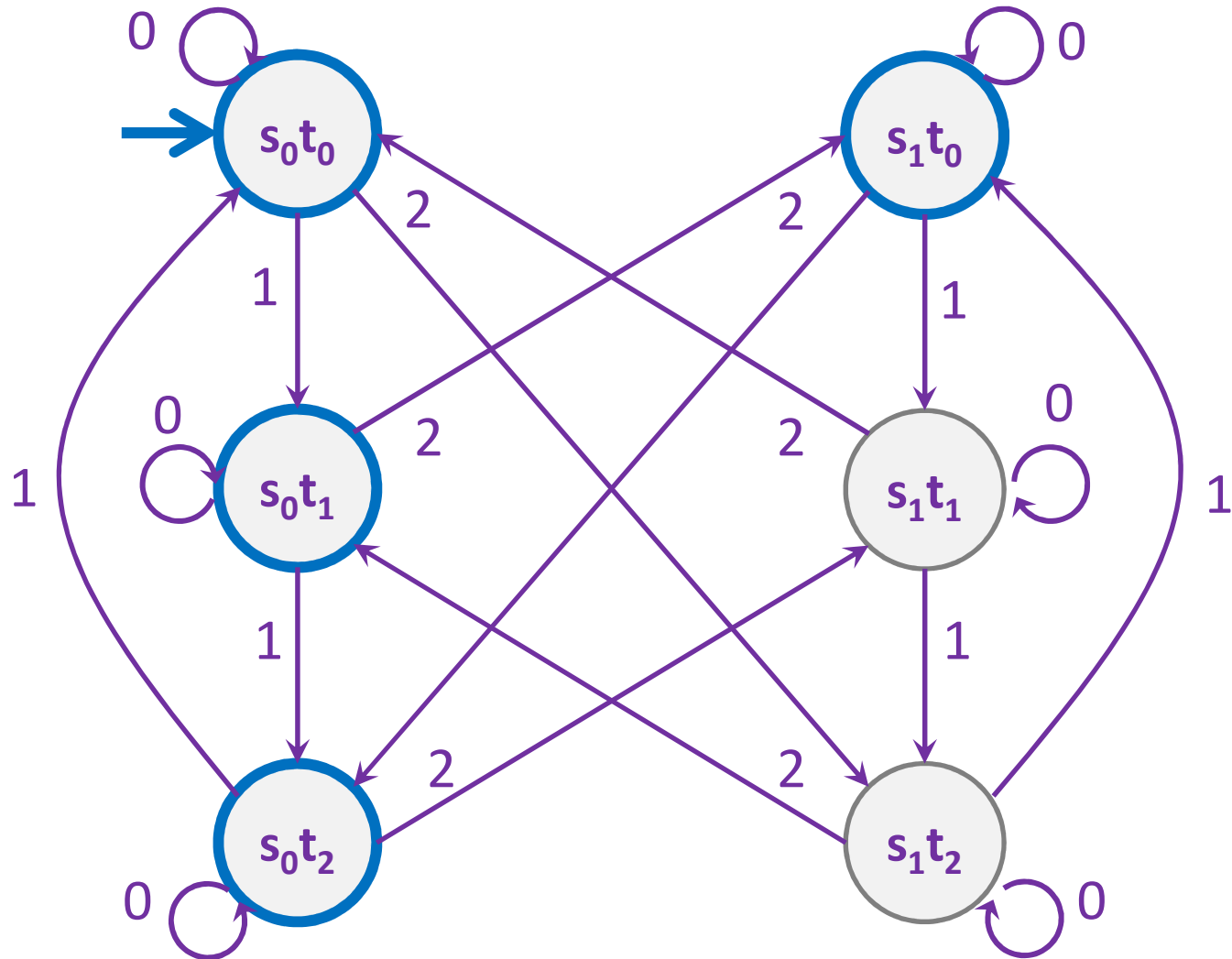
# Strings over $\{0,1,2\}$ w/ even number of 2's and mod 3 sum 0

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Strings over  $\{0,1,2\}$  w/ even number of 2's **OR** mod 3 sum 0

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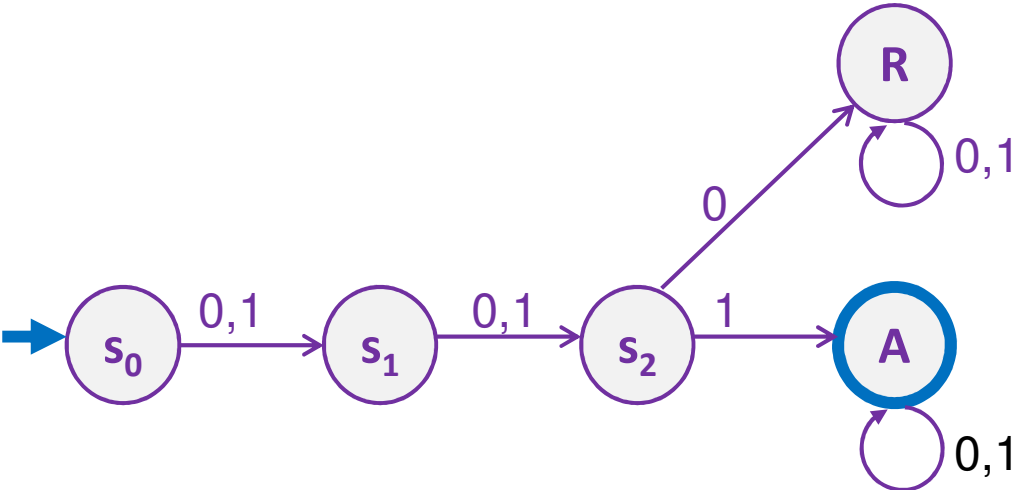


The set of binary strings with a **1** in the **3<sup>rd</sup>** position from the start

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The set of binary strings with a 1 in the 3<sup>rd</sup> position from the start

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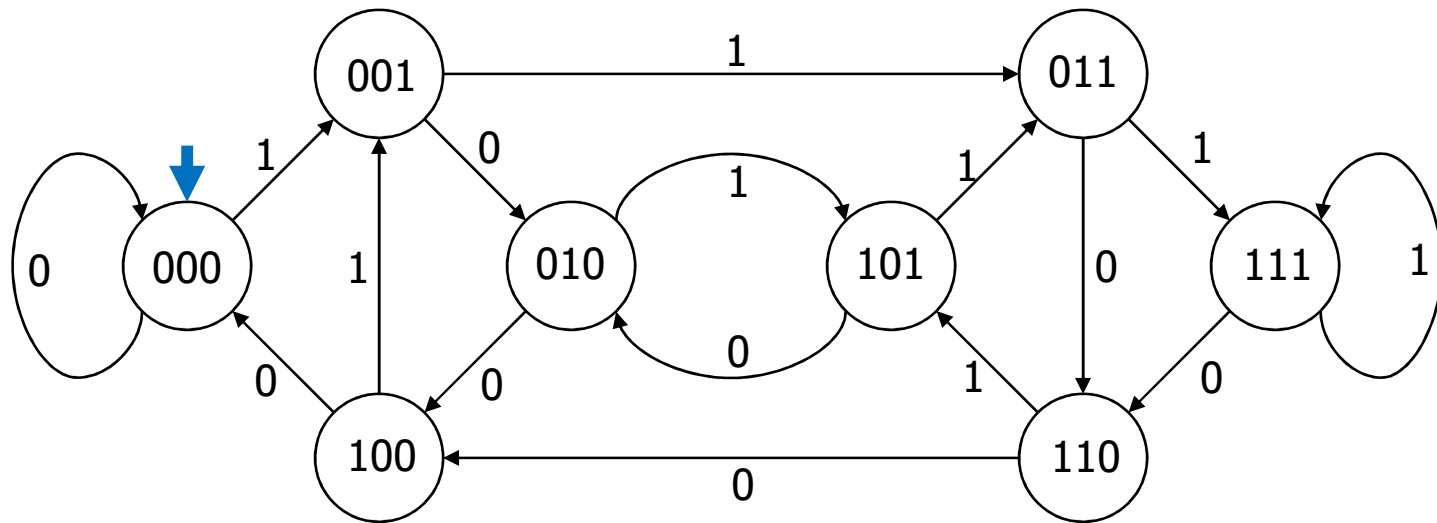


The set of binary strings with a **1** in the **3<sup>rd</sup>** position from the end

---

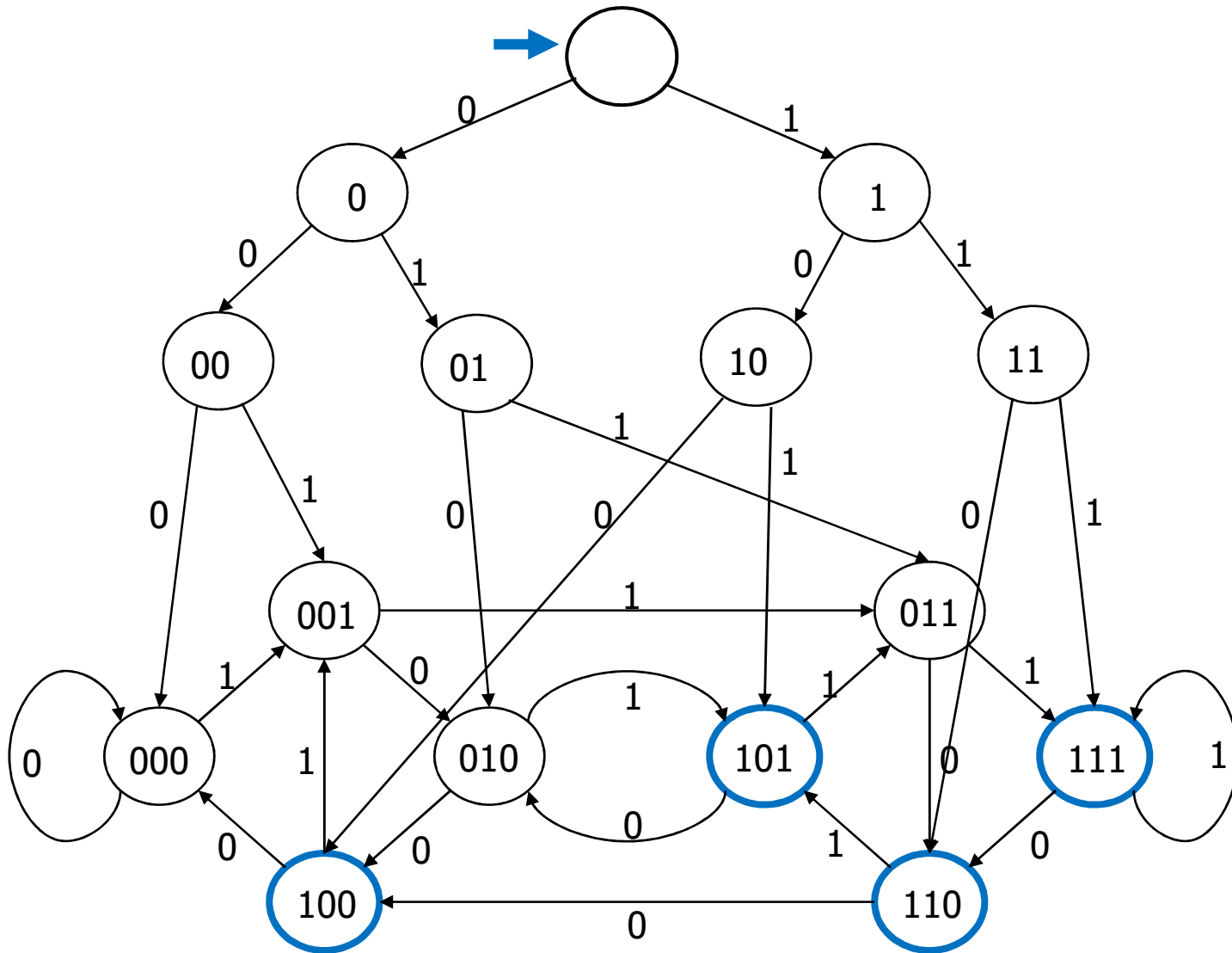
# 3 bit shift register “Remember the last three bits”

---



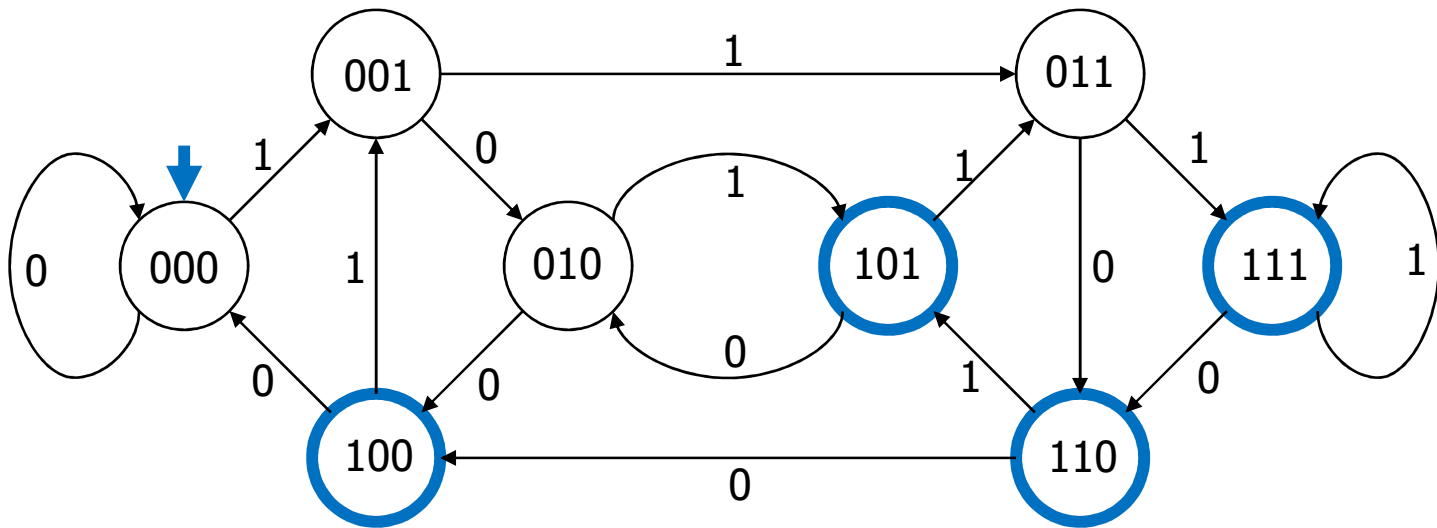
The set of binary strings with a 1 in the 3<sup>rd</sup> position from the end

---



The set of binary strings with a 1 in the 3<sup>rd</sup> position from the end

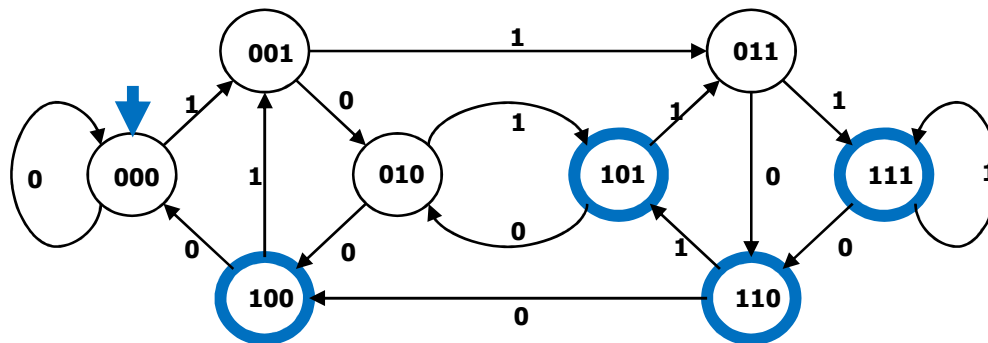
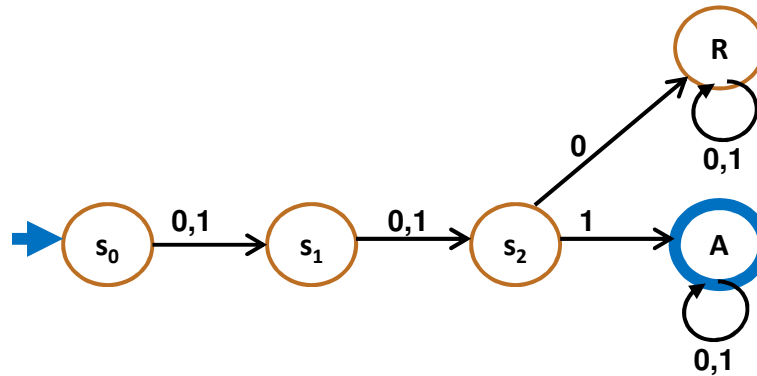
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# The beginning versus the end

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# Adding Output to Finite State Machines

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- **So far, we have considered finite state machines that just accept/reject strings**
  - called “Deterministic Finite Automata” or DFAs
- **Now we consider finite state machines *with output***
  - These are the kinds used as controllers



# Vending Machine

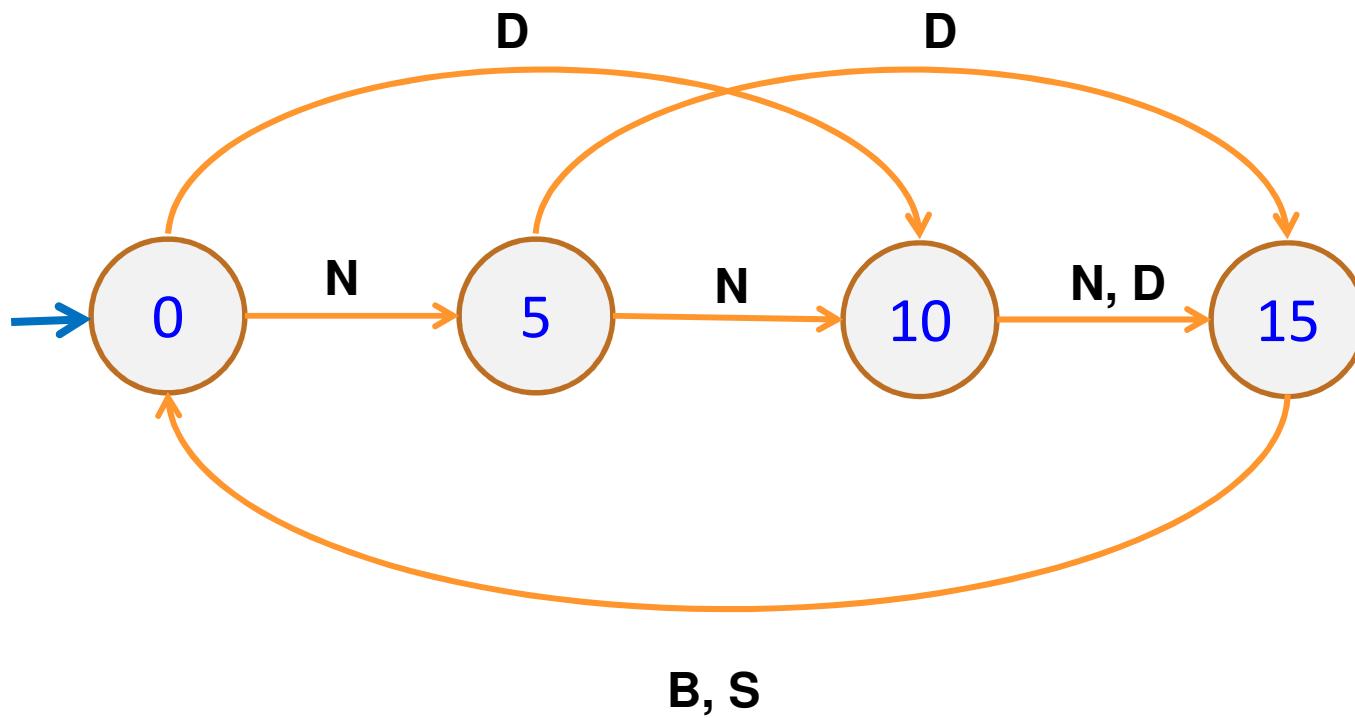


Enter 15 cents in dimes or nickels  
Press S or B for a candy bar



# Vending Machine, v0.1

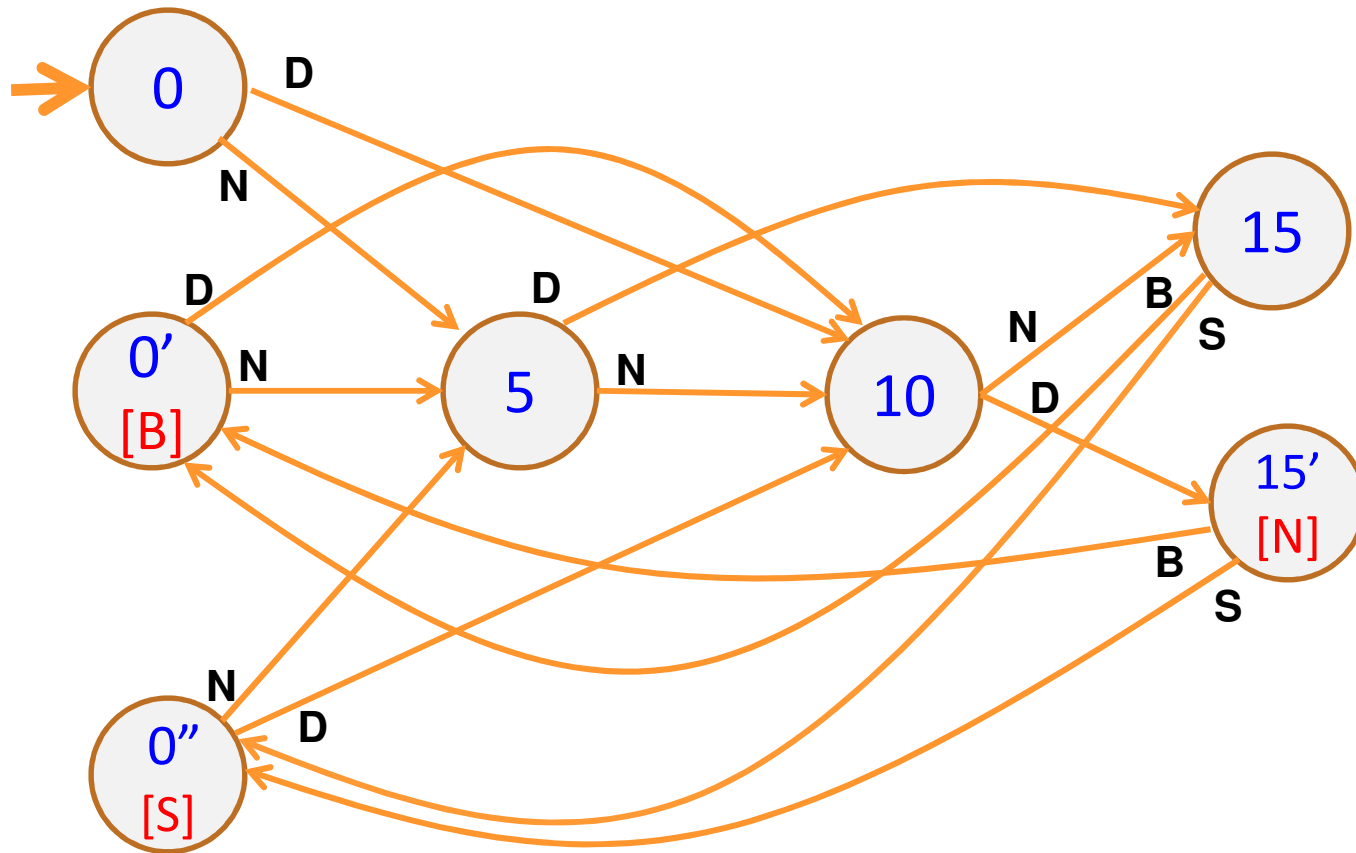
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Basic transitions on **N** (nickel), **D** (dime), **B** (butterfinger), **S** (snickers)

# Vending Machine, v0.2

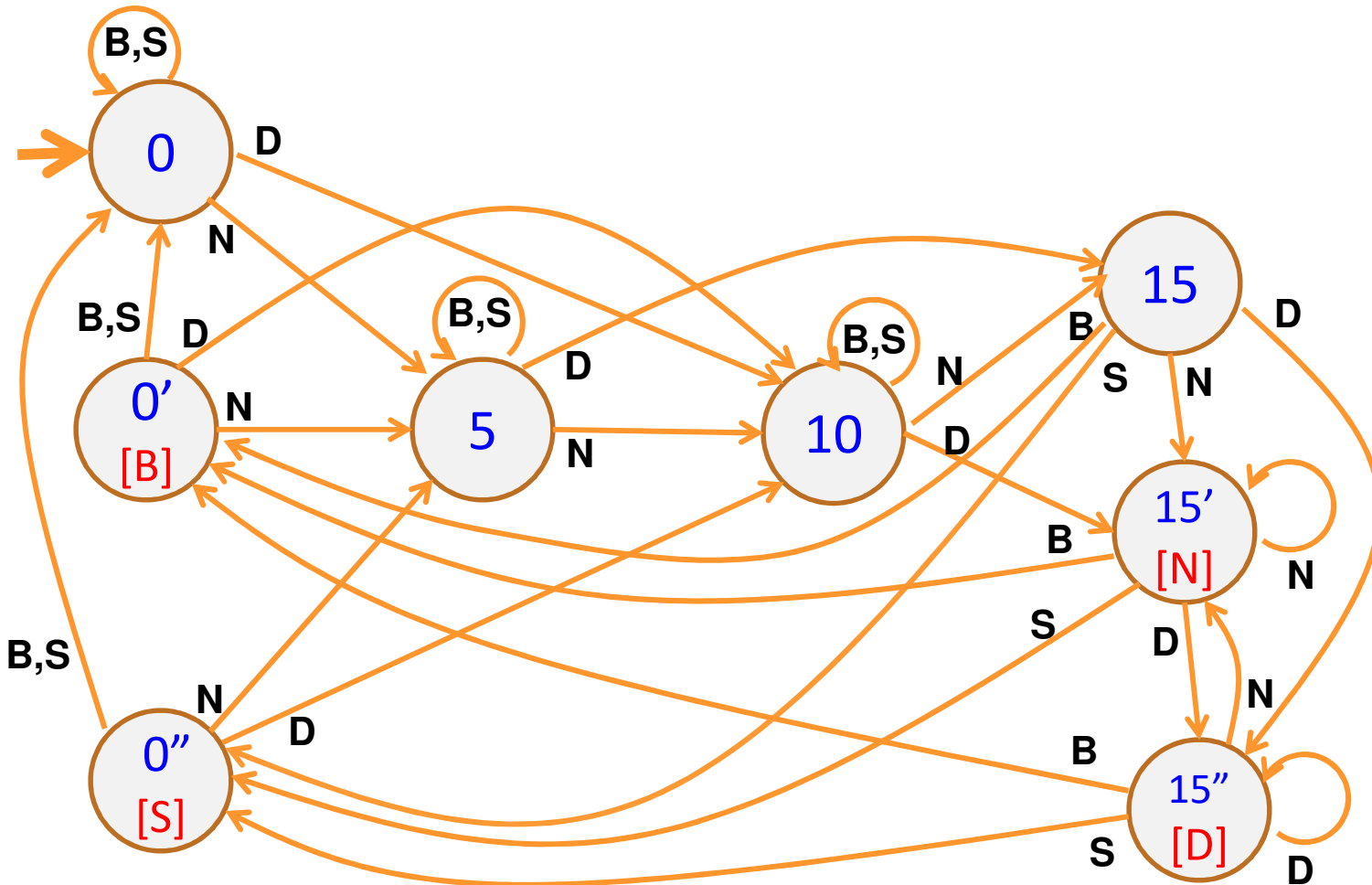
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Adding output to states: **N** – Nickel, **S** – Snickers, **B** – Butterfinger

# Vending Machine, v1.0

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Adding additional “unexpected” transitions to cover all symbols for each state