

cor.

$$a = \prod p_i^{a_i}, b = \prod p_i^{b_i}$$

$$\text{gcd}(a, b) = \prod p_i^{\min(a_i, b_i)}$$

An algorithm for gcd  
factor a (integers)  
factor b  
take min of exponents

2^n bit #

Factor n(a)

For d = 2, 3, 4, 5 ...  $\frac{\sqrt{a}}{2}$

see if d|a

if so continue on  $\frac{a}{d}$  starting with d

2^n trial divisions

n = 32 → 4 billion  
n = 64 → 16 billion billion  
10^25

if  $a = x \cdot y$  then either  $x \leq \sqrt{a}$   
or  $y \leq \sqrt{a}$

(if not  $x \cdot y > \sqrt{a} \cdot \sqrt{a} = a$ )

With  $\sqrt{a}$  optimized on

$$\sim 2^{1/2} \quad 2^{1000/2} = 2^{500} = 10^{150}$$

GCD(a, b)

while b  $\neq$  0 {

$$r = a \bmod b$$

$$a = b$$

$$b = r \quad \}$$

return a

Euclid's algorithm < 300 < 440

a = 440    b = 300

$$440 = 1 \cdot 300 + 140$$

$$300 = 2 \cdot 140 + 20$$

$$140 = 7 \cdot 20 + 0$$

20  
= 0  
→ gcd = 20

$$\text{let } a = qb + r \quad 0 \leq r < b$$

$\uparrow$  quotient      remainder

Claim  $\gcd(a, b) = \gcd(b, r)$

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if  $d|a$  &  $d|b$   
 then  $d|a+b$   
 and  $d|sa+tb \quad \forall s, t \in \mathbb{Z}$

$$d|a \Rightarrow \exists u \text{ st } a = du$$

$$d|b \Rightarrow \exists v \text{ st } b = dv$$

$$sa + tb = sdu + tdv$$

$$= d(su + tv)$$

$$\therefore d|sa + tb$$


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$$a = qb + r$$

if  $d|a$  &  $d|b$  then  $d|r$  since  $r = a - qb$

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