

Closed Book, Closed Notes Time Limit: 1 hour 50 minutes

1. Answer True or False to the following questions and briefly JUSTIFY each answer.

- (a) If a PDA M actually pushes a symbol on its stack then $L(M)$ is not regular.
- (b) For every infinite regular language L over Σ there are strings $x, y, z \in \Sigma^*$ such that $|y| \geq 1$ and for every $k \geq 0$, $xy^kz \in L$.
- (c) There is no algorithm to tell, given an arbitrary string P whether or not P is the ASCII for a syntactically correct C program.
- (d) If L is a CFL and $L = K \cap R$ for a regular language R then K is a CFL.
- (e) If L is not a CFL and $L = K \cap R$ for a regular language R then K is not a CFL.
- (f) There is no algorithm to tell, given an arbitrary program P and an input x , whether or not P runs forever on input x .
- (g) If the stack in a PDA M only has capacity for 100 characters on any input then $L(M)$ is regular.
- (h) If L is accepted by some PDA M then L^R is accepted by some PDA M' .
- (i) There is no algorithm to tell, given an arbitrary CFG G and input x whether or not $x \in L(G)$.
- (j) Context-free languages are more interesting than regular languages.

2. Classify each of the following sets as:

- A Both Regular and Context-Free,
- B Context-Free, but not regular,
- C Neither context-free nor regular but membership in the language can be decided by an algorithm.
- D Undecidable

You do not need to justify your answers.

- (a) $\{a^n b^m \mid m = 2n + 1\}$ A B C D
- (b) $\{a^n a^m \mid m = 2n + 1\}$ A B C D
- (c) $\{a^n b^m \mid m \equiv n \pmod{3}\}$ A B C D
- (d) $\{x c x^R \in \{a, b\}^* \mid x \text{ has an even number of a's}\}$ A B C D
- (e) $\{a^i b^j c^k \mid k = i + j\}$ A B C D
- (f) $\{a^i b^j a^k \mid k \neq i \text{ or } k \neq j\}$ A B C D
- (g) $\{a^i b^j a^k \mid k \neq i \text{ and } k \neq j\}$ A B C D
- (h) $\{a^n b^n a^m b^m \mid m, n \geq 0\}$ A B C D
- (i) $\{a^n b^m a^n b^m \mid m, n \geq 0\}$ A B C D
- (j) $\{xy \in \{a, b\}^* \mid x \neq y\}$ A B C D

3. Let $L = \{0^n 1^m \mid m \text{ is an integer multiple of } n\}$.
Use the Myhill-Nerode theorem to show that L is not regular.
4. Let $G = (V, \Sigma, R, S)$ be the context-free grammar with the following set of rules:

$$S \rightarrow aSaSb \mid aSbSa \mid bSaSa \mid SS \mid \epsilon$$

- (a) Use one of the general constructions that convert a CFG to a PDA to give a PDA M such that $L(M) = L(G)$.
- (b) Show each step of a computation of your PDA (list the stack contents, state, and amount of input remaining) that accepts input $aabbbaaab$.
- (c) (Bonus) What language does this grammar generate?
5. Let $\#_a(x)$ be the number of a's in x and $\#_b(x)$ be the number of b's in x .
Let $L = \{x \in \{a, b\}^* \mid \#_a(x) = \#_b(x) \text{ and both are even}\}$.
- (a) Give the formal definition of a context-free grammar G such that $L(G) = L$.
- (b) Give a parse tree for the string: $aabbbaab$
- (c) Write out the leftmost derivation corresponding to the parse tree in part (b).
6. Let $L = \{a^m b^n c^p \mid 0 \leq m < n < p\}$. Prove that L is not a context-free language.
7. For any context-free grammar $G = (V, \Sigma, R, S)$, we say that a nonterminal $A \in V$ is *useful* if there is a derivation $S \Rightarrow^* xAy \Rightarrow^* w$ for $w \in \Sigma^*$ and $x, y \in (V \cup \Sigma)^*$; otherwise it is *useless*. Suppose that G has the following rules:

$$\begin{aligned} S &\rightarrow AC \mid BS \mid B \\ A &\rightarrow aA \mid aT \\ B &\rightarrow CF \mid b \\ C &\rightarrow cC \mid D \\ D &\rightarrow aD \mid BD \mid C \\ E &\rightarrow aA \mid BSA \\ T &\rightarrow bB \mid b \\ U &\rightarrow bA \mid a \mid Wb \\ W &\rightarrow Ub \mid a \mid Bb \end{aligned}$$

- (a) Which non-terminals of G are useful?
- (b) Modify G to get an equivalent grammar G' whose nonterminals consist of precisely the nonterminals of G that are useful. (Don't remove any other non-terminals.)
- (c) Describe a reasonably efficient algorithm that will remove all useless non-terminals from a grammar and thus show that any grammar is equivalent to one with only useful non-terminals.