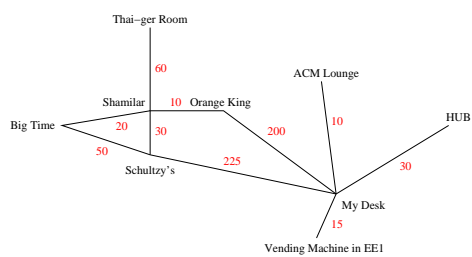


16—Shortest Paths and Dijkstra's Algorithm

May 17, 2002

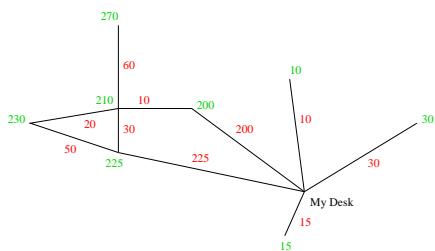
Weighted Graphs



Quest for Food
(all distances in yards)

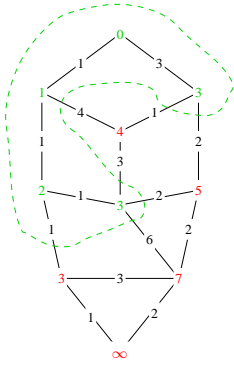
Serious About My Lunch

Single-Source Shortest Path



Compute shortest distance to all nodes from my desk

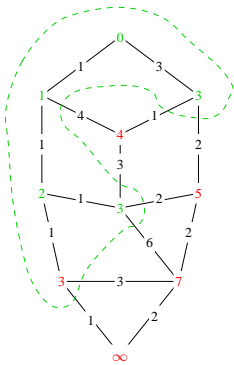
Dijkstra's Algorithm



Two kinds of vertices

- **Finished** vertices
Shortest distance computed
- **Unknown** vertices
Have *tentative* distance

Dijkstra's Algorithm



At each step:

1. Pick *closest* unknown vertex
2. Add it to finished vertices
3. Update distances

Dijkstra's vs. BFS

At each step:

1. Pick closest unknown vertex
2. Add it to finished vertices
3. Update distances

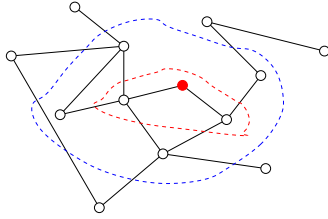
Dijkstra's Algorithm

At each step:

1. Pick vertex from queue
2. Add it to visited vertices
3. Update queue with neighbors

BFS

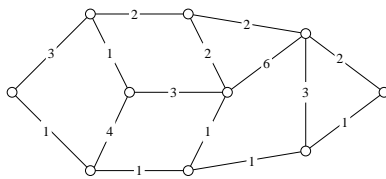
Grudge Match



1. *Finished* vertices in the middle
2. Vectors from fringe added at each step

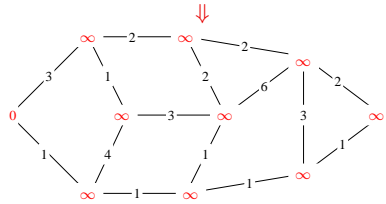
Dijkstra's and BFS

Yay! Example!

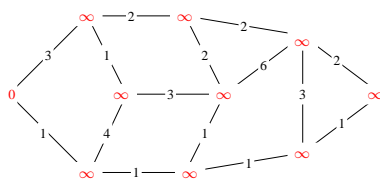


- All vertices unknown
- Start vertex distance 0
- All other vertices at ∞

Initialize



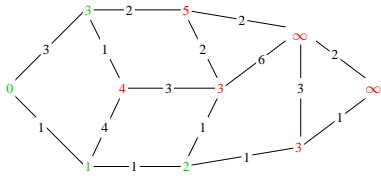
Exemplo-rific



At each step:

1. Pick closest unknown vertex
2. Add it to finished vertices
3. Update distances

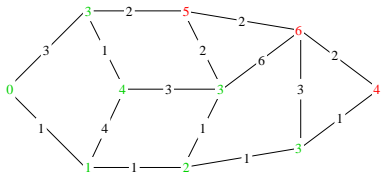
Exemplarama



At each step:

1. Pick closest unknown vertex
2. Add it to finished vertices
3. Update distances

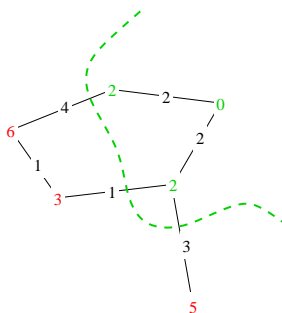
Examplophobia



At each step:

1. Pick closest unknown vertex
2. Add it to finished vertices
3. Update distances

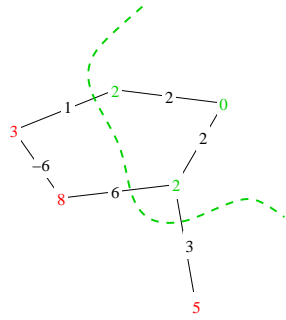
Would I Lie to You?



Why does this work?

Why are the distances of finished vertices correct?

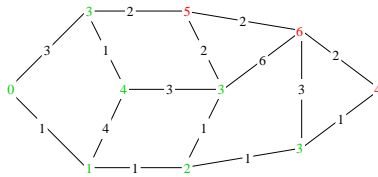
When Dijkstra's Wouldn't Work



Positive-only edge weights is essential

Key Lemma

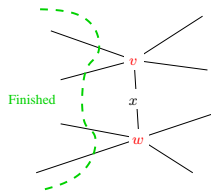
We finish with the closest vertices first



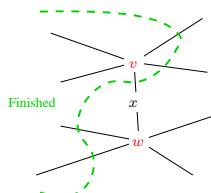
If w finished immediately after v , then
 $\text{FinalDistance}(w) \geq \text{FinalDistance}(v)$

Inductive Proof

If w finished immediately after v , then
 $\text{FinalDistance}(w) \geq \text{FinalDistance}(v)$



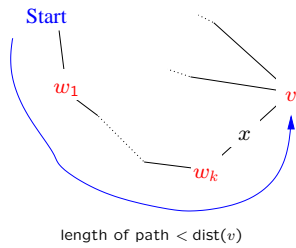
Tentative distance of w no bigger than v 's



Updates from v may *shrink* distance of w ,
 but can't get smaller than v 's distance

The Big Kahuna

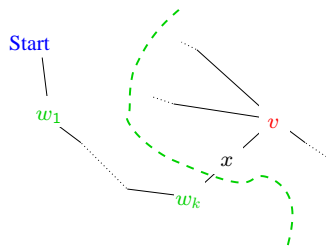
Suppose v 's final distance is *not* correct



Then exists a shorter path to v

- $\text{dist}(w_k) + x < \text{dist}(v)$
- Final distances on path are correct
- v finalized *after* all other nodes on the path

Gotcha!

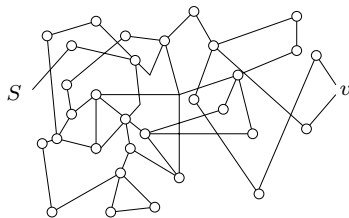


When w_k finalized, $\text{dist}(v) \leftarrow \text{dist}(w) + x$

- Final $\text{dist}(v)$ *greater* than this
- $\text{dist}(v)$ only *decreases* during the algorithm
- Contradiction!

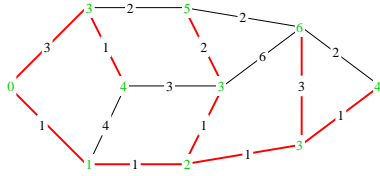
A Number is Not a Path

Sure, we know how far away v is. . .



. . . but how do we *get* there?

Shortest Path Tree

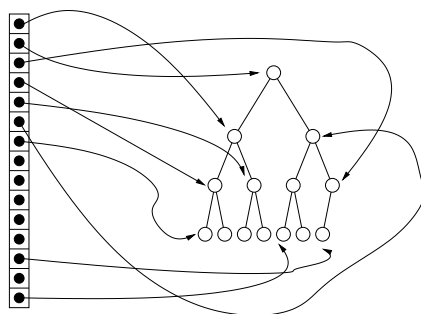


Analogue of BFS tree

In All Its Glory

```
Dijkstra(Graph G, Vertex s)
{
    PQ pq;
    for each (v in G) pq.Insert(v, INFINITY);
    pq.Change(s, 0);
    while (!pq.Empty()) {
        v = pq.DeleteMin();
        for each (w in v.Neighbors())
            if (dist(w) > dist(v) + w(v,w))
                pq.Change(w, dist(v)+w(v,w));
    }
}
```

Implementing the PQ



The List of Vertices

The PQ

PQ as a Heap



vtx.loc



heap

```
PQ::SwapUp(i)
{
  int p = (i-1)/2;
  while (p >= 0
        && heap[p].key
          > heap[i].key) {

    swap(heap[p], heap[i]);

    i = p;
    p = (i-1)/2;
  }
}
```