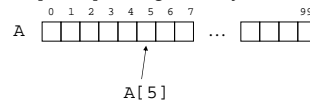


Pointers and Lists

CSE 326
Data Structures
Lecture 5

Basic Types and Arrays

- Basic Types
 - › integer, real (floating point), boolean (0,1), character
- Arrays
 - › A[0..99] : integer array

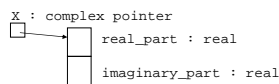


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Records and Pointers

- Record (also called a struct or data object)
 - › Group data together that are related



- › To access the fields we use “dot” notation.

X.real_part
X.imaginary_part

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Record Definition

- Record definition creates a new type

Definition

```
record complex : (  
  real_part : real,  
  imaginary_part : real  
)
```

Use in a declaration

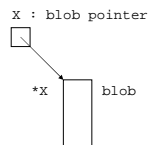
```
X : complex
```

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Pointer

- A pointer is a reference to a variable or record (or object in Java world).



- In C, if X is of type pointer to Y then *X is of type Y

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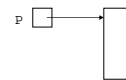
Creating a Record

- We use the “new” operator to create a record.

P : pointer to blob;

P □ (null pointer)

P := new blob;

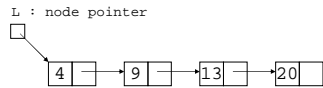


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Simple Linked List

- A linked list
 - › Group data together in a flexible, dynamic way.
 - › We'll describe several list ADTs later.



```
record node : (
  data : integer
  next : node pointer
)
```

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Memory Management – Global Allocator

- Global Allocator's store – always get and return blocks to global allocator
 - + Necessary for dynamic memory.
 - + Blocks of various sizes can be merged if they reside in contiguous memory.
 - Allocator may not handle blocks of different sizes well.
 - Allocator may be slower than a private store.

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Memory Management – Garbage Collection

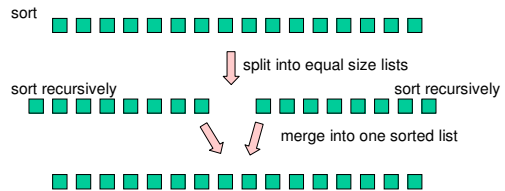
- Garbage collection – run time system recovers inaccessible blocks from time-to-time. Used in Lisp, Smalltalk, Java.
 - + No need to return blocks to an allocator or keep them in a private store.
 - Care must be taken to make unneeded blocks inaccessible.
 - When garbage collection kicks in there may be undesirable response time.

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List Mergesort

- Overall sorting plan



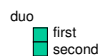
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Mergesort pseudocode

```
Mergesort(p : node pointer) : node pointer {
  Case {
    p = null : return p; //no elements
    p.next = null : return p; //one element
    else
      d : duo pointer; // duo has two fields first,second
      d := Split(p);
      return Merge(Mergesort(d.first),Mergesort(d.second));
  }
}
```

Note: Mergesort is destructive.



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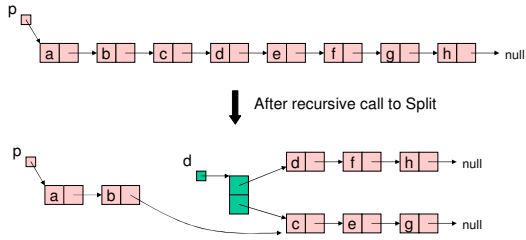
Split

```
Split(p : node pointer) : duo pointer {
  d : duo pointer;
  Case {
    p = null : d := new duo; return d
    p.next = null : d := new duo; d.first := p ; return d
    else :
      d := Split(p.next.next);
      p.next.next := d.first;
      d.first := p.next;
      p.next := d.second;
      d.second := p;
      return d;
  }
}
```

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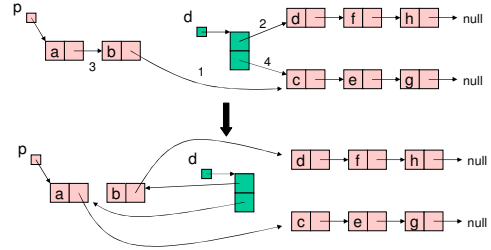
Split Example



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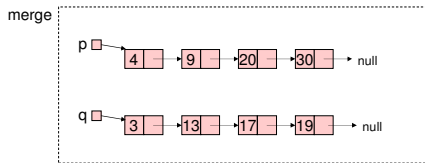
Split Example



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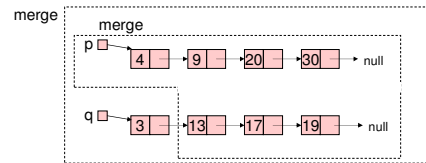
Merge Example



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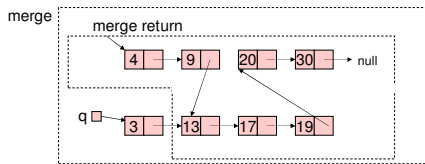
Merge Example



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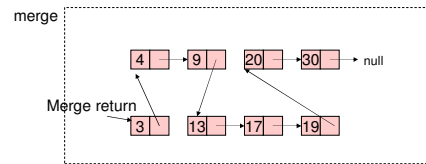
Merge Example



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Merge Example



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Merge Pseudocode Exercise

```

Merge(p,q : node pointer): node pointer{
  \precondition p and q point to sorted lists
  if p = null return q;
  if q = null return p;
  if p.data < q.data then
    p.next := merge(p.next,q);
    return p;
  else
    q.next := merge(p,q.next);
    return q
}

```

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Implementing Pointers in Arrays – “Cursor Implementation”


- This is needed in languages like Fortran, Basic, and assembly language
- Easiest when number of records is known ahead of time.
- Each record field of a basic type is associated with an array.
- A pointer field is an unsigned integer indicating an array index.

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Idea

Pointer World

n nodes
data next

data : basic type
next : node pointer

Nonpointer World

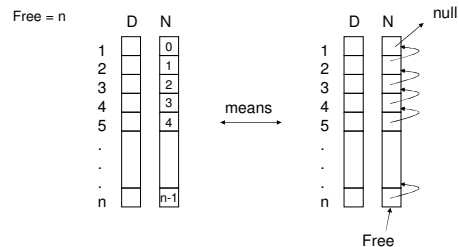
	D	N
1		
2		
3		
4		
5		
.		
.		
.		
n		

- D[] : basic type array
- N[] : integer array
- Pointer is an integer
- null is 0
- p.data is D[p]
- p.next is N[p]
- Free list needed for node allocation

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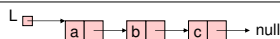
Initialization



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Example of Use



n = 8
L = 4
Free = 7

	D	N
1		3
2	c	0
3		0
4	a	6
5		8
6	b	2
7		5
8		1

```

InsertFront(L : integer, x : basic type) {
  q : integer;
  if not(Free = 0) then q := Free
  else return "overflow";
  Free := N[Free];
  D[q] := x;
  N[q] := L;
  L := q;
}

```

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Try DeleteFront

- Class Participation
- Define the cursor implementation of DeleteFront which removes the first member of the list when there is one.
 - › Remember to add garbage to free list.

```

DeleteFront(L : integer) {
  ???
}

```

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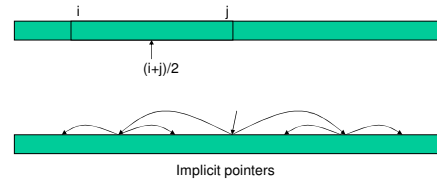
Pointer Summary

- Pointers can be implemented in several ways
 - › Explicit using “dot” notation in most high level languages
 - › Explicit using arrays of indices in all high level languages. (Cursor implementation)
 - › Implicit using calculation instead of storage. (Commonly used in nice structures like trees)

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Implicit Pointers in Binary Search



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DeleteFront Solution

```
DeleteFront(L : integer) {
  q : integer;
  if L = 0 then return "underflow"
  else {
    q := L;
    L := N[L];
    N[q] := Free;
    Free := q;
  }
}
```

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