Implementing an ADT: Representation invariants

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A data abstraction is defined by a specification

An ADT is a collection of procedural abstractions

Not a collection of procedures

Together, these procedural abstractions provide:

- A set of values
- <u>**All</u>** the ways of directly using that set of values</u>
 - Creating
 - Manipulating
 - Observing

Creators and producers: make new values Mutators: change the value (but don't affect ==) Observers: allow one to tell values apart

ADTs and specifications

Specification: only in terms of the abstraction Never mentions the representation An ADT is more than just a data structure data structure + a set of conventions

Why do we need to relate the specification to the representation?

Connecting specifications and implementations

Representation invariant: Object → boolean
Indicates whether a data structure is well-formed
Only well-formed representations are meaningful
Defines the set of valid values of the data structure
Abstraction function: Object → abstract value
What the data structure means (as an abstract value)
How the data structure is to be interpreted
How do you compute the inverse, abstract value → Object ?

Implementation of an ADT is provided by a class

To implement a data abstraction:

- Select the representation of instances, the rep
- Implement operations in terms of that rep

Choose a representation so that

- It is possible to implement operations
- The most frequently used operations are efficient
 - But which will these be?
 - Abstraction allows the rep to change later

CharSet Abstraction

// Overview: A CharSet is a finite mutable set of Characters

```
// effects: creates a fresh, empty CharSet
public CharSet ( )
```

```
// modifies: this
// effects: this
post = this
pre U {c}
public void insert (Character c);
```

```
// modifies: this
// effects: this
post = this
pre - {c}
public void delete (Character c);
```

// <u>returns</u>: ($c \in this$) public boolean member (Character c);

```
// returns: cardinality of this
public int size ( );
```

A CharSet implementation. What client code will expose the error?

```
class CharSet {
 private List<Character> elts
    = new ArrayList<Character>();
 public void insert(Character c) {
    elts.add(c);
  public void delete(Character c) {
    elts.remove(c);
                             CharSet s = new CharSet();
  public boolean member (Char
                             s.insert('a');
    return elts.contains(c)
                             s.insert('a');
                             s.delete('a');
  public int size() {
    return elts.size();
                             if (s.member('a'))
                                 // print "wrong";
}
                             else
   Where is the error?
                                  // print "right";
```

Where Is the Error?

The answer to this question tells you what needs to be fixed

Perhaps delete is wrong

It should remove all occurrences

Perhaps insert is wrong

It should not insert a character that is already there How can we know?

The representation invariant tells us

The representation invariant

- States data structure well-formedness
- Holds before and after every CharSet operation
- Operation implementations (methods) may depend on it
 Write it this way:

```
class CharSet {
   // Rep invariant: elts has no nulls and no duplicates
   private List<Character> elts;
```

```
Or, if you are the pedantic sort:
```

```
∀ indices i of elts . elts.elementAt(i) ≠ null
∀ indices i, j of elts .
i ≠ j ⇒ ¬ elts.elementAt(i).equals(elts.elementAt(j))
```

Now, we can locate the error

```
// Rep invariant:
// elts has no nulls and no duplicates
public void insert(Character c) {
   elts.add(c);
}
public void delete(Character c) {
   elts.remove(c);
}
```

Another rep invariant example

```
class Account {
   private int balance;
   // history of all transactions
   private List<Transaction> transactions;
   ...
}
```

```
// real-world constraints:
balance \ge 0
balance = \Sigma_i transactions.get(i).amount
// implementation-related constraints:
transactions \neq null
no nulls in transactions
```

Listing the elements of a CharSet

Consider adding the following method to CharSet: // <u>returns</u>: a List containing the members of this public List<Character> getElts();

Consider this implementation:

// Rep invariant: elts has no nulls and no duplicates
public List<Character> getElts() { return elts; }

Does the implementation of getElts preserve the rep invariant?

... sort of

Representation exposure

Consider this client code (outside the CharSet implementation):

```
CharSet s = new CharSet();
s.insert('a');
s.getElts().add('a');
s.delete('a');
if (s.member('a')) ...
```

Representation exposure is external access to the rep Representation exposure is almost always EVIL Enables violation of abstraction boundaries and the rep invariant If you do it, document why and how And feel guilty about it!

How can we avoid/prevent rep exposure?

Ways to avoid rep exposure

1. Exploit immutability

Character is immutable.

```
Character choose() {
   return elts.elementAt(0);
}
```

Defining fields as **private** is **not sufficient** to hide the representation

```
2. Make a copy
```

```
List<Character> getElts() {
   return new ArrayList<Character>(elts);
   // or: return (ArrayList<Character>) elts.clone();
}
Mutating a copy doesn't affect the original.
```

Don't forget to make a copy on the way in!

3. Make an immutable copy

```
List<Character> getElts() {
   return Collections.unmodifiableList<Character>(elts);
}
Client cannot mutate
Still need to make a copy on the way in
```

Checking rep invariants

Should code check that the rep invariant holds?

- Yes, if it's inexpensive
- Yes, for debugging (even when it's expensive)
- It's quite hard to justify turning the checking off
- Some private methods need not check (Why?)

Checking the rep invariant

Rule of thumb: check on entry and on exit (why?)

```
public void delete(Character c) {
    checkRep();
    elts.remove(c)
    // Is this guaranteed to get called?
    // See handouts for a less error-prone way to check at exit.
    checkRep();
}
...
/** Verify that elts contains no duplicates. */
private void checkRep() {
    for (int i = 0; i < elts.size(); i++) {
        assert elts.indexOf(elts.elementAt(i)) == i;
    }
}</pre>
```

Practice defensive programming

Assume that you will make mistakes

Write and incorporate code designed to catch them

On entry:

Check rep invariant

Check preconditions (requires clause)

On exit:

Check rep invariant

Check postconditions

Checking the rep invariant helps you discover errors

Reasoning about the rep invariant helps you avoid errors

Or prove that they do not exist!

We will discuss such reasoning, later in the term

The rep invariant constrains structure, not meaning

New implementation of insert that preserves the rep invariant:

```
public void insert(Character c) {
   Character cc = new Character(encrypt(c));
   if (!elts.contains(cc))
      elts.addElement(cc);
   }
  public boolean member(Character c) {
    return elts.contains(c);
  }
```

```
The program is still wrong
Clients observe incorrect behavior
What client code exposes the error?
Where is the error?
We must consider the meaning
The abstraction function helps us
```

```
CharSet s = new CharSet();
s.insert('a');
if (s.member('a'))
        // print "right";
else
        // print "wrong";
```