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# CSE 331

# Software Design & Implementation

Spring 2022

Section 6 – HW6 and Midterm Review

# Administrivia

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- Done with HW5!
- HW6 (ADT implementation) due next week (Thurs. 5/12)
- Midterm tomorrow during lecture!
- Any questions?

# Agenda

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- Walk-through of the test-script driver (to run `.test` files)
- Managing an expensive `checkRep`
- Midterm review

# Refresher: Format of script tests

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Each script test is expressed as text-based script `foo.test`

- One command per line, of the form: **Command** *arg<sub>1</sub>* *arg<sub>2</sub>* ...
- Script's output compared against `foo.expected`
- Precise details specified in the homework
- Match format **exactly**, including whitespace!

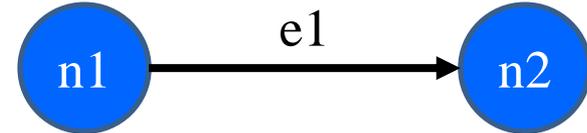
Command (in <code>foo.test</code> )	Output (in <code>foo.expected</code> )
<code>CreateGraph</code> <i>name</i>	<code>created graph</code> <i>name</i>
<code>AddNode</code> <i>graph label</i>	<code>added node</code> <i>label</i> <code>to graph</code>
<code>AddEdge</code> <i>graph parent child label</i>	<code>added edge</code> <i>label</i> <code>from parent to child in graph</code>
<code>ListNodes</code> <i>graph</i>	<code>graph contains:</code> <i>label</i> <sub><i>node</i></sub> ...
<code>ListChildren</code> <i>graph parent</i>	<code>the children of parent in graph are:</code> <i>child</i> ( <i>label</i> <sub><i>edge</i></sub> ) ...
<code>#</code> <i>This is comment text ...</i>	<code>#</code> <i>This is comment text ...</i>

# Refresher: example.test

---

```
# Create a graph  
CreateGraph graph1
```

```
# Add a pair of nodes  
AddNode graph1 n1  
AddNode graph1 n2
```



```
# Add an edge  
AddEdge graph1 n1 n2 e1
```

```
# Print all nodes in the graph  
ListNodes graph1
```

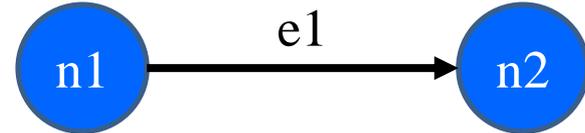
```
# Print all child nodes of n1 with outgoing edge  
ListChildren graph1 n1
```

# Refresher: example.expected

---

```
# Create a graph  
created graph graph1
```

```
# Add a pair of nodes  
added node n1 to graph1  
added node n2 to graph1
```



```
# Add an edge  
added edge e1 from n1 to n2 in graph1
```

```
# Print all nodes in the graph  
graph1 contains: n1 n2
```

```
# Print all child nodes of n1 with outgoing edge  
the children of n1 in graph1 are: n2(e1)
```

# How the script tests work

---

- In HW5, you wrote script tests in the form of `.test` files
  - As well as an `.expected` file for each test's expected outcome
- The JUnit class `ScriptFileTests` runs all these tests
  - Looks for all the `.test` files in the `src/test/resources/testScripts` folder
  - Compares test output against corresponding `.expected` file
- `ScriptFileTests` needs a bridge to your graph implementation
  - That's exactly what the `GraphTestDriver` class is for

# Graph Test Driver

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- **GraphTestDriver** knows how to read these test scripts
- **GraphTestDriver** calls a method to “do” each verb
  - **CreateGraph**, **AddNode**, **AddEdge** ...
  - One method stub per script command for you to fill with calls to your graph code
- Note: Completed test driver should sort lists before printing for **ListNodes** and **ListChildren**
  - Just to ensure predictable, deterministic output
  - Your graph implementation itself should not worry about sorting

# Graph Test Driver Output

---

- The Graph Test Driver is a client of our graph...
  - ...but not the only client.
  - Your graph should not be designed to be exclusively used for the test driver.
- ListChildren in the test driver should print out: “**the children of parent in graph are: *child*(*label<sub>edge</sub>*) ...**”
- This does **not** mean that you should have a method on your graph called ListChildren that returns this String
  - Because that isn't useful for other clients

# Sorting with the driver

---

- **Use the test driver appropriately!**
  - From before: “Completed test driver should sort lists before printing.”
- Script test output for hw5 needs to be sorted so we can mechanically check it.
- This means sorted output for tests does ***NOT*** mean sorted internal storage in graph.
  - If sorting behavior is needed, Graph ADT clients (including the test driver) can sort those labels.

# In other words...

---

The Graph ADT in general should **NOT** assume that node or edge labels are sorted or even comparable(!).  
(of course they can be tested for equality with equals() )

# Demo

---

Here's a quick tour of the `GraphTestDriver`!

# Expensive checkReps

---

- A complicated rep. invariant can be expensive to check
  - Especially iterating over internal collection(s)
  - For example, examining every edge in a graph
- A slow `checkRep` could cause our grading scripts to time-out
  - Can be really useful during testing/debugging, but
  - Need to disable the really slow checks before submitting
- We have a tension between two goals:
  - Thorough, possibly slow checking for development
  - Essential, necessarily fast checking for production/grading
- What to do?

# Use a debug flag to tune `checkRep`

---

- Repeatedly (un)commenting sections of code is a poor solution
- Instead, use a class-level constant as a toggle
  - Ex.: `private static final boolean DEBUG = ...;`
    - `false` for only the fast, essential checks
    - `true` for all the slow, thorough checks
  - Real-world code often has several such “debug levels”

```
private void checkRep() {  
    assert fast_checks();  
    if (DEBUG)  
        assert slow_checks();  
}
```

---

# Midterm Review

# intToString()

---

Fill in the implementation of a method that converts a **positive integer** to its **string representation in decimal** (invariant given on next slide).

```
{ { P: x > 0 } }
```

```
String intToString(int x)
```

Useful facts to recall:

1. Convert **char** **ch** that is one of `'0'`, `'1'`, ..., `'9'` to a corresponding **int** by doing `ch - '0'`
2. Convert **int** **x** that is one of `0`, `1`, ..., `9` to a corresponding **char** by doing `(char) (x + '0')`

# intToString()

---

```
{ { P: x > 0 } }
```

```
String intToString(int x) {
```

```
    StringBuilder buf =
```

```
    int k = , y = ;
```

```
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and  $y = x / 10^k$  } }
```

```
    while (y != 0) {
```

```
        k = k + 1;
```

```
    }
```

```
    return buf.reverse().toString();
```

```
}
```

How do we get the invariant to hold initially?

# intToString()

---

```
{ { P: x > 0 } }  
String intToString(int x) {  
    StringBuilder buf = new StringBuilder();  
    int k = 0, y = x;  
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and y = x / 10^k } }  
    while (y != 0) {  
  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

How do we fill out the loop body?

# intToString()

---

```
{ { P: x > 0 } }  
String intToString(int x) {  
    StringBuilder buf = new StringBuilder();  
    int k = 0, y = x;  
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and y = x / 10^k } }  
    while (y != 0) {  
  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

Inv changes  $k$  to  $k+1$ , so

- $y$  becomes  $x / 10^{k+1}$
- $y_{\text{post}} = x / 10^{k+1} = y_{\text{pre}} / 10$

# intToString()

---

```
{ { P: x > 0 } }  
String intToString(int x) {  
    StringBuilder buf = new StringBuilder();  
    int k = 0, y = x;  
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and y = x / 10^k } }  
    while (y != 0) {  
  
        y = y / 10;  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

Inv changes  $k$  to  $k+1$ , so

- $y$  becomes  $x / 10^{k+1}$
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# intToString()

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{ { P: x > 0 } }  
String intToString(int x) {  
    StringBuilder buf = new StringBuilder();  
    int k = 0, y = x;  
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and y = x / 10^k } }  
    while (y != 0) {  
  
        y = y / 10;  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

- Inv changes  $k$  to  $k+1$ , so
- buf stores lowest  $k+1$  digits

# intToString()

---

```
{ { P: x > 0 } }  
String intToString(int x) {  
    StringBuilder buf = new StringBuilder();  
    int k = 0, y = x;  
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and y = x / 10^k } }  
    while (y != 0) {  
  
        y = y / 10;  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

Inv changes  $k$  to  $k+1$ , so

- buf stores lowest  $k+1$  digits

$(k+1)$ -st lowest digit goes at end  
since buf stores them reversed

# intToString()

---

```
{ { P: x > 0 } }
```

```
String intToString(int x) {
```

```
    StringBuilder buf = new StringBuilder();
```

```
    int k = 0, y = x;
```

```
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and  $y = x / 10^k$  } }
```

```
    while (y != 0) {
```

```
        char ch = ?
```

```
        buf.append(ch);
```

```
        y = y / 10;
```

```
        k = k + 1;
```

```
    }
```

```
    return buf.reverse().toString();
```

```
}
```

How can we get the (k+1)-st lowest digit of x? And make it a char?

# intToString()

---

```
{ { P: x > 0 } }  
String intToString(int x) {  
    StringBuilder buf = new StringBuilder();  
    int k = 0, y = x;  
    { { Inv: P and buf stores the lowest k digits of x  
        in reverse order and y = x / 10^k } }  
    while (y != 0) {  
        char ch = (char) (y % 10 + '0');  
        buf.append(ch);  
        y = y / 10;  
        k = k + 1;  
    }  
  
    return buf.reverse().toString();  
}
```

# intToString() solution

---

```
{ { P: x > 0 } }
String intToString(int x) {
    StringBuilder buf = new StringBuilder();
    int k = 0, y = x;
    { { Inv: P and buf stores the lowest k digits of x
        in reverse order and y = x / 10^k } }
    while (y != 0) {
        char ch = (char) (y % 10 + '0');
        buf.append(ch);
        y = y / 10;
        k = k + 1;
    }
    { { buf stores the digits of x in reverse order } }
    return buf.reverse().toString();
}
```

Why does this hold?

# intToString() solution

---

```
{ { P: x > 0 } }
String intToString(int x) {
    StringBuilder buf = new StringBuilder();
    int k = 0, y = x;
    { { Inv: P and buf stores the lowest k digits of x
        in reverse order and y = x / 10^k } }
    while (y != 0) {
        char ch = (char) (y % 10 + '0');
        buf.append(ch);
        y = y / 10;
        k = k + 1;
    }
    { { buf stores the digits of x in reverse order } }
    return buf.reverse().toString();
}
```

Why does this hold?  
 $y = 0 \Rightarrow x < 10^k$  so x only  
has k digits

# Specifications

---

Suppose we have a **BankAccount** class with instance variable `balance`. Consider the following specifications (ignore `@param`):

- A. `@effects` decreases `balance` by `amount` 
- B. `@requires` `amount >= 0` and `amount <= balance`  
`@effects` decreases `balance` by `amount` 
- C. `@throws` `InsufficientFundsException` if `balance < amount`  
`@effects` decreases `balance` by `amount` 

Which specifications does this implementation meet?

```
void withdraw(int amount) {  
    balance -= amount;  
}
```

# Specifications

---

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`@effects` decreases `balance` by `amount` 
- C. `@throws` `InsufficientFundsException` if `balance < amount`  
`@effects` decreases `balance` by `amount` 

Which specifications does this implementation meet?

```
void withdraw(int amount) {  
    if (balance >= amount) balance-=amount;  
}
```

# Specifications

---

Suppose we have a **BankAccount** class with instance variable `balance`. Consider the following specifications (ignore `@param`):

- A. `@effects` decreases `balance` by `amount` 
- B. `@requires` `amount >= 0` and `amount <= balance`  
`@effects` decreases `balance` by `amount` 
- C. `@throws` `InsufficientFundsException` if `balance < amount`  
`@effects` decreases `balance` by `amount` 

Which specifications does this implementation meet?

```
void withdraw(int amount) {  
    if (amount < 0) throw new IllegalArgumentException();  
    balance -= amount;  
}
```

# Specifications

---

Suppose we have a **BankAccount** class with instance variable `balance`. Consider the following specifications (ignore `@param`):

- A. `@effects` decreases `balance` by `amount` 
- B. `@requires` `amount >= 0` and `amount <= balance`  
`@effects` decreases `balance` by `amount` 
- C. `@throws` `InsufficientFundsException` if `balance < amount`  
`@effects` decreases `balance` by `amount` 

Which specifications does this implementation meet?

```
void withdraw(int amount) throws InsufficientFundsException {  
    if (balance < amount) throw new InsufficientFundsException();  
    balance -= amount;  
}
```

# Testing

---

Consider the `BankAccount` class again. What are some good test cases?

```
public class BankAccount {
    /** @return current balance of account */
    public void balance() { ... }

    /**
     * @param amount to withdraw
     * @requires amount >= 0
     * @throws InsufficientFundsException
     *         if balance < amount
     * @effects decreases balance by amount
     */
    public void withdraw(int amount) { ... }
}
```

Specification test heuristic:

- amount  $\leq$  balance
- amount  $>$  balance

Boundary test heuristic:

- amount = balance
- amount  $>$  balance

Others?

Should we test amount  $<$  0?

# More Reasoning

---

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ ? }}
        xk = xk * x;
        {{ ? }}
        val = val + a[k+1]*xk;
        {{ ? }}

        k = k + 1;
        {{ ? }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

# More Reasoning

---

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ ? }}
        xk = xk * x;
        {{ ? }}
        val = val + a[k+1]*xk;
        {{ ? }}

        k = k + 1;
        {{ ? }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

Does the invariant hold before the loop?

# More Reasoning

---

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ ? }}
        xk = xk * x;
        {{ ? }}
        val = val + a[k+1]*xk;
        {{ ? }}

        k = k + 1;
        {{ ? }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

# More Reasoning

---

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ inv && k != n }}
        xk = xk * x;
        {{ ? }}
        val = val + a[k+1]*xk;
        {{ ? }}

        k = k + 1;
        {{ ? }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

# More Reasoning

---

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ inv && k != n }}
        xk = xk * x;
        {{ xk = x^(k+1) && val = a[0] + a[1]*x + ... + a[k]*x^k && k != n }}
        val = val + a[k+1]*xk;
        {{ ? }}

        k = k + 1;
        {{ ? }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

# More Reasoning

---

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ inv && k != n }}
        xk = xk * x;
        {{ xk = x^(k+1) && val = a[0] + a[1]*x + ... + a[k]*x^k && k != n }}
        val = val + a[k+1]*xk;
        {{ xk = x^(k+1) && val = a[0] + a[1]*x + ... + a[k+1]*x^(k+1)
            && k != n }}
        k = k + 1;
        {{ ? }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

# More Reasoning

---

Let's check that this method is correct.

```
/** Return the value of this IntPoly at point x */
public int valueAt(int x) {
    int val = a[0];
    int xk = 1;
    int k = 0;
    int n = a.length - 1;
    {{ inv: xk = x^k && val = a[0] + a[1]*x + ... + a[k]*x^k }}
    while (k != n) {
        {{ inv && k != n }}
        xk = xk * x;
        {{ xk = x^(k+1) && val = a[0] + a[1]*x + ... + a[k]*x^k && k != n }}
        val = val + a[k+1]*xk;
        {{ xk = x^(k+1) && val = a[0] + a[1]*x + ... + a[k+1]*x^(k+1)
            && k != n }}

        k = k + 1;
        {{ inv && k-1 != n }} -> {{ inv }}
    }
    {{ val = a[0] + a[1]*x + ... + a[n]*x^n }}
    return val;
}
```

Do we reach the postcondition?

---

Good luck on the midterm!