
CSE 331

Software Design & Implementation

Section: Dijkstra's Algorithm; MVC; HW7

Reminders

- On HW7, it is ok to go back and modify your HW6 Graph

Upcoming Deadlines

- HW6 due 11pm tonight (7/28)
- Prep. Quiz: HW7 due 11pm Monday (8/01)

Last Time...

- Subtyping
- Generics
- Event-driven programming

Today's Agenda

- HW7 Overview
- Dijkstra's Algorithm
- Model-View-Controller (MVC)
- Campus Dataset

HW7 – Overview

- HW7 includes 2 folders:
 - **hw-tasks/**
 - **hw-pathfinder/**
- When done, attach the tag **hw7-final**
 - Reminder: commit/push everything, and **then** create/push the tag in a **separate transaction!**
 - Remember to check **Repository > Graph** on GitLab to verify that your tag is on the correct commit!

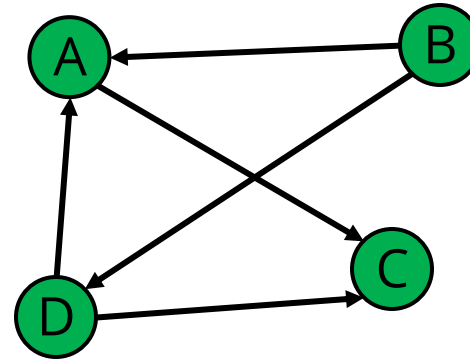
HW7 – Tasks

- You will first need to make your graph class **generic** to take other types for node and edge labels that are not Strings.
 - a. Update HW5/6 to use the generic graph ADT
 - b. Make sure all the HW5/6 tests pass!
- You will need to implement some of **TaskSorter**
 - Tasks can be dependent on other tasks (i.e. one needs to be completed before the other)
 - What's a natural way to represent this? A graph!
 - Given a set of tasks and dependencies, can we find an ordering of tasks that satisfies the dependencies?
 - This algorithm is already written for you (we suggest you take a look)

HW7 – Tasks

- Tasks are nodes, dependencies are edges
- Let's take a look at a visual:
 - If $X \rightarrow Y$, task X must be done before task Y .
 - What order can we complete these tasks in?

B -> D -> A -> C



HW7 – Pathfinder

Next part: a program to find the shortest walking routes through campus *ca.* 2006

- Network of walkways in campus constitutes a graph!

Pathfinder progresses through 3 steps:

1. Implement Dijkstra's algorithm
 - Starter code gives a path ADT to store search result: `pathfinder.datastructures.Path`
2. Run tests for your implementation of Dijkstra's algorithm
3. Complete starter code for the Pathfinder application

Dijkstra's algorithm

- Named for its inventor, Edsger Dijkstra (1930–2002)
 - Truly one of the “founders” of computer science
 - Just one of his many contributions
- Key idea: find shortest path based on numeric edge weights:
 - Track the path to each node with least-yet-seen cost
 - Shrink a set of pending nodes as they are visited
- *A priority queue* makes handling weights efficient and convenient
 - Helps track which node to process next
- **Note:** Dijkstra's algorithm requires all edge weights be **nonnegative**
 - (Other graph search algorithms can handle negative weights – see Bellman-Ford algorithm)

Priority queue

- A queue-like ADT that reorders elements by associated *priority*
 - Whichever element has the least value dequeues next (not FIFO)
 - Priority of an element traditionally given as a separate integer
- Java provides a standard implementation, **PriorityQueue<E>**
 - Implements the **Queue<E>** interface but has distinct semantics
 - Enqueue (add) with the **add** method
 - Dequeue (remove highest priority) with the **remove** method
- **PriorityQueue<E>** uses comparison order for priority order
 - Default: class **E** implements **Comparable<E>**
 - May configure otherwise with a **Comparator<E>**

Priority queue – example

```
q = new PriorityQueue<Double>();
```

| | | |
|--|--|--|
| | | |
|--|--|--|

```
q.add(5.1);
```

| | | |
|-----|--|--|
| 5.1 | | |
|-----|--|--|

```
q.add(4.2);
```

| | | |
|-----|-----|--|
| 4.2 | 5.1 | |
|-----|-----|--|

```
q.add(0.3);
```

| | | |
|-----|-----|-----|
| 0.3 | 4.2 | 5.1 |
|-----|-----|-----|

```
q.remove(); // 0.3
```

| | | |
|-----|-----|--|
| 4.2 | 5.1 | |
|-----|-----|--|

```
q.add(0.8);
```

| | | |
|-----|-----|-----|
| 0.8 | 4.2 | 5.1 |
|-----|-----|-----|

```
q.remove(); // 0.8
```

| | | |
|-----|-----|--|
| 4.2 | 5.1 | |
|-----|-----|--|

```
q.add(20.4);
```

| | | |
|-----|-----|------|
| 4.2 | 5.1 | 20.4 |
|-----|-----|------|

```
q.remove(); // 4.2
```

| | | |
|-----|------|--|
| 5.1 | 20.4 | |
|-----|------|--|

Finding the “shortest” path

- In HW7, edge labels are numbers, called *weights*
 - Labeled graphs like that are called *weighted graphs*
 - An edge’s weight is considered its *cost* (think time, distance, price, ...)
- HW7 measures the “shortest” path by the total weight of its edges
 - So really, the path with the least cost
 - Find using *Dijkstra’s algorithm*
 - Edge weights crucially relevant
- There are other definitions of “shortest” path that we will not consider

Aside: `break` vs. `continue`

- `break` exits the loop, while `continue` skips the rest of this iteration

```
for (int i = 0; i < 5; i++) {  
    if (i == 3) { break; }  
    System.out.println(i + " ");  
}
```

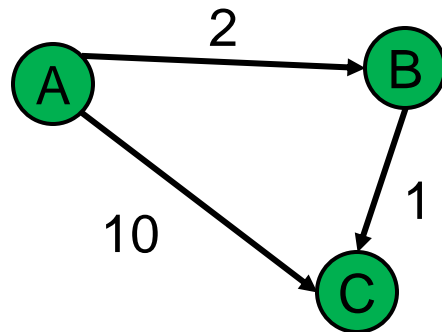
```
// out: 0 1 2
```

```
for (int i = 0; i < 5; i++) {  
    if (i == 3) { continue; }  
    System.out.println(i + " ");  
}
```

```
// out: 0 1 2 4
```

Dijkstra's algorithm

- **Main idea:** Start at the source node and find the shortest path to all reachable nodes.
 - This will include the shortest path to your destination!
- What is the shortest path from A to C for the given graph using Dijkstra's algorithm? Using BFS?



Dijkstra's algorithm – pseudocode

active = priority queue of paths.

finished = empty set of nodes.

add a path from start to itself to active

<inv ???> What would be a good invariant for this loop?

while active is non-empty:

minPath = active.removeMin()

 minDest = destination node in minPath

 if minDest is dest:

 return minPath

 if minDest is in finished:

 continue

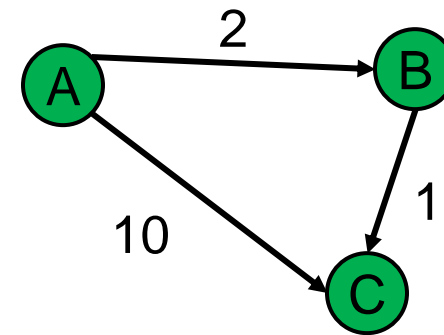
 for each edge e = (minDest, child):

 if child is not in finished:

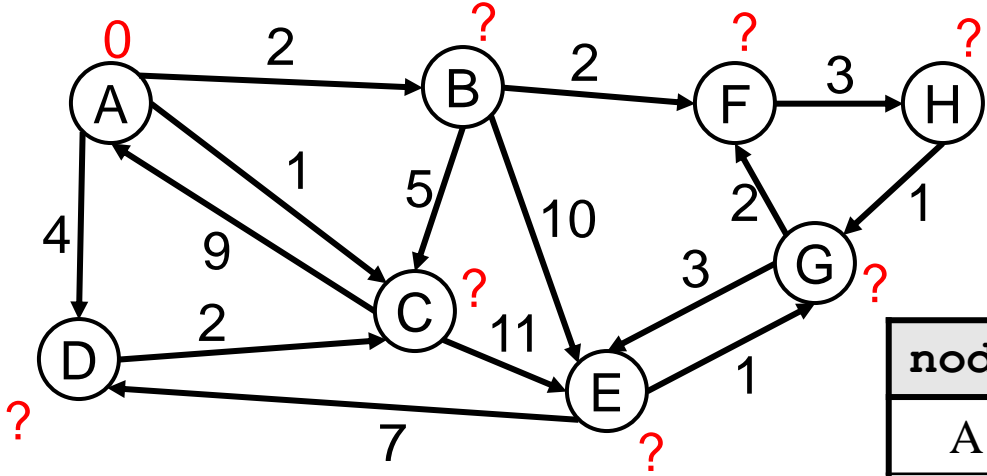
 newPath = minPath + e

 add newPath to active

 add minDest to finished



Dijkstra's algorithm – paths from A

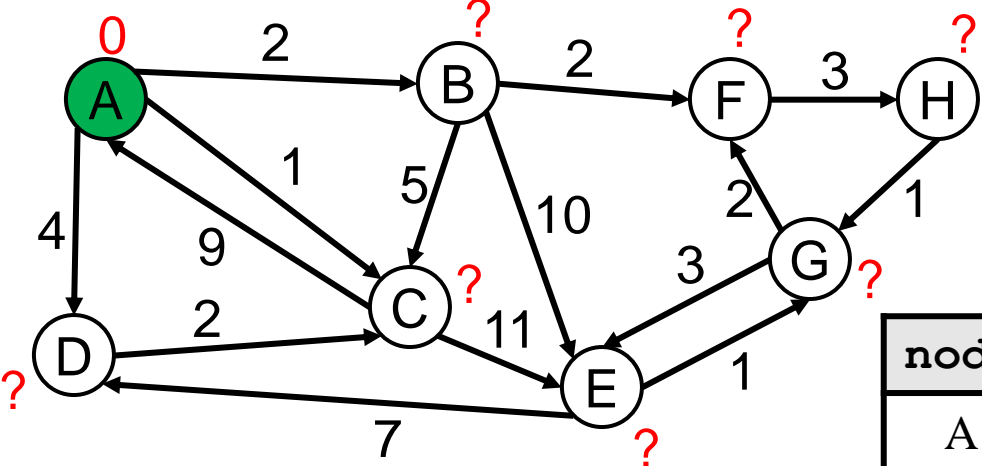


priority queue

| path | cost |
|------|------|
| [A] | 0 |
| | |
| | |
| | |

| node | finished | cost | prev |
|------|----------|------|------|
| A | | 0 | - |
| B | | | |
| C | | | |
| D | | | |
| E | | | |
| F | | | |
| G | | | |
| H | | | |

Dijkstra's algorithm – paths from A

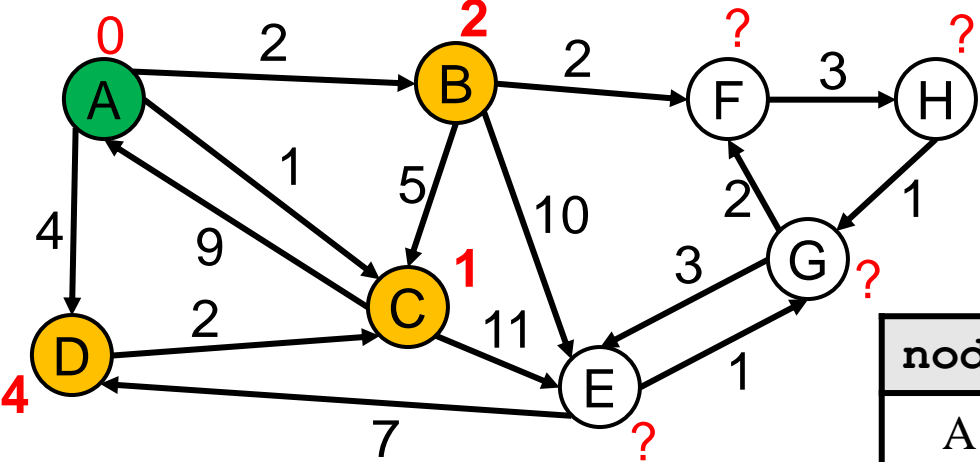


priority queue

| path | cost |
|------|------|
| | |
| | |
| | |
| | |
| | |

| node | finished | cost | prev |
|------|----------|------|------|
| A | Y | 0 | - |
| B | | | |
| C | | | |
| D | | | |
| E | | | |
| F | | | |
| G | | | |
| H | | | |

Dijkstra's algorithm – paths from A

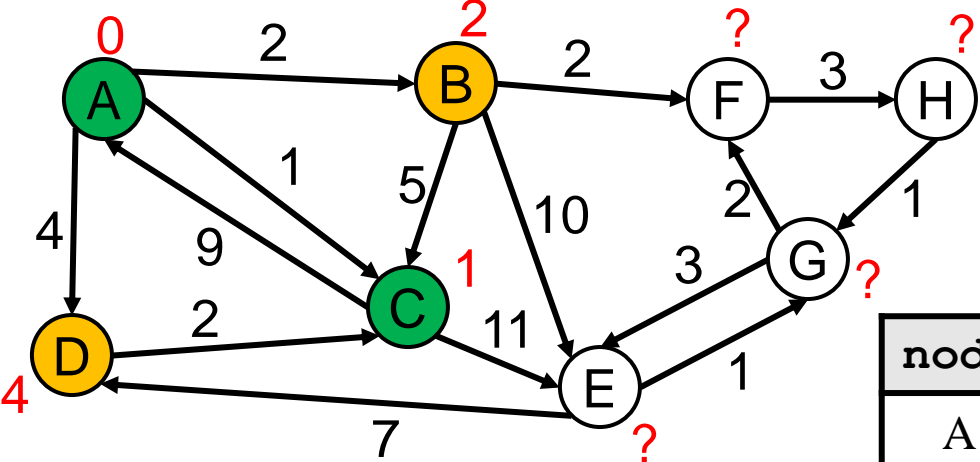


priority queue

| path | cost |
|--------|------|
| [A, C] | 1 |
| [A, B] | 2 |
| [A, D] | 4 |
| | |

| node | finished | cost | prev |
|------|----------|----------|------|
| A | Y | 0 | - |
| B | | ≤ 2 | A |
| C | | ≤ 1 | A |
| D | | ≤ 4 | A |
| E | | | |
| F | | | |
| G | | | |
| H | | | |

Dijkstra's algorithm – paths from A

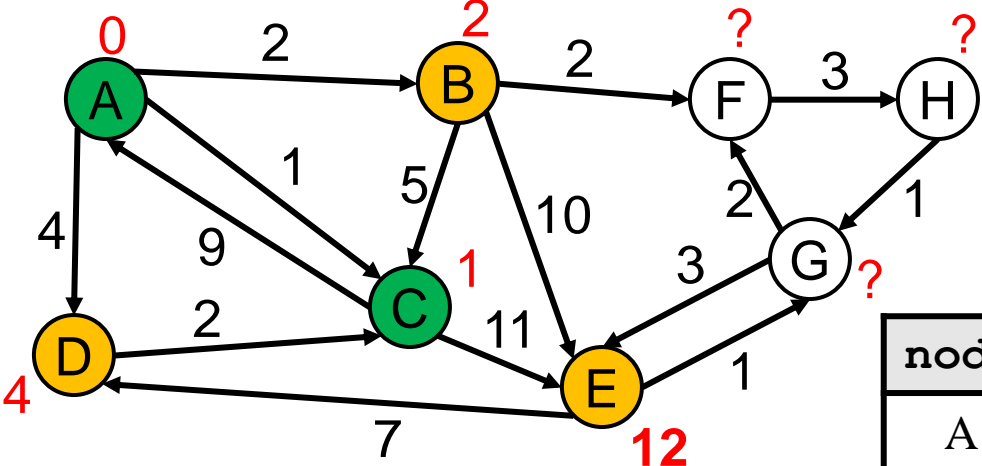


priority queue

| path | cost |
|--------|------|
| [A, B] | 2 |
| [A, D] | 4 |
| | |
| | |

| node | finished | cost | prev |
|------|----------|----------|------|
| A | Y | 0 | - |
| B | | ≤ 2 | A |
| C | Y | 1 | A |
| D | | ≤ 4 | A |
| E | | | |
| F | | | |
| G | | | |
| H | | | |

Dijkstra's algorithm – paths from A

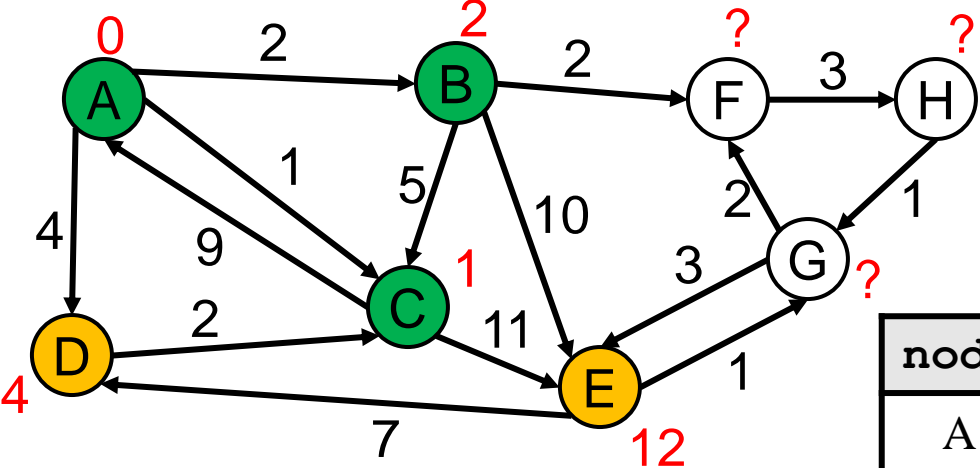


priority queue

| path | cost |
|-----------|------|
| [A, B] | 2 |
| [A, D] | 4 |
| [A, C, E] | 12 |
| | |

| node | finished | cost | prev |
|------|----------|-----------|------|
| A | Y | 0 | - |
| B | | ≤ 2 | A |
| C | Y | 1 | A |
| D | | ≤ 4 | A |
| E | | ≤ 12 | C |
| F | | | |
| G | | | |
| H | | | |

Dijkstra's algorithm – paths from A

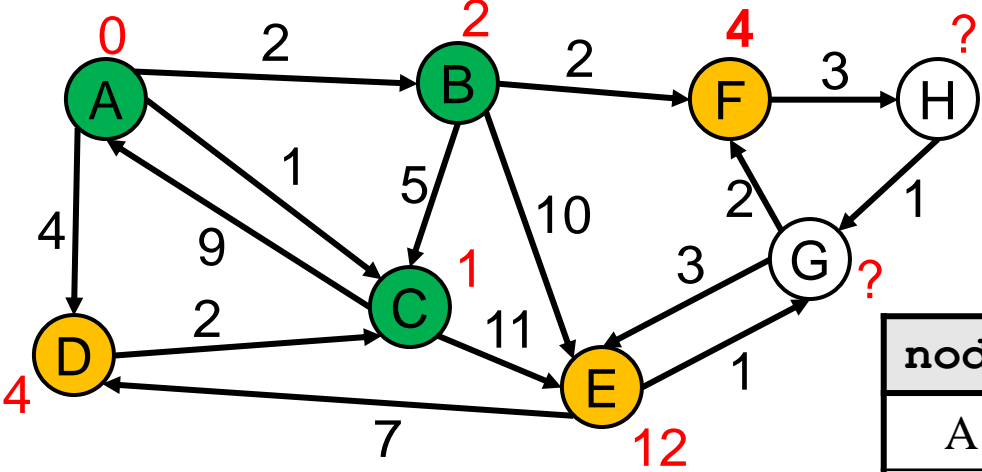


priority queue

| path | cost |
|-----------|------|
| [A, D] | 4 |
| [A, C, E] | 12 |
| | |
| | |

| node | finished | cost | prev |
|------|----------|-----------|------|
| A | Y | 0 | - |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | | ≤ 4 | A |
| E | | ≤ 12 | C |
| F | | | |
| G | | | |
| H | | | |

Dijkstra's algorithm – paths from A

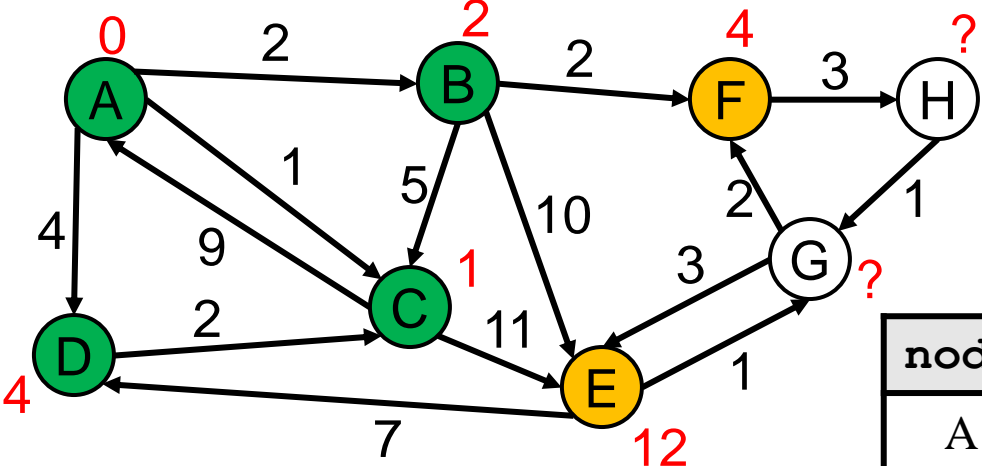


priority queue

| path | cost |
|-----------|------|
| [A, D] | 4 |
| [A, B, F] | 4 |
| [A, C, E] | 12 |
| [A, B, E] | 12 |

| node | finished | cost | prev |
|------|----------|------|----------|
| A | Y | 0 | - |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | | ≤ 4 | A |
| E | | ≤ 12 | C |
| F | | ≤ 4 | B |
| G | | | |
| H | | | |

Dijkstra's algorithm – paths from A

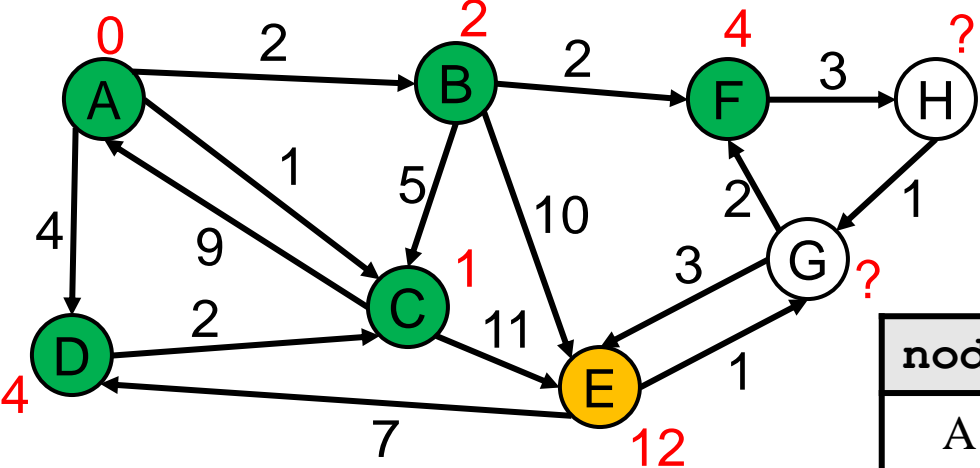


priority queue

| path | cost |
|-----------|------|
| [A, B, F] | 4 |
| [A, C, E] | 12 |
| [A, B, E] | 12 |
| | |

| node | finished | cost | prev |
|------|----------|-----------|------|
| A | Y | 0 | - |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | | ≤ 12 | C |
| F | | ≤ 4 | B |
| G | | | |
| H | | | |

Dijkstra's algorithm – paths from A

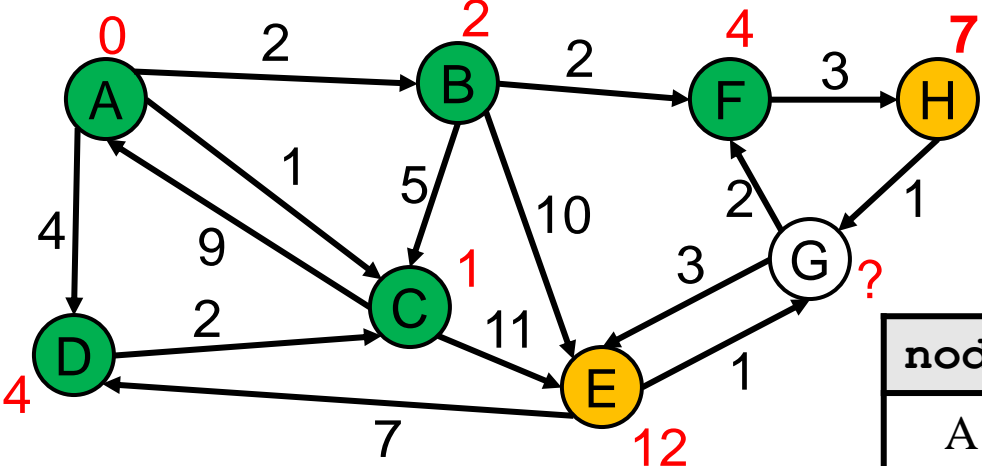


priority queue

| path | cost |
|-----------|------|
| [A, C, E] | 12 |
| [A, B, E] | 12 |
| | |
| | |

| node | finished | cost | prev |
|------|----------|-----------|------|
| A | Y | 0 | - |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | | ≤ 12 | C |
| F | Y | 4 | B |
| G | | | |
| H | | | |

Dijkstra's algorithm – paths from A

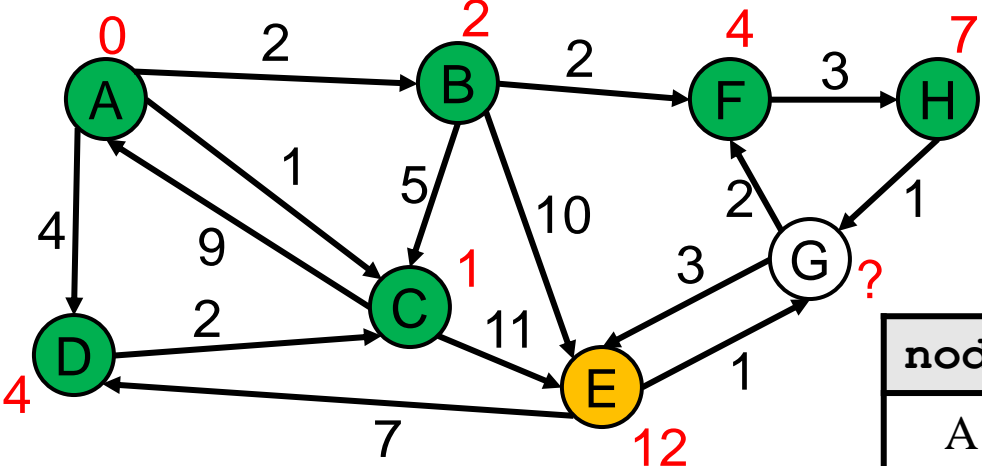


priority queue

| path | cost |
|--------------|------|
| [A, B, F, H] | 7 |
| [A, C, E] | 12 |
| [A, B, E] | 12 |
| | |

| node | finished | cost | prev |
|------|----------|-----------|------|
| A | Y | 0 | - |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | | ≤ 12 | C |
| F | Y | 4 | B |
| G | | | |
| H | | ≤ 7 | F |

Dijkstra's algorithm – paths from A

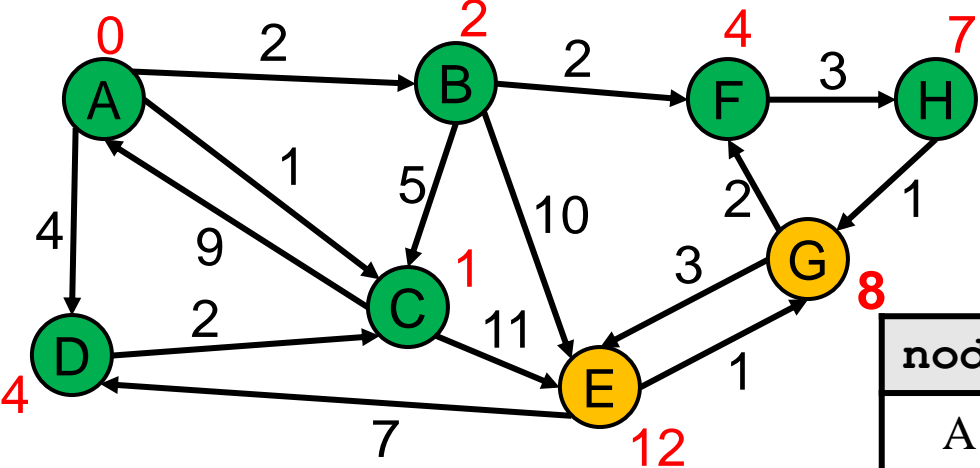


priority queue

| path | cost |
|-----------|------|
| [A, C, E] | 12 |
| [A, B, E] | 12 |
| | |
| | |

| node | finished | cost | prev |
|------|----------|-----------|------|
| A | Y | 0 | - |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | | ≤ 12 | C |
| F | Y | 4 | B |
| G | | | |
| H | Y | 7 | F |

Dijkstra's algorithm – paths from A

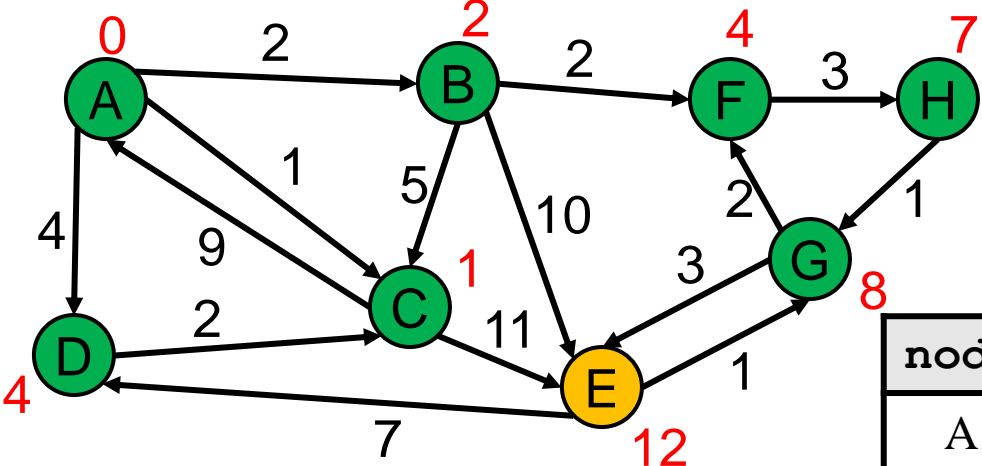


priority queue

| path | cost |
|-----------------|------|
| [A, B, F, H, G] | 8 |
| [A, C, E] | 12 |
| [A, B, E] | 12 |
| | |

| node | finished | cost | prev |
|------|----------|-----------|------|
| A | Y | 0 | - |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | | ≤ 12 | C |
| F | Y | 4 | B |
| G | | ≤ 8 | H |
| H | Y | 7 | F |

Dijkstra's algorithm – paths from A

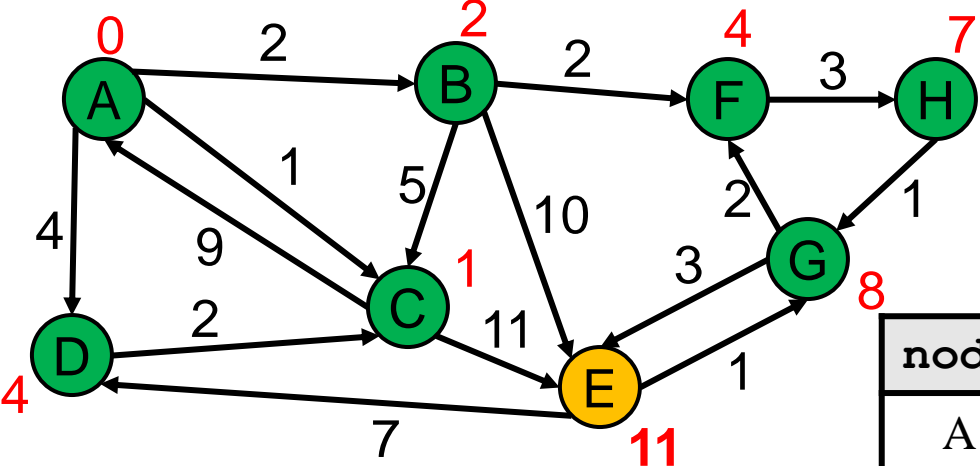


priority queue

| path | cost |
|-----------|------|
| [A, C, E] | 12 |
| [A, B, E] | 12 |
| | |
| | |

| node | finished | cost | prev |
|------|----------|------|------|
| A | Y | 0 | - |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | | ≤ 12 | C |
| F | Y | 4 | B |
| G | Y | 8 | H |
| H | Y | 7 | F |

Dijkstra's algorithm – paths from A

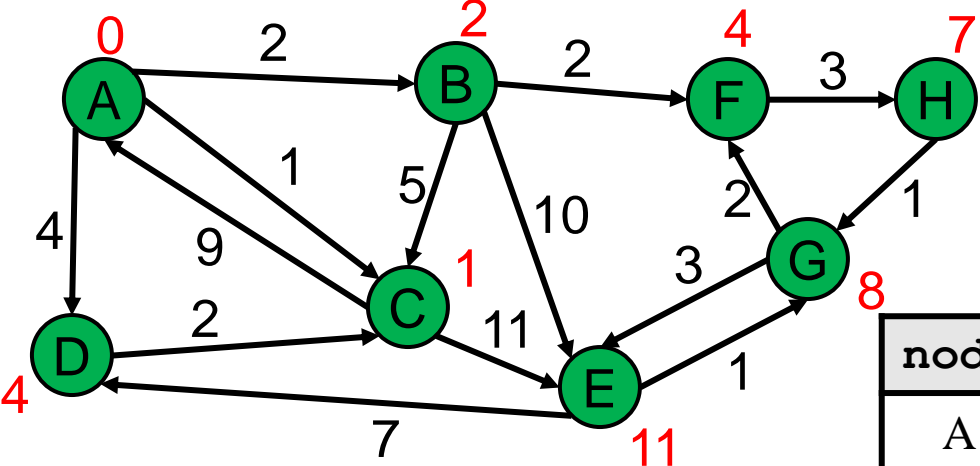


priority queue

| path | cost |
|--------------------|------|
| [A, B, F, H, G, E] | 11 |
| [A, C, E] | 12 |
| [A, B, E] | 12 |
| | |

| node | finished | cost | prev |
|------|----------|-----------|----------|
| A | Y | 0 | - |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | | ≤ 11 | G |
| F | Y | 4 | B |
| G | Y | 8 | H |
| H | Y | 7 | F |

Dijkstra's algorithm – paths from A

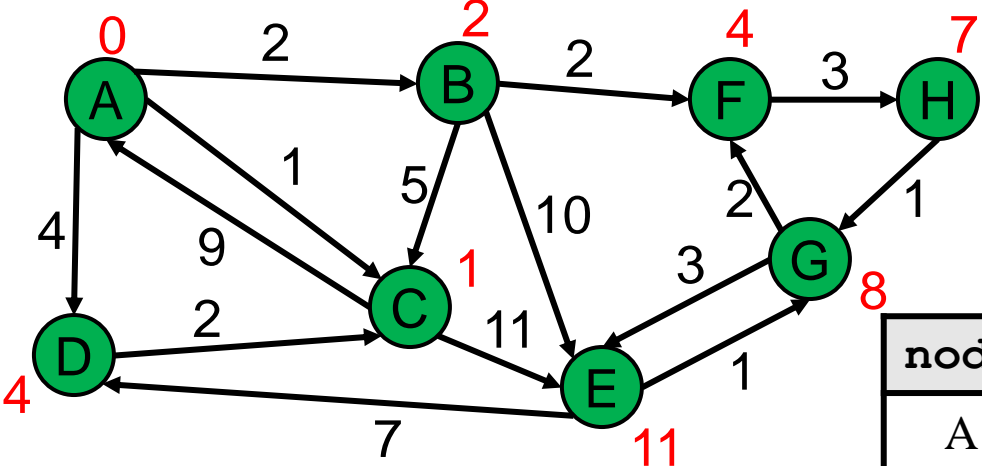


priority queue

| path | cost |
|-----------|------|
| [A, C, E] | 12 |
| [A, B, E] | 12 |
| | |
| | |

| node | finished | cost | prev |
|------|----------|------|------|
| A | Y | 0 | - |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | Y | 11 | G |
| F | Y | 4 | B |
| G | Y | 8 | H |
| H | Y | 7 | F |

Dijkstra's algorithm – paths from A

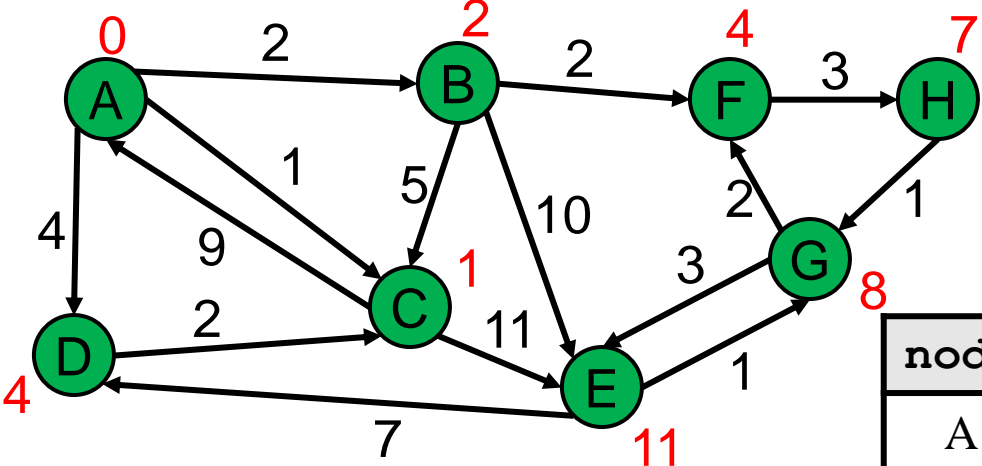


priority queue

| path | cost |
|-----------|------|
| [A, B, E] | 12 |
| | |
| | |
| | |

| node | finished | cost | prev |
|------|----------|------|------|
| A | Y | 0 | - |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | Y | 11 | G |
| F | Y | 4 | B |
| G | Y | 8 | H |
| H | Y | 7 | F |

Dijkstra's algorithm – paths from A



Now we know the cost and path to every single node by looking at the table!

priority queue

| path | cost |
|------|------|
| | |
| | |
| | |
| | |
| | |

| node | finished | cost | prev |
|------|----------|------|------|
| A | Y | 0 | - |
| B | Y | 2 | A |
| C | Y | 1 | A |
| D | Y | 4 | A |
| E | Y | 11 | G |
| F | Y | 4 | B |
| G | Y | 8 | H |
| H | Y | 7 | F |

Dijkstra's algorithm - Worksheet

- Now it's your turn!

Dijkstra's algorithm – pseudocode

```
active = priority queue of paths.  
finished = empty set of nodes.  
add a path from start to itself to active  
<inv: All paths found so far are shortest paths>  
while active is non-empty:  
    minPath = active.removeMin()  
    minDest = destination node in minPath  
    if minDest is dest:  
        return minPath  
    if minDest is in finished:  
        continue  
    for each edge e = (minDest, child):  
        if child is not in finished:  
            newPath = minPath + e  
            add newPath to active  
    add minDest to finished
```

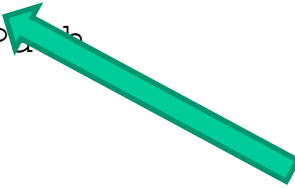
Dijkstra's algorithm – pseudocode

```
active = priority queue of paths.  
finished = empty set of nodes.  
add a path from start to itself to active  
<inv: All paths found so far are shortest paths>  
while active is non-empty:  
    minPath = active.removeMin()  
    minDest = destination node in minPath  
    if minDest is dest:  
        return minPath  
    if minDest is in finished:  
        continue  
    for each edge e = (minDest, child):  
        if child is not in finished:  
            newPath = minPath + e  
            add newPath to active  
    add minDest to finished
```

What else?

Dijkstra's algorithm – pseudocode

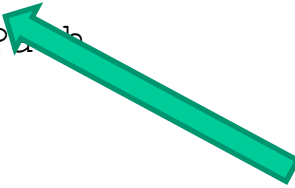
```
active = priority queue of paths.  
finished = empty set of nodes.  
add a path from start to itself to active  
<inv: All paths found so far are shortest paths>  
while active is non-empty:  
    minPath = active.removeMin()  
    minDest = destination node in minPath  
    if minDest is dest:  
        return minPath  
    if minDest is in finished:  
        continue  
    for each edge e = (minDest, child):  
        if child is not in finished:  
            newPath = minPath + e  
            add newPath to active  
    add minDest to finished
```




All nodes not reached yet are farther away than those reached so far

Dijkstra's algorithm – pseudocode

```
active = priority queue of paths.  
finished = empty set of nodes.  
add a path from start to itself to active  
<inv: All paths found so far are shortest paths>  
while active is non-empty:  
    minPath = active.removeMin()  
    minDest = destination node in minPath  
    if minDest is dest:  
        return minPath  
    if minDest is in finished:  
        continue  
    for each edge e = (minDest, child):  
        if child is not in finished:  
            newPath = minPath + e  
            add newPath to active  
    add minDest to finished
```



All nodes not reached yet are farther away than those reached so far



The queue contains all paths formed by adding 1 more edge to a node we already reached.

Dijkstra's algorithm – pseudocode

```
active = priority queue of paths.  
finished = empty set of nodes.  
add a path from start to itself to active  
<inv: All paths found so far are shortest paths & ...>  
while active is non-empty:  
    minPath = active.removeMin()  
    minDest = destination node in minPath  
    if minDest is dest:  
        return minPath  
    if minDest is in finished:  
        continue  
    for each edge e = (minDest, child):  
        if child is not in finished:  
            newPath = minPath + e  
            add newPath to active  
    add minDest to finished
```



Let's take a
moment to think
what else is true
here?

Dijkstra's algorithm – pseudocode

active = priority queue of paths.

finished = empty set of nodes.

add a path from start to itself to active

<inv: All paths found so far are shortest paths & ...>

while active is non-empty:

minPath = **active.removeMin()**

 minDest = destination node in minPath

 if minDest is dest:

 return minPath

 if minDest is in finished:

 continue

 for each edge e = (minDest, child):

 if child is not in finished:

 newPath = minPath + e

 add newPath to active

 add minDest to finished



It follows from our updated invariant that this path is the shortest path (assuming node is not in finished)

Script testing in HW7

- Extends the test-script mechanism from HW5/6
 - Using numeric weights instead of string labels on edges
 - New command **FindPath** to find shortest path with Dijkstra's algorithm
- Must write the test driver (**PathfinderTestDriver**) yourself
 - Feel free to copy pieces from **GraphTestDriver** in HW5/6

| Command (in <i>foo.test</i>) | Output (in <i>foo.expected</i>) |
|--|--|
| FindPath <i>graph node₁ node_n</i> | path from node₁ to node_n: <i>node₁ to node₂ with weight $w_{1,2}$</i> <i>node₂ to node₃ with weight $w_{2,3}$</i> ... <i>node_{n-1} to node_n with weight $w_{n-1,n}$</i> total cost: <i>w</i> |
| ... | ... |

Model-View-Controller

Model-View-Controller

- Model-View-Controller (MVC) is a ubiquitous design pattern:
 - The **model** abstracts + represents the application's data.
 - The **view** provides a user interface to display the application data.
 - The **controller** handles user input to affect the application.

Model-View-Controller: Example

Accessing my Google Drive files through my laptop and my phone

| Laptop | Phone |
|---|--|
| View: The screen displays options for me to select files | |
| Control: Get input selection from mouse/keyboard | Control: Get input selection from touch sensor |
| Control: Request the selected file from Google Drive | |
| Model: Google Drive sends back the request file to my device | |
| Control: Receive the file and pass it to View | |
| View: The screen displays the file | |

HW7: text-based View-Controller

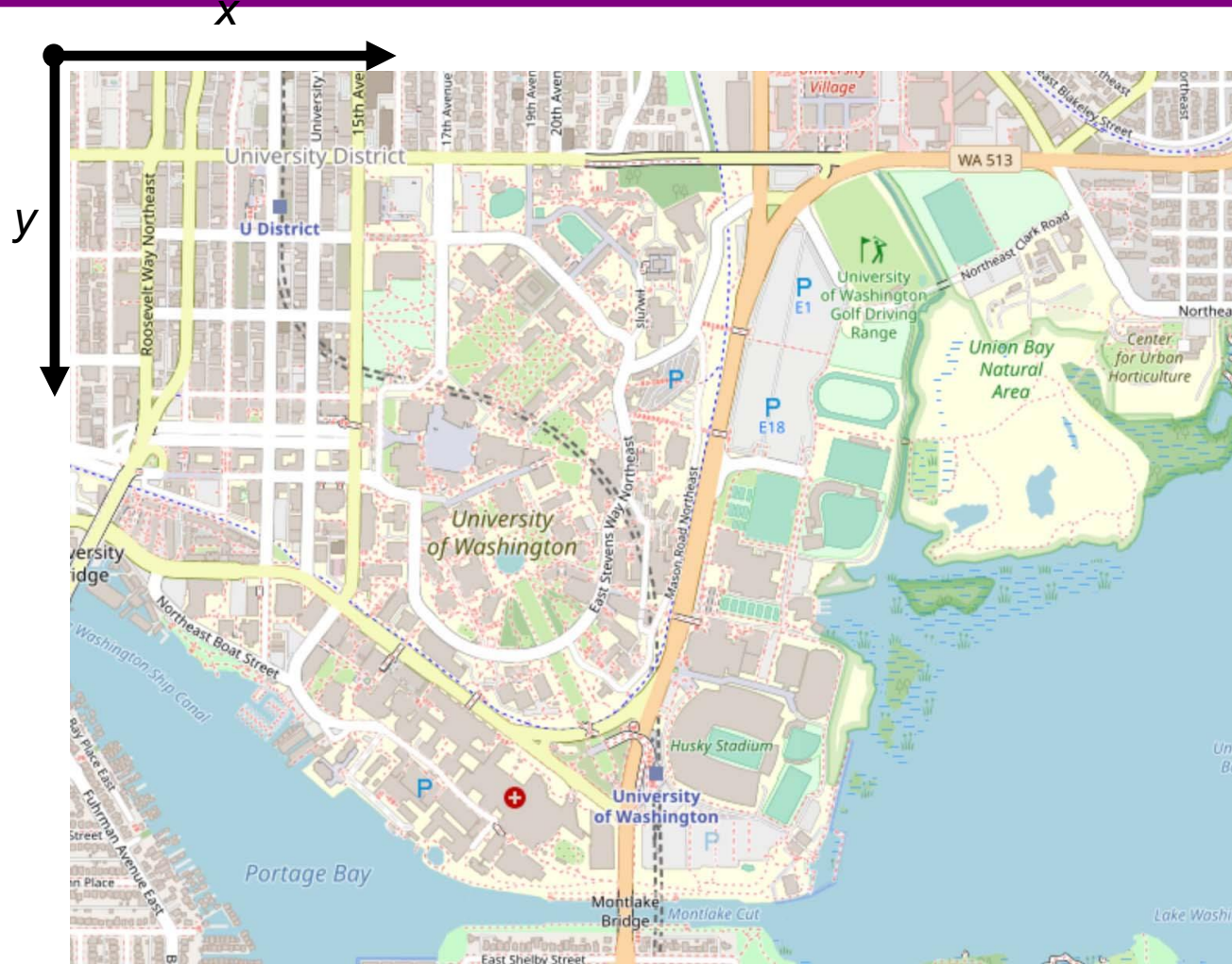
- **TextInterfaceView**
 - Displays output to users from the result received from **TextInterfaceController**.
 - Receives input from users.
 - Does not process anything; directly pass the input to the **TextInterfaceController** to process.
- **TextInterfaceController**
 - Process the passed input from the **TextInterfaceView**
 - Include talking to the **Model** (the graph & supporting code)
 - Give the processed result back to the **TextInterfaceView** to display to users.

* HW9 will be using the same **Model** but different and more sophisticated View and Controller

Campus dataset

- Two CSV files in `src/main/resources/data`:
 - `campus_buildings.csv` – building entrances on campus
 - `campus_paths.csv` – straight-line walkways on campus
- Exact points on campus identified with (x, y) coordinates
 - Pixels on a map of campus (`campus_map.jpg`, next to CSV files)
 - Position $(0, 0)$, the origin, is the top left corner of the map
- Parser in starter code: `pathfinder.parser.CampusPathsParser`
 - `CampusBuilding` object for each entry of `campus_buildings.csv`
 - `CampusPath` object for each entry of `campus_paths.csv`

Campus dataset – coordinate plane



campus_map.jpg

Campus dataset – sample

- **campus_buildings.CSV** has entries like the following:

| <i>shortName</i> | <i>longName</i> | <i>x</i> | <i>y</i> |
|------------------|-----------------|------------|-----------|
| BGR, | By George, | 1671.5499, | 1258.4333 |
| MOR, | Moore Hall, | 2317.1749, | 1859.502 |

- **campus_paths.CSV** has entries like the following:

| <i>x1</i> | <i>y1</i> | <i>x2</i> | <i>y2</i> | <i>distance</i> |
|-----------|-----------|------------|------------|-----------------|
| 1810.0, | 431.5, | 1804.6429, | 437.92857, | 17.956615... |
| 1810.0, | 431.5, | 1829.2857, | 409.35714, | 60.251364... |

- See **campus_routes.jpg** for nice visual rendering of **campus_paths.csv**

Campus dataset – demo

- Let's go through the starter files of HW 7.

HW 7 – Model-View-Controller

- HW7 is an MVC application, with much given as starter code.
 - View: `pathfinder.textInterface.TextInterfaceView`
 - Controller: `pathfinder.textInterface.TextInterfaceController`
- You will need to fill out the code in `pathfinder.CampusMap`.
 - Since your code implements the model functionality

Before next lecture...

1. Do [HW6](#) by tonight!
 - No written portion
 - Coding portion (push and tag on GitLab)
2. Feel free to add additional JUnit tests or script tests!

HW 7 – Model-View-Controller

- HW7 is an MVC application, with much given as starter code.
 - View: `pathfinder.textInterface.TextInterfaceView`
 - Controller: `pathfinder.textInterface.TextInterfaceController`
- You will need to fill out the code in `pathfinder.CampusMap`.
 - Since your code implements the model functionality