



# CSE332: Data Abstractions Lecture 1: Introduction; ADTs; Stacks/Queues

Dan Grossman Spring 2012

# Welcome!

We have 10 weeks to learn fundamental data structures and algorithms for organizing and processing information - "Classic" data structures / algorithms and how to analyze rigorously their efficiency and when to use them - Queues, dictionaries, graphs, sorting, etc. Parallelism and concurrency (!) Today in class: · Course mechanics · What this course is about And how it fits into the CSE curriculum · Start (finish?) ADTs, stacks, and queues - Largely review Spring 2012 CSE332: Data Abstractions 2

### Concise to-do list

### In next 24-48 hours:

- Adjust class email-list settings
- Email homework 0 (worth 0 points) to me
- · Read all course policies
- Read/skim Chapters 1 and 3 of Weiss book
  - Relevant to Project 1, due next week
  - Will start Chapter 2 on Wednesday

### Possibly:

- Set up your Eclipse / Java environment for Project 1
  - Thursday's section will help

http://www.cs.washington.edu/education/courses/cse332/12sp/

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# Communication

- Course email list: cse332a sp12@u
  - Students and staff already subscribed
  - You must get announcements sent there
  - Fairly low traffic
- Course staff: cse332-staff@cs plus individual emails
- Discussion board
  - For appropriate discussions; TAs will monitor
  - Optional, won't use for important announcements
- Anonymous feedback link •
  - For good and bad: if you don't tell me, I don't know

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3

### Course staff

Dan Grossman

Tyler Robison







Stanley Wang

Dan: Faculty, "341 guy", loves 332 too, did parallelism/concurrency part Tyler: Grad student, TAed 332 many times, taught it Summer 2010 Stanley: Took 332 last quarter

### Office hours, email, etc. on course web-page

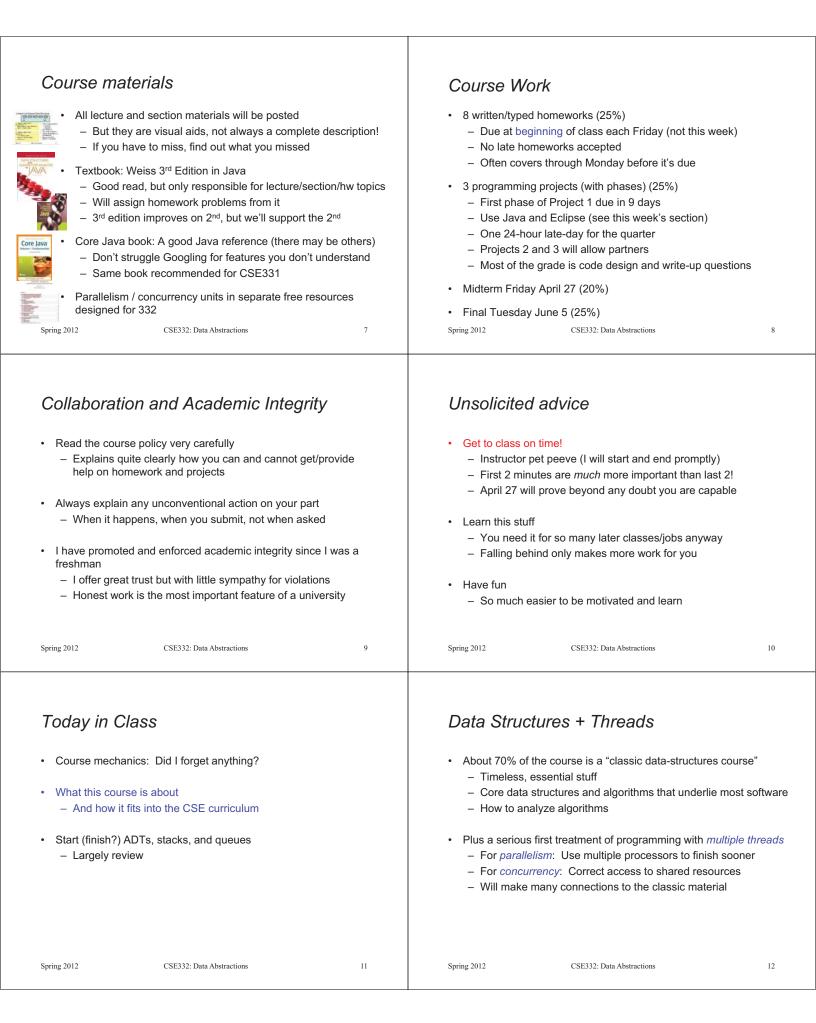
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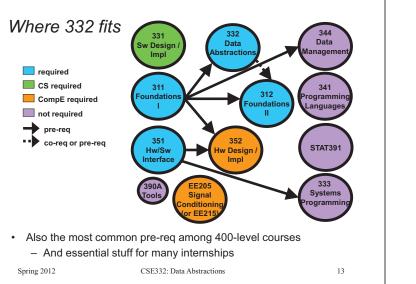
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4

# Course meetings

- · Lecture (Dan)
  - Materials posted (sometimes afterwards), but take notes
  - Ask questions, focus on key ideas (rarely coding details)
- Section (Tyler)
  - Often focus on software (Java features, tools, project issues)
  - Reinforce key issues from lecture
  - Answer homework questions, etc.
  - An important part of the course (not optional)
- Office hours
  - Use them: please visit me
  - Ideally not just for homework questions (but that's great too)





# What is 332 is about

- · Deeply understand the basic structures used in all software
  - Understand the data structures and their trade-offs
  - Rigorously analyze the algorithms that use them (math!)
  - Learn how to pick "the right thing for the job"
- · Experience the purposes and headaches of multithreading
- · Practice design, analysis, and implementation
  - The elegant interplay of "theory" and "engineering" at the core of computer science

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#### 14

### Goals

- Be able to make good design choices as a developer, project manager, etc.
  - Reason in terms of the general abstractions that come up in all non-trivial software (and many non-software) systems
- Be able to justify and communicate your design decisions

### Dan's take:

- 3 years from now this course will seem like it was a waste of your time because you can't imagine not "just knowing" every main concept in it
- Key abstractions computer scientists and engineers use almost every day
- A big piece of what separates us from others

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15

# Trade-offs

A data structure strives to provide many useful, efficient operations

### But there are unavoidable trade-offs:

- Time vs. space
- One operation more efficient if another less efficient
- Generality vs. simplicity vs. performance

That is why there are many data structures and educated CSEers internalize their main trade-offs and techniques

- And recognize logarithmic < linear < quadratic < exponential

### Data structures

- (Often highly *non-obvious*) ways to organize information to enable *efficient* computation over that information
  - Key goal over the next week is introducing asymptotic analysis to precisely and generally describe efficient use of time and space

A data structure supports certain operations, each with a:

- Meaning: what does the operation do/return
- Performance: how efficient is the operation

### Examples:

- List with operations insert and delete
- Stack with operations push and pop

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16

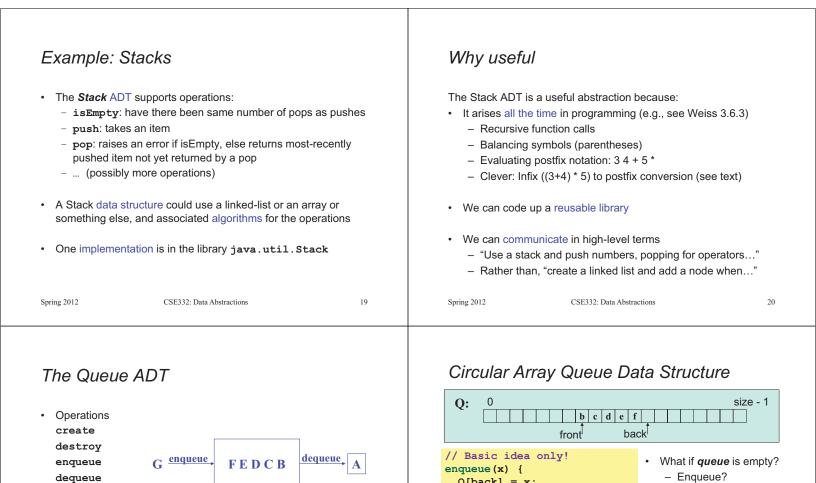
# Terminology

- Abstract Data Type (ADT)
  - Mathematical description of a "thing" with set of operations
- Algorithm

 A high level, language-independent description of a step-bystep process

- Data structure
  - A specific family of algorithms for implementing an ADT
- Implementation of a data structure
  - A specific implementation in a specific language

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# Linked List Queue Data Structure

is\_empty

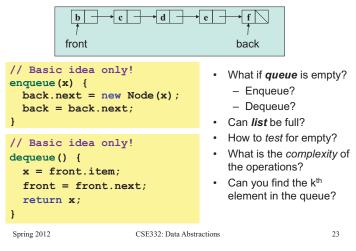
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Just like a stack except:

Stack: LIFO (last-in-first-out)

Just as useful and ubiquitous

- Queue: FIFO (first-in-first-out)



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# Circular Array vs. Linked List

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- Dequeue?

the operations?

Can you find the kth

element in the queue?

22

What if array is full?

How to *test* for empty?

What is the complexity of

	back	Array: – May waste unneeded space or run out of space	List: – Always just enough space – But more space per element
);	<ul> <li>What if <i>queue</i> is empty?</li> <li>Enqueue?</li> <li>Dequeue?</li> <li>Can <i>list</i> be full?</li> </ul>	<ul> <li>Space per element excellen</li> <li>Operations very simple / fas</li> <li>Constant-time access to k<sup>th</sup> element</li> </ul>	t – Operations very simple / fast
	<ul> <li>How to <i>test</i> for empty?</li> <li>What is the <i>complexity</i> of the operations?</li> <li>Can you find the k<sup>th</sup> element in the queue?</li> </ul>	<ul> <li>For operation insertAtPosition must shift all later elements</li> <li>Not in Queue ADT</li> <li>This is something every train sleep – it's like knowing here</li> </ul>	must traverse all earlier elements – Not in Queue ADT ed computer scientist knows in his/her
Data Abstra	ctions 23		32: Data Abstractions 24

Q[back] = x;

dequeue() {

return x;

// Basic idea only!

x = Q[front];

}

}

21

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back = (back + 1) % size

front = (front + 1) % size;

