# CSE 333

## Lecture 21 -- fork, pthread_create

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# Administrivia

pthreads exercise due Monday morning

HW4 is due Thursday night

- **<panic>** if you haven't started yet **</panic>**
- Usual late days apply ***if*** you have any left

Final exam Wed., Dec. 13, 2:30 pm

- Some review in section next week; last-minute review Q&A Tue. 12/12, 4:30 pm, EEB 045 (bring questions!)
- Topic list and old exams on web now

# Administrivia (Monday)

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Please fill out course evals this week

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# Some common HW4 bugs

Your server works, but is really really slow

- check the 2nd argument to the QueryProcessor constructor

Funny things happen after the first request

- make sure you're not destroying the HTTPConnection object too early (e.g., falling out of scope in a while loop)

Server crashes on blank request

- make sure you handle the case that read( ) [or WrappedRead] returns 0

# Previously

We implemented a searchserver, but it was sequential

- it processed requests one at a time, in spite of client interactions blocking for arbitrarily long periods of time
    - this led to terrible performance

Servers should be concurrent

- process multiple requests simultaneously
    - issue multiple I/O requests simultaneously
    - overlap the I/O of one request with computation of another
    - utilize multiple CPUs / cores

# Today

We'll go over three versions of searchserver

- sequential
- concurrent
  - processes [ **fork( )** ]
  - threads [ **pthread_create( )** ]

Alternative (which we won't get to): non-blocking, event driven version

  - non-blocking I/O [ **select( )** ]

Reference: *Computer Systems: A Programmer's Perspective*

- 351 textbook: good source for process/thread/OS concepts

# Sequential

pseudocode:

```
listen_fd = Listen(port);

while(1) {
  client_fd = accept(listen_fd);
  buf = read(client_fd);
  resp = ProcessQuery(buf);
  write(client_fd, resp);
  close(client_fd);
}
```

*look at **searchserver_sequential/***

# Whither sequential?

Benefits

- super simple to build

Disadvantages

- incredibly poorly performing
  - one slow client causes **all** others to block
  - poor utilization of network, CPU, disks

# fork( )

`pid_t fork(void);`

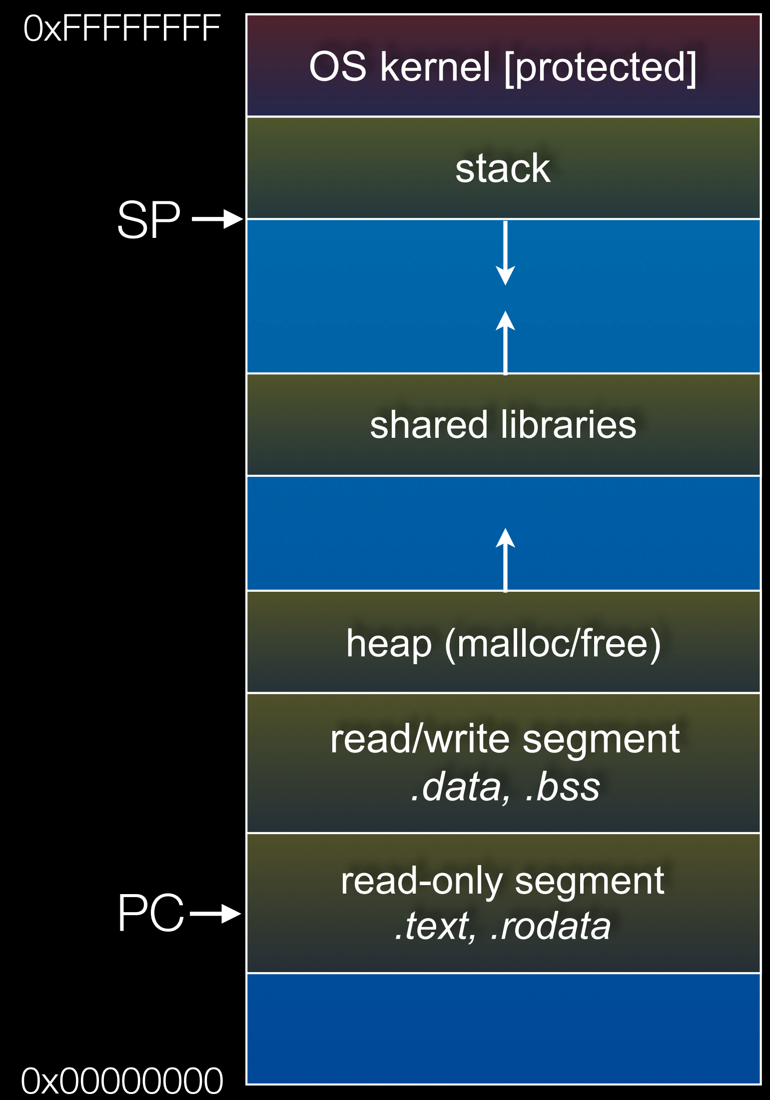
Fork is used to create a new process (the "child") that is an exact clone of the current process (the "parent")

- everything is cloned (except threads)
  - all variables, file descriptors, open sockets, etc.
  - the heap, the stack, etc.
- primarily used in two patterns
  - servers: fork a child to handle a connection
  - shells: fork a child, which then exec's a new program

# fork( ) and address spaces

Remember this picture...?

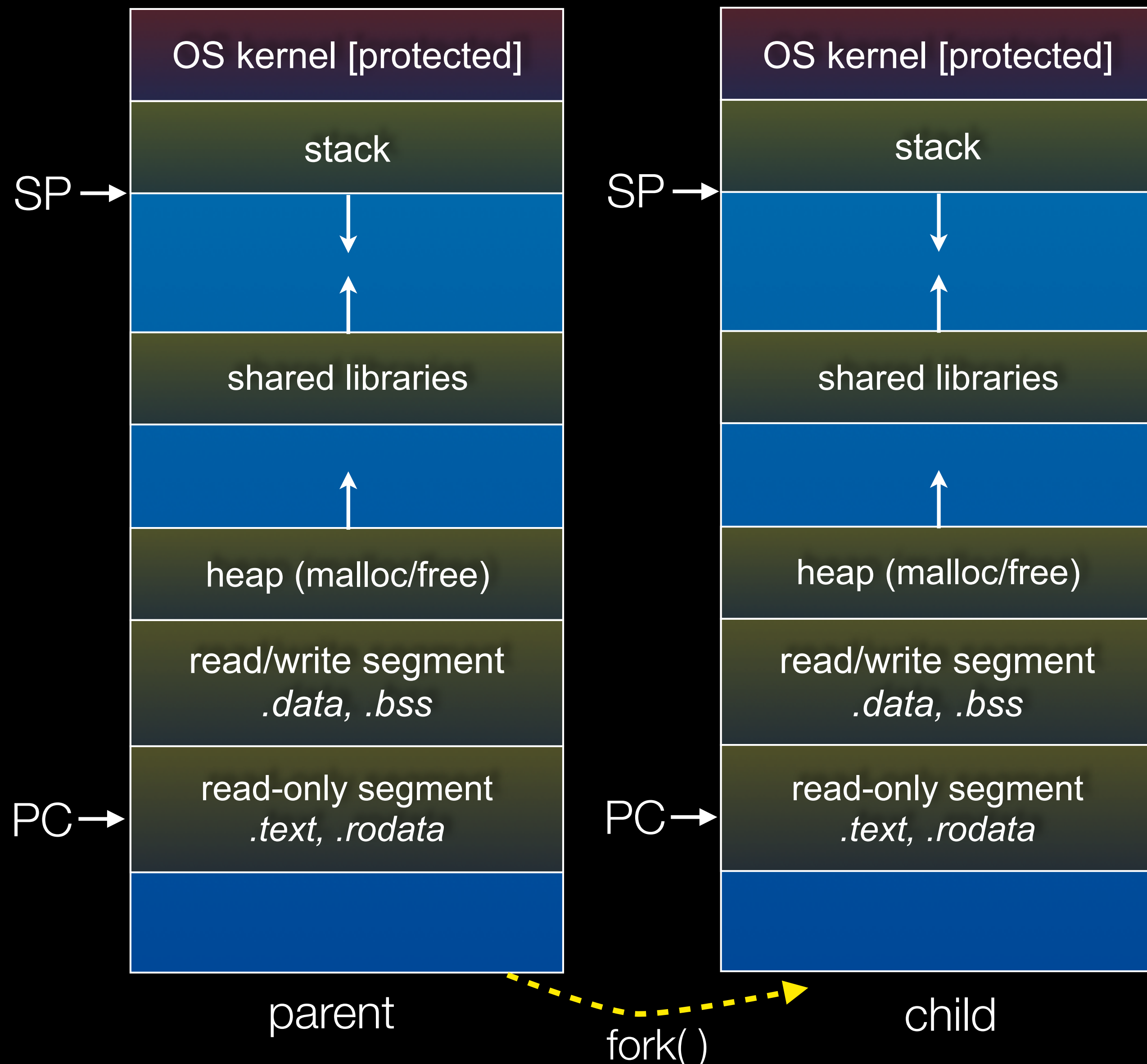
- a process executes within an ***address space***
- the address space includes:
    - a stack (for stack frames)
    - heap (for dynamically allocated data)
    - text segment (containing code)
    - etc.



# fork( ) and address spaces

Fork causes the OS to clone the address space, creating a brand new process

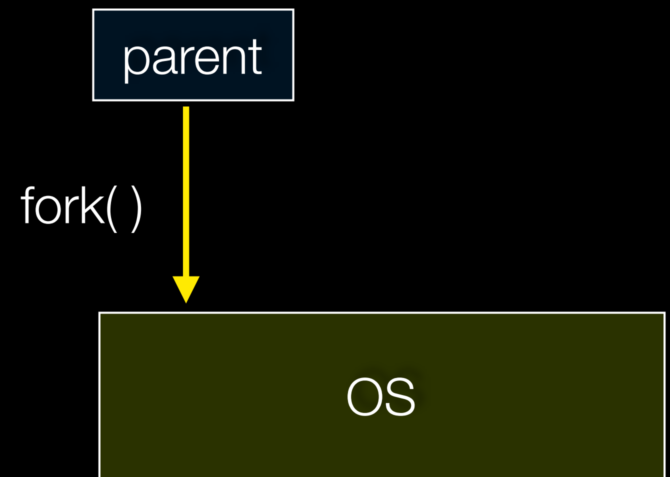
- the new process starts life as a **copy** the old process in (nearly) every way
- the **copies** of the heap, stack, text segment, etc. are (nearly) identical
- the new process has **copies** of the parent's data structures, stack-allocated variables, open file descriptors, and so on



# fork( )

## fork( ) has peculiar semantics

- the parent invokes fork( )
- the operating system clones the parent
- **both** the parent and the child return from fork
  - parent receives child's pid
  - child receives a "0" as pid



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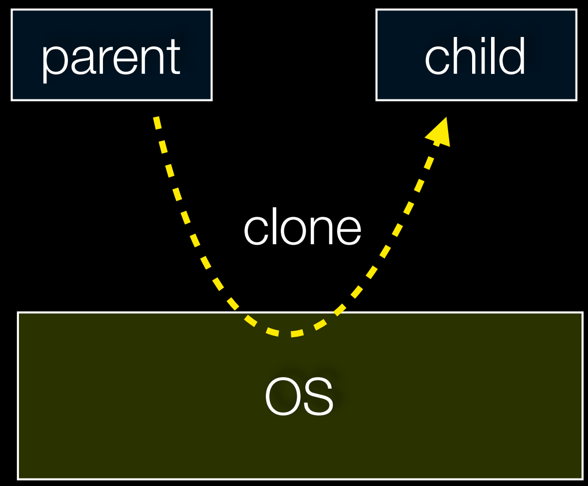

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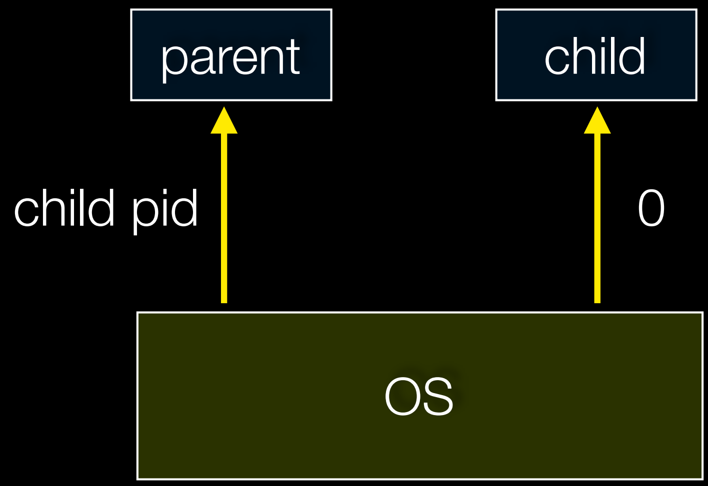

# fork( )

*fork_example.cc*

# Concurrency with processes

The ***parent*** process blocks on **accept( )**, waiting for a new client to connect

- when a new connection arrives, the parent calls **fork( )** to create a ***child*** process
- the child process handles that new connection, and **exit( )**'s when the connection terminates

Remember that children become "zombies" after death

- option a) parent calls wait( ) to "reap" children
- option b) use the double-fork trick

# Graphically

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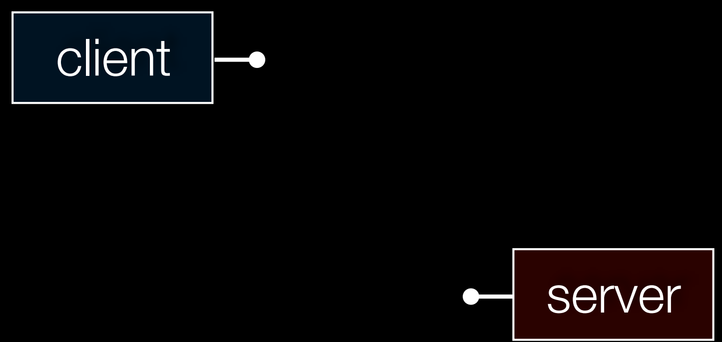

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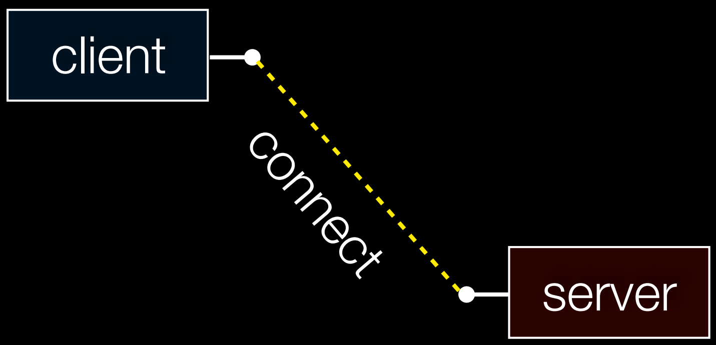

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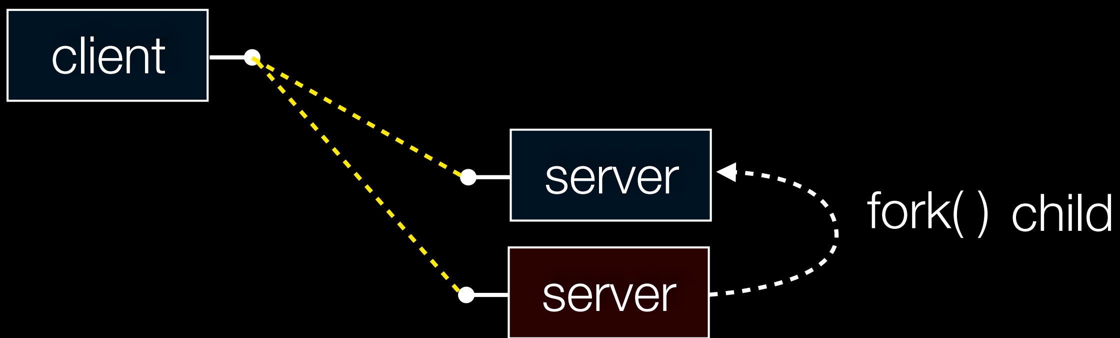

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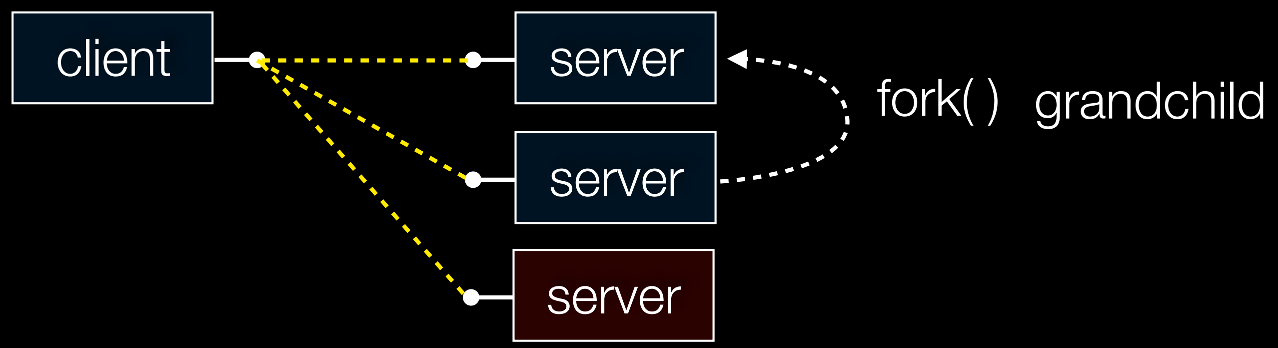

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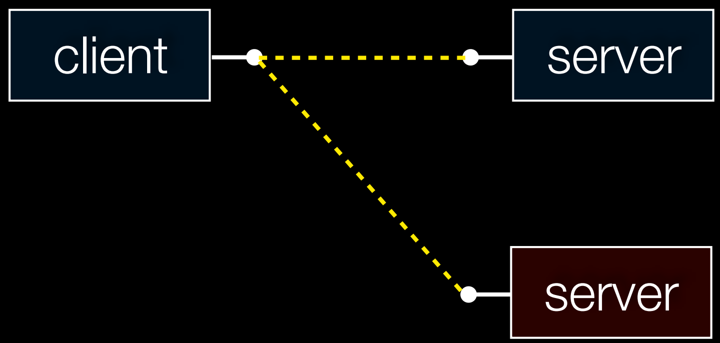
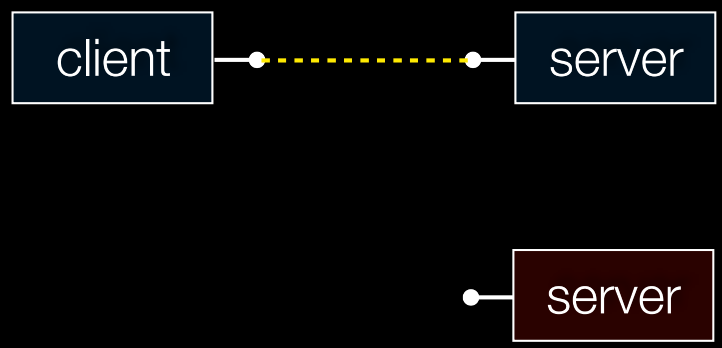


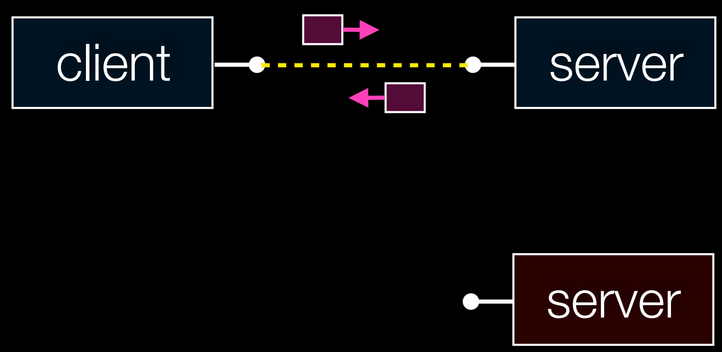

child exit( )'s / parent wait( )'s

# Graphically




parent closes its
client connection

# Graphically



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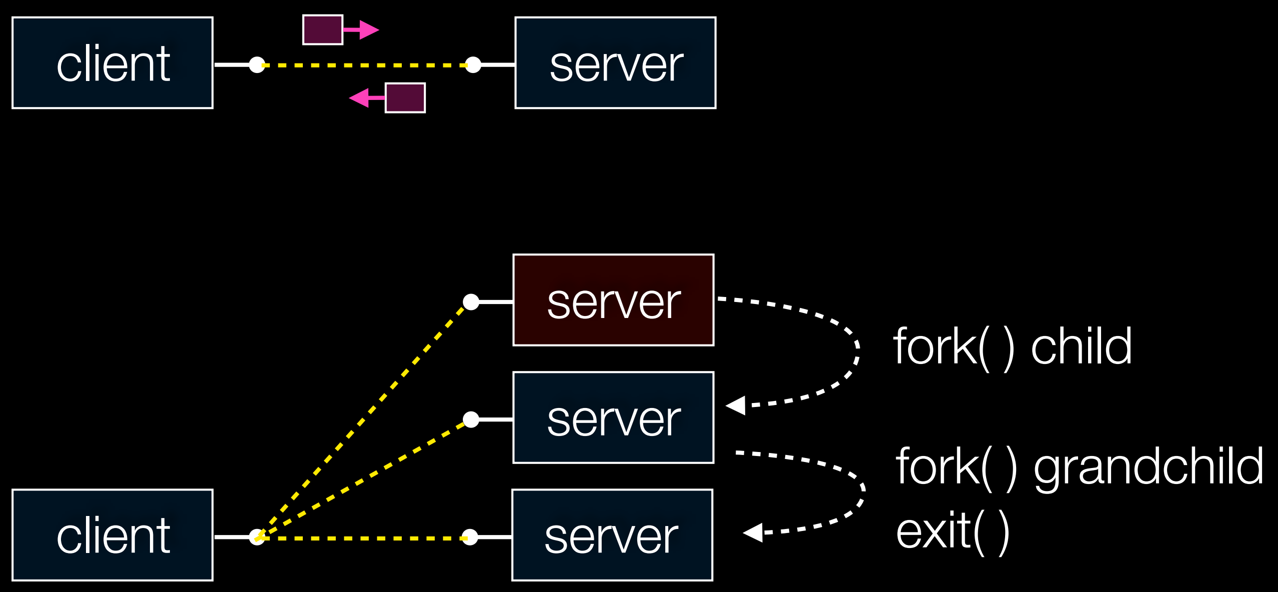

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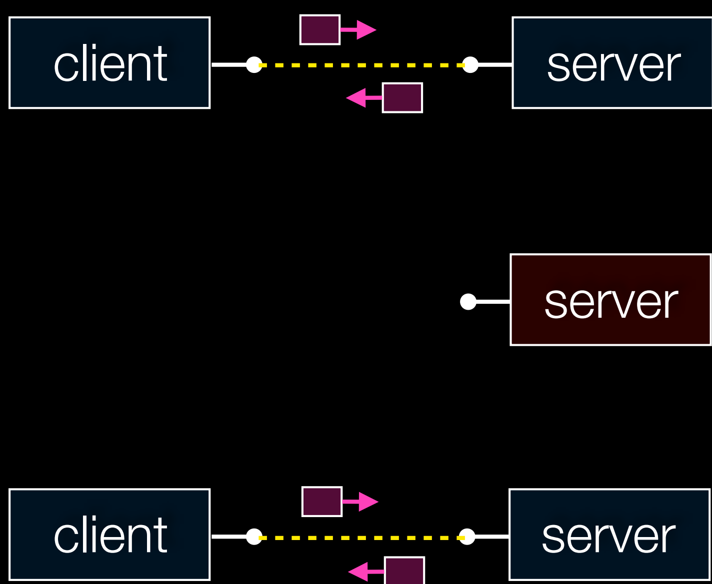

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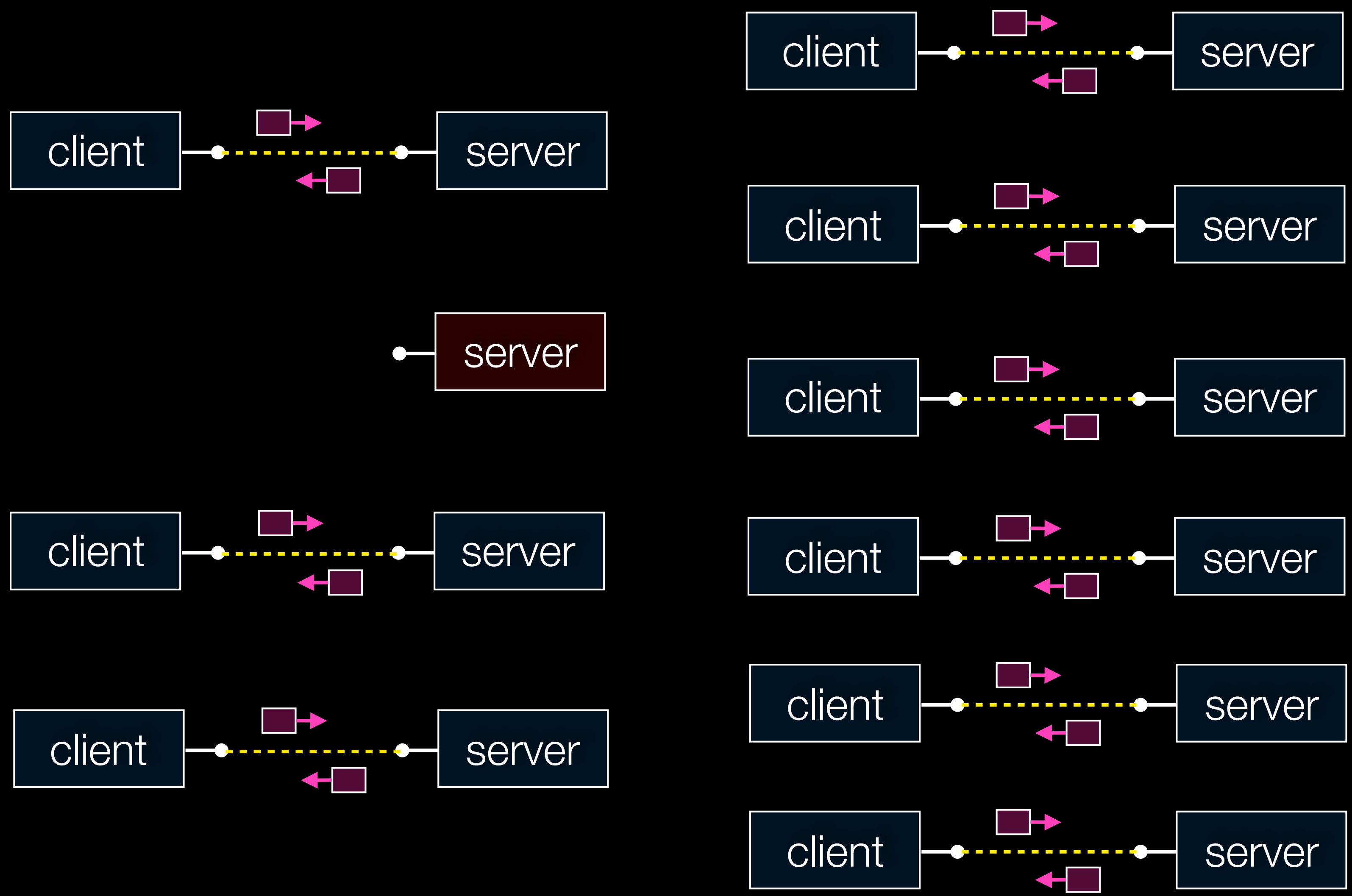


# Concurrent with processes

*look at **searchserver_processes***

# Whither concurrent processes?

## Benefits

- almost as simple as sequential
  - in fact, most of the code is identical!
- parallel execution; good CPU, network utilization

## Disadvantages

- processes are heavyweight
  - relatively slow to fork
  - context switching latency is high
- communication between processes is complicated

# How slow is fork?

*run* ***forklatency.cc***

# Implications?

**0.25 ms** per fork

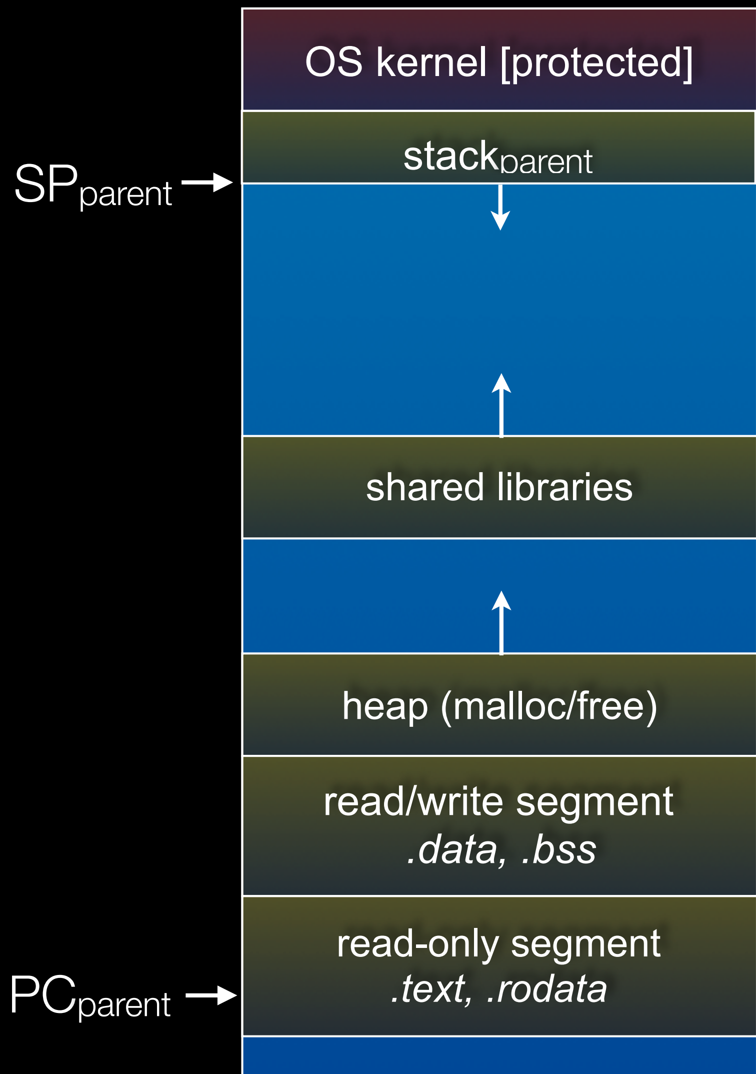
- maximum of (1000 / 0.25) = 4,000 connections per second per core
- ~0.5 billion connections per day per core
  - fine for most servers
  - too slow for a few super-high-traffic front-line web services
    - Facebook served *O*(750 billion) page views per day 4 years ago!
    - would need 3,000 -- 6,000 cores just to handle fork( ), i.e., without doing any work for each connection!

# threads

Threads are like lightweight processes

- like processes, they execute concurrently
    - multiple threads can run simultaneously on multiple cores/CPUs
- unlike processes, threads cohabit the same address space
    - the threads within a process see the same heap and globals
        - threads can communicate with each other through variables
        - but, threads can interfere with each other: need synchronization
    - each thread has its own stack

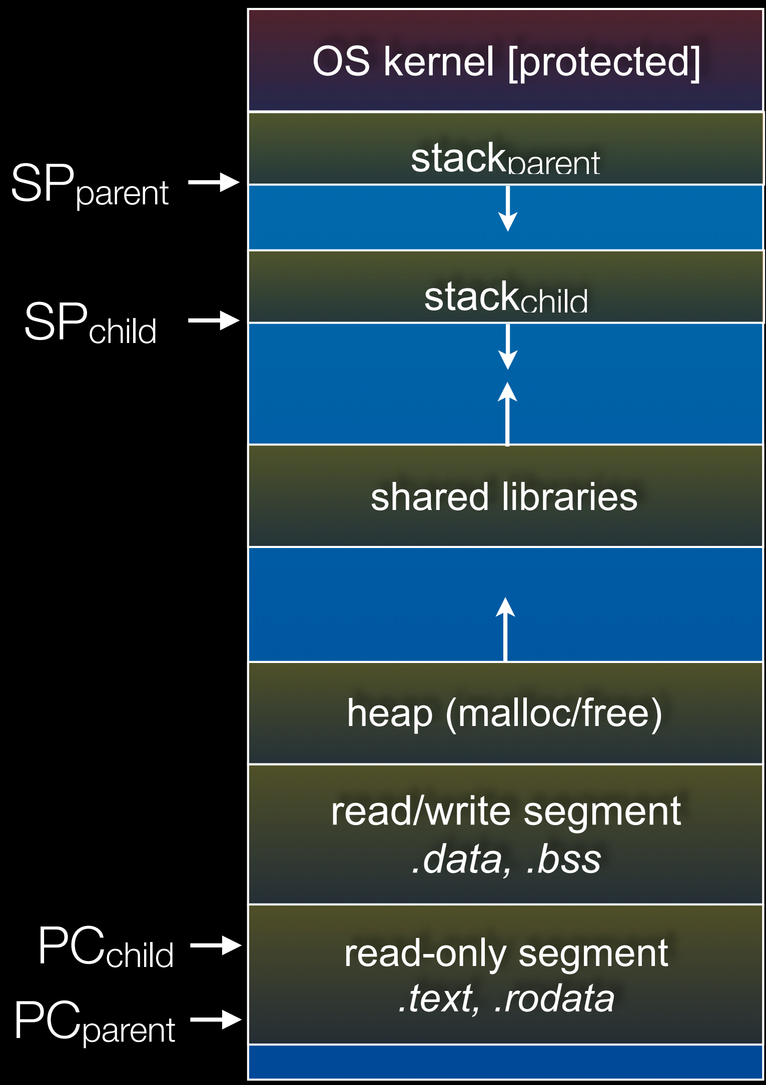
# threads and the address space




Pre- thread create

- one thread of execution running in the address space
  - the “main” thread
  - therefore, one stack, SP, PC
- that main thread invokes a function to create a new thread
  - typically “pthread_create( )”

# threads and the address space




## Post- thread create

- two threads of execution running in the address space
  - the "main" thread (parent)
  - the child thread
  - thus, two stacks, SPs, PCs
- both threads share the heap and text segment (globals)
  - they can cooperatively modify shared data

# threads

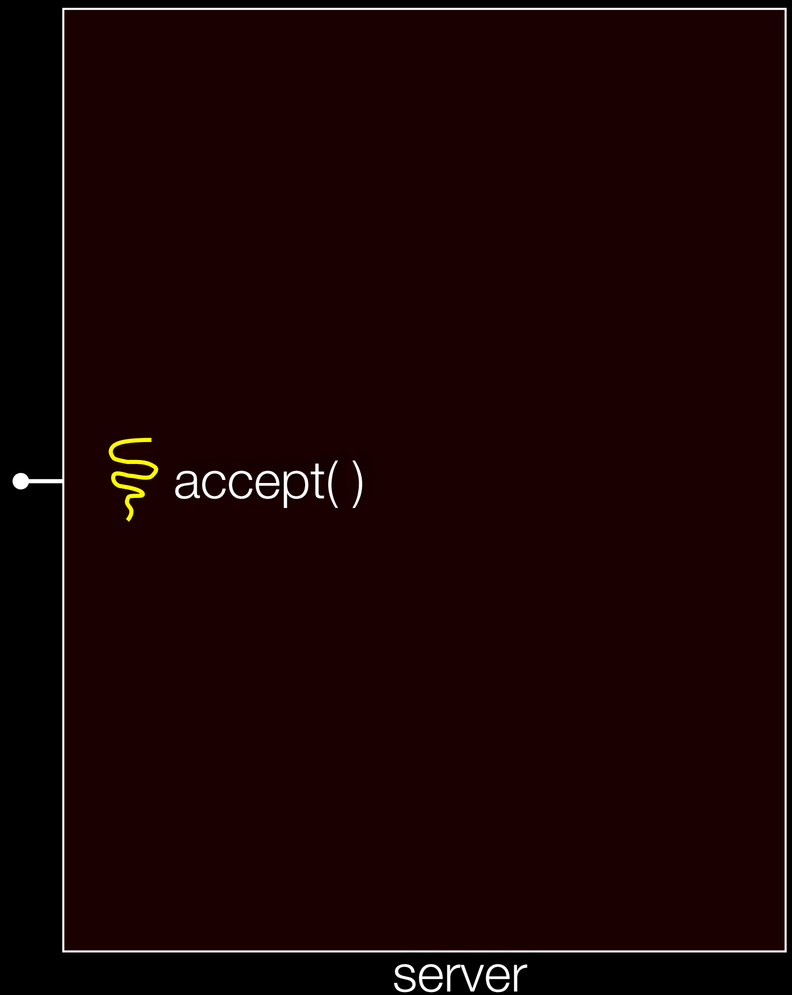
*see thread_example.cc*

# Concurrent server with threads

A single ***process*** handles all of the connections

- but, a parent ***thread*** forks (or dispatches) a new thread to handle each connection
- the child thread:
  - handles the new connection
  - exits when the connection terminates

# Graphically



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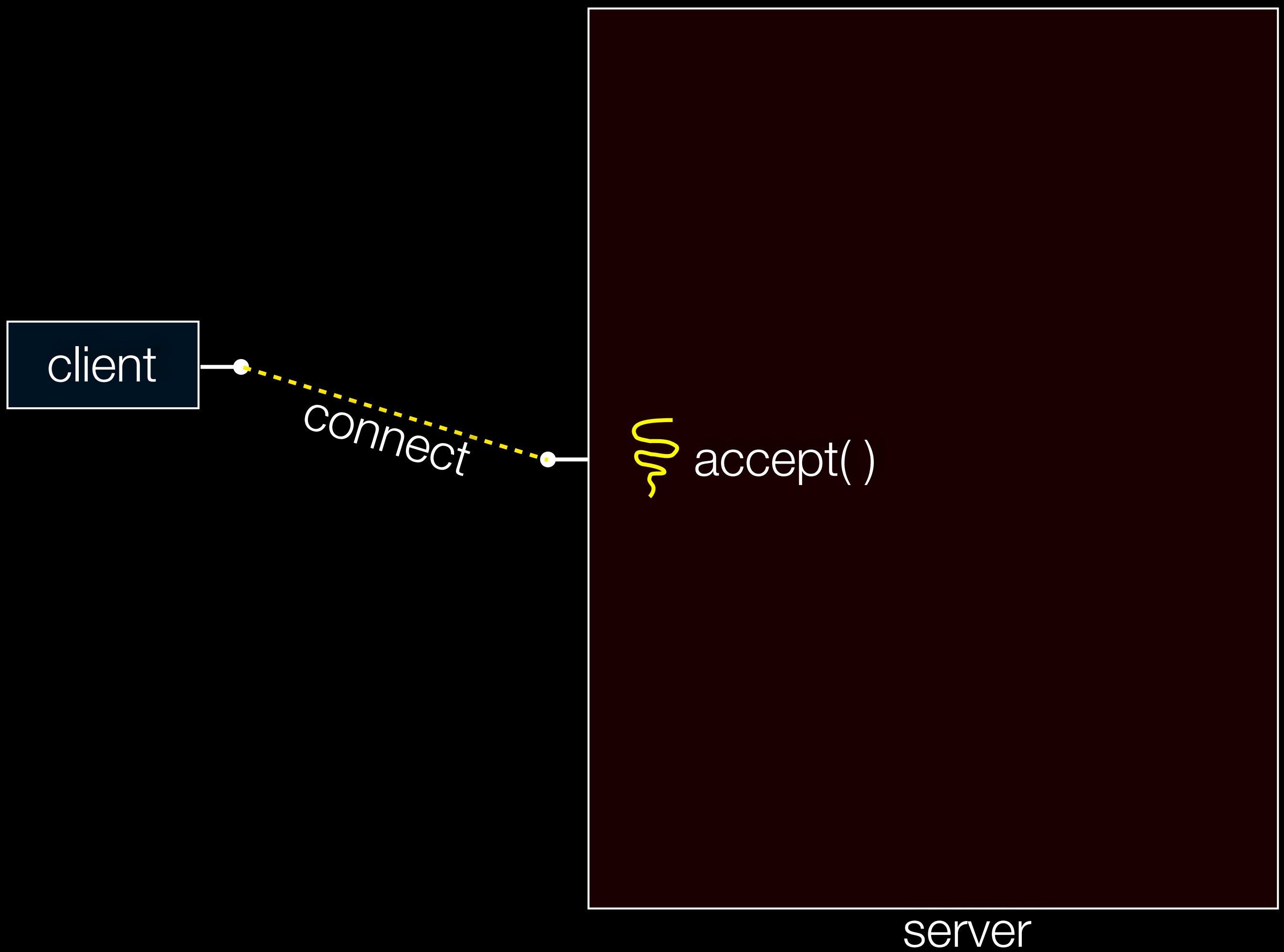

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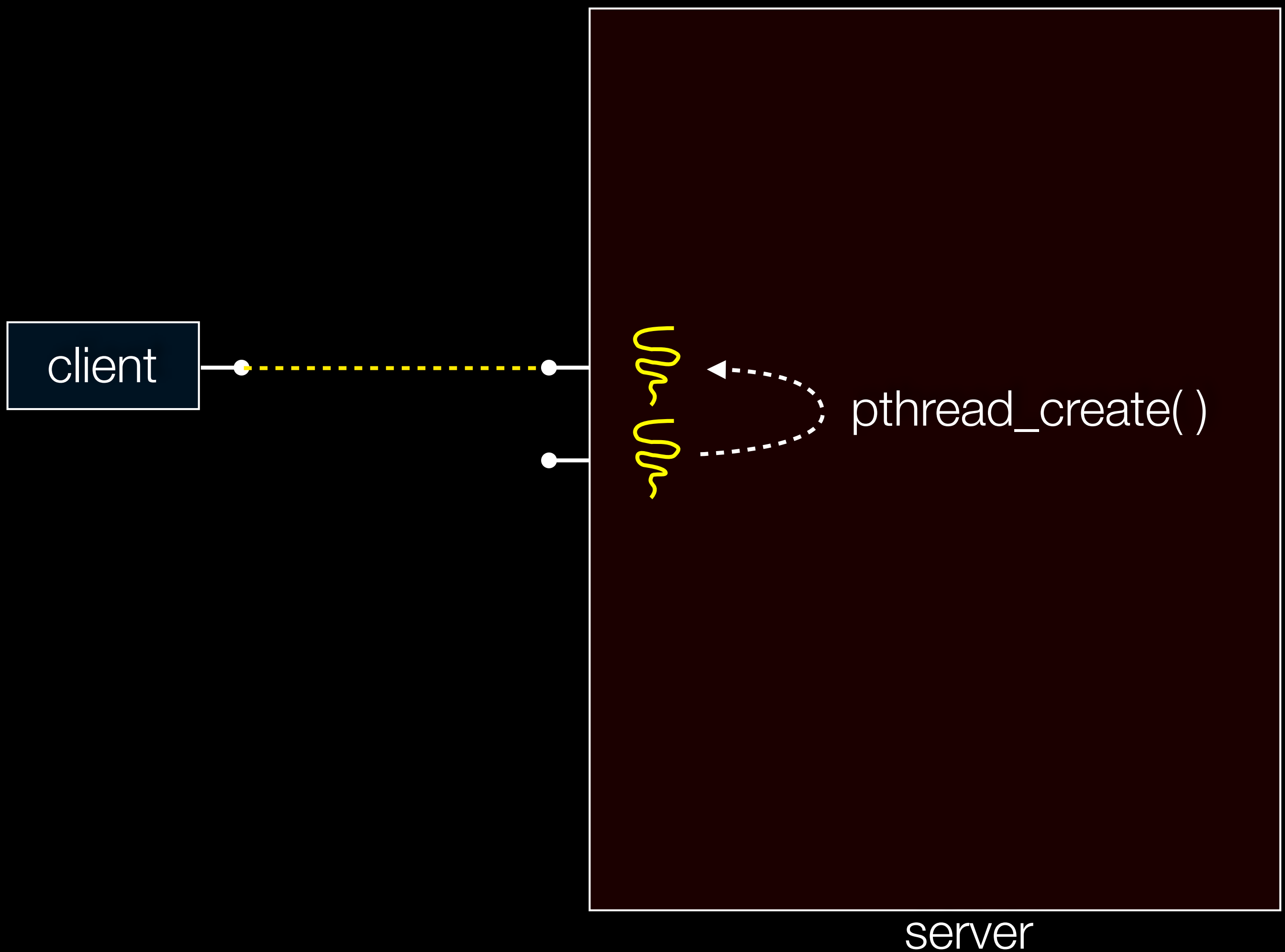

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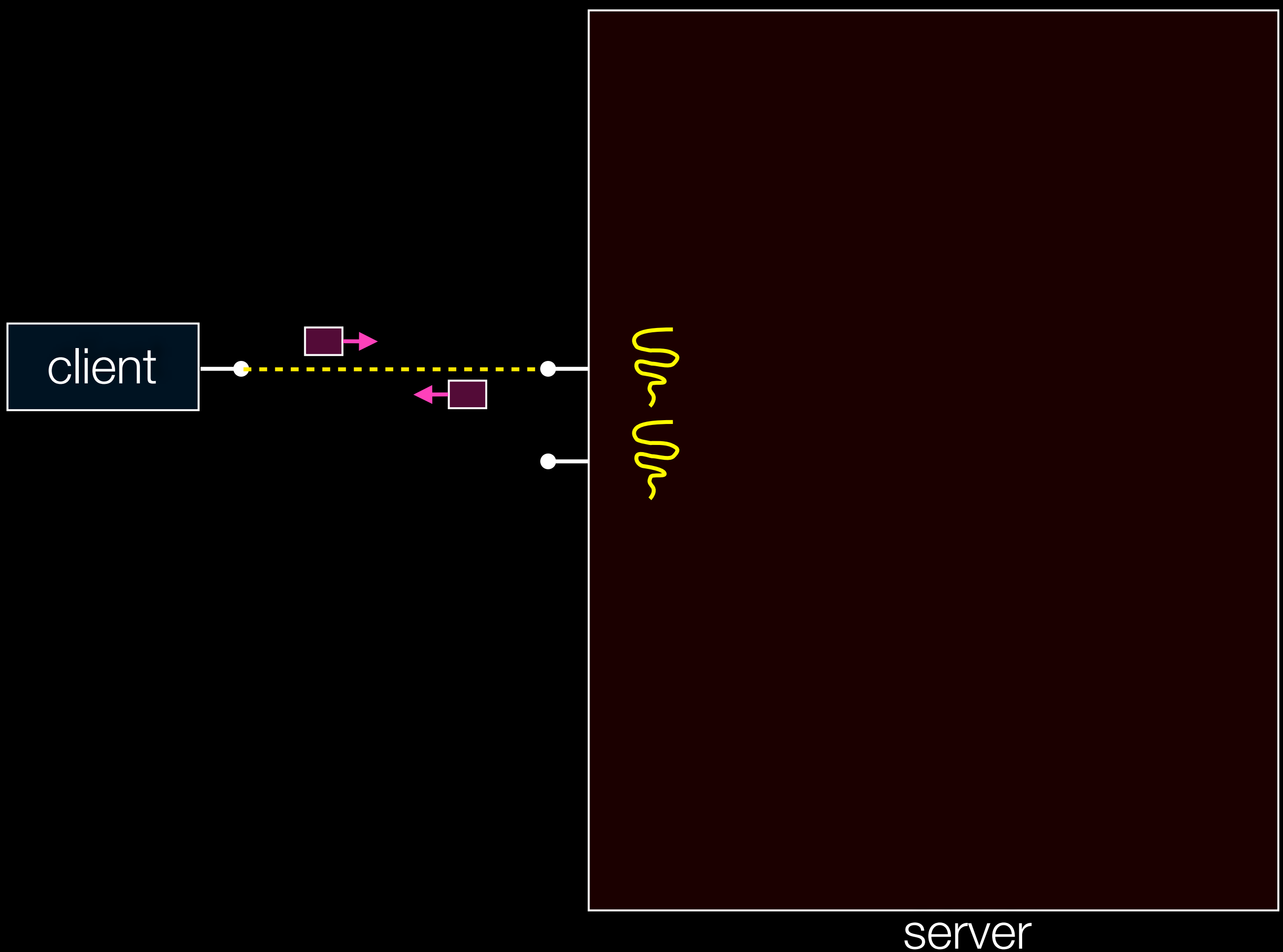

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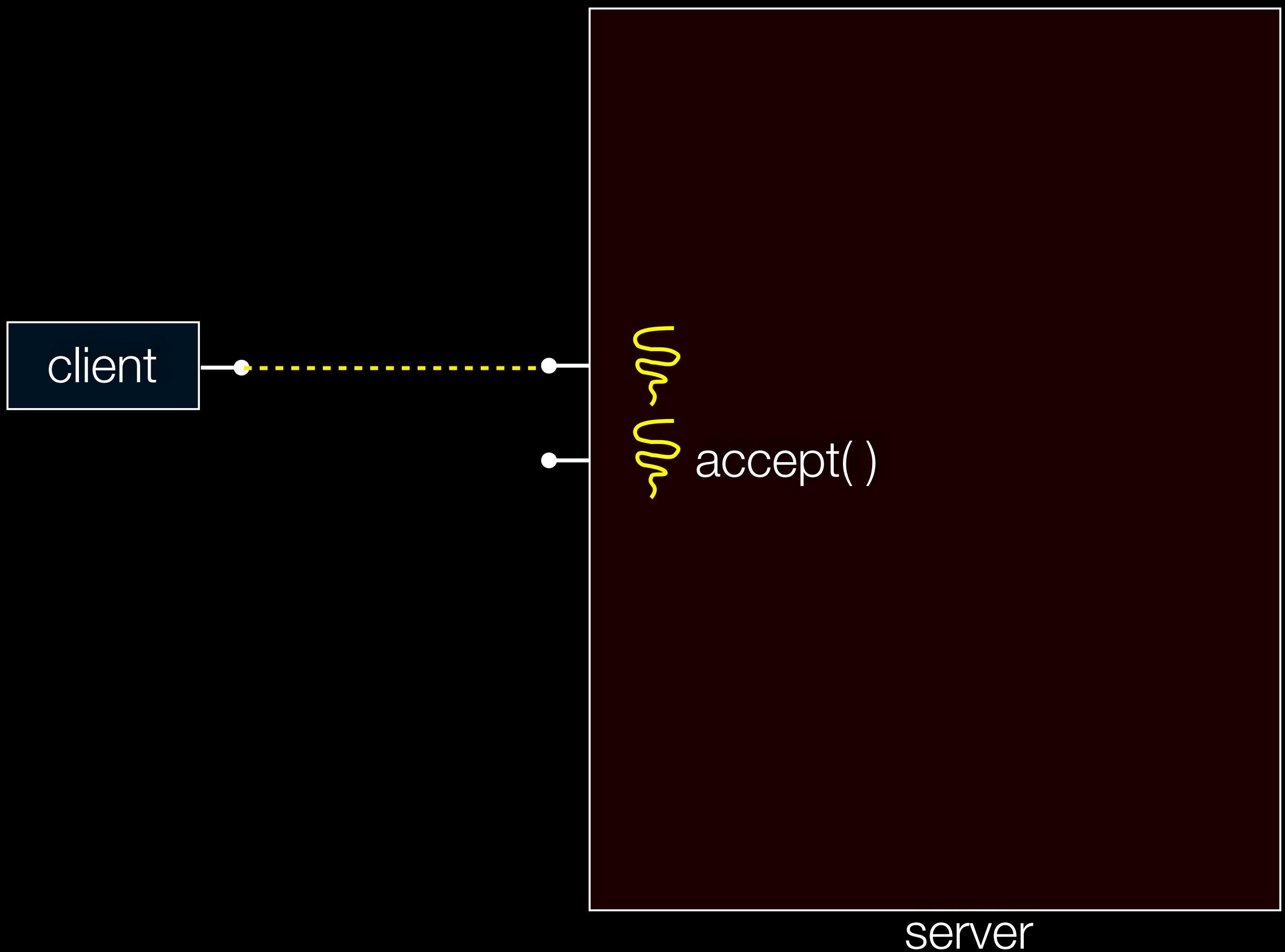

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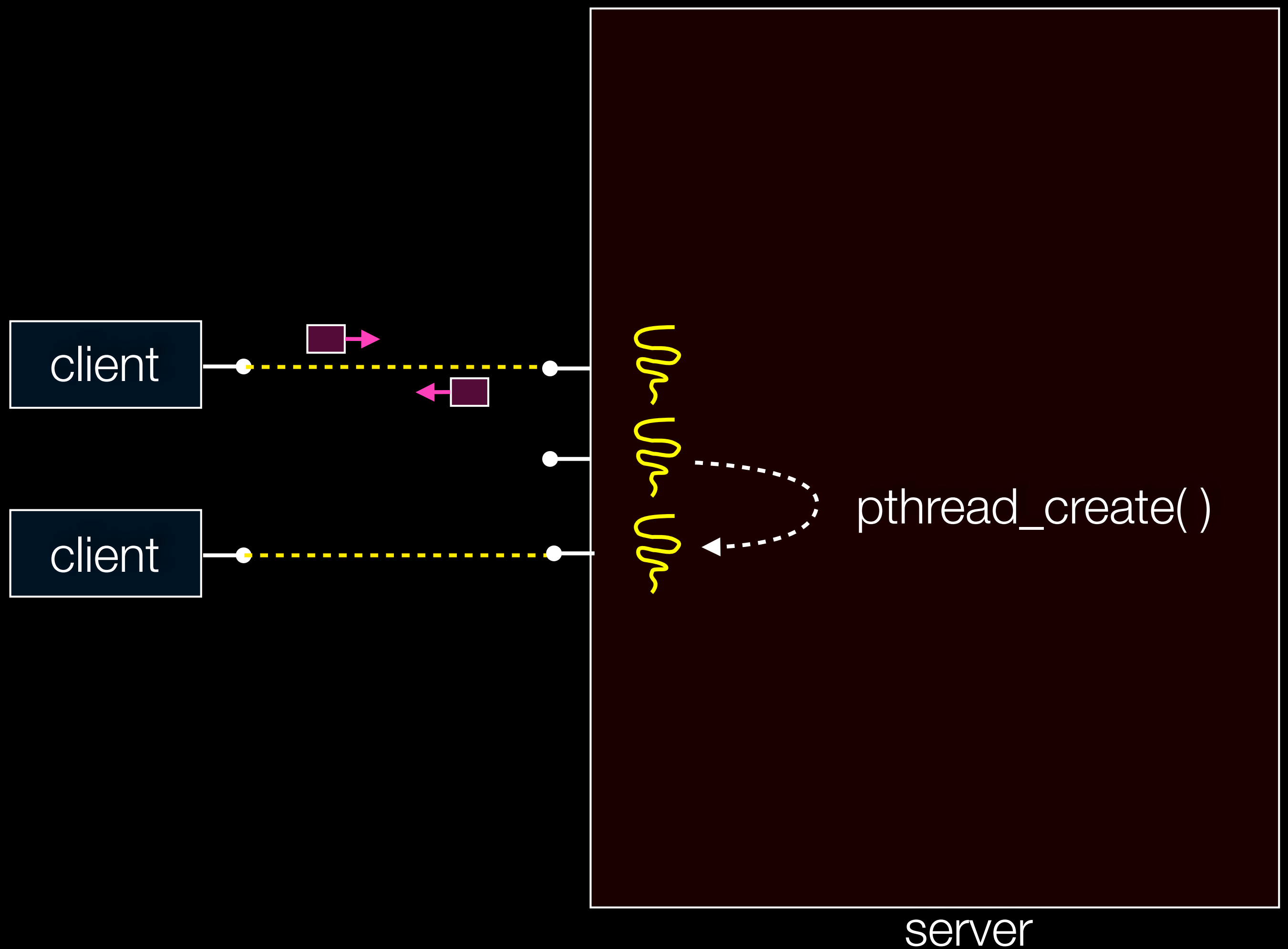

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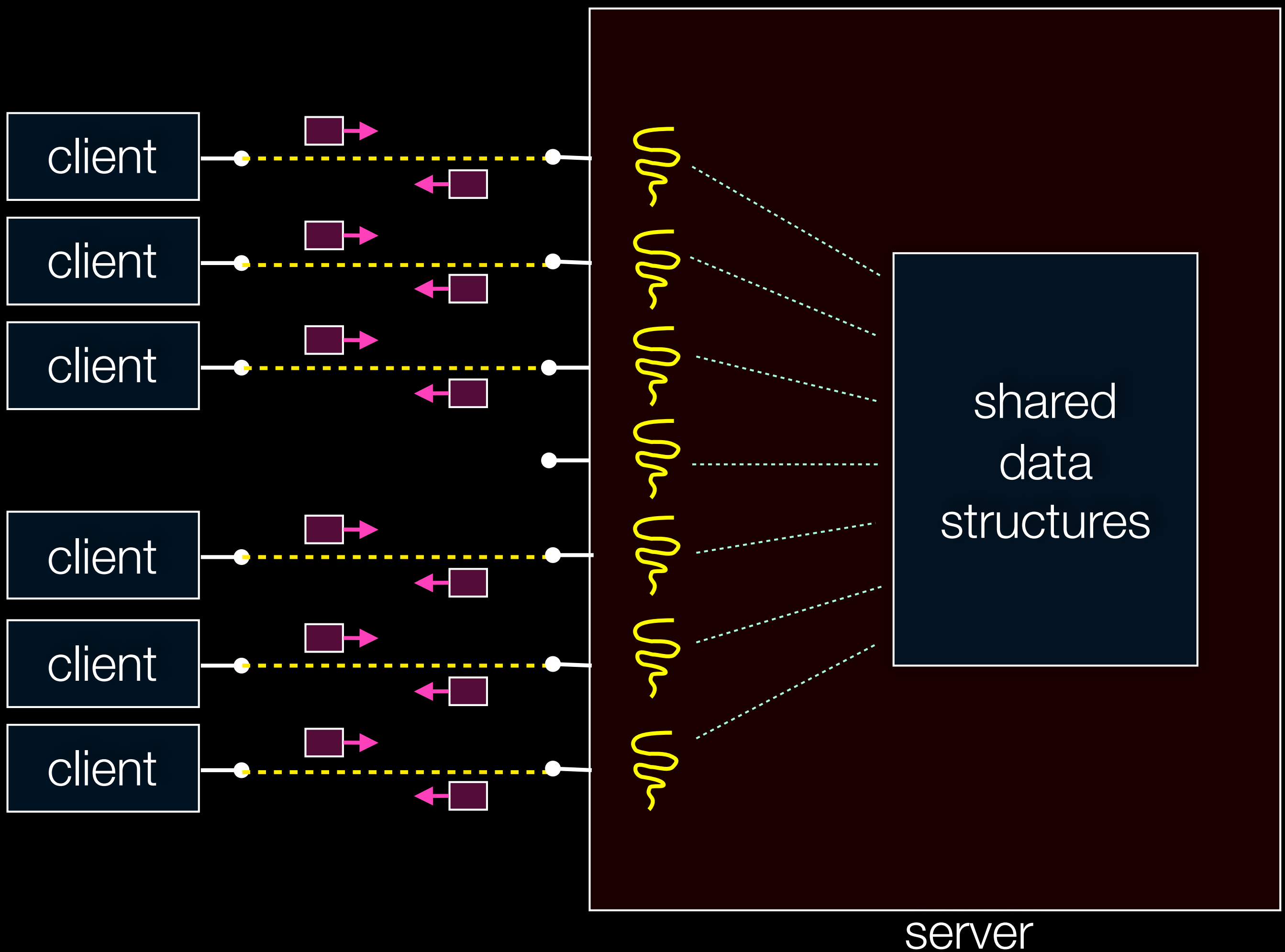

# Concurrent with threads

*look at **searchserver_threads/***

# Whither concurrent threads?

## Benefits

- straight-line code
  - still the case that much of the code is identical to sequential!
- parallel execution; good CPU, network utilization
  - lower overhead than processes
- shared-memory communication is possible

## Disadvantages

- ***synchronization*** is complicated
- shared fate within a process; one rogue thread can hurt you badly

# How fast is pthread_create?

*run* ***threadlatency.cc***

# Implications?

**0.036 ms** per thread create; ~10x faster than process forking

- maximum of (1000 / 0.018) = ~60,000 connections per second
- ~10 billion connections per day per core
    - much better

But, writing safe multithreaded code can be serious voodoo

# Thread Pools

In real servers we'd like to avoid overhead needed to create a new thread or process for every request

Idea: thread pools

- Create a set of worker threads or processes on server startup, put them in a queue
- When a request arrives, remove the first worker thread from the queue and assign it to handle the request
- When a worker is done it places itself back on the queue and then sleeps until dequeued and handed a new request

# Threads and races

What happens if two threads try to mutate the same data structure?

- they might interfere in painful, non-obvious ways, depending on the specifics of the data structure
  - imagine if two threads try to push an item onto the head of the linked list at the same time
  - depending on how the threads interleave, you might end up with a correct answer, or you might break the data structure altogether

# Simple “race” example

If no milk, buy some more

- liveness: if out, somebody buys
- safety: at most one person buys

What happens with multiple threads?

```
if (!milk) {
  buy milk
}
```

# Simple “race” example

Does this fix the problem?

```
if (!note) {
  if (!milk) {
    leave note
    buy milk
    remove note
  }
}
```

# Synchronization

Synchronization is the act of preventing two (or more) concurrently running threads from interfering with each other when operating on shared data

- need some mechanism to coordinate the threads
  - "let me go first, then you go"
- many different coordination mechanisms have been invented
  - take cse451 for details

# Locks

## lock acquire

- wait until the lock is free, then take it

## lock release

- release the lock
- if other threads are waiting for it
  - wake up exactly one of them
  - give it the lock

## simplifies concurrent code

- prevents more than one thread from entering a *critical section*

```
... non-critical code ...

lock.acquire();
  critical section
lock.release();

... non-critical code ...
```

# Simple “race” solution

What is the critical section?

- checking for milk
- buying more milk if out

These two steps must be uninterrupted, i.e., ***atomic***

- solution: protect the critical section with a lock

```
milk_lock.lock()

if (!milk) {
  buy milk
}

milk_lock.unlock()
```

# pthreads and locks

pthread_mutex_init( )

- creates a mutex (a.k.a. a lock)

pthread_mutex_lock( )

- grabs the lock

pthread_mutex_unlock( )

- releases the lock

*see **lock_example.cc***

# C++ 11 Threads

C++ 11 added threads and concurrency to the libraries

- <thread> - thread objects
- <mutex> - locks to handle critical sections
- <condition_variable> - used to block objects until notified to resume
- <atomic> - indivisible, atomic operations
- <future> - asynchronous access to data

Might be built on top of <pthread.h>, maybe not

Definitely use in C++ 11 code, but pthreads will still be around for a long, long time (and use pthreads in current exercise)

# Exercise 1

Write a simple "proxy" server

- forks a process for each connection
- reads an HTTP request from the client
  - relays that request to www.cs.washington.edu
- reads the response from www.cs.washington.edu
  - relays the response to the client, closes the connection

Try visiting your proxy using a web browser :)

# Exercise 2

Write a client program that:

- loops, doing “requests” in a loop.  Each request must:
  - connect to one of the echo servers from the lecture
  - do a network exchange with the server
  - close the connection
- keeps track of the latency (time to do a request) distribution
- keeps track of the throughput (requests / s)
- prints these out

See you on Wednesday !