Functions

Some function definitions:

```
- fun square(x:int):int = x * x;
val square = fn : int -> int
- fun swap(a:int, b:string):string*int = (b,a);
val swap = fn : int * string -> string * int
```

Some function types:

```
int -> int
int*string -> string*int
• in general: T<sub>arg(s)</sub> -> T<sub>result(s)</sub>
• * binds tighter than ->
```

Some function calls:

```
- square(3);
val it = 9 : int
- swap(3 * 4, "billy" ^ "bob");
val it = ("billybob",12) : string * int
```

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Expression-orientation

Function body is a single expression

```
fun square(x:int):int = x * x
```

- · not a statement list
- no return keyword needed

Like equality in math

 a call to a function is equivalent to its body, after substituting its formals for the actuals in the call

```
square(3) \Leftrightarrow (x*x)[x\rightarrow 3] \Leftrightarrow 3*3
```

There are no statements in ML, only expressions

· simplicity, regularity, and orthogonality in action

What would be statements in other languages are recast as expressions in ML

Expression Orientation: A Big Idea

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If expression

General form:

```
if test then e1 else e2
```

- return value of either e1 or e2,
 based on whether test is true or false
- cannot omit else part

```
- fun max(x:int, y:int):int =
=     if x >= y then x else y;
val max = fn : int * int -> int
```

Like ?: operator in C

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- don't need the separate ${\tt if}$ statement

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Static typechecking of if expression

What are the rules for typechecking an if expression? What's the type of the result of if?

Some basic principles of typechecking:

- · values are members of types
- the type of an expression must include all the values that might possibly result from evaluating that expression at run-time

Requirements on each if expression:

- the type of the <code>test</code> expression must be <code>bool</code>
- the type of the result of the if must include whatever values might be returned from the if
 - the if might return the result of either e1 or e2

A solution: e1 and e2 must have the same type, and that type is the type of the result of the if expression

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Let expression

An expression that introduces a new nested scope with local variable declarations

• unlike { ... } statements in C, which don't compute results

General form:

```
\begin{array}{lll} \mathbf{let} & \mathbf{val} & id_1 \colon type_1 = \mathbf{e}_1 \\ & \dots \\ & \mathbf{val} & id_n \colon type_n = \mathbf{e}_n \\ \mathbf{in} & & & \\ & & e_{body} \\ \mathbf{end} & & & \end{array}
```

• $type_i$ are optional; they'll be inferred

Evaluates each e_i and binds it to id_i , in turn

• each e_i can refer to the previous $id_1..id_{i-1}$ bindings

Evaluates $e_{\ensuremath{\textit{body}}}$ and returns it as the result of the let expression

• can refer to all the $id_1..id_n$ bindings

The id_i bindings disappear after e_{body} is evaluated

• they're in a nested, local scope

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Example scopes

```
- val x = 3;
val x = 3 : int
- fun f(y:int):int =
   let
    val z = x + y
    val x = 4
   in
     (let
        val y = z + x
        x + y + z
      end)
     + x + y + z
   end;
val f = fn : int -> int
- val x = 5;
val\ x = 5 : int
- f(x);
???
```

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