

# The Hardware/Software Interface

CSE351 Spring 2011

1<sup>st</sup> Lecture, March 28

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## **Teaching Assistants:**

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# Who is Luis?



PhD in architecture,  
multiprocessors, parallelism,  
compilers.



# Who are you?

- **55+ students (wow!)**
- **Who has written programs in assembly before?**
- **Written a threaded program before?**
  
- **What is an interface?**
  
- **Why do we need a hardware/software interface?**

# C vs. Assembler vs. Machine Programs

```
if ( x != 0 ) y = (y+z) / x;
```

```

cmpl  $0, -4(%ebp)
je    .L2
movl  -12(%ebp), %eax
movl  -8(%ebp), %edx
leal  (%edx,%eax), %eax
movl  %eax, %edx
sarl  $31, %edx
idivl -4(%ebp)
movl  %eax, -8(%ebp)
.L2:

```

```

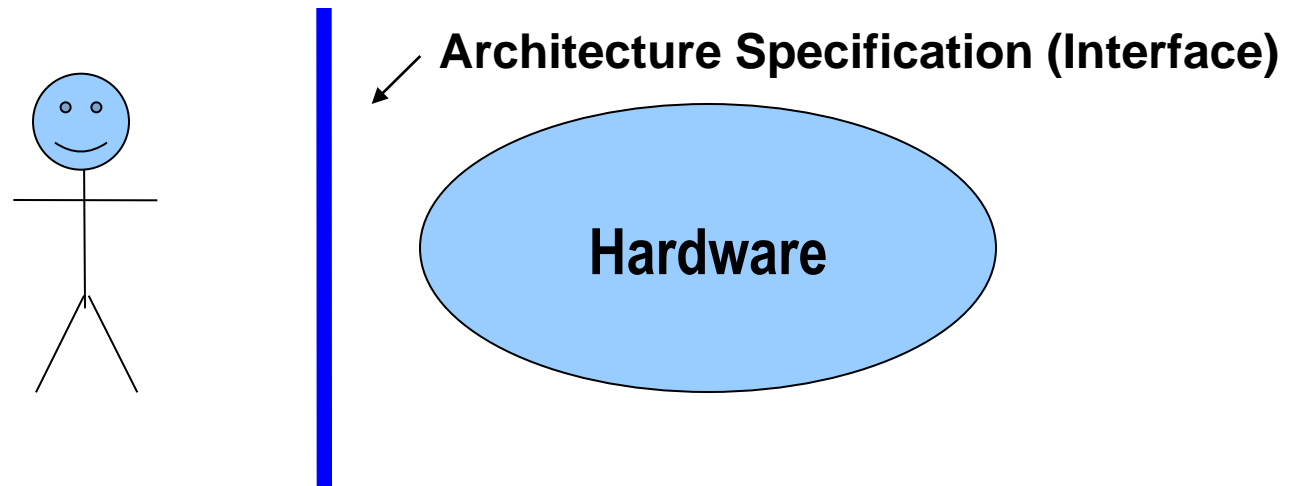
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```

- The three program fragments are equivalent
- You'd rather write C!
- The hardware likes bit strings!
  - The machine instructions are actually much shorter than the bits required to represent the characters of the assembler code

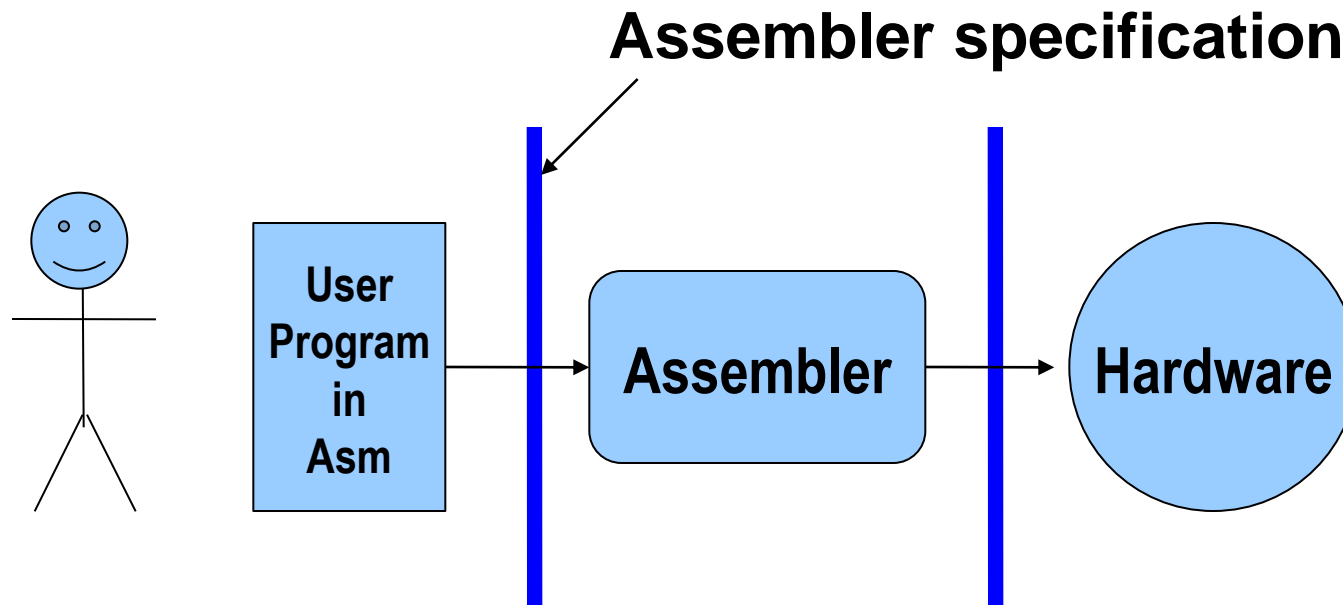
# HW/SW Interface: The Historical Perspective

- **Hardware started out quite primitive**
  - Design was expensive  $\Rightarrow$  the instruction set was very simple
    - E.g., a single instruction can add two integers
- **Software was also very primitive**



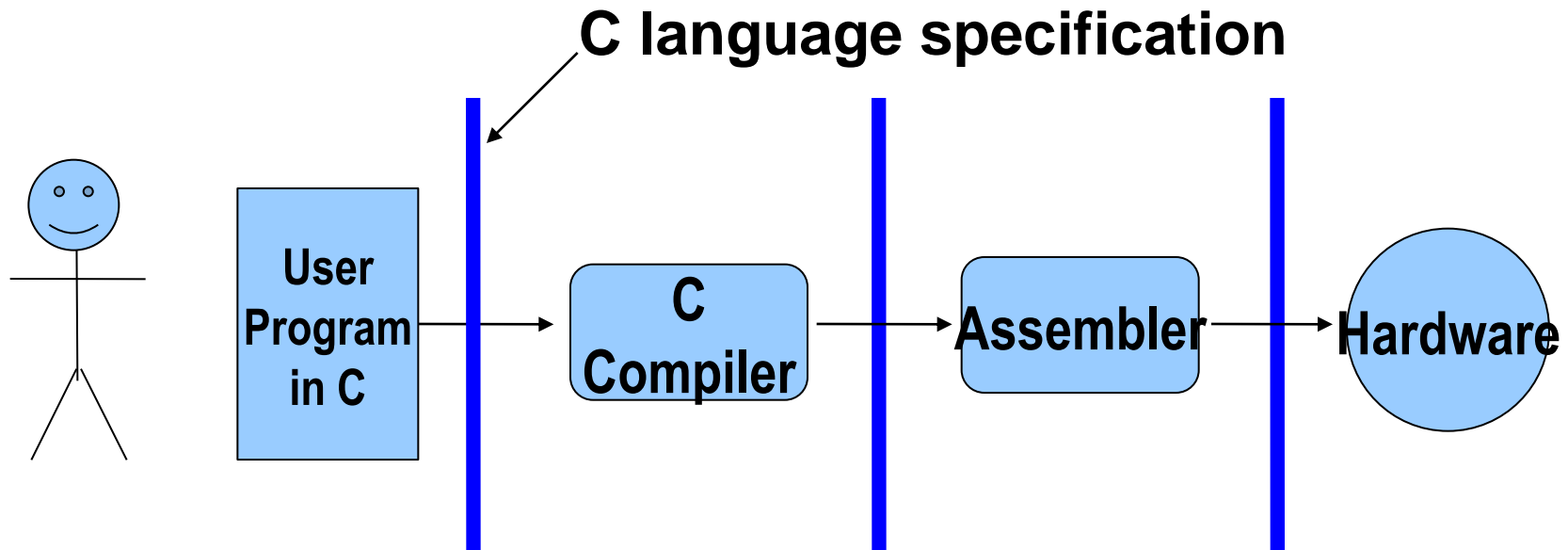
# HW/SW Interface: Assemblers

- Life was made a lot better by assemblers
  - 1 assembly instruction = 1 machine instruction, but...
  - different syntax: assembly instructions are character strings, not bit strings

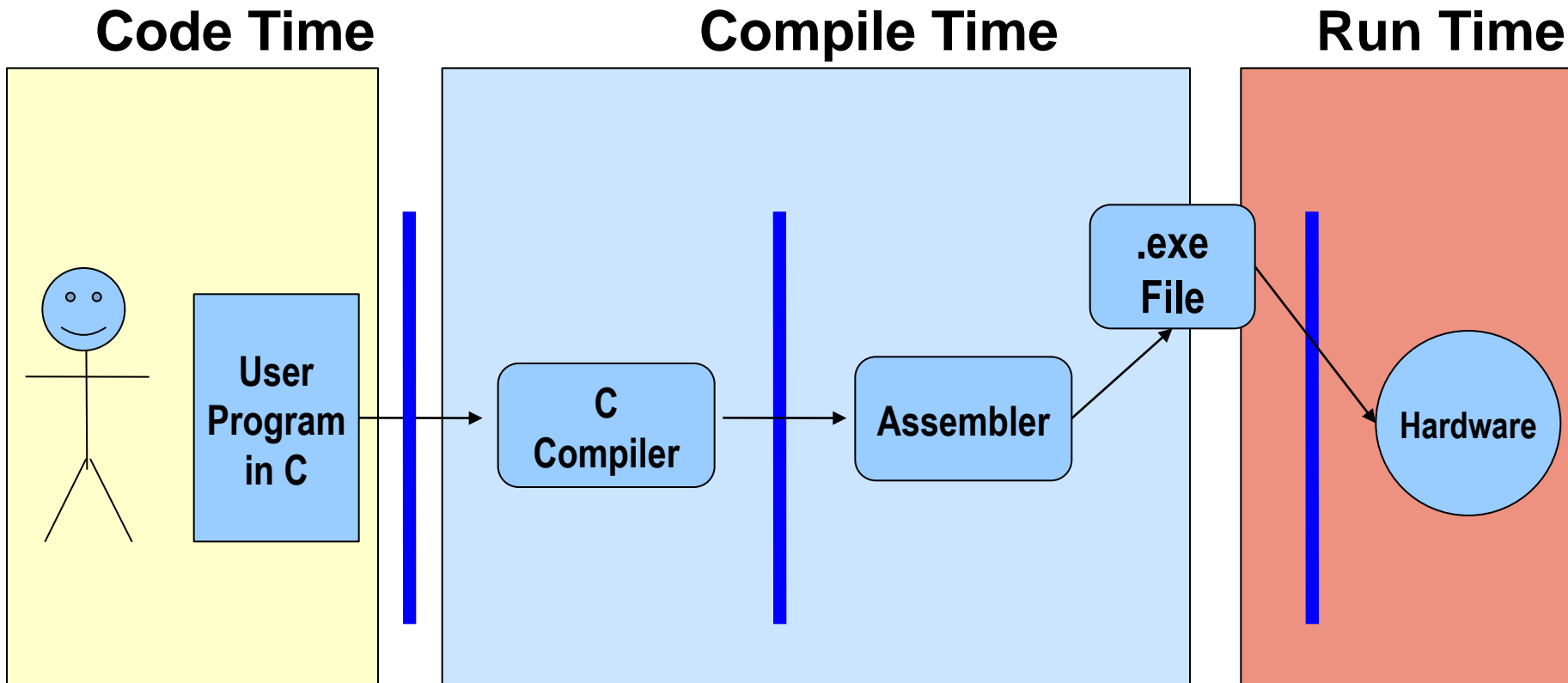


# HW/SW Interface: Higher Level Languages (HLL's)

- Higher level of abstraction:
  - 1 HLL line is compiled into many (many) assembler lines



# HW/SW Interface: Code / Compile / Run Times



**Note:** *The compiler and assembler are just programs, developed using this same process.*



# Overview

- **Course themes: big and little**
  - **Four important realities**
  - **How the course fits into the CSE curriculum**
  - **Logistics**
- 
- **(ready? 😊)**

# The Big Theme

- **THE HARDWARE/SOFTWARE INTERFACE**
- **How does the hardware (0s and 1s, processor executing instructions) relate to the software (Java programs)?**
- **Computing is about abstractions (but don't forget reality)**
- **What are the abstractions that we use?**
- **What do YOU need to know about them?**
  - When do they break down and you have to peek under the hood?
  - What assumptions are being made that may or may not hold in a new context or for a new technology?
  - What bugs can they cause and how do you find them?
- **Become a better programmer and begin to understand the thought processes that go into building computer systems**

# Little Theme 1: Representation

- **All digital systems represent everything as 0s and 1s**
- **Everything includes:**
  - Numbers – integers and floating point
  - Characters – the building blocks of strings
  - Instructions – the directives to the CPU that make up a program
  - Pointers – addresses of data objects in memory
- **These encodings are stored in registers, caches, memories, disks, etc.**
- **They all need addresses**
  - A way to find them
  - Find a new place to put a new item
  - Reclaim the place in memory when data no longer needed

# Little Theme 2: Translation

- **There is a big gap between how we think about programs and data and the 0s and 1s of computers**
- **Need languages to describe what we mean**
- **Languages need to be translated one step at a time**
  - Word-by-word
  - Phrase structures
  - Grammar
- **We know Java as a programming language**
  - Have to work our way down to the 0s and 1s of computers
  - Try not to lose anything in translation!
  - We'll encounter Java byte-codes, C language, assembly language, and machine code (for the X86 family of CPU architectures)

# Little Theme 3: Control Flow

- **How do computers orchestrate the many things they are doing – seemingly in parallel**
- **What do we have to keep track of when we call a method, and then another, and then another, and so on**
- **How do we know what to do upon “return”**
- **User programs and operating systems**
  - Multiple user programs
  - Operating system has to orchestrate them all
    - Each gets a share of computing cycles
    - They may need to share system resources (memory, I/O, disks)
  - Yielding and taking control of the processor
    - Voluntary or by force?

# Course Outcomes

- **Foundation: basics of high-level programming (Java)**
- **Understanding of some of the abstractions that exist between programs and the hardware they run on, why they exist, and how they build upon each other**
- **Knowledge of some of the details of underlying implementations**
- **Become more effective programmers**
  - More efficient at finding and eliminating bugs
  - Understand the many factors that influence program performance
  - Facility with some of the many languages that we use to describe programs and data
- **Prepare for later classes in CSE**

# Reality 1: Ints $\neq$ Integers & Floats $\neq$ Reals

- Representations are finite
- Example 1: Is  $x^2 \geq 0$ ?
  - Floats: Yes!
  - Ints:
    - $40000 * 40000 \rightarrow 1600000000$
    - $50000 * 50000 \rightarrow ??$
- Example 2: Is  $(x + y) + z = x + (y + z)$ ?
  - Unsigned & Signed Ints: Yes!
  - Floats:
    - $(1e20 + -1e20) + 3.14 \rightarrow 3.14$
    - $1e20 + (-1e20 + 3.14) \rightarrow ??$

# Code Security Example

```
/* Kernel memory region holding user-accessible data */
#define KSIZE 1024
char kbuf[KSIZE]; int len = KSIZE;

/* Copy at most maxlen bytes from kernel region to user buffer */
int copy_from_kernel(void *user_dest, int maxlen) {
    /* Byte count len is minimum of buffer size and maxlen */
    if (KSIZE > maxlen) len = maxlen;
    memcpy(user_dest, kbuf, len);
    return len;
}
```

- Similar to code found in FreeBSD's implementation of `getpeername`
- There are legions of smart people trying to find vulnerabilities in programs



# Typical Usage

```
/* Kernel memory region holding user-accessible data */
#define KSIZE 1024
char kbuf[KSIZE]; int len = KSIZE;

/* Copy at most maxlen bytes from kernel region to user buffer */
int copy_from_kernel(void *user_dest, int maxlen) {
    /* Byte count len is minimum of buffer size and maxlen */
    if (KSIZE > maxlen) len = maxlen;
    memcpy(user_dest, kbuf, len);
    return len;
}
```

```
#define MSIZE 528

void getstuff() {
    char mybuf[MSIZE];
    copy_from_kernel(mybuf, MSIZE);
    printf("%s\n", mybuf);
}
```

# Malicious Usage

```
/* Kernel memory region holding user-accessible data */
#define KSIZE 1024
char kbuf[KSIZE]; int len = KSIZE;

/* Copy at most maxlen bytes from kernel region to user buffer */
int copy_from_kernel(void *user_dest, int maxlen) {
    /* Byte count len is minimum of buffer size and maxlen */
    if (KSIZE > maxlen) len = maxlen;
    memcpy(user_dest, kbuf, len);
    return len;
}
```

```
#define MSIZE 528

void getstuff() {
    char mybuf[MSIZE];
    copy_from_kernel(mybuf, -MSIZE);
    . . .
}
```

# Reality #2: You've Got to Know Assembly

- **Chances are, you'll never write a program in assembly code**
  - Compilers are much better and more patient than you are
- **But: Understanding assembly is the key to the machine-level execution model**
  - Behavior of programs in presence of bugs
    - High-level language model breaks down
  - Tuning program performance
    - Understand optimizations done/not done by the compiler
    - Understanding sources of program inefficiency
  - Implementing system software
    - Operating systems must manage process state
  - Creating / fighting malware
  - x86 assembly is the language of choice

# Assembly Code Example

## ■ Time Stamp Counter

- Special 64-bit register in Intel-compatible machines
- Incremented every clock cycle
- Read with rdtsc instruction

## ■ Application

- Measure time (in clock cycles) required by procedure

```
double t;  
start_counter();  
P();  
t = get_counter();  
printf("P required %f clock cycles\n", t);
```

# Code to Read Counter

- Write small amount of assembly code using GCC's asm facility
- Inserts assembly code into machine code generated by compiler

```
/* Set *hi and *lo to the high and low order bits
   of the cycle counter.
*/

void access_counter(unsigned *hi, unsigned *lo)
{
    asm("rdtsc; movl %%edx,%0; movl %%eax,%1"
        : "=r" (*hi), "=r" (*lo) /* output */
        : /* input */
        : "%edx", "%eax"); /* clobbered */
}
```

# Reality #3: Memory Matters

- **Memory is not unbounded**
  - It must be allocated and managed
  - Many applications are memory-dominated
- **Memory referencing bugs are especially pernicious**
  - Effects are distant in both time and space
- **Memory performance is not uniform**
  - Cache and virtual memory effects can greatly affect program performance
  - Adapting program to characteristics of memory system can lead to major speed improvements

# Memory Referencing Bug Example

```
double fun(int i)
{
    volatile double d[1] = {3.14};
    volatile long int a[2];
    a[i] = 1073741824; /* Possibly out of bounds */
    return d[0];
}
```

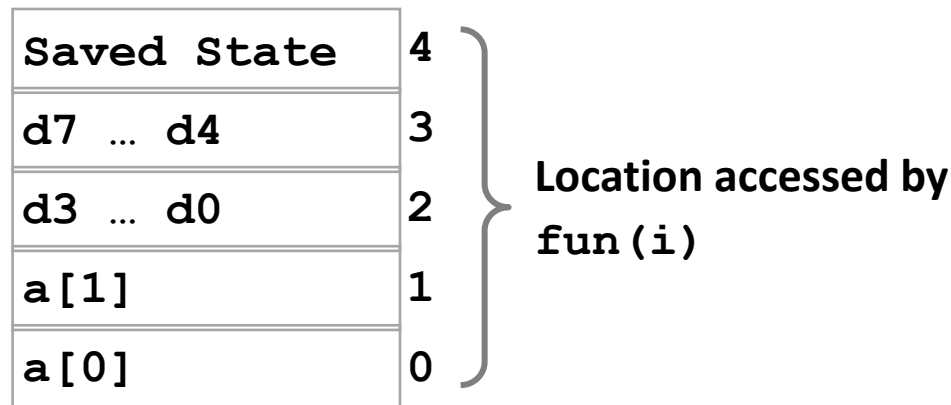
```
fun(0)    ->    3.14
fun(1)    ->    3.14
fun(2)    ->    3.1399998664856
fun(3)    ->    2.00000061035156
fun(4)    ->    3.14, then segmentation fault
```

# Memory Referencing Bug Example

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```

## Explanation:





# Memory Referencing Errors


- **C (and C++) do not provide any memory protection**
  - Out of bounds array references
  - Invalid pointer values
  - Abuses of malloc/free
- **Can lead to nasty bugs**
  - Whether or not bug has any effect depends on system and compiler
  - Action at a distance
    - Corrupted object logically unrelated to one being accessed
    - Effect of bug may be first observed long after it is generated
- **How can I deal with this?**
  - Program in Java (or C#, or ML, or ...)
  - Understand what possible interactions may occur
  - Use or develop tools to detect referencing errors

# Memory System Performance Example

- Hierarchical memory organization
- Performance depends on access patterns
  - Including how program steps through multi-dimensional array

```
void copyij(int src[2048][2048],
            int dst[2048][2048])
{
    int i,j;
    for (i = 0; i < 2048; i++)
        for (j = 0; j < 2048; j++)
            dst[i][j] = src[i][j];
}
```

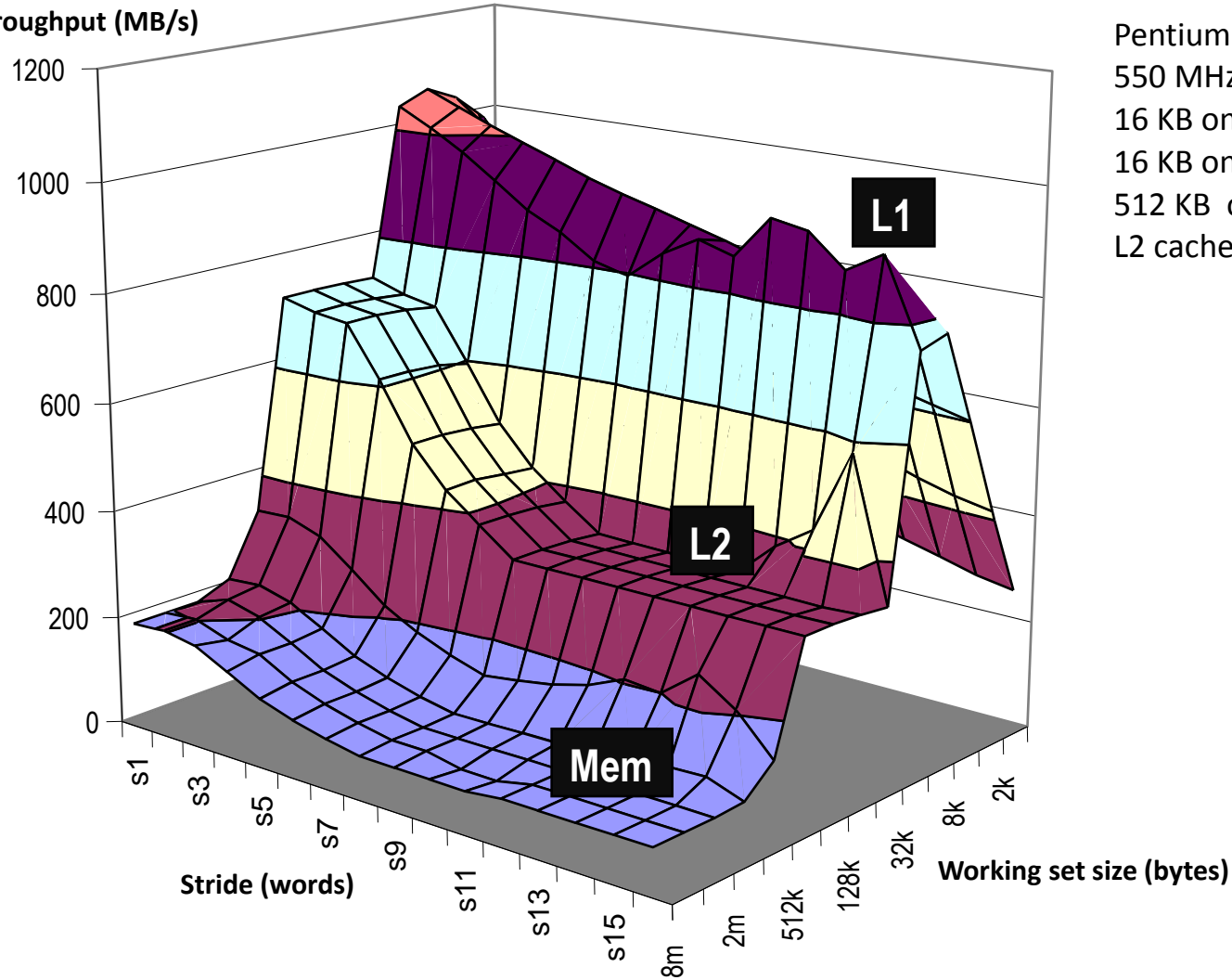
```
void copyji(int src[2048][2048],
            int dst[2048][2048])
{
    int i,j;
    for (j = 0; j < 2048; j++)
        for (i = 0; i < 2048; i++)
            dst[i][j] = src[i][j];
}
```



**21 times slower  
(Pentium 4)**

# The Memory Mountain

Read throughput (MB/s)



Pentium III Xeon

550 MHz

16 KB on-chip L1 d-cache

16 KB on-chip L1 i-cache

512 KB off-chip unified

L2 cache

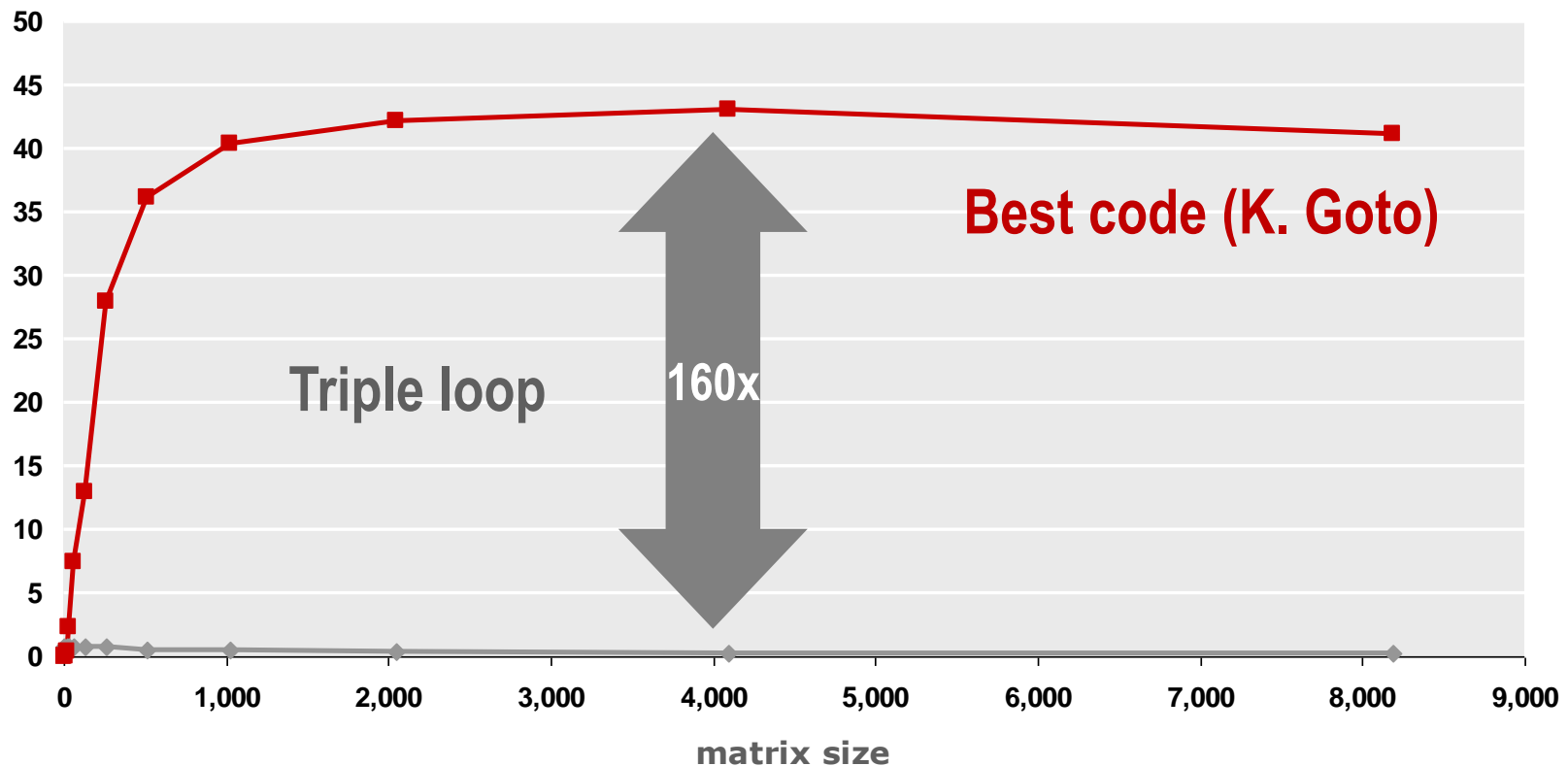
# Reality #4: Performance isn't counting ops

- **Exact op count does not predict performance**
  - Easily see 10:1 performance range depending on how code written
  - Must optimize at multiple levels: algorithm, data representations, procedures, and loops
- **Must understand system to optimize performance**
  - How programs compiled and executed
  - How memory system is organized
  - How to measure program performance and identify bottlenecks
  - How to improve performance without destroying code modularity and generality

# Example Matrix Multiplication

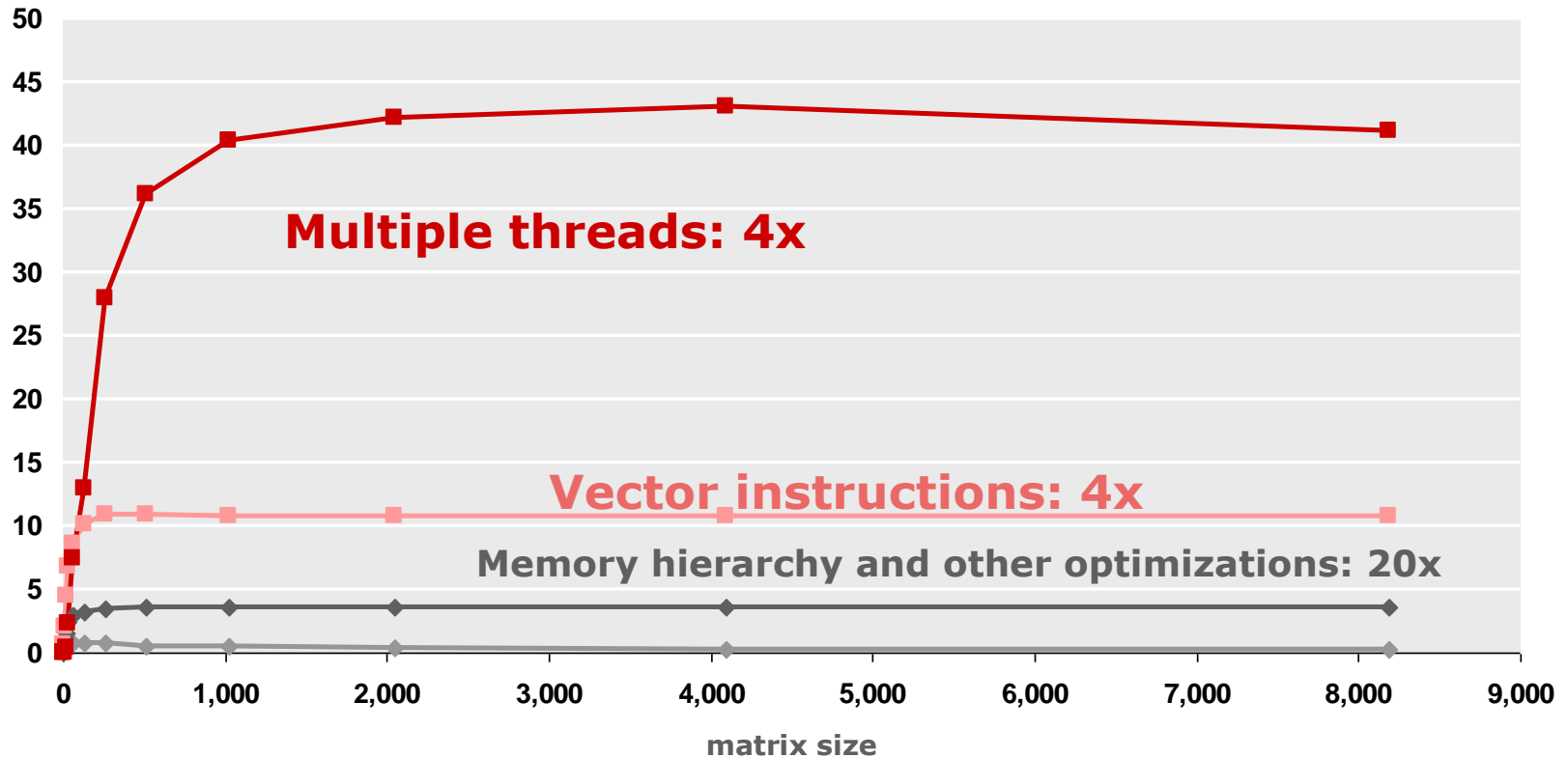
- Standard desktop computer, vendor compiler, using optimization flags
- Both implementations have **exactly** the same operations count ( $2n^3$ )

**Matrix-Matrix Multiplication (MMM) on 2 x Core 2 Duo 3 GHz (double precision)**  
Gflop/s



# MMM Plot: Analysis

Matrix-Matrix Multiplication (MMM) on 2 x Core 2 Duo 3 GHz  
Gflop/s



- Reason for 20x: blocking or tiling, loop unrolling, array scalarization, instruction scheduling, search to find best choice
- *Effect: less register spills, less L1/L2 cache misses, less TLB misses*

# CSE351's role in new CSE Curriculum

## ■ Pre-requisites

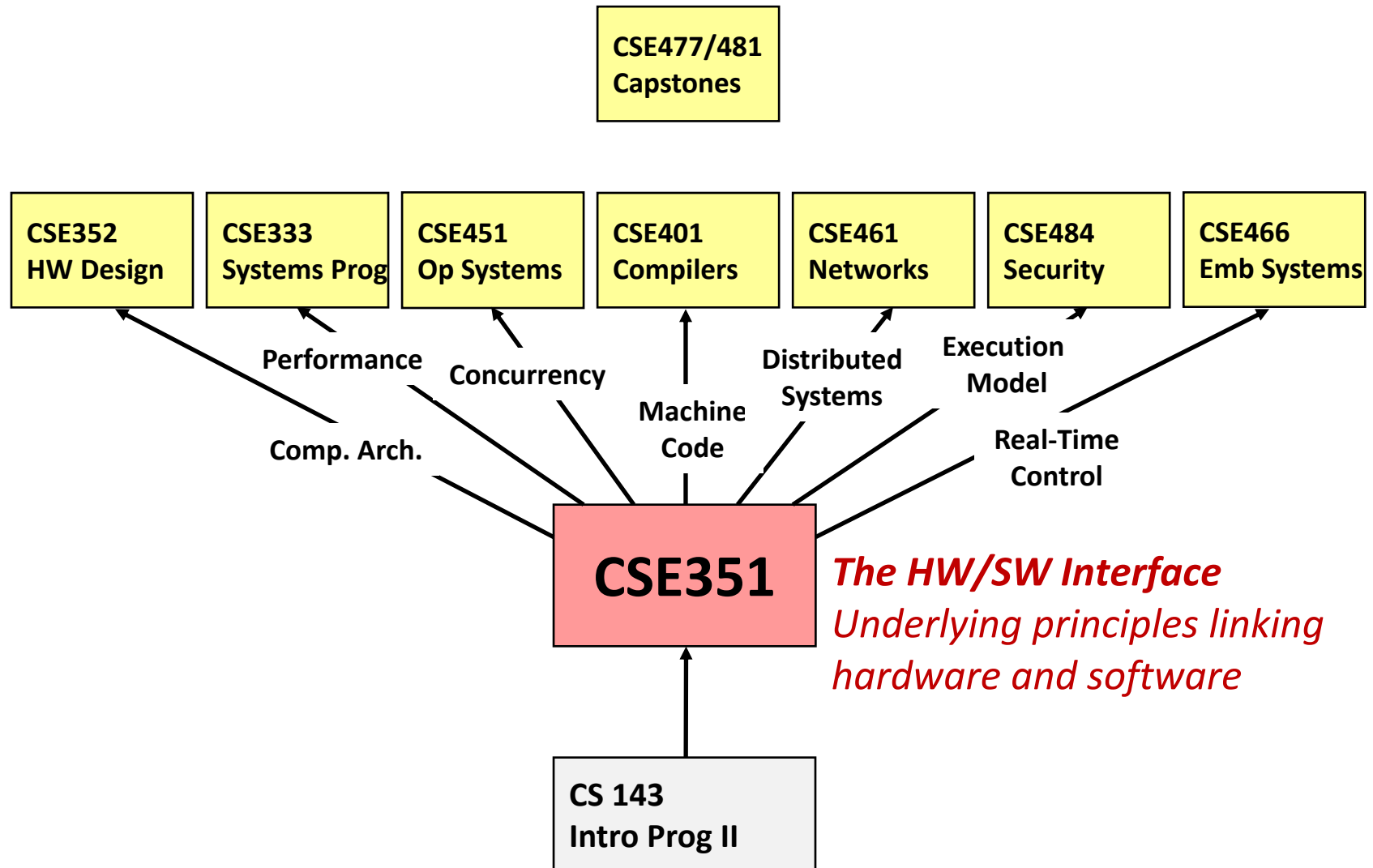
- 142 and 143: Intro Programming I and II

## ■ One of 6 core courses

- 311: Foundations I
- 312: Foundations II
- 331: SW Design and Implementation
- 332: Data Abstractions
- 351: HW/SW Interface
- 352: HW Design and Implementation

## ■ 351 sets the context for many follow-on courses

# CSE351's place in new CSE Curriculum





# Course Perspective

## ■ Most systems courses are Builder-Centric

- Computer Architecture
  - Design pipelined processor in Verilog
- Operating Systems
  - Implement large portions of operating system
- Compilers
  - Write compiler for simple language
- Networking
  - Implement and simulate network protocols

# Course Perspective (Cont.)

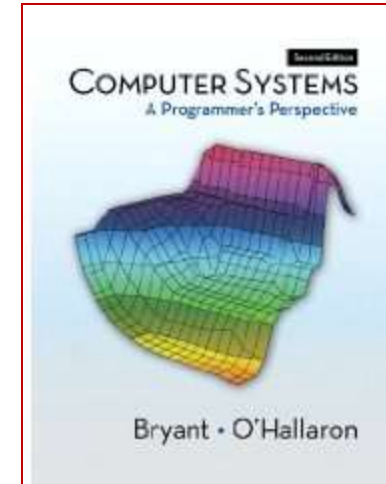
## ■ This course is Programmer-Centric

- Purpose is to show how software really works
- By understanding the underlying system, one can be more effective as a programmer
  - Better debugging
  - Better basis for evaluating performance
  - How multiple activities work in concert (e.g., OS and user programs)
- Not just a course for dedicated hackers
  - What every CSE major needs to know
- Provide a context in which to place the other CSE courses you'll take

# Textbooks

## ■ **Computer Systems: A Programmer's Perspective, 2<sup>nd</sup> Edition**

- Randal E. Bryant and David R. O'Hallaron
- Prentice-Hall, 2010
- <http://csapp.cs.cmu.edu>
- This book really matters for the course!
  - How to solve labs
  - Practice problems typical of exam problems



## ■ **A good C book.**

- C: A Reference Manual (Harbison and Steele)
- The C Programming Language (Kernighan and Ritchie)

# Course Components

- **Lectures (~30)**
  - Higher-level concepts – I'll assume you've done the reading in the text
- **Sections (~10)**
  - Applied concepts, important tools and skills for labs, clarification of lectures, exam review and preparation
- **Written assignments (4)**
  - Problems from text to solidify understanding
- **Labs (4)**
  - Provide in-depth understanding (via practice) of an aspect of systems
- **Exams (midterm + final)**
  - Test your understanding of concepts and principles

# Resources

## ■ Course Web Page

- <http://www.cse.washington.edu/351>
- Copies of lectures, assignments, exams

## ■ Course Discussion Board

- Keep in touch outside of class – help each other
- Staff will monitor and contribute

## ■ Course Mailing List

- Low traffic – mostly announcements; you are already subscribed

## ■ Staff email

- Things that are not appropriate for discussion board or better offline

## ■ Anonymous Feedback (will be linked from homepage)

- Any comments about anything related to the course where you would feel better not attaching your name

# Policies: Grading

- **Exams: weighted 1/3 (midterm), 2/3 (final)**
- **Written assignments: weighted according to effort**
  - We'll try to make these about the same
- **Labs assignments: weighted according to effort**
  - These will likely increase in weight as the quarter progresses
- **Grading:**
  - 25% written assignments
  - 35% lab assignments
  - 40% exams

# Welcome to CSE351!

- **Let's have fun**
- **Let's learn – together**
- **Let's communicate**
- **Let's set the bar for a useful and interesting class**
  
- **Many thanks to the many instructors who have shared their lecture notes – I will be borrowing liberally through the qtr – they deserve all the credit, the errors are all mine**
  - UW: Gaetano Borriello (Inaugural edition of CSE 351, Spring 2010)
  - CMU: Randy Bryant, David O'Halloran, Gregory Kesden, Markus Püschel
  - Harvard: Matt Welsh
  - UW: Tom Anderson, Luis Ceze, John Zahorjan