

Lecture 5

◆ Logistics

- HW1 was due before lecture
- HW2 posted today, due in one week
- Lab2 ongoing
- Question on how to post lab grades/check list
- Final exam scheduled: 6/9/ 8:30am here EEB 105

◆ Last lecture

- Logic gates and truth tables
- Implementing logic functions

◆ Today's lecture

- Canonical forms
- NAND and NOR

de Morgan's theorem

◆ Replace

- • with +, + with •, 0 with 1, and 1 with 0
- All variables with their complements

◆ Example 1: $Z = A'B' + A'C'$

$$Z' = (A'B' + A'C')'$$

$$= (A+B) \cdot (A+C)$$

◆ Example 2: $Z = A'B'C + A'BC + AB'C + ABC'$

$$Z' = (A'B'C + A'BC + AB'C + ABC')'$$

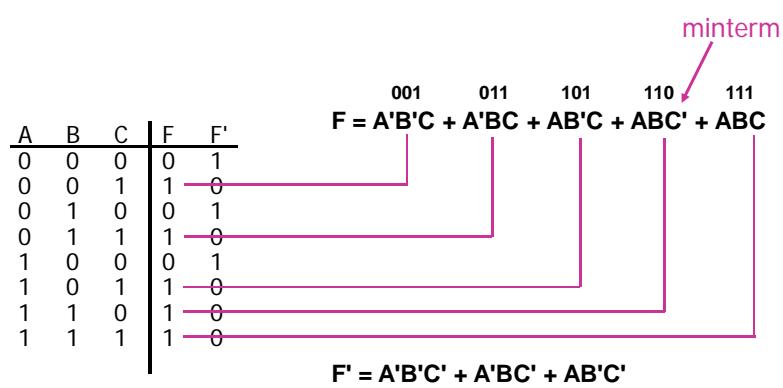
$$= (A+B+C') \cdot (A+B'+C) \cdot (A'+B+C) \cdot (A'+B'+C)$$

Canonical forms

- ◆ Canonical forms
 - Standard forms for Boolean expressions
 - Generally not the simplest forms
 - ↳ Can be minimized
 - Derived from truth table
- ◆ Two canonical forms
 - Sum-of-products (minterms)
 - Product-of-sum (maxterms)

Sum-of-products canonical form (SOP)

- ◆ Also called disjunctive normal form (DNF)
 - Commonly called a **minterm expansion**



Minterms

- ◆ Variables appears exactly once in each minterm
 - In true or inverted form (but not both)

A	B	C	minterms
0	0	0	$A'B'C'$ m0
0	0	1	$A'B'C$ m1
0	1	0	$A'BC$ m2
0	1	1	$A'BC$ m3
1	0	0	$AB'C$ m4
1	0	1	$AB'C$ m5
1	1	0	ABC' m6
1	1	1	ABC m7

↑
short-hand notation

F in canonical form:

$$\begin{aligned} F(A,B,C) &= \Sigma m(1,3,5,6,7) \\ &= m1 + m3 + m5 + m6 + m7 \\ &= A'B'C + A'BC + AB'C + ABC' + ABC \end{aligned}$$

canonical form → minimal form

$$\begin{aligned} F(A,B,C) &= A'B'C + A'BC + AB'C + ABC + ABC' \\ &= (A'B' + A'B + AB' + AB)C + ABC' \\ &= ((A' + A)(B' + B))C + ABC' \\ &= ABC' + C \\ &= AB + C \end{aligned}$$

Product-of-sums canonical form (POS)

- ◆ Also called conjunctive normal form (CNF)
 - Commonly called a **maxterm expansion**

A	B	C	F	F'
0	0	0	0	1
0	0	1	1	0
0	1	0	0	1
0	1	1	1	0
1	0	0	0	1
1	0	1	1	0
1	1	0	1	0
1	1	1	1	0

$F = (A + B + C)(A + B' + C)(A' + B + C)$

↑
maxterm

$F' = (A+B+C')(A+B'+C')(A'+B+C')(A'+B'+C)(A'+B'+C')$

Maxterms

- ◆ Variables appears exactly once in each maxterm
 - In true or inverted form (but not both)

A	B	C	maxterms
0	0	0	$A+B+C$ M0
0	0	1	$A+B+C'$ M1
0	1	0	$A+B'+C$ M2
0	1	1	$A+B'+C'$ M3
1	0	0	$A'+B+C$ M4
1	0	1	$A'+B+C'$ M5
1	1	0	$A'+B'+C$ M6
1	1	1	$A'+B'+C'$ M7

short-hand notation

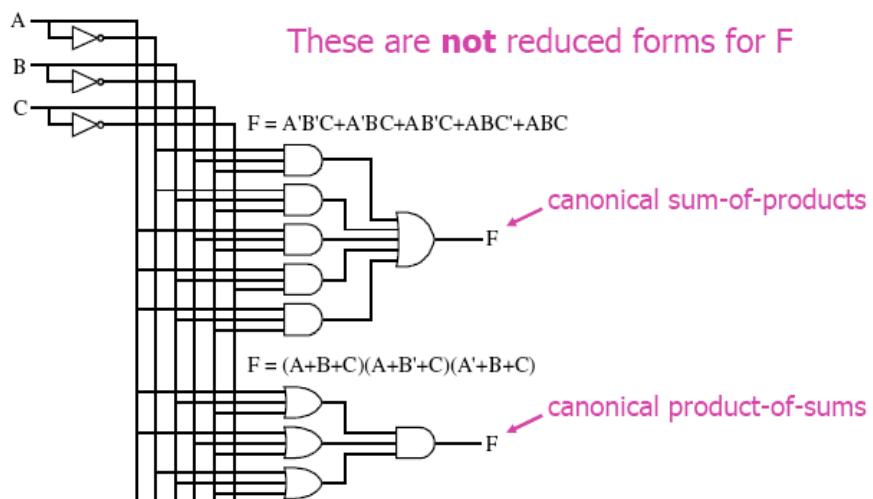
F in canonical form:

$$\begin{aligned} F(A,B,C) &= \prod_{i=0}^2 M_i \\ &= M_0 \cdot M_2 \cdot M_4 \\ &= (A+B+C)(A+B'+C)(A'+B+C) \end{aligned}$$

canonical form → minimal form

$$\begin{aligned} F(A,B,C) &= (A+B+C)(A+B'+C)(A'+B+C) \\ &= (A+B+C)(A+B'+C) \cdot \\ &\quad (A+B+C)(A'+B+C) \\ &= (A+C)(B+C) \end{aligned}$$

Canonical implementations of $F = AB + C$



Conversion between canonical forms

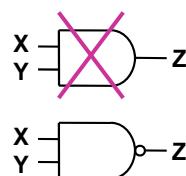
- ◆ Minterm to maxterm
 - Use maxterms that aren't in minterm expansion
 - $F(A,B,C) = \sum m(1,3,5,6,7) = \prod M(0,2,4)$
- ◆ Maxterm to minterm
 - Use minterms that aren't in maxterm expansion
 - $F(A,B,C) = \prod M(0,2,4) = \sum m(1,3,5,6,7)$
- ◆ Minterm of F to minterm of F'
 - Use minterms that don't appear
 - $F(A,B,C) = \sum m(1,3,5,6,7) \quad F'(A,B,C) = \sum m(0,2,4)$
- ◆ Maxterm of F to maxterm of F'
 - Use maxterms that don't appear
 - $F(A,B,C) = \prod M(0,2,4) \quad F'(A,B,C) = \prod M(1,3,5,6,7)$

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NAND/NOR more common/efficient

- ◆ CMOS logic gates are more common and efficient in the inverted forms
 - NAND, NOR, NOT
 - Even though Canonical forms discussed so far used AND/OR, NAND/NOR preferred for real hardware implementation



X	Y	Z
0	0	1
0	1	1
1	0	1
1	1	0

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NAND and NOR (truth table)

$(X + Y)' = X' \cdot Y'$
 NOR is equivalent to AND
 with inputs complemented

X	Y	X'	Y'	$(X + Y)'$	$X' \cdot Y'$
0	0	1	1	1	1
0	1	1	0	0	0
1	0	0	1	0	0
1	1	0	0	0	0

$(X \cdot Y)' = X' + Y'$
 NAND is equivalent to OR
 with inputs complemented

X	Y	X'	Y'	$(X \cdot Y)'$	$X' + Y'$
0	0	1	1	1	1
0	1	1	0	1	1
1	0	0	1	1	1
1	1	0	0	0	0

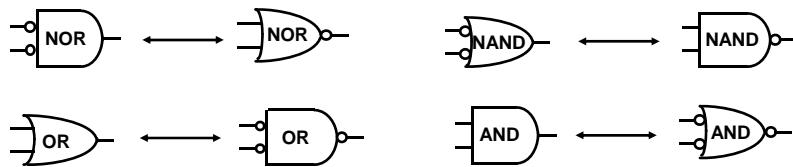
NAND and NOR (logic gates)

◆ de Morgan's

- Standard form: $A'B' = (A + B)'$ $A' + B' = (AB)'$
- Inverted: $A + B = (A'B')'$ $(AB) = (A' + B')'$

- AND with complemented inputs \equiv NOR
- OR with complemented inputs \equiv NAND
- OR \equiv NAND with complemented inputs
- AND \equiv NOR with complemented inputs

pushing
the
bubble



Converting to use NAND/NOR

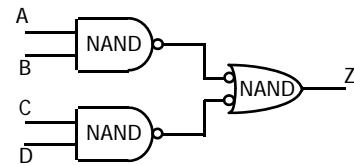
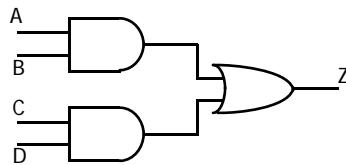
◆ Introduce inversions ("bubbles")

- Introduce bubbles in pairs
 - ↳ Conserve inversions
 - ↳ Do not alter logic function

◆ Example

- AND/OR to NAND/NAND

$$\begin{aligned} Z &= AB + CD \\ &= (A' + B')' + (C' + D')' \\ &= [(A' + B')(C' + D')]' \\ &= [(AB)'(CD)']' \end{aligned}$$



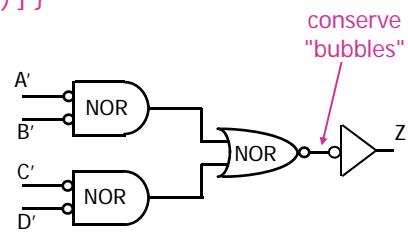
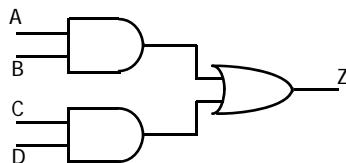
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Converting to use NAND/NOR (con't)

◆ Example: AND/OR network to NOR/NOR

$$\begin{aligned} Z &= AB + CD \\ &= (A' + B')' + (C' + D')' \\ &= [(A' + B')' + (C' + D')']'' \\ &= \{\{(A' + B')' + (C' + D')'\}'\} \end{aligned}$$

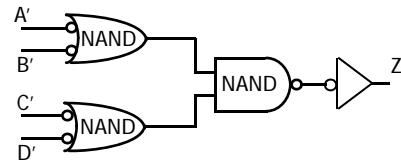
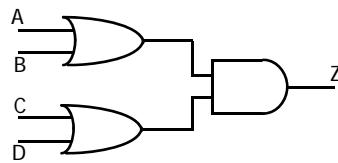


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Converting to use NAND/NOR (con't)

- ◆ Example: OR/AND to NAND/NAND

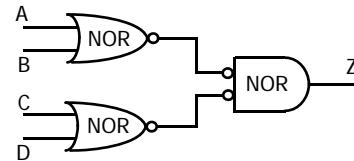
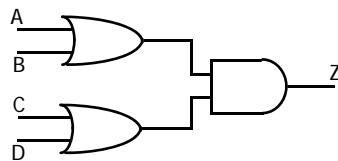


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Converting between forms (con't)

- ◆ Example: OR/AND to NOR/NOR



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