

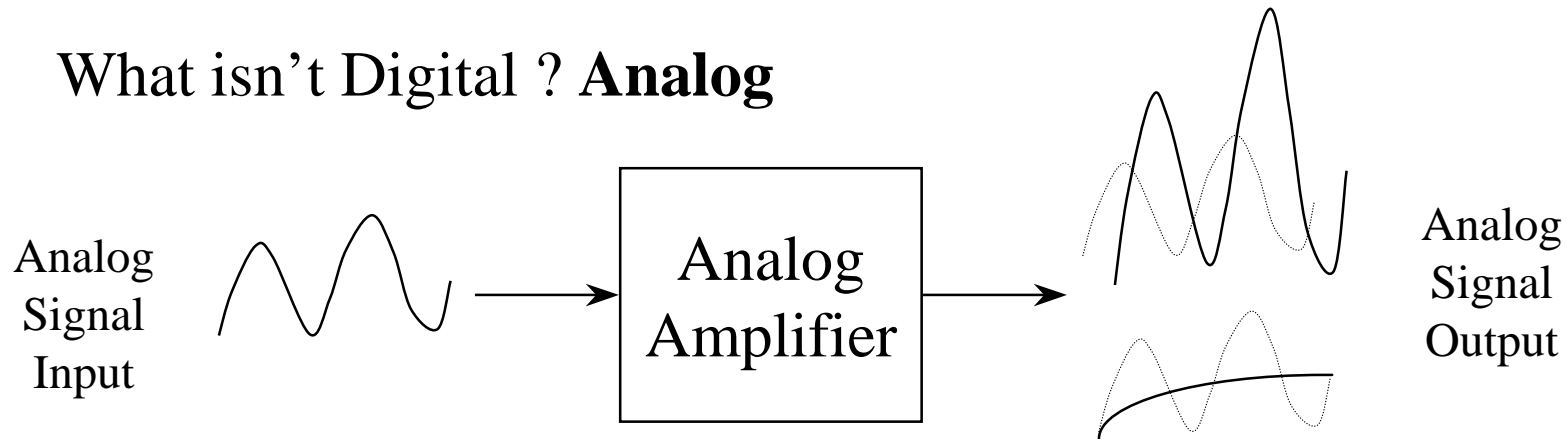
# CSE370: Introduction to Digital Design

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- ❑ Course staff
  - Larry Arnstein, ...
- ❑ Course web
  - [www.cs.washington.edu/education/courses/370/99au/](http://www.cs.washington.edu/education/courses/370/99au/)
- ❑ This week: Introduction
  - Keywords "digital" and "design"
  - Examples and basics
- ❑ Content and Grading
  - Weekly homeworks = 40%
  - Class participation = 10%
  - Midterm = 25%
  - Final (comprehensive) = 25%
- ❑ Software: DesignWorks

# What is Digital?

## What isn't Digital ? **Analog**



### Key Design Issue:

Frequency response

### Examples:

- Bullhorn
- Noise Filter

### Strengths

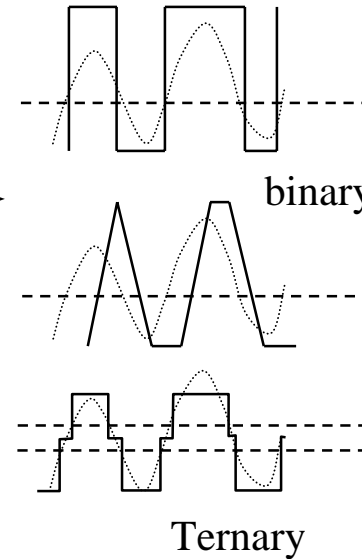
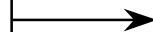
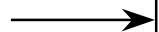
**Compact, and Error Tolerant**  
**One wire can carry a Symphony**

### Weaknesses

**“Lossy”-- Subject to distortion,**  
**good for media, not computation**

# What's Digital

Input signal  
"regarded" as  
Digital



Output  
signal is  
restored or  
"quantized"

Key issues:

Switching Speed  
and Delay

Weaknesses

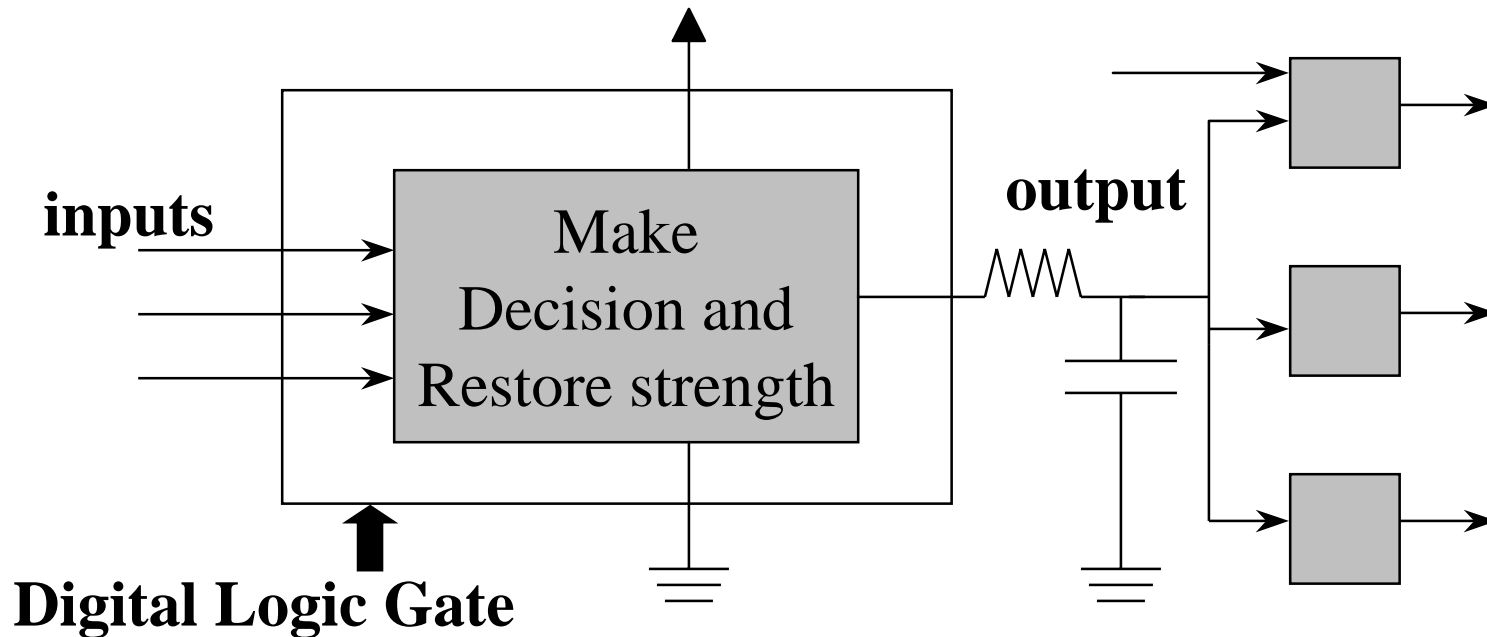
- Inefficient
- Rely on parallelism to get things done
- (Sega Dreamcast is a 128 Bit Machine)
- Error intolerant

Strengths

- Lossless -- Only loss is quantization.
- Great for computation!
- Highly reproducible and shrinkable

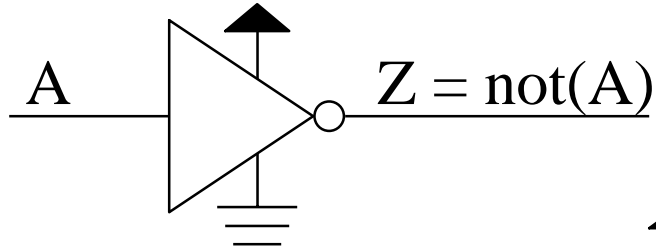
# The Digital Gate (Electronic)

- Value of output determined by values of inputs
- The output strength (voltage or current) must be high enough to activate several similar devices.



**Allows complex computation with no degradation of information**

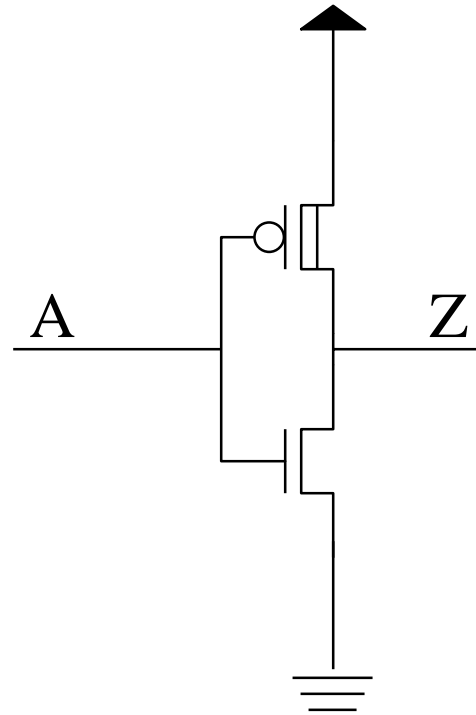
# An Example: The CMOS Inverter



pMOS  
Transistor



nMOS  
Transistor



nMOS + pMOS = CMOS

## MOS Transistors

The resistance of the “channel” is determined by the voltage of the gate.

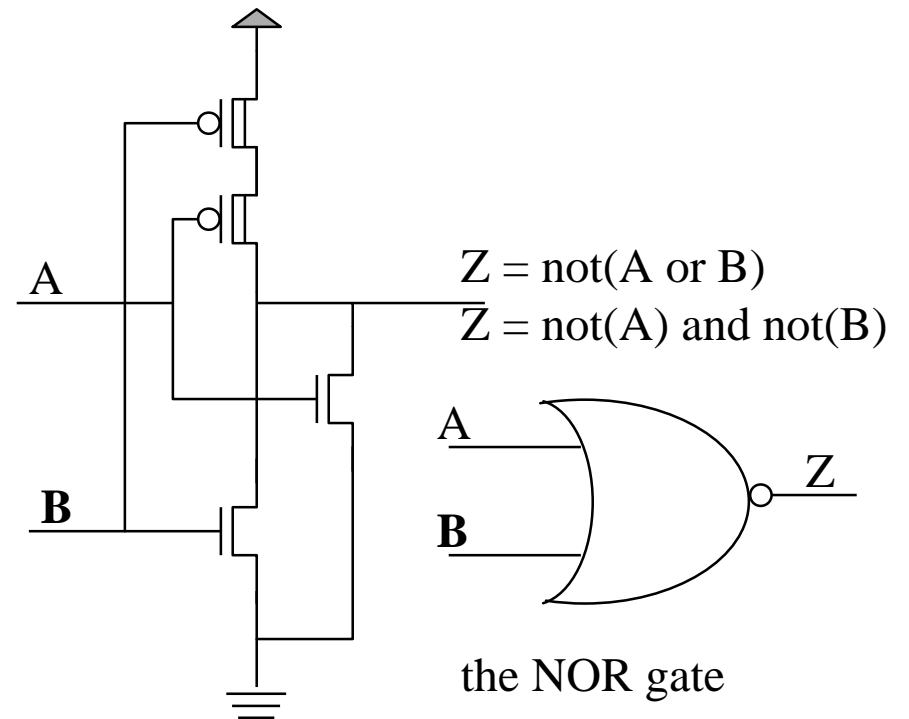
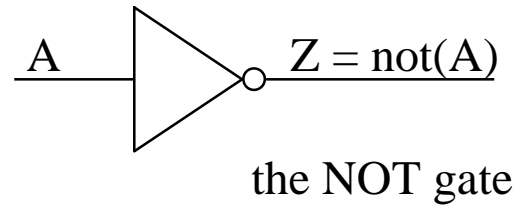
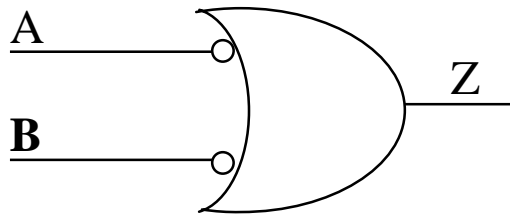
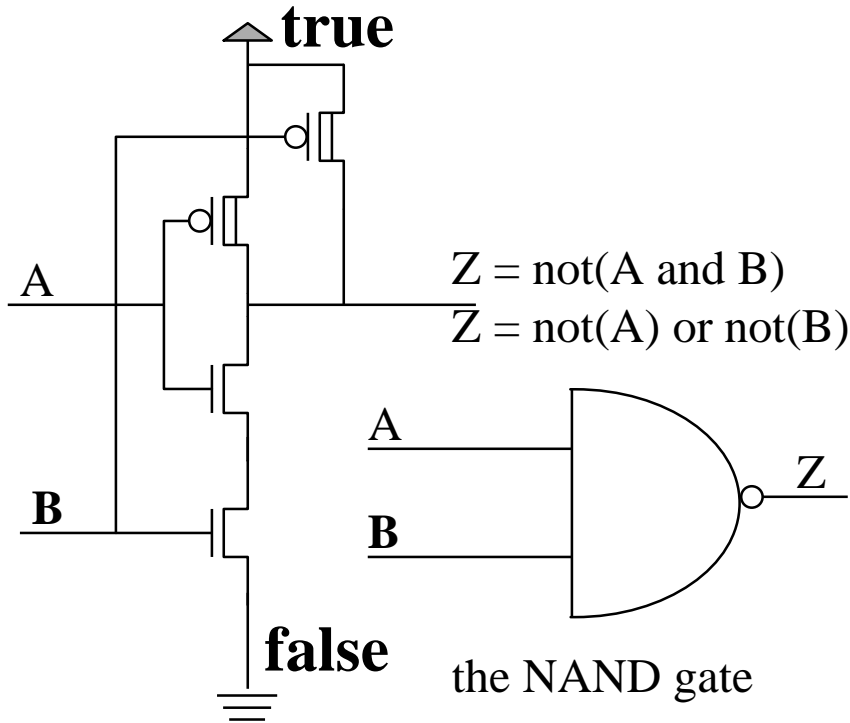
A positive voltage on the input reduces the resistance of the nMOS channel, and increases the resistance of the pMOS channel.

## Key properties:

Very small input current (leakage only), very high “gain” amplifier

**Weak input → Strong Output**

# Basic Units of Digital Computation in CMOS



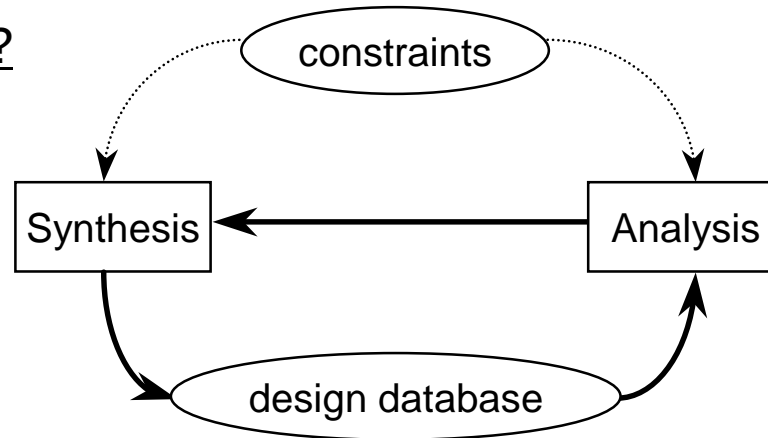
## Other forms of digital hardware

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- ❑ First Digital System: Molecular Structure such as DNA
- ❑ Beads -- the abacus (~500 BC)
- ❑ Mechanisms -- the position of a gear or shaft (1600's -- Babbage 1800's)
- ❑ Electromechanisms -- relays (Zuse 1900's)
- ❑ Electrical Currents (Tubes, Bipolar Transistors 1930's-1970's)
- ❑ **Voltages (CMOS 1980's-present)**
- ❑ Fiber Optics -- Light or Light off
- ❑ Dynamic RAM -- Charged or Discharged capacitor (on MOS gate)
- ❑ Nonvolatile memory (erasable) -- Trapped electrons or no trapped electrons
- ❑ Programmable ROM -- Fuses
- ❑ Bubble memory -- Presence or absence of a magnetic bubble
- ❑ Magnetic disk -- Direction of magnetic flux
- ❑ Compact disc -- the Presence or absence of a pit

# What is design (cont'd)?

## □ What is engineering design?



## □ What is logic design?

### ➤ Constraints =

Function, Cost, Performance, Power, Servicability, Safety, Project Schedule, etc.

### ➤ Database =

Schematic, Layout, Algorithm, Mechanical Drawing -- whatever the next person in line needs (mfg'r, technician).

### ➤ Analysis = S

simulate it, Cost it out, Total up the power, Build a prototype, Prove that it works (formal or otherwise)

### ➤ Synthesis = Making engineering decisions based on available information in a reasonable amount of time + Lots of creativity!



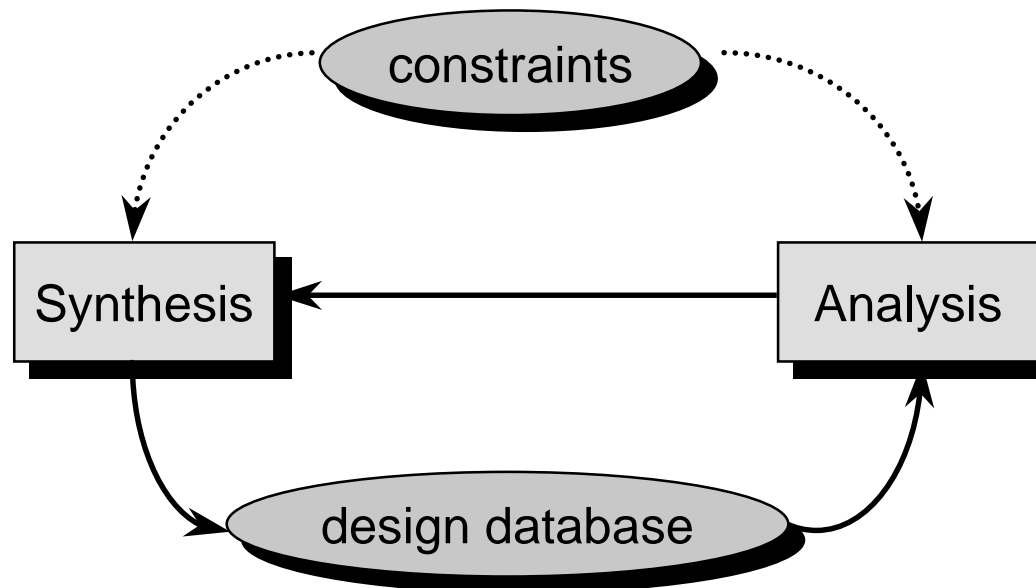
# What can go wrong?

Inefficiency due to:

- *lack of automation*
- *lack of design re-use*

Iteration due to:

- *inexperience (expertise intensive)*
- *sub-optimal tool performance*



Failure due to:

- *Unreasonable constraints*
- *Limited access to the design space*

# What we will learn in 370

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This is *the watershed* course for understanding how all computers work.

And...

Lots of practical stuff including design tools, hardware technologies, etc.

The Future?

The Modern Era -- Computers

The **Tertiary** -- Sequential Logic

The **Jurassic** -- Math Hardware

The **Cambrian** -- Combination Logic

The **Primordial Soup** -- True, False, One, Zero  
(Binary Number Systems)

# Applications of logic design

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- ❑ Conventional computer design
  - CPUs, busses, peripherals
- ❑ Networking and communications
  - phones, modems, routers
- ❑ Embedded products
  - in cars, toys, appliances, entertainment devices
- ❑ Scientific equipment
  - testing, sensing, reporting
- ❑ The world of computing is much much bigger than just PCs!

# CSE 370: concepts/skills/abilities

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- ❑ Understanding the basics of logic design (concepts)
- ❑ Understanding sound design methodologies (concepts)
- ❑ Modern specification methods (concepts)
- ❑ Familiarity with a full set of CAD tools (skills)
- ❑ Appreciation for the differences and similarities (abilities) in hardware and software design

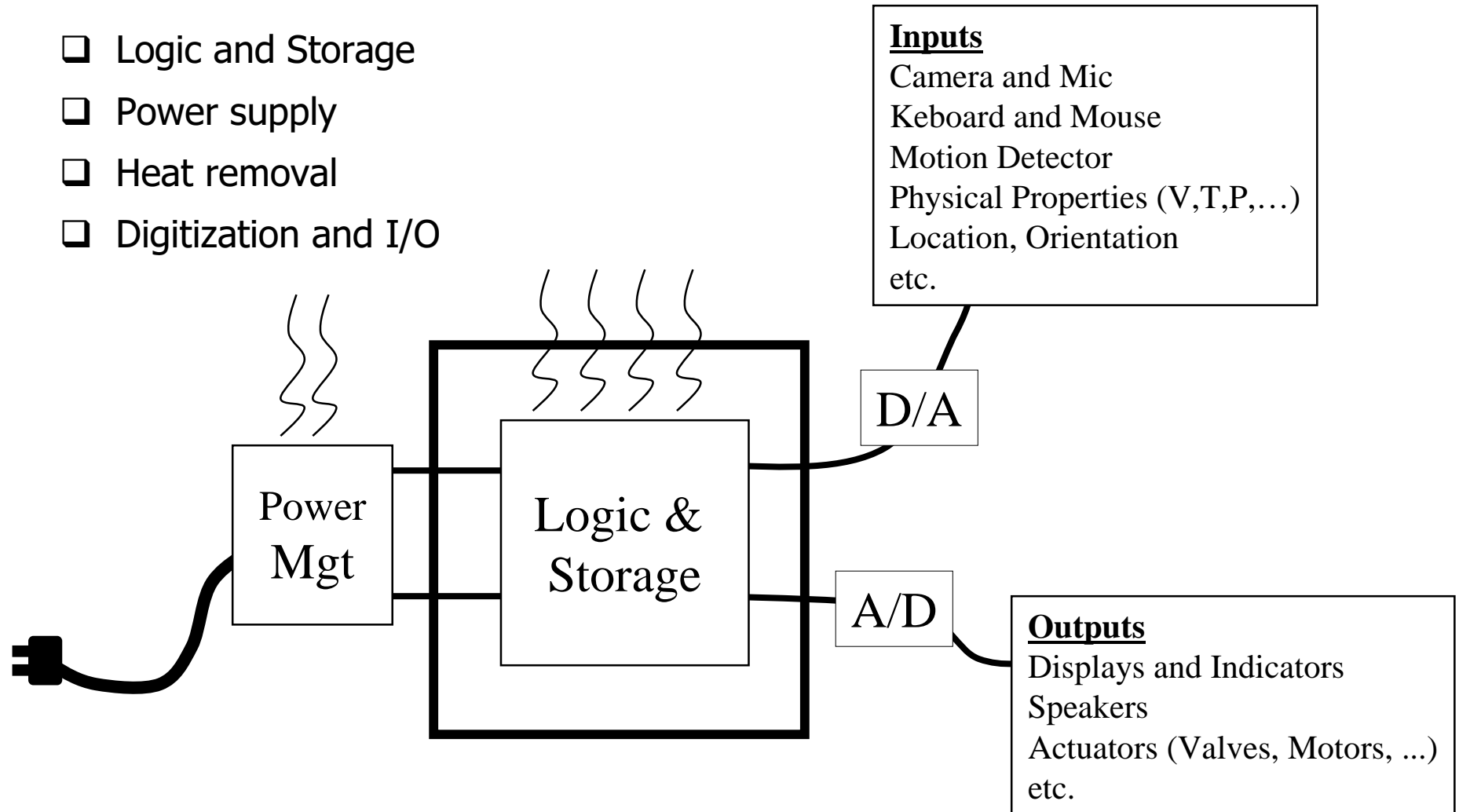
New ability: Understand the basic tools and techniques for designing digital systems, including special purpose hardware and general purpose computers, using a variety of implementation technologies such as LSI and programmable logic

# Computation: abstract vs. implementation

- ❑ Up to now, computation has been a mental exercise (paper, programs)
- ❑ This class is about physically implementing computation using physical devices that use voltages to represent logical values dealing with constraints
- ❑ Basic units of computation are:
  - representation: "0", "1" or True/False
  - assignment:  $x = y$
  - data operations:  $x + y - 5$
  - control:
    - sequential statements:  $A; B; C$
    - conditionals:  $\text{if } x == 1 \text{ then } y$
    - loops:  $\text{for } ( i = 1 ; i == 10, i++)$
    - procedures:  $A; \text{proc}(\dots); B;$
- ❑ We will study how each of these are implemented in hardware and composed into computational structures

# System Components

- ❑ Logic and Storage
- ❑ Power supply
- ❑ Heat removal
- ❑ Digitization and I/O



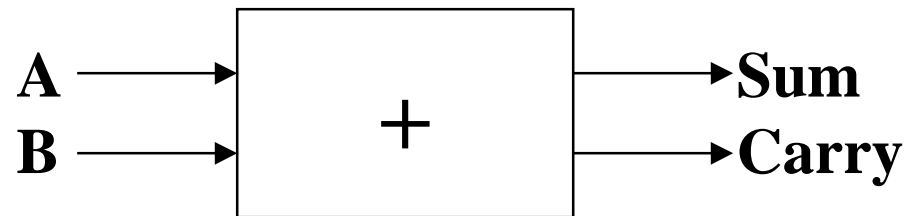
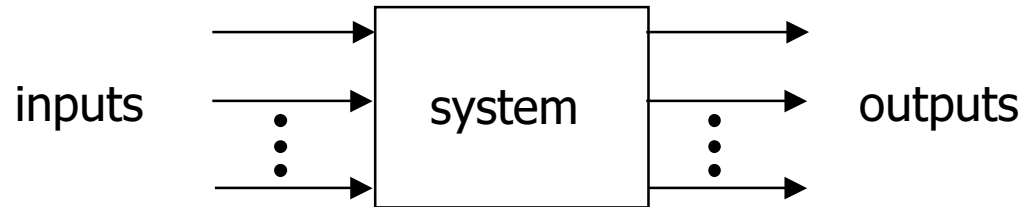
# Representation of digital designs

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- Physical devices (transistors, relays)
  - Switches
  - Truth tables
  - Boolean algebra
  - Gates
  - Waveforms
  - Finite state behavior
  - Register-transfer behavior
  - Concurrent abstract specifications
- scope of CSE 370
-

# Combinational vs. sequential digital circuits

- A simple model of a digital system is a unit with inputs and outputs:



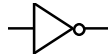
- Combinational means "memory-less"
  - a digital circuit is combinational if its output values only depend on its input values
  - A given input will also have the same steady state output, no matter what has come before



# Combinational logic symbols

- Common combinational logic systems have standard symbols called logic gates

- Buffer, NOT



- AND, NAND



- OR, NOR



easy to implement  
with CMOS transistors  
(the switches we have  
available and use most)

# Sequential logic

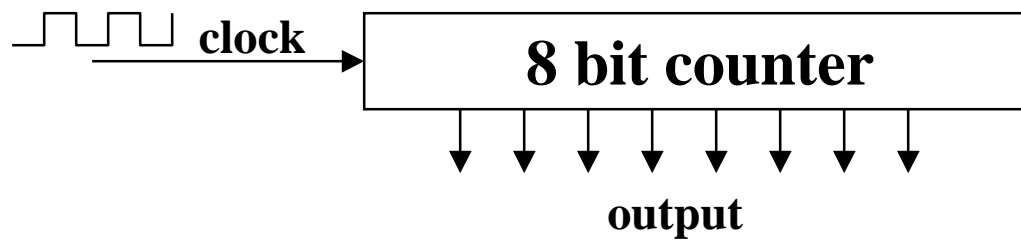
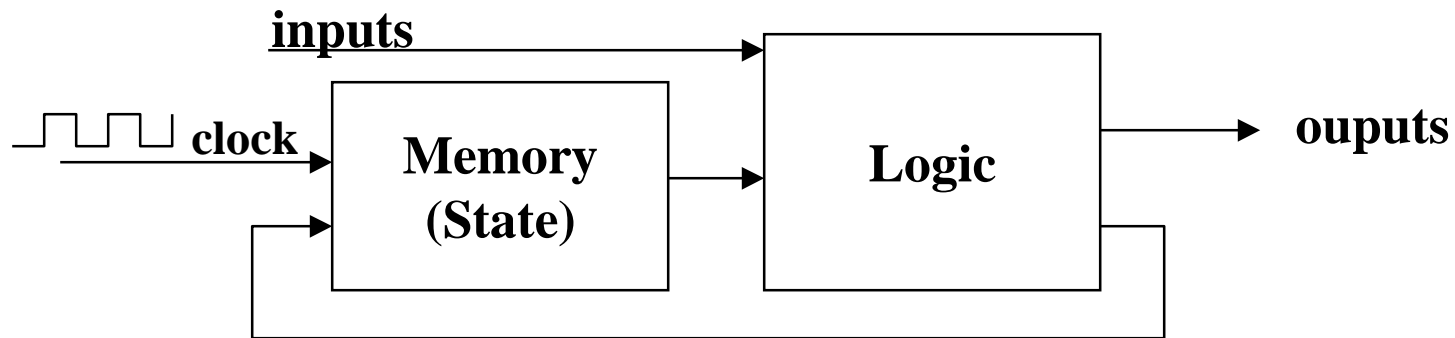
- ❑ Sequential systems
  - exhibit behaviors (output values) that depend not only on the current input values, but also on previous (sequence of) input values
- ❑ In reality, all real circuits are sequential
  - because the outputs do not change instantaneously after an input change why not, and why is it then sequential?
- ❑ A fundamental abstraction of digital design is to reason (mostly) about steady-state behaviors
  - look at the outputs only after sufficient time has elapsed for the system to make its required changes and settle down



**Output depends on who got there first -- sequence!**  
**Problem: How to guarantee GrantA & GrantB == 0**

# One at a Time...Synchronous Sequential

- ❑ Ignore changes on inputs except during a specific time window
  - Special signal -- the clock determines when to consider new inputs
- ❑ Isolate "memory" from "logic"
  - during a clock transition, compute new state and outputs using inputs and old state.
  - Assume inputs and state are stable during the clock transition



# Summary

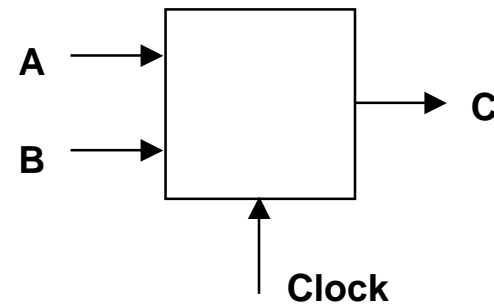
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## □ Combinational:

- input A, B
- wait for clock edge
- observe C
- wait for another clock edge
- observe C again: will stay the same

## □ Sequential:

- input A, B
- wait for clock edge
- observe C
- wait for another clock edge
- observe C again: may be different



# Abstractions and Simplifications

- ❑ How we make life easy for ourselves
  - Quantization and Amplification -- Lossless propagation of data in exchange for low bandwidth
  - Generally stick to binary representation
  - Basic logic gates as building blocks in exchange for complex transistor networks (density)
  - Clearly defined state transitions on clock edges. Assume inputs stable during clock transistions.
  
- ❑ Tools for managing complexity that we will see
  - Truth tables and Boolean algebra to represent combinational logic
  - Binary arithmetic to manipulate groups of signals
  - State diagrams to represent sequential logic
  - Hardware description languages to represent digital logic
  - Waveforms to represent temporal behavior

## An example

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- ❑ Calendar subsystem: number of days in a month (to control watch display)
  - used in controlling the display of a wrist-watch LCD screen
  
  - inputs: month, leap year flag
  - outputs: number of days
  - No need to remember previous entries or results to give correct output

## Implementation in software

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```
integer number_of_days ( month, leap_year_flag) {
    switch (month) {
        case 1: return (31);
        case 2: if (leap_year_flag == 1) then return (29)
                else return (28);
        case 3: return (31);
        ...
        case 12: return (31);
        default: return (0);
    }
}
```

**Clue: Return value is unrelated to previous result call**

# Implementation as a combinational system

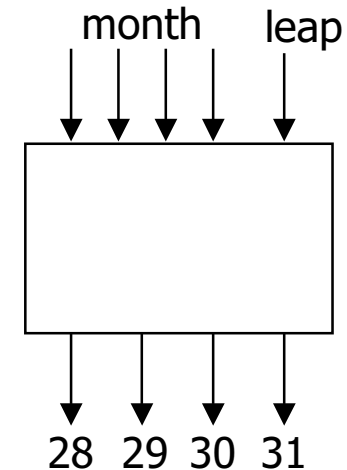
## □ Encoding:

- how many bits for each input/output?
- binary number for month
- four wires for 28, 29, 30, and 31

## □ Behavior:

- combinational
- truth table specification

month	leap	28	29	30	31
0001	-	0	0	0	1
0010	0	1	0	0	0
0010	1	0	1	0	0
0011	-	0	0	0	1
0100	-	0	0	1	0
...					
1100	-	0	0	0	1
1101	-	0	0	0	0
111-	-	0	0	0	0
0000	-	0	0	0	0





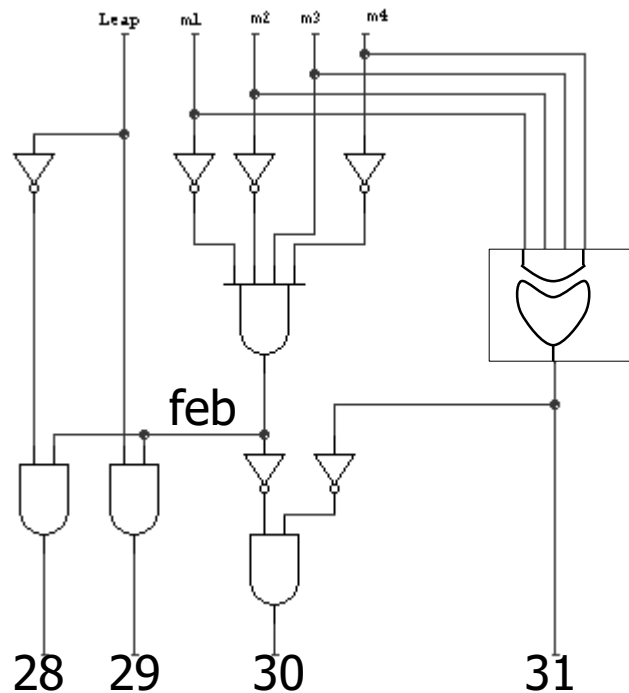
# Combinational example (cont'd)

## □ Truth-table to logic

- 28 = 1 when month=0010 and leap=0
- $28 = m1' \cdot m2' \cdot m3 \cdot m4' \cdot leap'$

- 31 = 1 when month=0001 or month=0011 or ... month=1100
- $31 = (m1' \cdot m2' \cdot m3' \cdot m4) + (m1' \cdot m2' \cdot m3 \cdot m4) + \dots (m1 \cdot m2 \cdot m3' \cdot m4')$
- 31 = can we simplify more?

symbol for or      symbol for and      symbol for not



month	leap	28	29	30	31
0001	-	0	0	0	1
0010	0	1	0	0	0
0010	1	0	1	0	0
0011	-	0	0	0	1
0100	-	0	0	1	0
...					
1100	-	0	0	0	1
1101	-	-	-	-	-
111-	-	-	-	-	-
0000	-	-	-	-	-

**Example of “multilevel logic”  
speed v. cost, parallel v. serial**

## Another example

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- ❑ Door combination lock:
  - Enter three values, if they match the combination then the door opens;
  - inputs: values V1, V2, V3
  - outputs: door open/close
- ❑ Unspecified
  - What happens if wrong combination is entered?
  - How do we know when a new combination is starting?
  - Are V1-V3 on different sets of wires or are they sequenced on the same set of wires? If sequenced how do we know when a new number has been entered?
  - How does it get locked again after being opened?

# Refinement of Specification

- ❑ All values come in on a single set of wires:

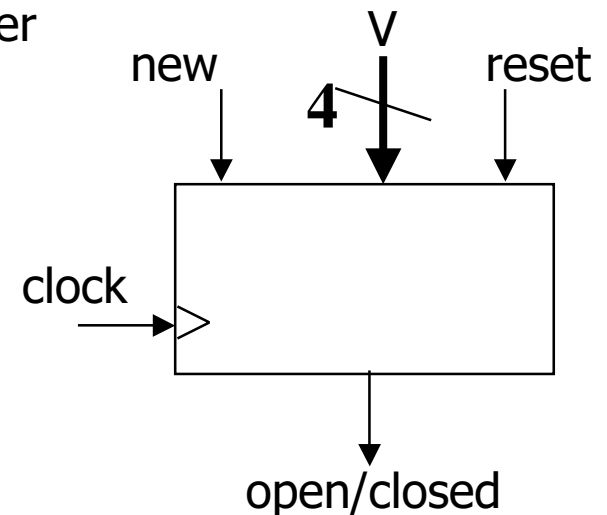
So its sequential so we need a clock to control the sequence.

- ❑ Any other problems?

Yes -- might not get new value every clock cycle – need another input NEW – we could use new for clock, but usually want our subsystem to be in synch with other sub-systems so we don't have user provided clocks.

- ❑ How to handle wrong combination and relocking?

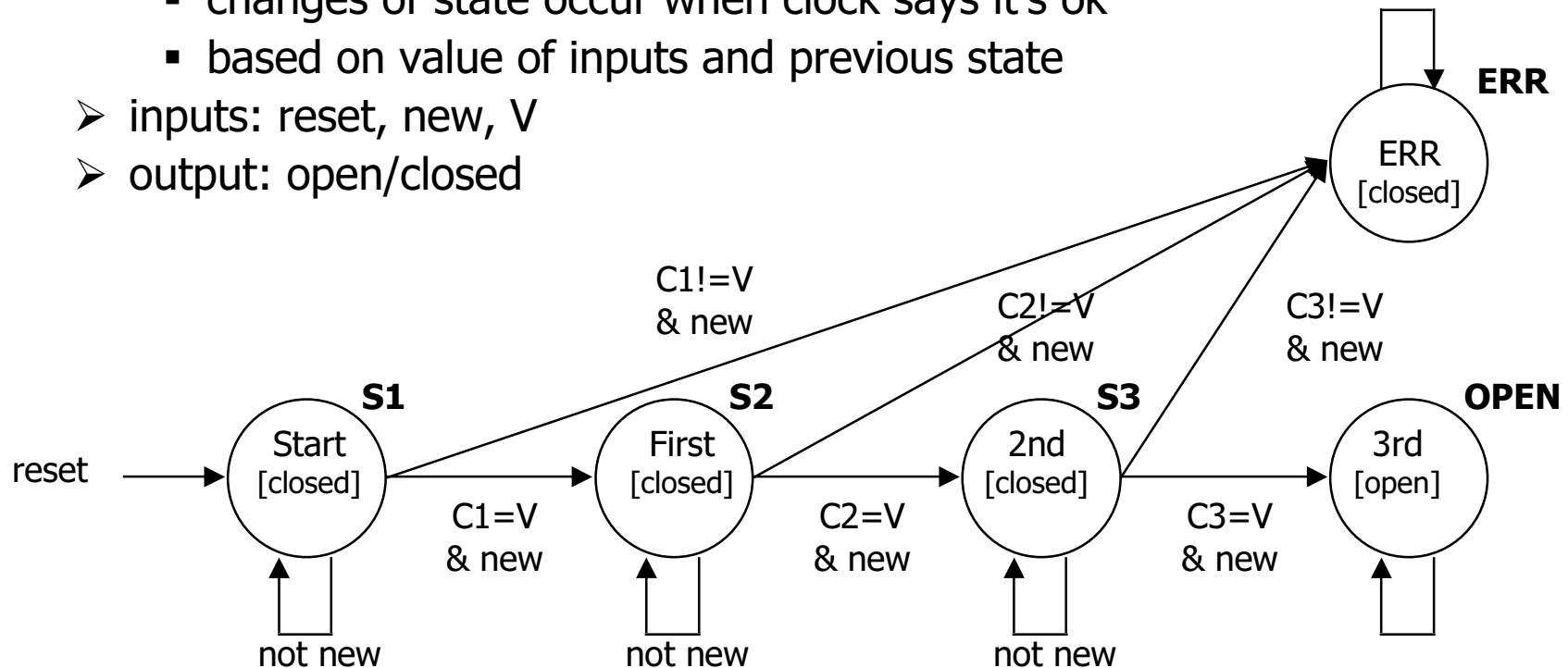
Add a reset input



# Sequential example (cont'd): abstract control

## □ Finite-state diagram

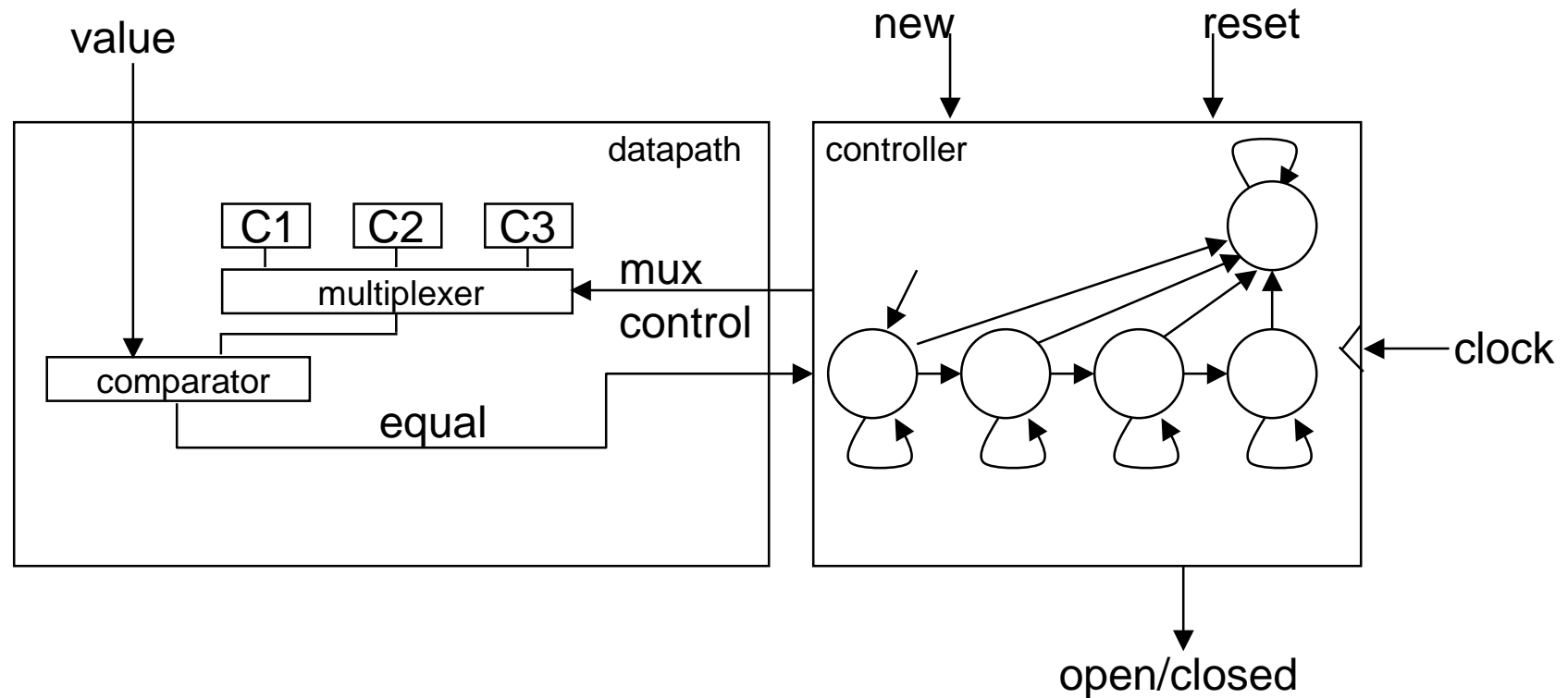
- states: 5 states: S1, S2, S3, Open, Err
  - represents synchronization points in execution of machine
  - each state has outputs
- transitions: 6 from state to state, 5 self transitions, 1 global
  - changes of state occur when clock says it's ok
  - based on value of inputs and previous state
- inputs: reset, new, V
- output: open/closed



# Sequential example implementation

## □ Internal structure: Datapath and Control

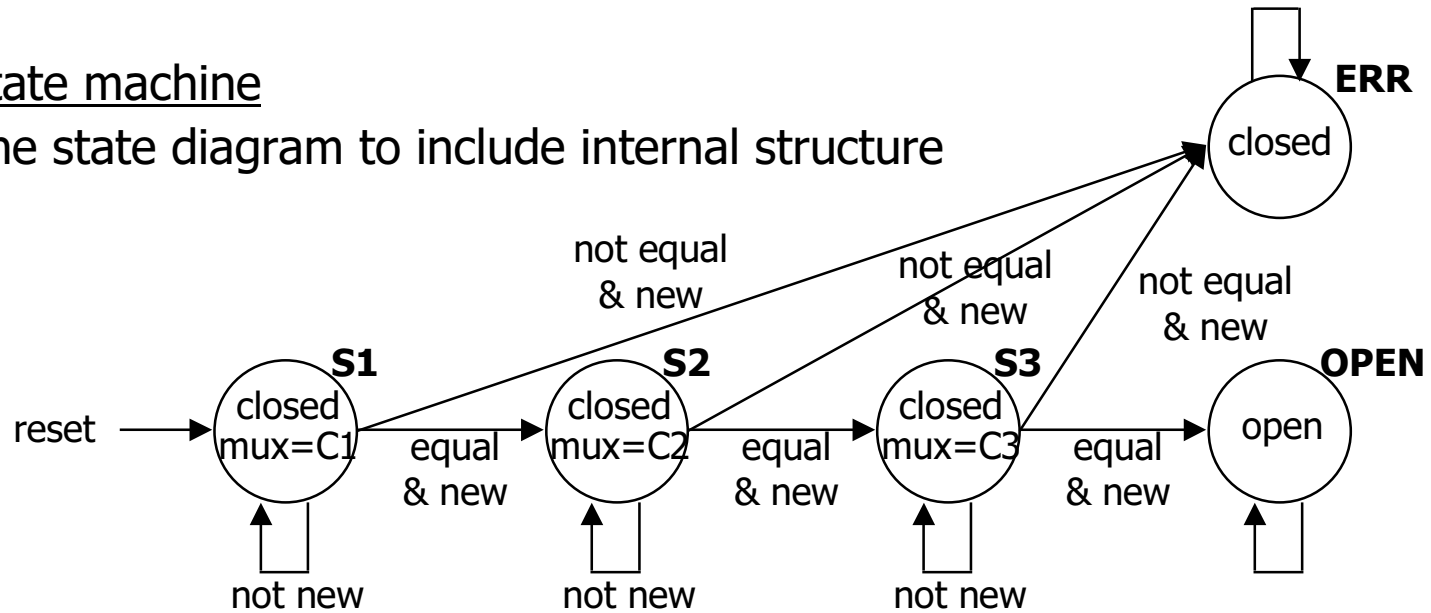
- data-path (combinational logic for multi-bit numbers)
  - storing and operating on n-bit numbers
- control
  - decide what to do next based on datapath results, state, and inputs



# Implementing the Controller

## □ Finite-state machine

➤ refine state diagram to include internal structure



➤ generate state table (much like a truth-table)

reset	new	equal	state	next state	mux	open/closed
1	-	-	-	S1	C1	closed
0	0	-	S1	S1	C1	closed
0	1	0	S1	ERR	-	closed
0	1	1	S1	S2	C2	closed
...						
0	1	1	S3	OPEN	-	open
...						

# Design of Controller

## □ Encode state table

- state can be: S1, S2, S3, OPEN, or ERR
  - needs at least 3 bits to encode: 000, 001, 010, 011, 100
  - and as many as 5: 00001, 00010, 00100, 01000, 10000
  - choose 4 bits: 0001, 0010, 0100, 1000, 0000
- output mux can be: C1, C2, or C3
  - needs 2 to 3 bits to encode
  - choose 3 bits: 001, 010, 100
- output open/closed can be: open or closed
  - needs 1 or 2 bits to encode
  - choose 1 bits: 1, 0

reset	new	equal	state	next state	mux	open/closed
1	–	–	–	0001	001	0
0	0	–	0001	0001	001	0
0	1	0	0001	0000	–	0
0	1	1	0001	0010	010	0
...						
0	1	1	0100	1000	–	1
...						

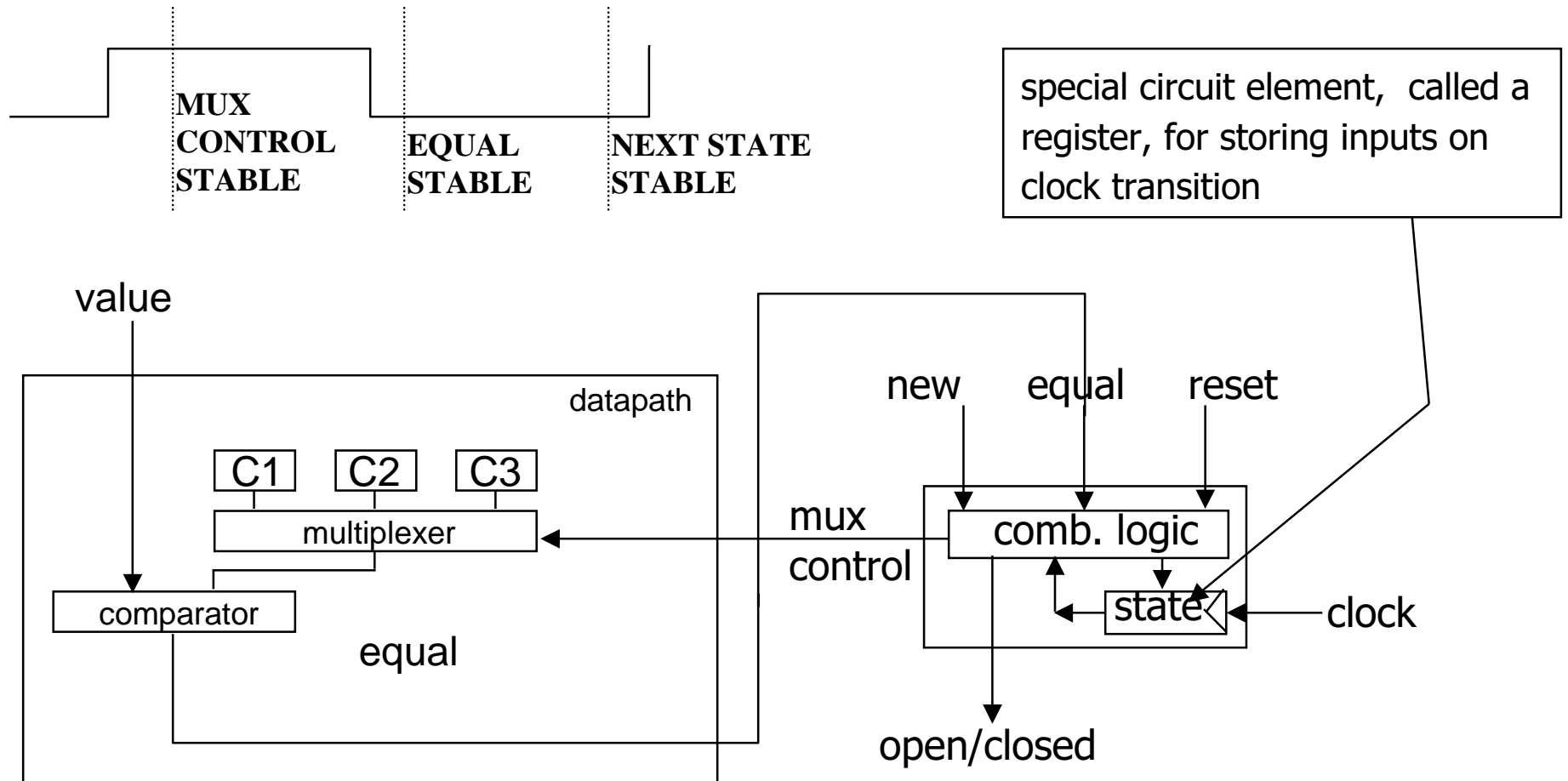
good choice of encoding!

mux is identical to last 3 bits of state

open/closed is identical to first bit of state

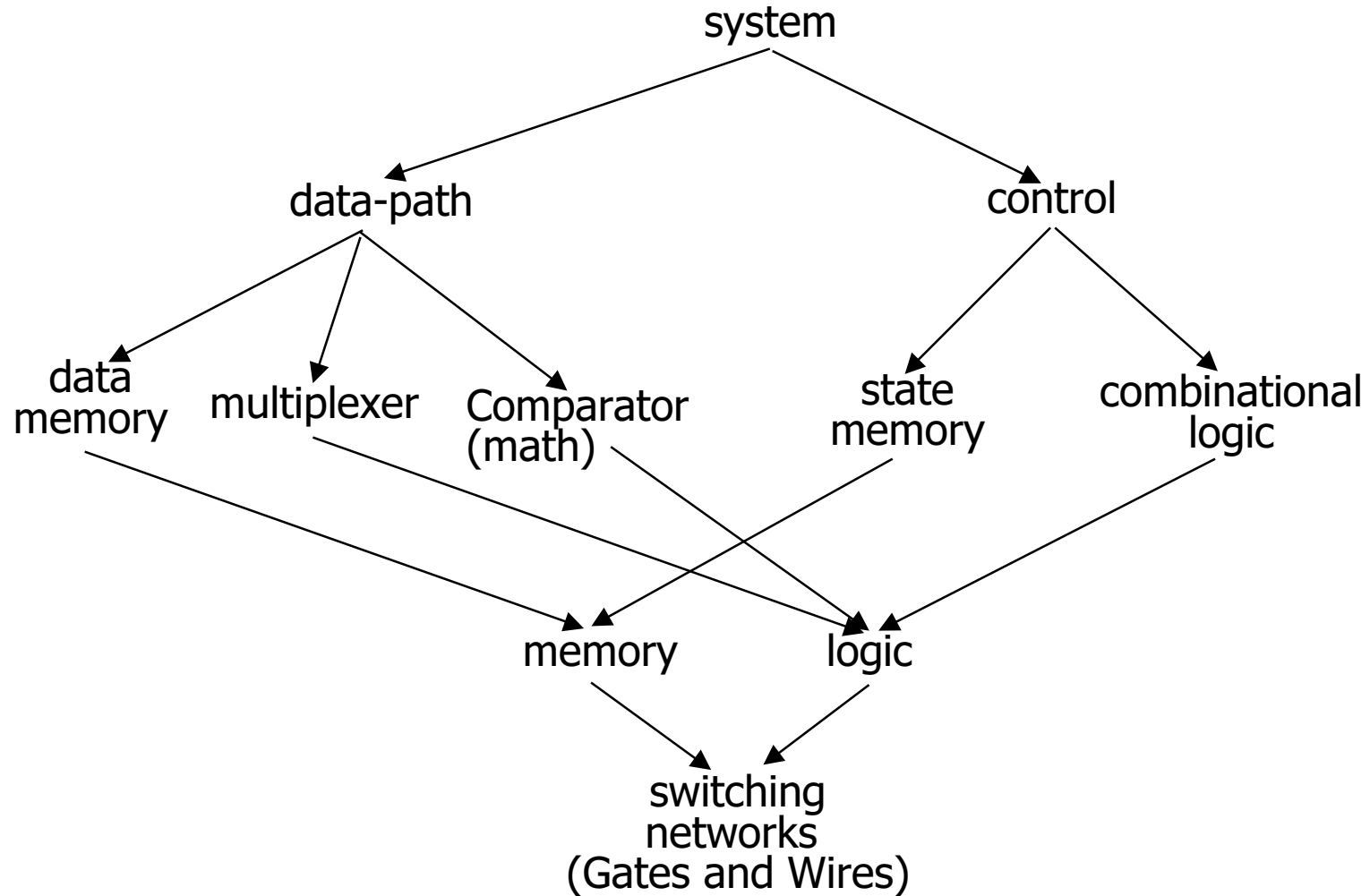
# Implementation of Controller w/ Datapath

- Implementation of the controller





# Design hierarchy



# Summary

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- ❑ That was what the entire course is about (Almost)
  - converting solutions to problems into combinational and sequential networks effectively organizing the design hierarchically
  - doing so with a modern set of design tools that lets us handle large designs effectively
  - taking advantage of optimization opportunities
  
- ❑ Learn where software and hardware come together
  - How does a software program get executed on a digital hardware system?

# What is happening now in digital design?

- ❑ Big change in the way industry does hardware design over last few years
  - larger and larger designs
  - shorter and shorter time to market
  - cheaper and cheaper products
- ❑ Scale
  - pervasive use of computer-aided design tools over hand methods
  - multiple levels of design representation
- ❑ Time
  - emphasis on abstract design representations
  - programmable rather than fixed function components
  - automatic synthesis techniques
  - importance of sound design methodologies
- ❑ Cost
  - higher levels of integration
  - use of simulation to debug designs

# Some Basics. Lecture Overview

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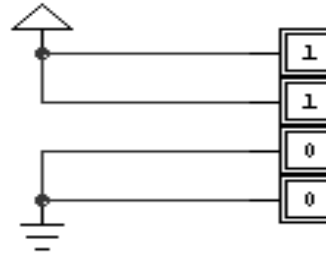
## □ Today

- The basics: Electronics, binary numbers, base conversion
- Number systems
  - 2's complement numbers
- Combinational logic
  - Logic functions and truth tables

# The basics → binary numbers

## □ Binary (base 2) Number Representation

- low V → 0
- hi V → 1



## □ Base conversion (binary, octal, decimal, hexadecimal)

- Positional number system
  - $101_2 = 1 \times 2^2 + 0 \times 2^1 + 1 \times 2^0 = 5_{10}$
  - $101_8 = 1 \times 8^2 + 0 \times 8^1 + 1 \times 8^0 = 65_{10}$
  - $101_{16} = 1 \times 16^2 + 0 \times 16^1 + 1 \times 16^0 = 257_{10}$
- Conversion between binary/octal/hex
  - Binary: 10011110001
  - Octal: 10 | 011 | 110 | 001 = 2361<sub>8</sub>
  - Hex: 100 | 1111 | 0001 = 4F1<sub>16</sub>

## □ Addition and subtraction are trivial, but worth practicing

- See Katz, appendix A

# Number systems

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- ❑ How do we write negative binary numbers?
- ❑ Historically: Three approaches
  - Sign and magnitude
  - One's complement
  - Two's complement
- ❑ Two's complement makes addition and subtraction easy
  - Used almost universally in present-day systems
- ❑ Fractional Numbers
- ❑ Floating Point Representations

# Sign-magnitude Representation

- ❑ The most-significant bit (msb) is the sign digit
  - 0 ≡ positive
  - 1 ≡ negative
- ❑ The remaining bits are the number's magnitude
- ❑ Benefit: easy to perform negation: just flip MSB
- ❑ Problem 1: Two representations for zero
  - 0 = 0000 *and also* -0 = 1000
- ❑ Problem 2: Arithmetic is cumbersome

$$\begin{array}{r} 1001 \\ + \underline{0011} \\ \neq 1100 \end{array}$$

Instead, must change order and perform subtraction

$$\begin{array}{r} 011 \\ - \underline{001} \\ = 0010 \end{array}$$

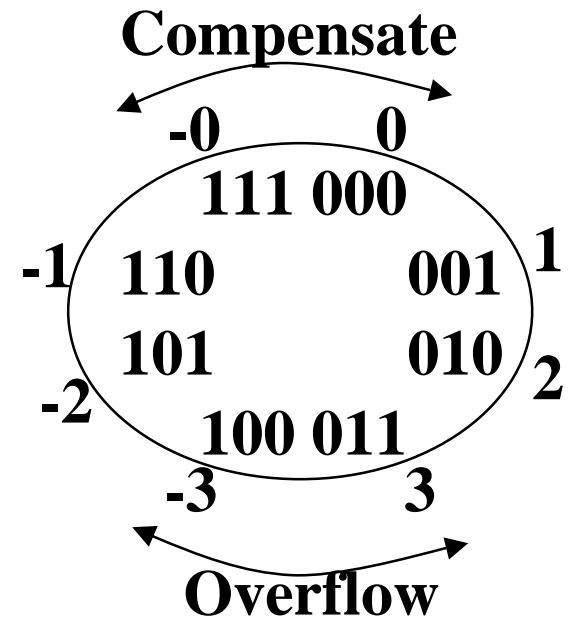
# One's complement

- ❑ Negative number: Bitwise complement of positive number

- $0011 \equiv 3_{10}$
- $1100 \equiv -3_{10}$

- ❑ Almost solves the arithmetic problem

$$\begin{array}{r}
 110 \quad (-1) \\
 + \underline{010} \quad (+2) \\
 \neq 000 \\
 + \quad \underline{1} \\
 = \quad 1
 \end{array}$$



- ❑ Remaining problem: Must compensate for two zeros

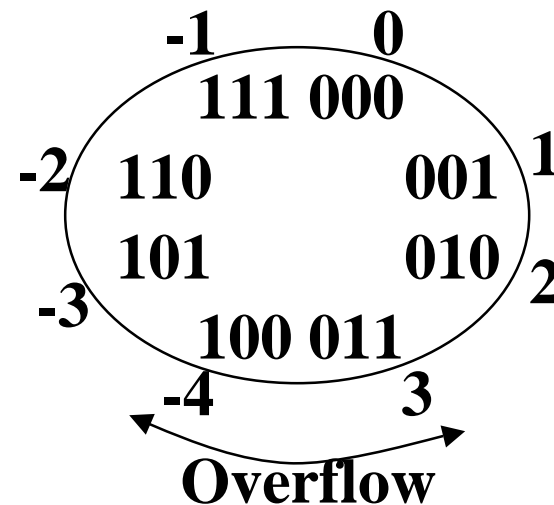
- Add (or subtract) carry when (MSB carry & opposite sign)
- Overflow when (MSB carry & same sign)



# Two's complement

- ❑ Negative number: Bitwise complement plus one
  - $0011 \equiv 3_{10}$
  - $1101 \equiv -3_{10}$
- ❑ Shift the 1's complement number wheel to eliminate -0

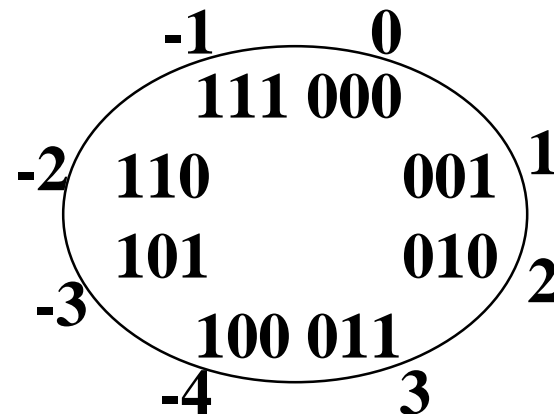
- ❑ msb is still the sign digit
  - $0 \equiv$  positive
  - $1 \equiv$  negative
- ❑ No need to compensate for +/-



# Twos complement (con't)

- ❑ Complementing a complement restores the original number
- ❑ Arithmetic is easy
  - We ignore the carry
    - Same as a full rotation around the wheel
  - Subtraction = negation and addition
    - Easy to implement in hardware

$$\begin{array}{r} 111 \quad (-1) \\ + \underline{010} \quad (+2) \\ = 001 \\ \text{(ignore carry)} \end{array}$$



# Where are we now?

