
CSE 373

Priority queue implementation; Intro to heaps
read: Weiss Ch. 6

slides created by Marty Stepp
<http://www.cs.washington.edu/373/>

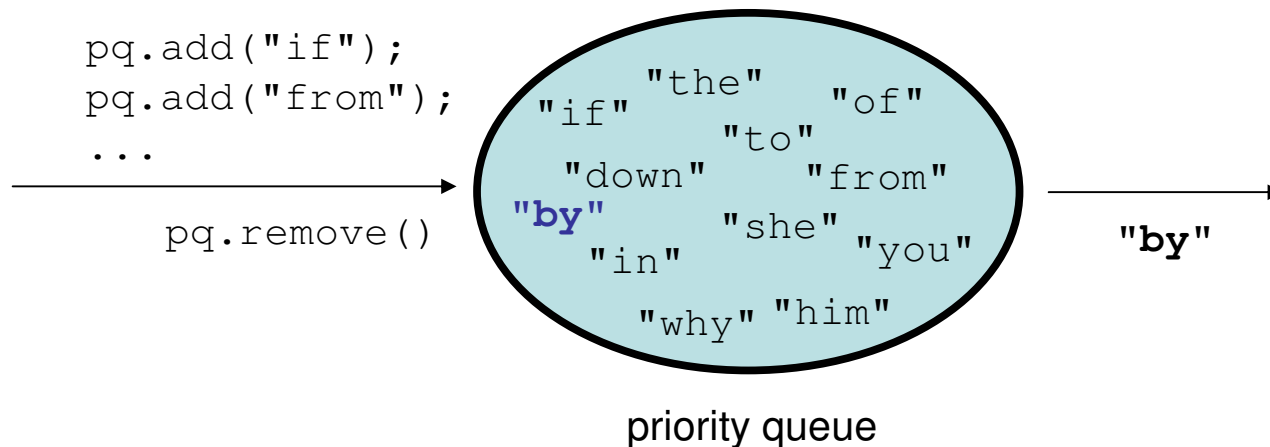
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Prioritization problems

- **print jobs:** CSE lab printers constantly accept and complete jobs from all over the building. We want to print faculty jobs before staff before student jobs, and grad students before undergrad, etc.
- **ER scheduling:** Scheduling patients for treatment in the ER. A gunshot victim should be treated sooner than a guy with a cold, regardless of arrival time. How do we always choose the most urgent case when new patients continue to arrive?
- *key operations we want:*
 - *add an element (print job, patient, etc.)*
 - *get/remove the most "important" or "urgent" element*

Priority Queue ADT

- **priority queue**: A collection of ordered elements that provides fast access to the minimum (or maximum) element.
 - `add` adds in order
 - `peek` returns **minimum** or "highest priority" value
 - `remove` removes/returns **minimum** value
 - `isEmpty, clear, size, iterator` $O(1)$



Unfilled array?

- Consider using an unfilled array to implement a priority queue.
 - `add`: Store it in the next available index, as in a list.
 - `peek`: Loop over elements to find minimum element.
 - `remove`: Loop over elements to find min. Shift to remove.

```
queue.add(9);  
queue.add(23);  
queue.add(8);  
queue.add(-3);  
queue.add(49);  
queue.add(12);  
queue.remove();
```

<i>index</i>	0	1	2	3	4	5	6	7	8	9
<i>value</i>	9	23	8	-3	49	12	0	0	0	0
<i>size</i>	6									

- How efficient is `add`? `peek`? `remove`?
 - $O(1)$, $O(N)$, $O(N)$
 - (`peek` must loop over the array; `remove` must shift elements)

Sorted array?

- Consider using a *sorted* array to implement a priority queue.
 - `add`: Store it in the proper index to maintain sorted order.
 - `peek`: Minimum element is in index [0].
 - `remove`: Shift elements to remove min from index [0].

```
queue.add(9);  
queue.add(23);  
queue.add(8);  
queue.add(-3);  
queue.add(49);  
queue.add(12);  
queue.remove();
```

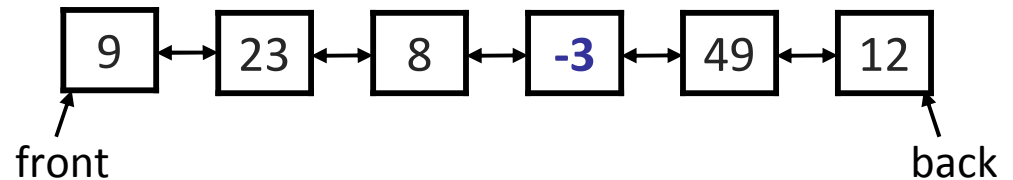
<i>index</i>	0	1	2	3	4	5	6	7	8	9
<i>value</i>	-3	8	9	12	23	49	0	0	0	0
<i>size</i>	6									

- How efficient is `add`? `peek`? `remove`?
 - $O(N)$, $O(1)$, $O(N)$
 - (add and remove must shift elements)

Linked list?

- Consider using a doubly linked list to implement a priority queue.
 - add: Store it at the end of the linked list.
 - peek: Loop over elements to find minimum element.
 - remove: Loop over elements to find min. Unlink to remove.

```
queue.add(9);  
queue.add(23);  
queue.add(8);  
queue.add(-3);  
queue.add(49);  
queue.add(12);  
queue.remove();
```

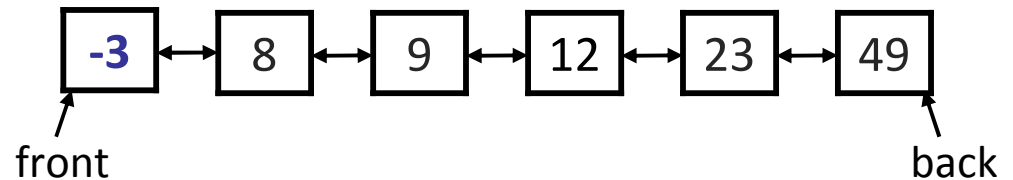


- How efficient is add? peek? remove?
 - $O(1)$, $O(N)$, $O(N)$
 - (peek and remove must loop over the linked list)

Sorted linked list?

- Consider using a *sorted* linked list to implement a priority queue.
 - add: Store it in the proper place to maintain sorted order.
 - peek: Minimum element is at the front.
 - remove: Unlink front element to remove.

```
queue.add(9);  
queue.add(23);  
queue.add(8);  
queue.add(-3);  
queue.add(49);  
queue.add(12);  
queue.remove();
```

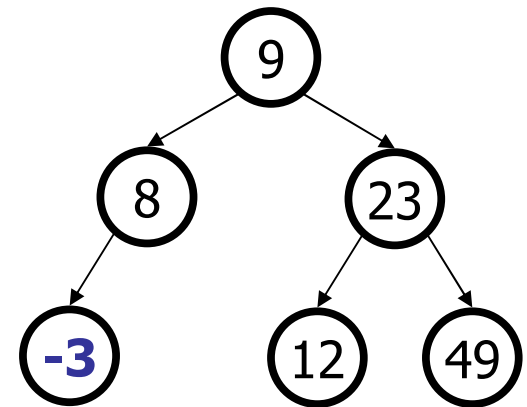


- How efficient is add? peek? remove?
 - $O(N)$, $O(1)$, $O(1)$
 - (add must loop over the linked list to find the proper insertion point)

Binary search tree?

- Consider using a binary search tree to implement a PQ.
 - add: Store it in the proper BST L/R - ordered spot.
 - peek: Minimum element is at the far left edge of the tree.
 - remove: Unlink far left element to remove.

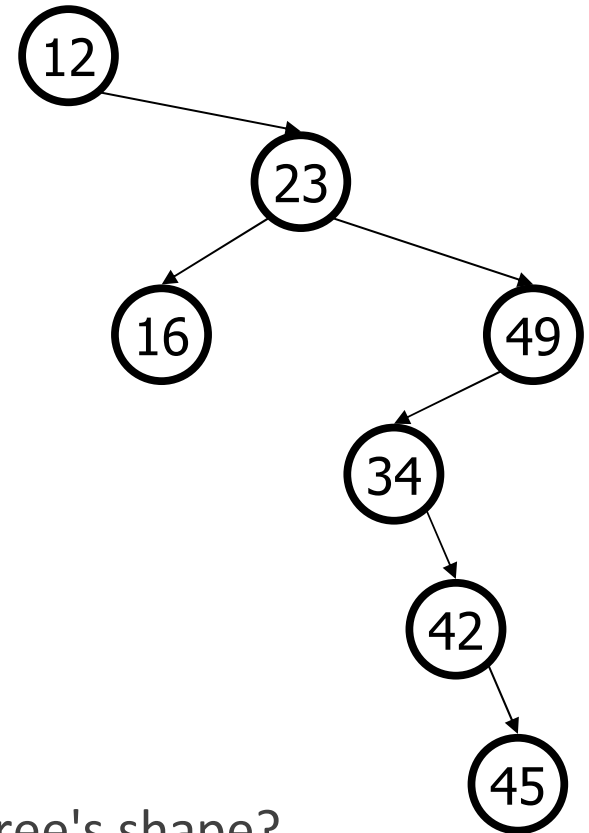
```
queue.add(9);  
queue.add(23);  
queue.add(8);  
queue.add(-3);  
queue.add(49);  
queue.add(12);  
queue.remove();
```



- How efficient is add? peek? remove?
 - $O(\log N)$, $O(\log N)$, $O(\log N)$...?
 - (good in theory, but the tree tends to become unbalanced to the right)

Unbalanced binary tree

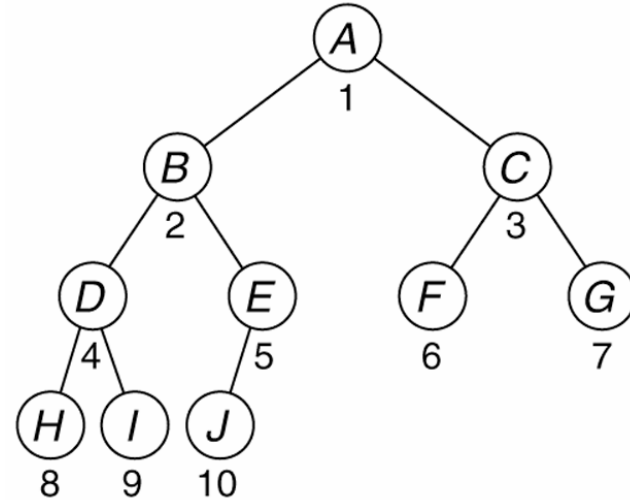
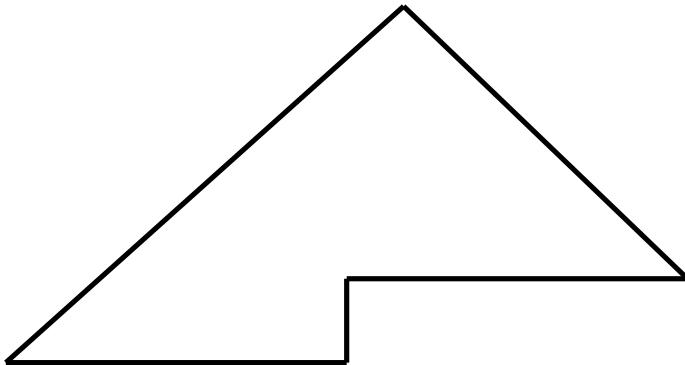
```
queue.add(9);  
queue.add(23);  
queue.add(8);  
queue.add(-3);  
queue.add(49);  
queue.add(12);  
queue.remove();  
queue.add(16);  
queue.add(34);  
queue.remove();  
queue.remove();  
queue.add(42);  
queue.add(45);  
queue.remove();
```



- Simulate these operations. What is the tree's shape?
- A tree that is *unbalanced* has a height close to N rather than $\log N$, which breaks the expected runtime of many operations.

Heaps

- **heap**: A *complete* binary tree with *vertical* ordering.
 - **complete tree**: Every level is full except possibly the lowest level, which must be filled from left to right
 - (i.e., a node may not have any children until all possible siblings exist)

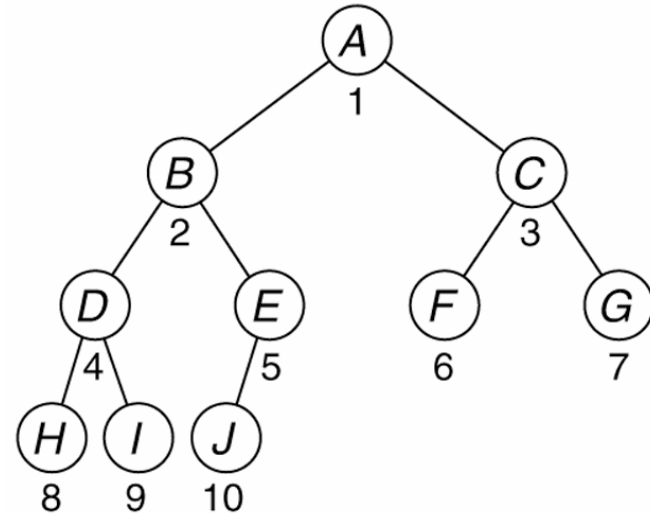


Heap ordering

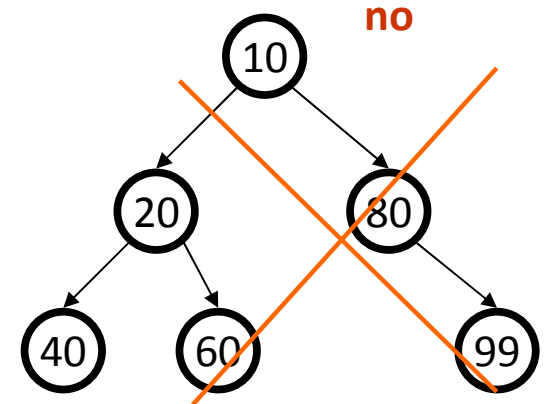
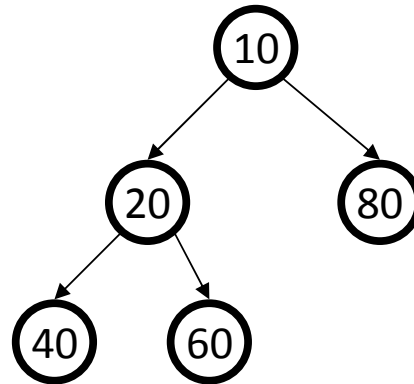
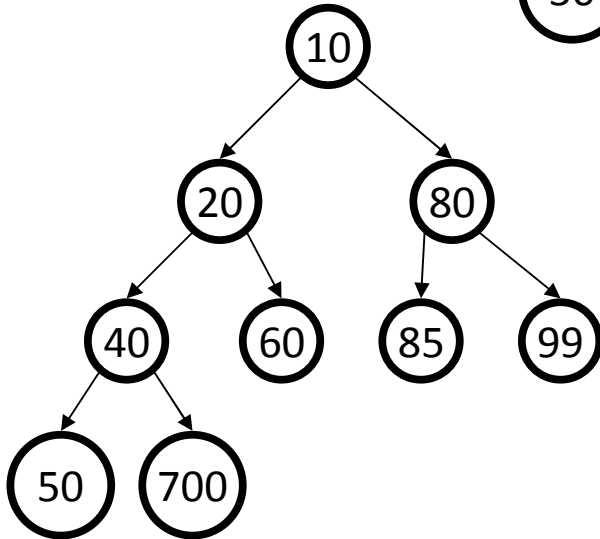
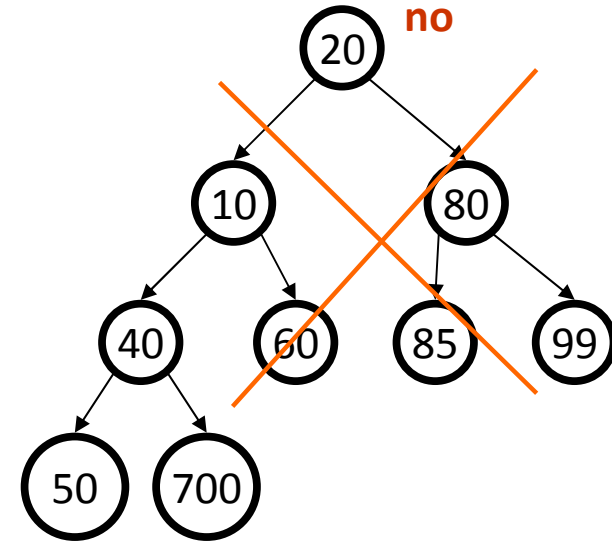
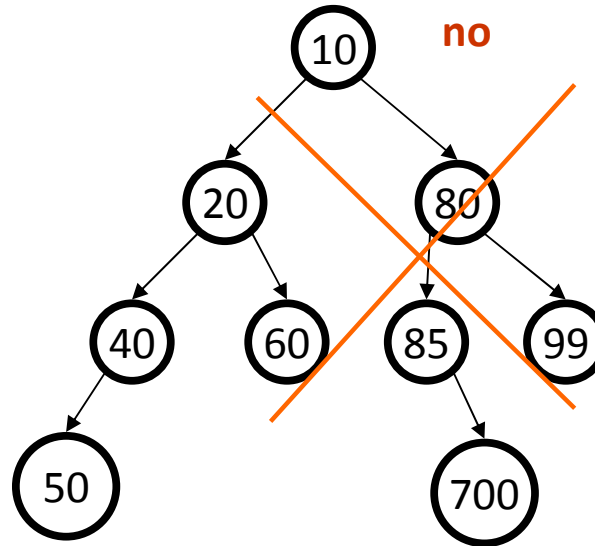
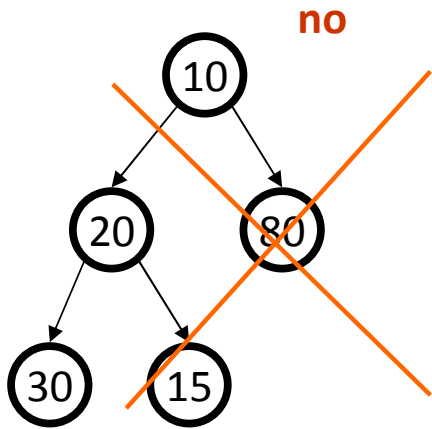
- **heap ordering:** If $P \leq X$ for every element X with parent P .
 - Parents' values are always smaller than those of their children.
 - Implies that minimum element is always the root (a "min-heap").
 - variation: "max-heap" stores largest element at root, reverses ordering
 - Is a heap a BST? How are they related?



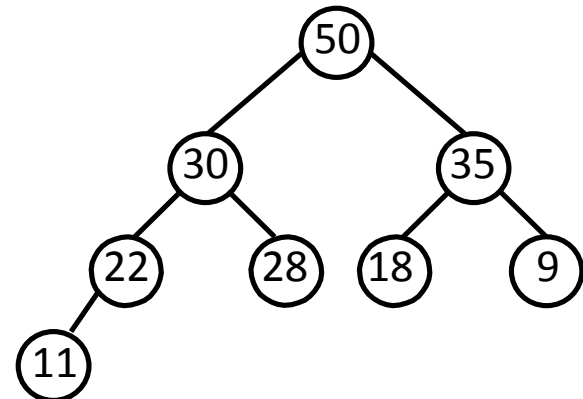
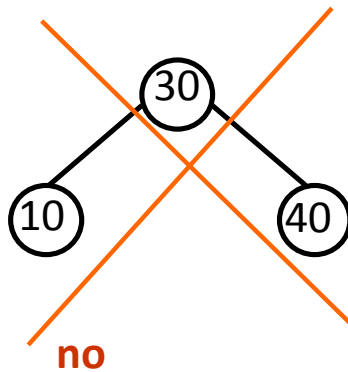
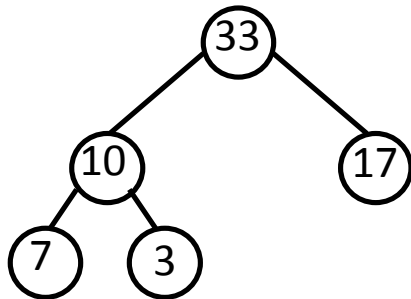
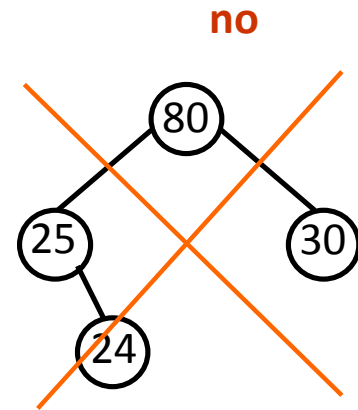
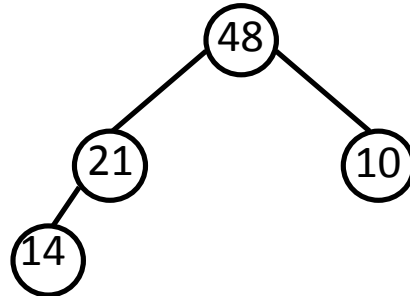
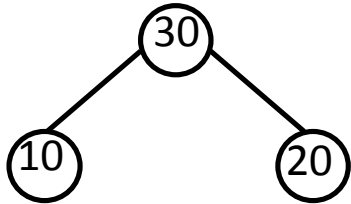
$$P \leq X$$



Which are min-heaps?

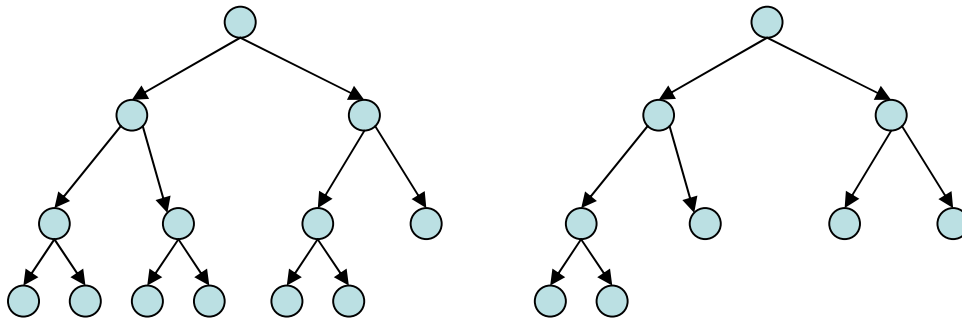


Which are max-heaps?



Heap height and runtime

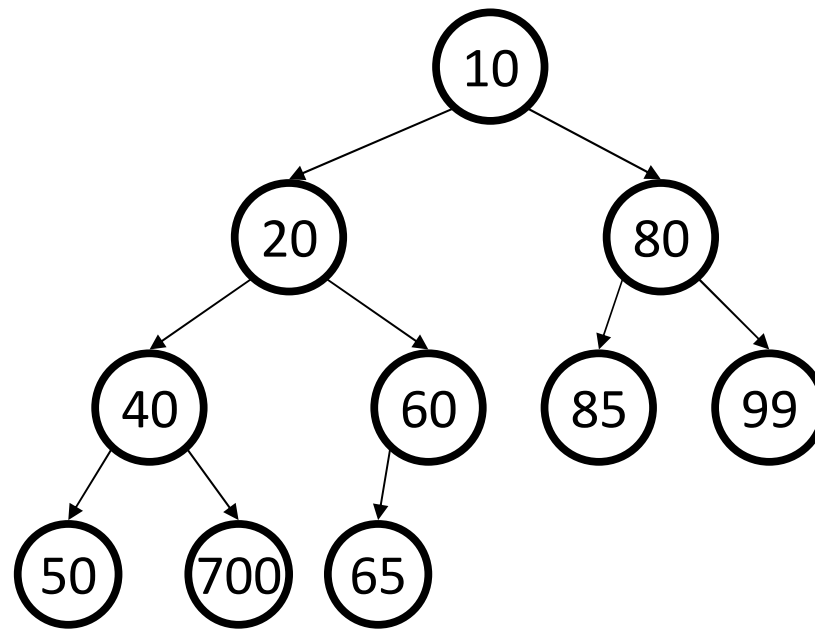
- The height of a complete tree is always $\log N$.
 - How do we know this for sure?
- Because of this, if we implement a priority queue using a heap, we can provide the following runtime guarantees:
 - add: $O(\log N)$
 - peek: $O(1)$
 - remove: $O(\log N)$



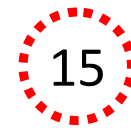
n -node complete tree
of height h :
 $2^h \leq n \leq 2^{h+1} - 1$
 $h = \lfloor \log n \rfloor$

The add operation

- When an element is added to a heap, where should it go?
 - Must insert a new node while maintaining heap properties.
 - `queue.add(15);`

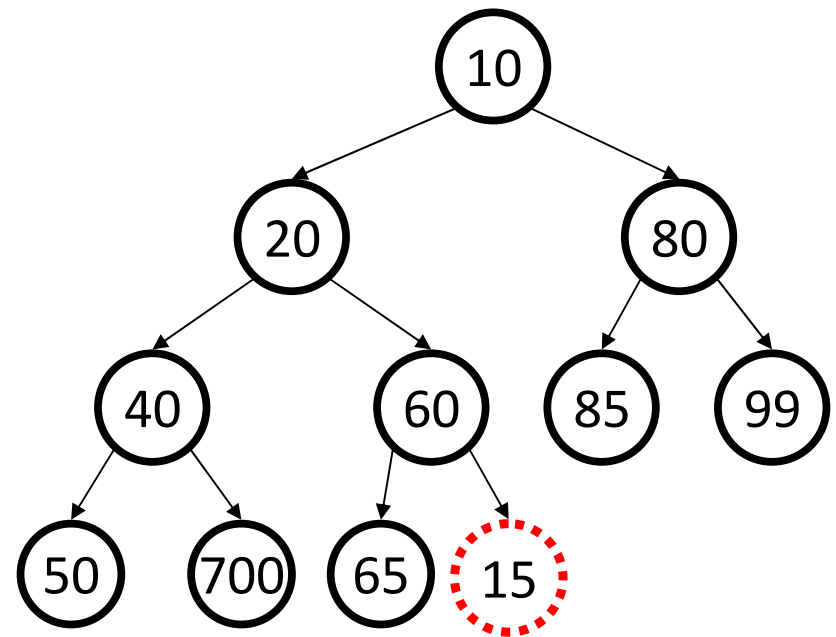
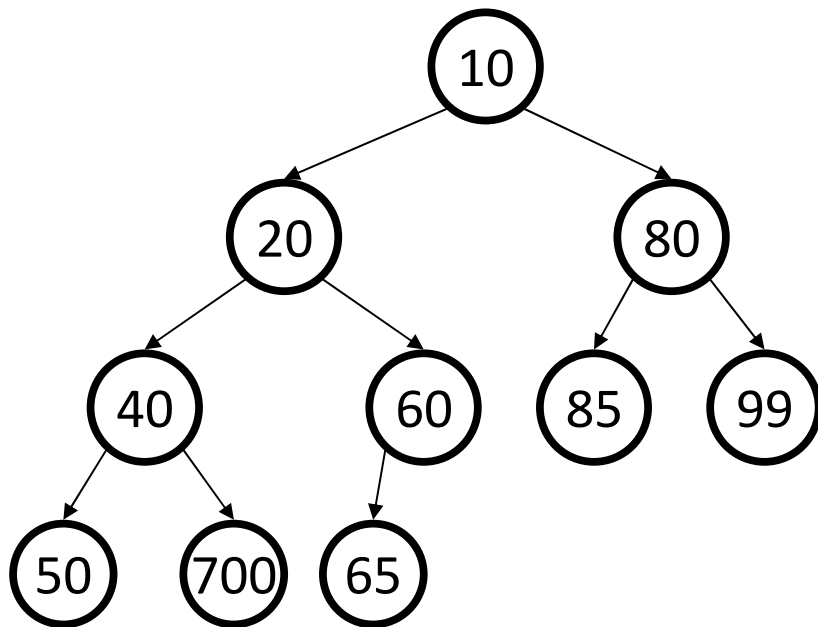


new node



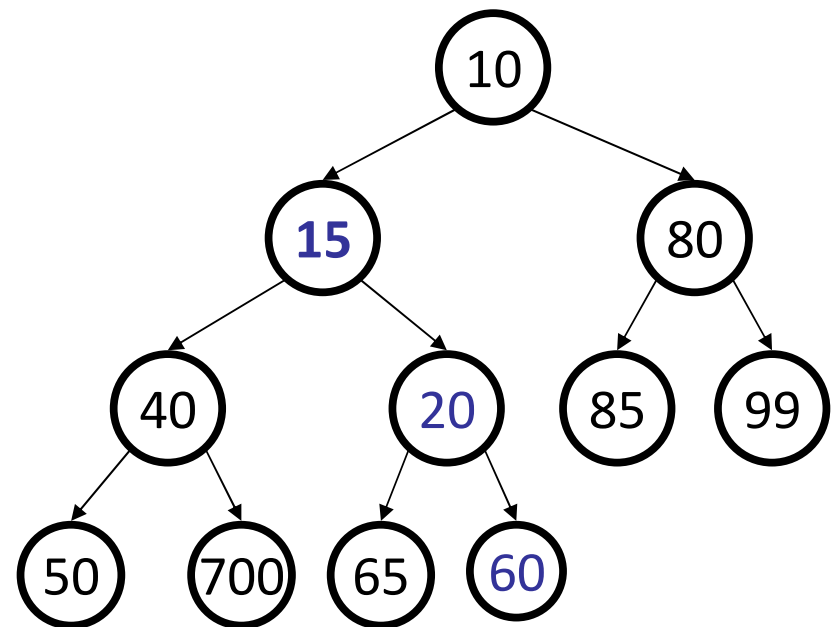
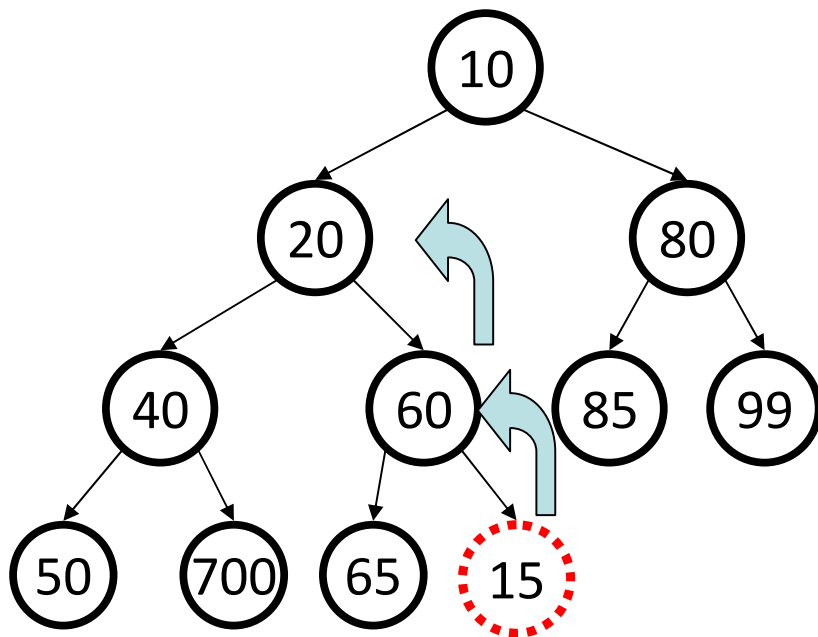
The add operation

- When an element is added to a heap, it should be initially placed as the *rightmost leaf* (to maintain the completeness property).
 - But the heap ordering property becomes broken!



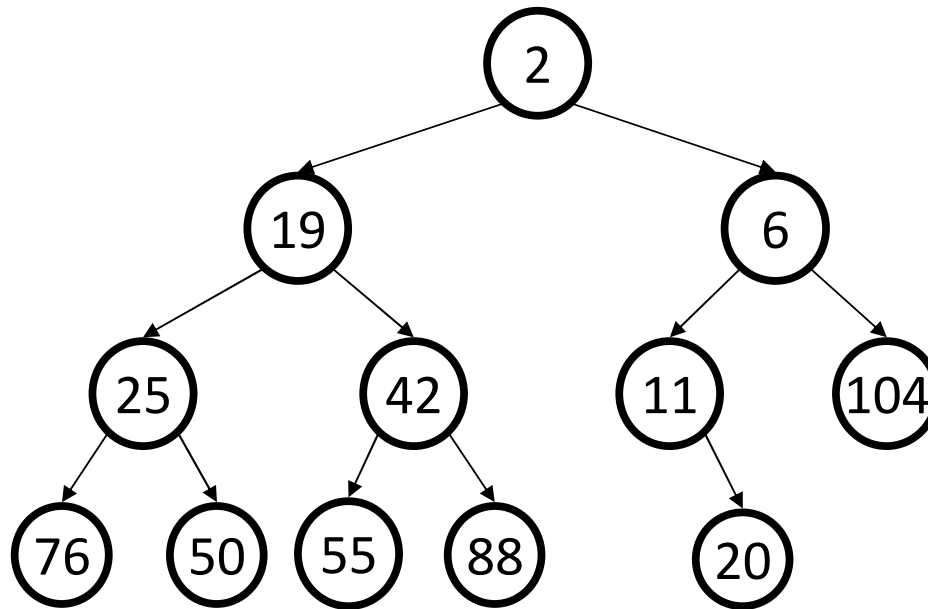
"Bubbling up" a node

- **bubble up:** To restore heap ordering, the newly added element is shifted ("bubbled") up the tree until it reaches its proper place.
 - Weiss: "*percolate up*" by swapping with its parent
 - How many bubble-ups are necessary, at most?



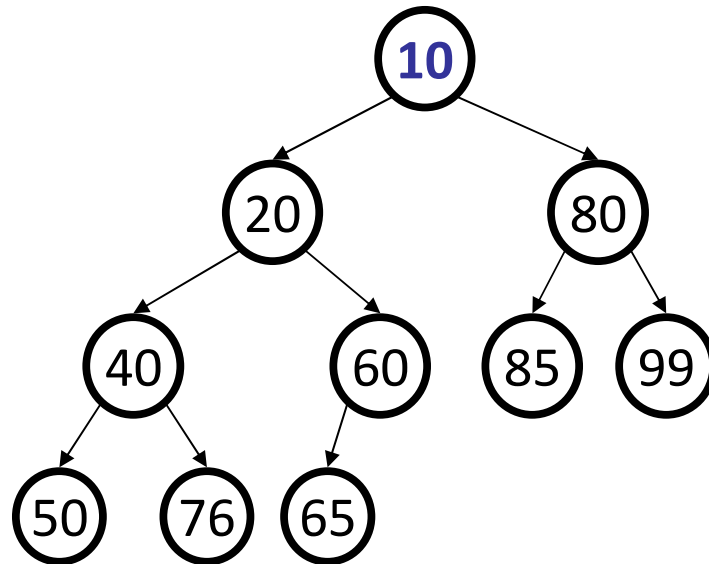
Bubble-up exercise

- Draw the tree state of a min-heap after adding these elements:
 - 6, 50, 11, 25, 42, 20, 104, 76, 19, 55, 88, 2



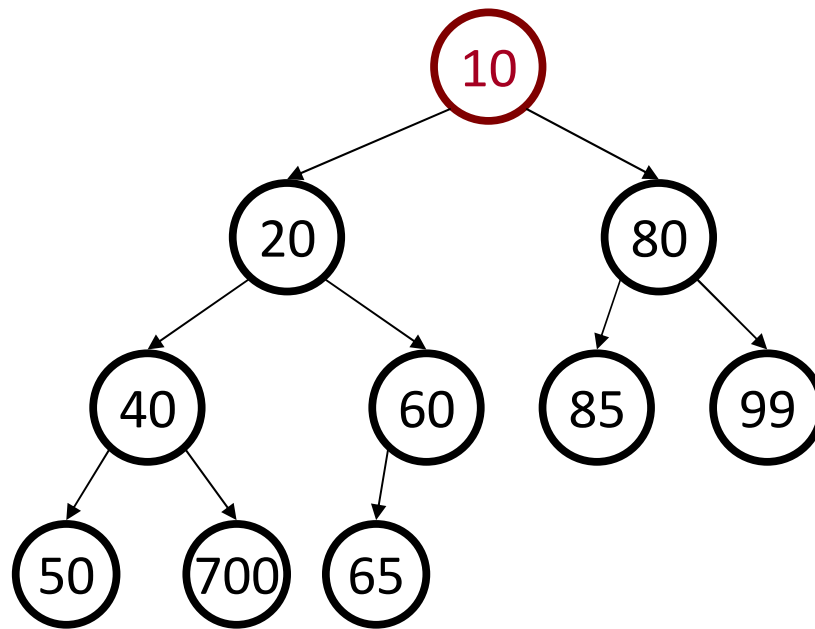
The peek operation

- A peek on a min-heap is trivial to perform.
 - because of heap properties, minimum element is always the root
 - $O(1)$ runtime
- Peek on a max-heap would be $O(1)$ as well (return max, not min)



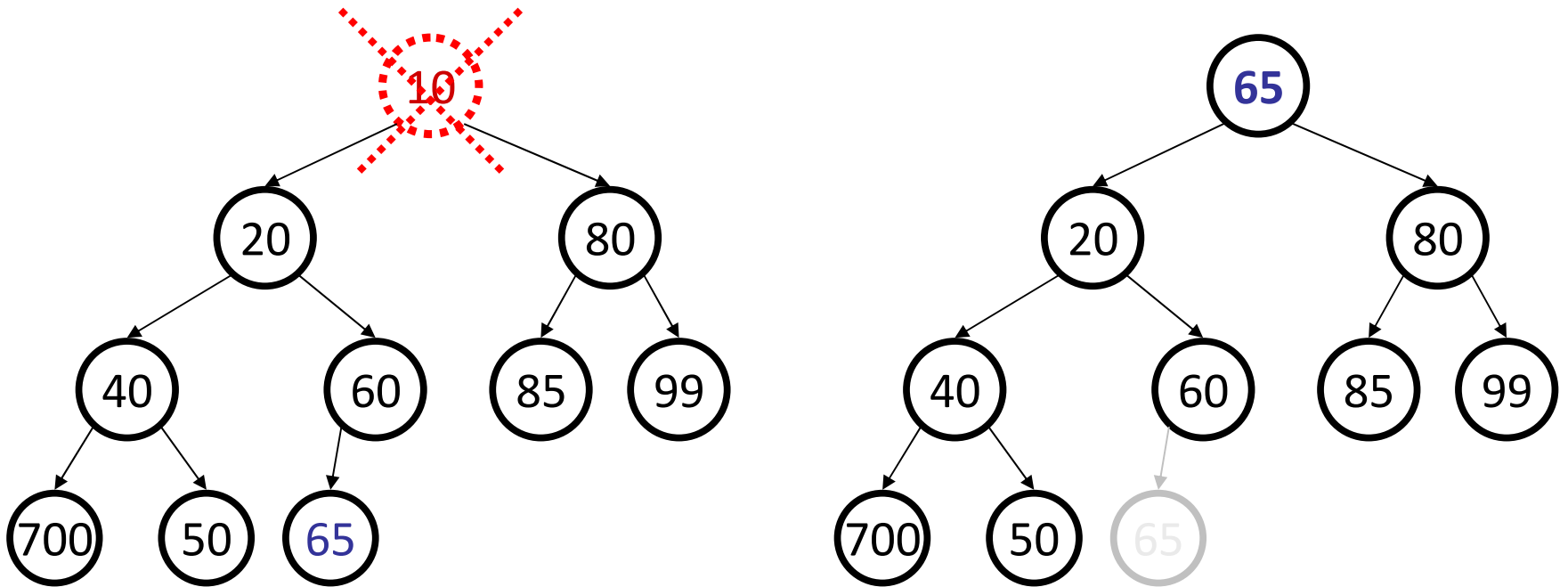
The remove operation

- When an element is removed from a heap, what should we do?
 - The root is the node to remove. How do we alter the tree?
 - `queue.remove()` ;



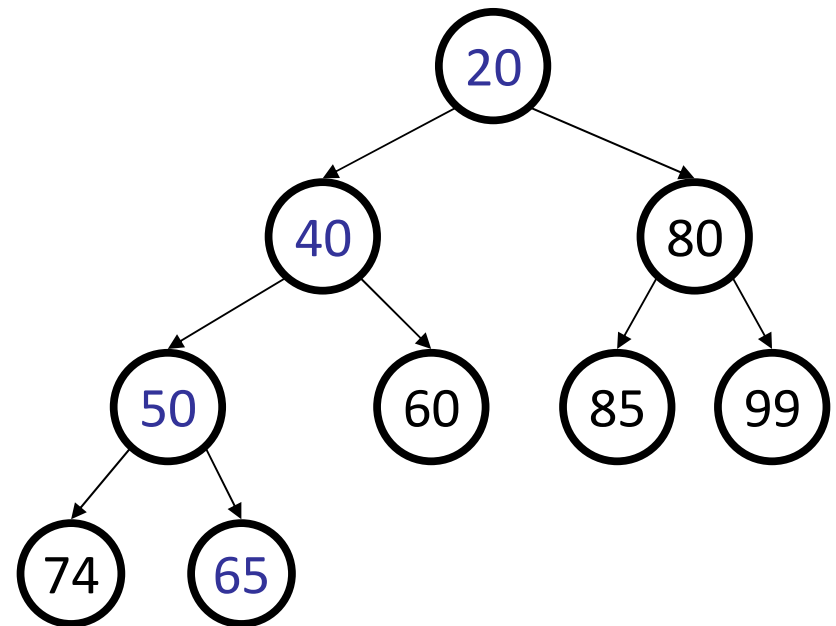
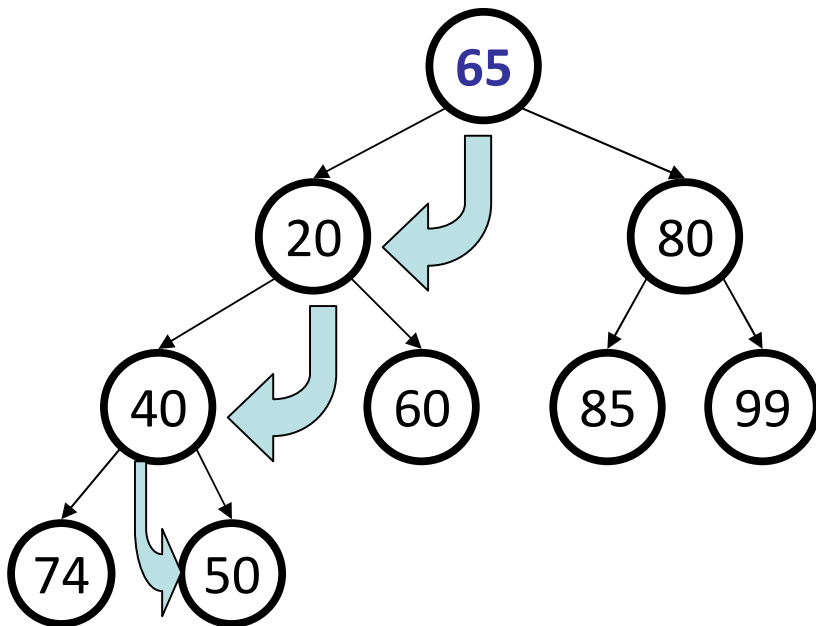
The remove operation

- When the root is removed from a heap, it should be initially replaced by the *rightmost leaf* (to maintain completeness).
 - But the heap ordering property becomes broken!



"Bubbling down" a node

- **bubble down:** To restore heap ordering, the new improper root is shifted ("bubbled") down the tree until it reaches its proper place.
 - Weiss: "*percolate down*" by swapping with its smaller child (why?)
 - How many bubble-down are necessary, at most?



Bubble-down exercise

- Suppose we have the min-heap shown below.
- Show the state of the heap tree after remove has been called 3 times, and which elements are returned by the removal.

