

---

# CSE 373

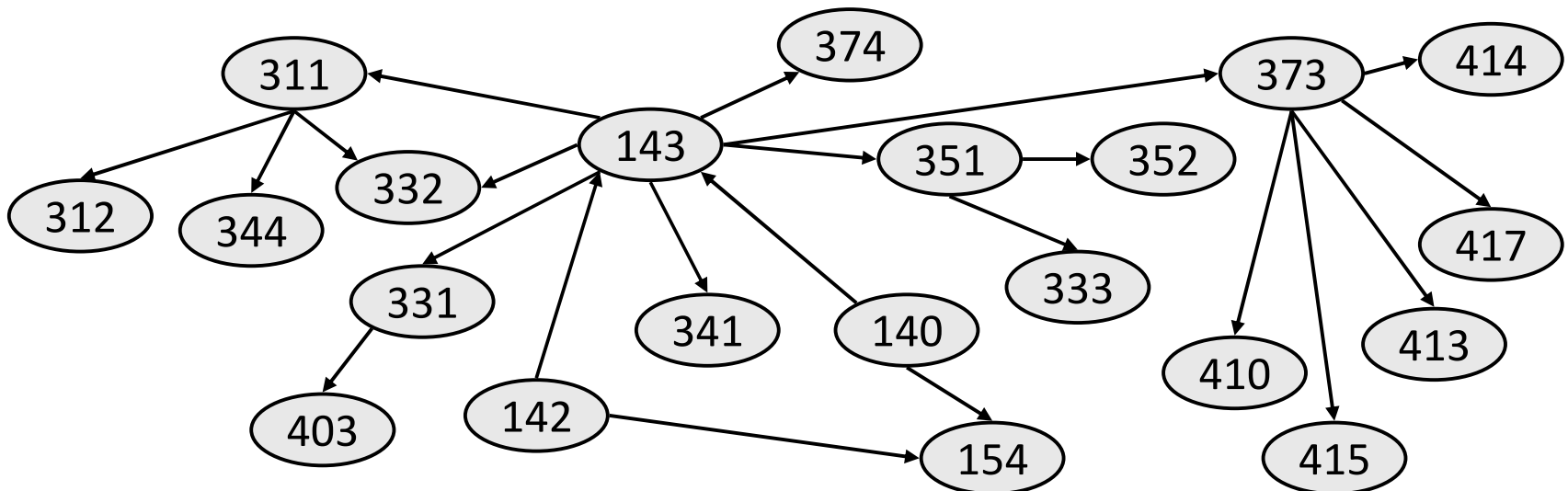
Graphs 4: Topological Sort  
reading: Weiss Ch. 9

slides created by Marty Stepp  
<http://www.cs.washington.edu/373/>

© University of Washington, all rights reserved.

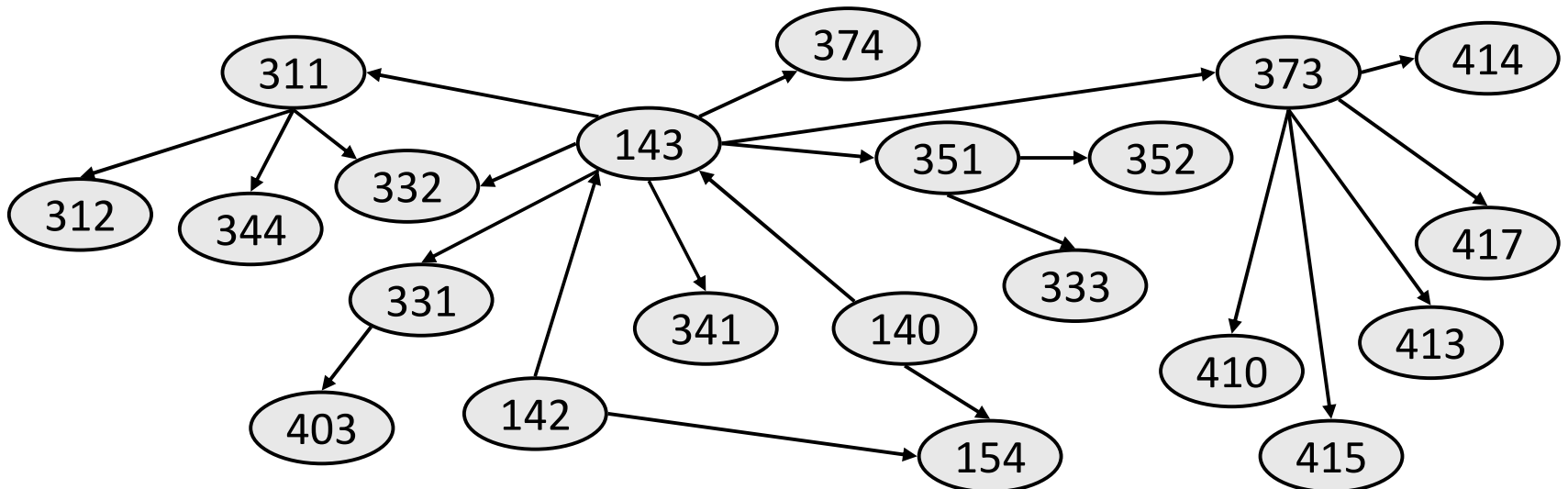
# Ordering a graph

- Suppose we have a directed acyclic graph (DAG) of courses, and we want to find an order in which the courses can be taken.
  - Must take all prereqs before you can take a given course. Example:
    - [142, 143, 140, 154, 341, 374, 331, 403, 311, 332, 344, 312, 351, 333, 352, 373, 414, 410, 417, 413, 415]
    - There might be more than one allowable ordering.
  - How can we find a valid ordering of the vertices?



# Topological Sort

- **topological sort:** Given a digraph  $G = (V, E)$ , a total ordering of  $G$ 's vertices such that for every edge  $(v, w)$  in  $E$ , vertex  $v$  precedes  $w$  in the ordering. Examples:
  - determining the order to recalculate updated cells in a spreadsheet
  - finding an order to recompile files that have dependencies
    - (any problem of finding an order to perform tasks with dependencies)

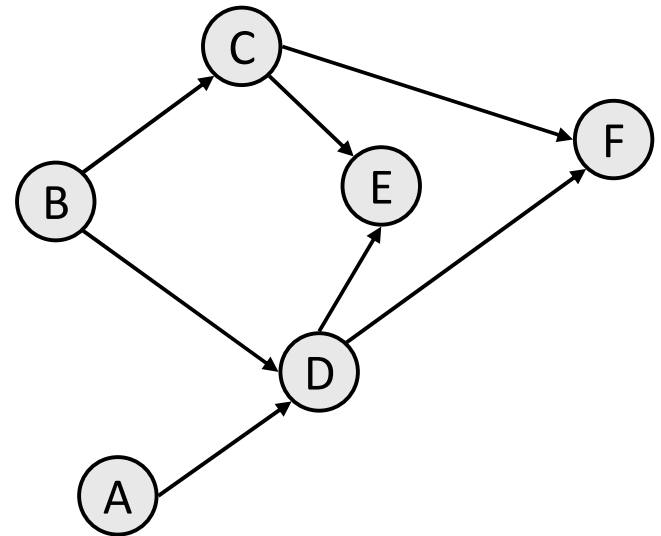


# Topo sort example

---

- How many valid topological sort orderings can you find for the vertices in the graph below?

- [A, B, C, D, E, F], [A, B, C, D, F, E],
- [A, B, D, C, E, F], [A, B, D, C, F, E],
- [B, A, C, D, E, F], [B, A, C, D, F, E],
- [B, A, D, C, E, F], [B, A, D, C, F, E],
- [B, C, A, D, E, F], [B, C, A, D, F, E],
- ...

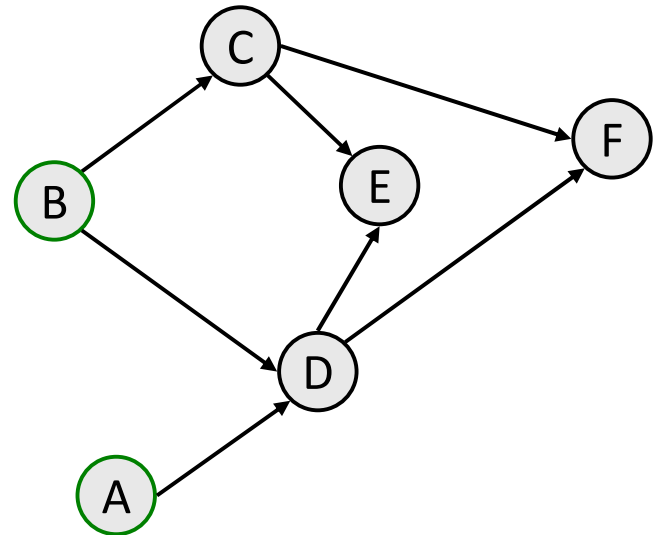


- What if there were a new vertex G unconnected to the others?

# Topo sort: Algorithm 1

---

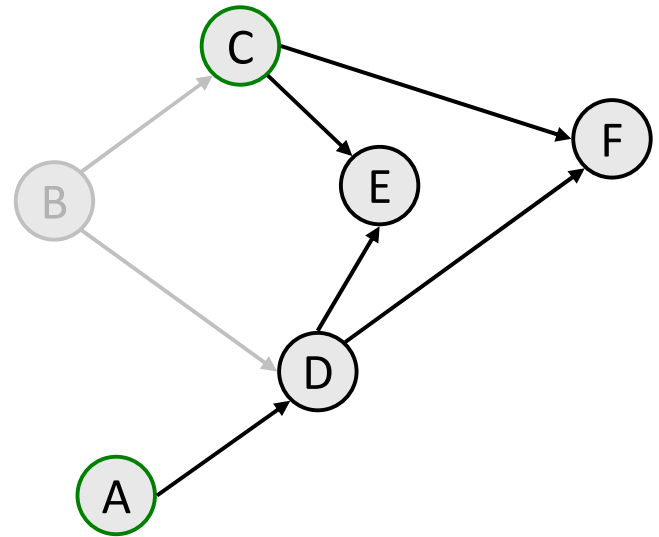
- function `topologicalSort()`:
  - *ordering* := {}.
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - *ordering* +=  $v$ .



# Topo sort example

---

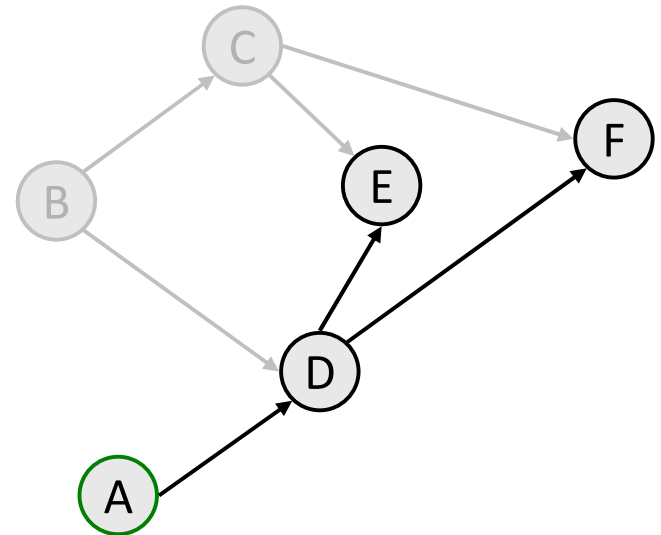
- function `topologicalSort()`:
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .
  - $ordering = \{ B \}$



# Topo sort example

---

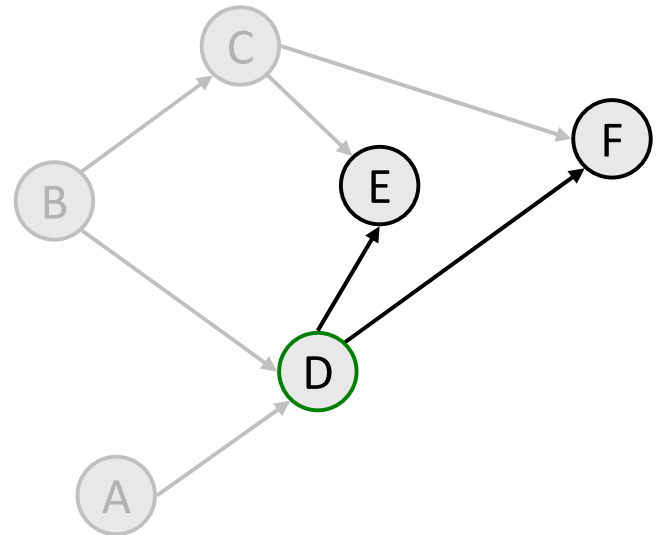
- function `topologicalSort()`:
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .
  - $ordering = \{ B, C \}$



# Topo sort example

---

- function `topologicalSort()`:
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .
  - $ordering = \{ B, C, A \}$

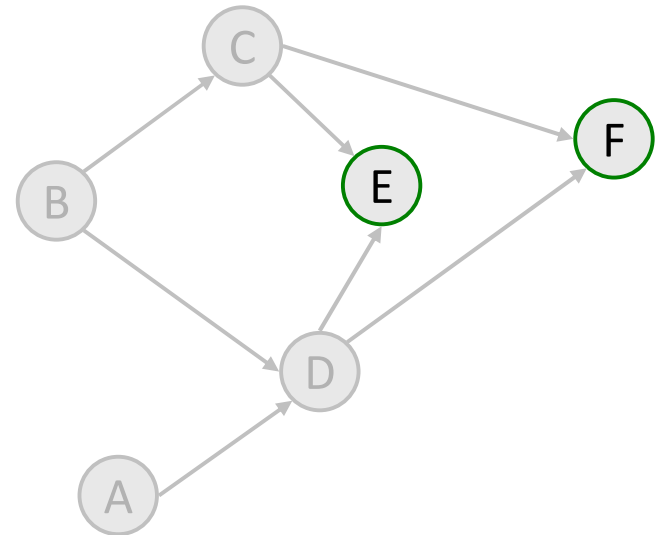




# Topo sort example

---

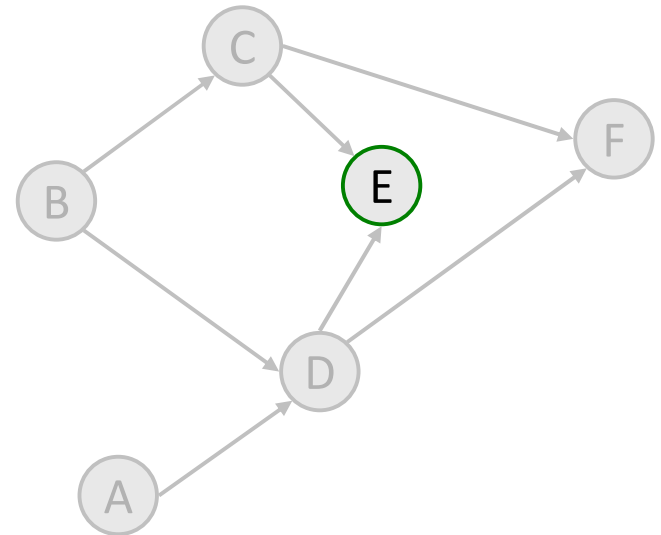
- function `topologicalSort()`:
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .
  - $ordering = \{ B, C, A, D \}$



# Topo sort example

---

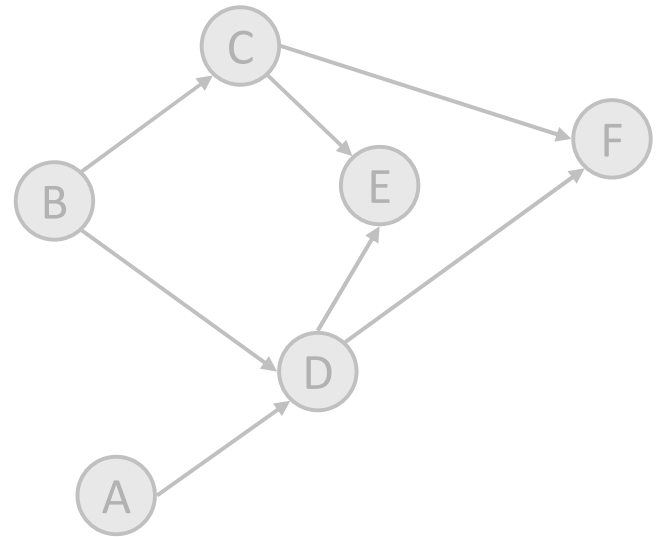
- function `topologicalSort()`:
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .
  - $ordering = \{ B, C, A, D, F \}$



# Topo sort example

---

- function `topologicalSort()`:
  - $ordering := \{ \}$ .
  - Repeat until graph is empty:
    - Find a vertex  $v$  with in-degree of 0 (no incoming edges).
      - (If there is no such vertex, the graph cannot be sorted; stop.)
    - Delete  $v$  and all of its outgoing edges from the graph.
    - $ordering += v$ .
  - $ordering = \{ B, C, A, D, F, E \}$



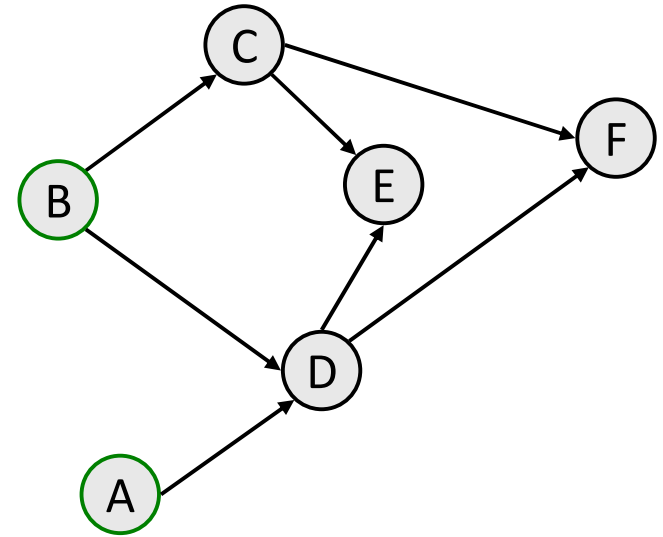
# Revised algorithm

---

- We don't want to literally delete vertices and edges from the graph while trying to topological sort it; so let's revise the algorithm:
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}$ .
  - $queue := \{\text{all vertices with in-degree} = 0\}$ .
  - $ordering := \{ \}$ .
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue.
    - $ordering += v$ .
    - Decrease the in-degree of all  $v$ 's neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}$ .
  - If all vertices are processed, success.  
Otherwise, there is a cycle.

# Topo sort example 2

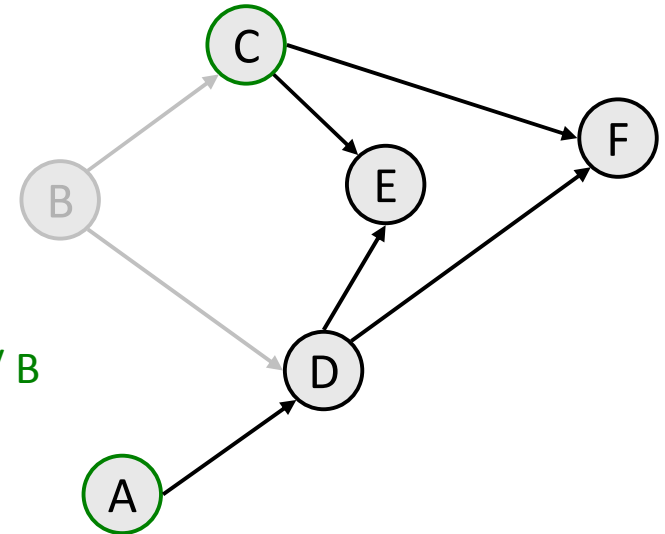
- function `topologicalSort()`:
  - *map* := {each vertex  $\rightarrow$  its in-degree}.
  - *queue* := {all vertices with in-degree = 0}.
  - *ordering* := { }.
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue.
    - *ordering* +=  $v$ .
    - Decrease the in-degree of all  $v$ 's neighbors by 1 in the *map*.
    - *queue* += {any neighbors whose in-degree is now 0}.
  
- *map* := { A=0, B=0, C=1, D=2, E=2, F=2 }
- *queue* := { B, A }
- *ordering* := { }



# Topo sort example 2

- function `topologicalSort()`:

- $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}$ .
- $queue := \{\text{all vertices with in-degree} = 0\}$ .
- $ordering := \{\}$ .
- Repeat until queue is empty:
  - Dequeue the first vertex  $v$  from the queue. // B
  - $ordering += v$ .
  - Decrease the in-degree of all  $v$ 's // C, D neighbors by 1 in the  $map$ .
  - $queue += \{\text{any neighbors whose in-degree is now } 0\}$ .

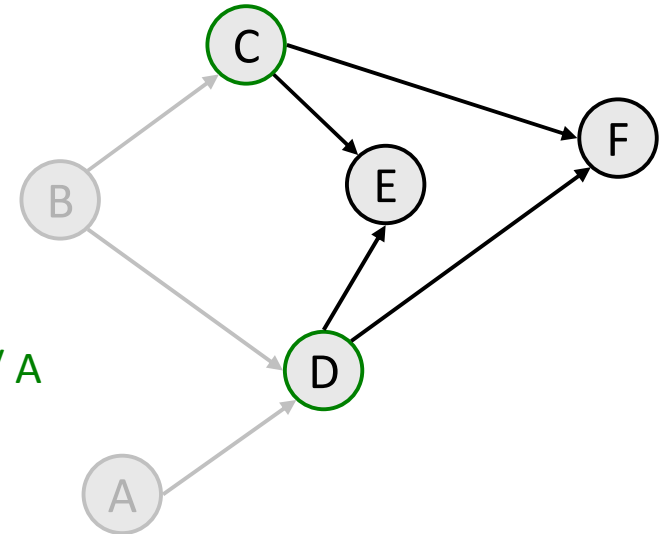


- $map := \{ A=0, B=0, C=0, D=1, E=2, F=2 \}$
- $queue := \{ A, C \}$
- $ordering := \{ B \}$

# Topo sort example 2

- function `topologicalSort()`:

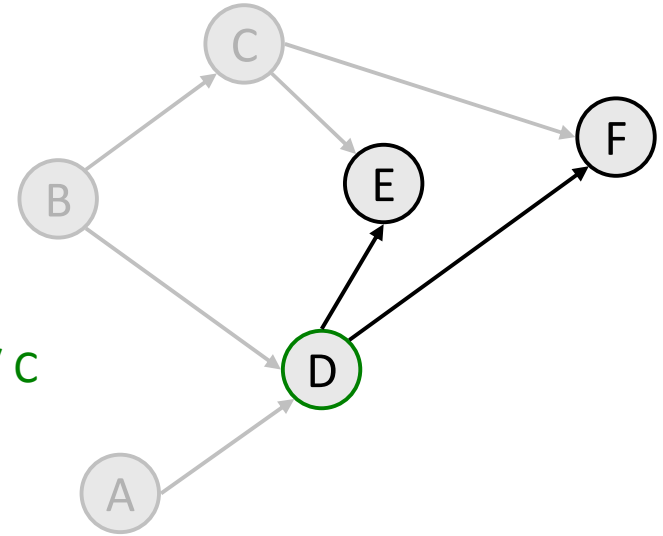
- $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}$ .
- $queue := \{\text{all vertices with in-degree} = 0\}$ .
- $ordering := \{\}$ .
- Repeat until queue is empty:
  - Dequeue the first vertex  $v$  from the queue. // A
  - $ordering += v$ .
  - Decrease the in-degree of all  $v$ 's // D neighbors by 1 in the  $map$ .
  - $queue += \{\text{any neighbors whose in-degree is now } 0\}$ .



- $map := \{ A=0, B=0, C=0, \mathbf{D=0}, E=2, F=2 \}$
- $queue := \{ C, \mathbf{D} \}$
- $ordering := \{ B, \mathbf{A} \}$

# Topo sort example 2

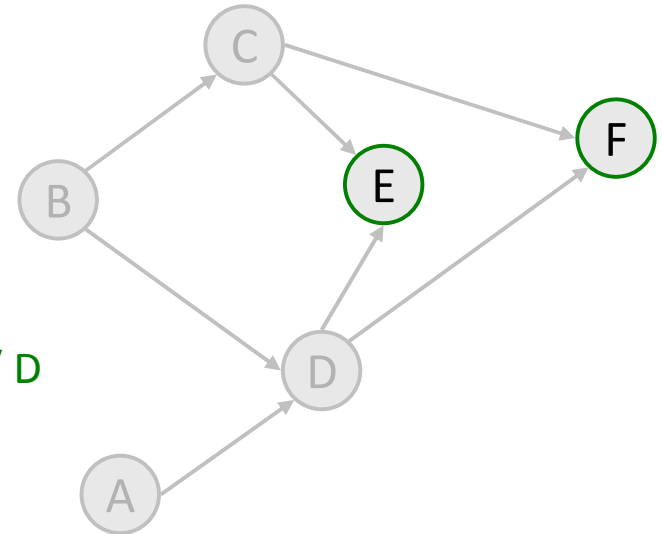
- function `topologicalSort()`:
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}$ .
  - $queue := \{\text{all vertices with in-degree} = 0\}$ .
  - $ordering := \{\}$ .
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. *// C*
    - $ordering += v$ .
    - Decrease the in-degree of all  $v$ 's neighbors *// E, F* by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}$ .
  
- $map := \{ A=0, B=0, C=0, D=0, E=1, F=1 \}$
- $queue := \{ D \}$
- $ordering := \{ B, A, C \}$





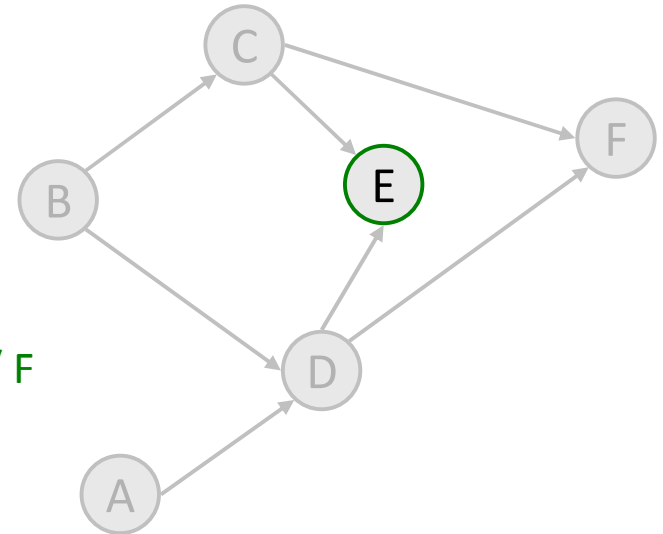
# Topo sort example 2

- function `topologicalSort()`:
  - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}$ .
  - $queue := \{\text{all vertices with in-degree} = 0\}$ .
  - $ordering := \{\}$ .
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. *// D*
    - $ordering += v$ .
    - Decrease the in-degree of all  $v$ 's *// F, E* neighbors by 1 in the  $map$ .
    - $queue += \{\text{any neighbors whose in-degree is now } 0\}$ .
  
- $map := \{ A=0, B=0, C=0, D=0, E=0, F=0 \}$
- $queue := \{ F, E \}$
- $ordering := \{ B, A, C, D \}$



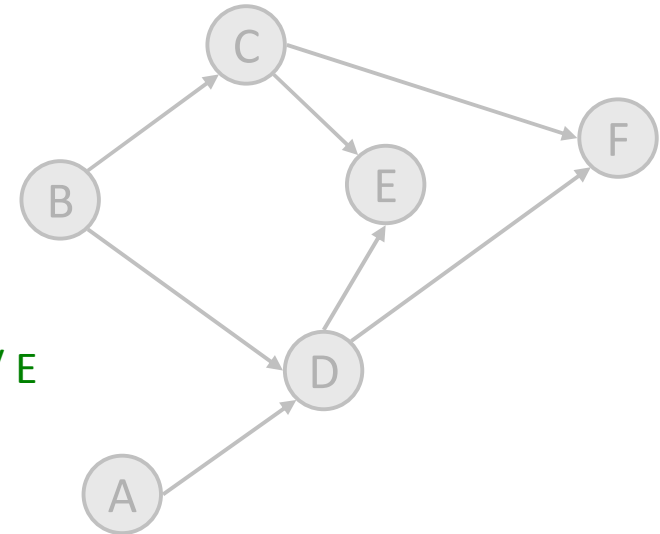
# Topo sort example 2

- function `topologicalSort()`:
  - `map` := {each vertex  $\rightarrow$  its in-degree}.
  - `queue` := {all vertices with in-degree = 0}.
  - `ordering` := { }.
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. // F
    - `ordering` +=  $v$ .
    - Decrease the in-degree of all  $v$ 's // none neighbors by 1 in the `map`.
    - `queue` += {any neighbors whose in-degree is now 0}.
  
- `map` := { A=0, B=0, C=0, D=0, E=0, F=0 }
- `queue` := { E }
- `ordering` := { B, A, C, D, F }



# Topo sort example 2

- function `topologicalSort()`:
  - `map` := {each vertex  $\rightarrow$  its in-degree}.
  - `queue` := {all vertices with in-degree = 0}.
  - `ordering` := { }.
  - Repeat until queue is empty:
    - Dequeue the first vertex  $v$  from the queue. // E
    - `ordering` +=  $v$ .
    - Decrease the in-degree of all  $v$ 's // none neighbors by 1 in the `map`.
    - `queue` += {any neighbors whose in-degree is now 0}.
  
- `map` := { A=0, B=0, C=0, D=0, E=0, F=0 }
- `queue` := { }
- `ordering` := { B, A, C, D, F, E }



# Topo sort runtime

---

- What is the runtime of our topological sort algorithm?
  - (with an "adjacency map" graph internal representation)
  - function `topologicalSort()`:
    - $map := \{\text{each vertex} \rightarrow \text{its in-degree}\}$ . //  $O(V)$
    - $queue := \{\text{all vertices with in-degree} = 0\}$ .
    - $ordering := \{\}$ .
    - Repeat until queue is empty: //  $O(V)$ 
      - Dequeue the first vertex  $v$  from the queue. //  $O(1)$
      - $ordering += v$ . //  $O(1)$
      - Decrease the in-degree of all  $v$ 's neighbors by 1 in the  $map$ . //  $O(E)$  for all passes
      - $queue += \{\text{any neighbors whose in-degree is now } 0\}$ .
  - Overall:  $O(V + E)$  ; essentially  $O(V)$  time on a sparse graph (fast!)