CSE 374 Programming Concepts & Tools

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Lecture 9 – C: Locals, Ivalues and rvalues, more pointers

The story so far...

- The low-level execution model of a process (one address space)
- Basics of C:
 - Language features: functions, pointers, arrays
 - Idioms: Array-lengths, strings with '\0' terminators
 - Control constructs and int guards
- Today, more features:
 - Local declarations
 - Storage duration and scope
 - Left vs. right expressions; more pointers
 - Dangling pointers
 - Stack arrays and implicit pointers (confusing)
- Later: structs; the heap and manual memory management (and some hacking)

Storage, lifetime, and scope

- At run-time, every variable needs space
 - When is the space allocated and deallocated?
- Every variable has scope
 - Where can the variable be used (unless another variable shadows it)?
- C has several answers (with inconsistent reuse of the word static)
- Some answers rarely used but understanding storage, lifetime, and scope is important
- Related: Allocating space is separate from initializing that space
 - Use uninitialized bits? Hopefully crash but who knows?
 - Unlike Java, which zeros out objects and complains about uninitialized locals

Storage, lifetime, and scope

- Global variables allocated before main, deallocated after main.
 Scope is entire program
 - Usually bad style, kind of like public static Java fields
 - But can be OK for truly global data like conversion tables, physical constants, etc.
- Static global variables like global variables but scope is just that source file, kind of like private static Java fields
 - Related: static functions cannot be called from other files
- Static local variables like global variables (!) but scope is just that function, rarely used (We won't use them)
- Local variables (often called automatic) allocated "when reached" deallocated "after that block", scope is that block
 - With recursion, multiple copies of same variable (one per stack frame/function activation)
 - Like local variables in Java

Ivalues vs rvalues

- In intro courses we are usually fairly sloppy about the difference between the left side of an assignment and the right. To "really get" C, it helps to get this straight:
 - Law #1: Left-expressions get evaluated to locations (addresses)
 - Law #2: Right-expressions get evaluated to values
 - Law #3: Values include numbers and pointers (addresses)
- The key difference is the "rule" for variables:
 - As a left-expression, a variable is a location and we are done
 - As a right-expression, a variable gets evaluated to its location's contents, and then we are done
 - Most things do not make sense as left expressions
- Note: This is true in Java too

Function arguments

- Storage and scope of arguments is like for local variables
- But initialized by the caller ("copying" the value)
- So assigning to an argument has no affect on the caller
- But assigning to the space pointed-to by an argument might

```
void f() {
    int i=17;
    int j=g(i);
    printf("%d %d",i,j);
}
int g(int x) {
    x = x+1;
    return x+1;
}
```

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```
void f() { int g(int* p) { int k = *p; int j=g(&i); int *q = &k; printf("%d %d",i,j); p = q; (*p) = (*q) + 1; return (*q) + 1;
```

Pointers to pointers to ...

- Any level of pointer makes sense:
 - Example: argv, *argv, **argv
 - Same example: argv, argv[0], argv[0][0]
- But &(&p) makes no sense (&p is not a left-expression, the value is an address but the value is in no-particularplace)
- This makes sense (well, at least it's legal C):

```
void f(int x) {
  int*p = &x;
  int**q = &p;
  ... can use x, p, *p, q, *q, **q, ...
}
```

 Note: When playing, you can print pointers (i.e., addresses) with %p (just numbers in hexadecimal)

Dangling pointers

```
int* f(int x) {
  int *p;
  if(x) {
    int y = 3;
    p = &y; /* ok */
  } /* ok, but p now dangling */
  /* y = 4 does not compile */
  *p = 7; /* could CRASH but probably not */
  return p; /* uh-oh, but no crash yet */
void g(int *p) { *p = 123; }
void h() {
  g(f(7)); /* HOPEFULLY YOU CRASH (but maybe not) */
```

Arrays and Pointers

- If p has type T* or type T[]:
 - *p has type T
 - If i is an int, p+i refers to the location of an item of type
 T that is i items past p (not +i storage locations unless each item of type T takes up exactly 1 unit of storage)
 - p[i] is defined to mean *(p+i)
 - if p is used in an expression (including as a function argument) it has type T*
 - Even if it is declared as having type T[]
 - One consequence: array arguments are always "passed by reference" (as a pointer), not "by value" (which would mean copying the entire array value)

Arrays revisited

 "Implicit array promotion": a variable of type T[] becomes a variable of type T* in an expression

```
void f1(int^* p) \{ *p = 5; \}
int* f2() {
  int x[3]; /* x on stack */
  x[0] = 5;
/* (&x)[0] = 5; wrong */
  *x = 5:
  *(x+0) = 5;
  f1(x);
/* f1(&x); wrong – watch types! */
/* x = &x[2]; wrong - x isn't really a pointer! */
  int *p = &x[2];
  return x; /* wrong – dangling pointer – but type correct */
```