Program and memory layout

- By convention in MIPS the layout is:
 - Note that only half of the addressing space is taken by user
 Other half is O.S.



Procedures

- Procedures/functions are the major program structuring mechanism
- Calling and returning form a procedure requires a protocol between *caller* and *callee*
- Protocol is based on conventions

Procedures/Functions -- Protocol

- Each machine (compiler?) has its own set of protocol(s)
- Protocol: combination of hardware/software
 - e.g., "jal" is hardware
 - use of register \$29 as \$sp is software
- Protocol: sequence of steps to be followed at each call and each return
 - controlled by hardware and/or software
- In RISC machines
 - hardware performs simple instructions
 - software (compiler/assembler) controls sequence of instructions

Program stack

- Each executing program (process) has a *stack*
- Stack = dynamic data structure accessed in a LIFO manner
- Program stack automatically allocated by O.S.
- At the start of the program, register \$sp (\$29 in MIPS) is automatically loaded to point to the first empty slot on top of stack
 - After that it will be your responsibility to manage \$sp
- By convention, stack grows towards lower addresses
 - to allocate new space (i.e., when you *push*), decrement \$sp
 - to free space on top of stack (*pop*), increment \$sp

Push operation

- *push* adds an item on top of stack
 - one instruction to manipulate the data, e.g. "sw \$6,0(\$sp)"
 - one instruction to adjust the stack pointer e.g., "addi \$sp,\$sp,-4" (the assembler will accept "subu \$sp,4")



Pop operation

- *pop* removes the item on top of stack and stores it in a register
 - one instruction to adjust the stack pointer e.g., "addiu \$sp,\$sp,4"





Procedure call requirements (caller/callee)

- Caller must pass the return address to the callee
- Caller must pass the parameters to the callee
- Caller must save what is *volatile* (registers) and could be used by the callee
- Callee must save the return address (in case it becomes a
- caller)
- Callee must provide (stack) storage for its own use
- Caller/callee should support recursive calls

CSE378 Procedures.

Mechanism

- Registers are used for
 - passing return address in \$ra
 - jal target
 - passing a small number of parameters (up to 4 in \$a0 to \$a3)
 - keeping track of the stack
 (\$sp)
 - returning function values (in \$v0 and \$v1)

- Stack is used for
 - saving temporary registers to be used by caller/callee
 - saving info about the caller (return address)
 - passing parameters if needed
 - allocating local data for the called procedure

Procedure calls and register conventions

Register	Name	Function	Comment
\$0	Zero	Always 0	No-op on write
\$1	\$at	Reserved for assembler	Don't use it
\$2-3	\$v0-v1	Expr. Eval/funct. Return	
\$4-7	\$a0-a3	Proc./func. Call parameters	
\$8-15	\$t0-t7	Temporaries; volatile	Not saved on proc. Calls
\$16-23	\$s0-s7	Temporaries	Should be saved on calls
\$24-25	\$t8-t9	Temporaries; volatile	Not saved on proc. Calls
\$26-27	\$k0-k1	Reserved for O.S.	Don't use them
\$28	\$gp	Pointer to global static memory	
\$29	\$sp	Stack pointer	
\$30	\$fp	Frame pointer	
\$31	\$ra	Proc./funct return address	

Who does what on a call (one sample protocol)

- Caller
 - Saves any volatile register
 (\$t0-\$t9) that has contents that need to be kept
 - Puts up to 4 arguments in \$a0-\$a3
 - If more than 4 arguments, pushes the rest on the stack
 - calls with jal instruction

- Callee
 - saves \$ra on stack
 - saves any non-volatile register
 (\$s0-s7) that it will use

Who does what on return

- Callee
 - restores any non-volatile
 register (\$s0-\$s7) it has used
 - restores \$ra
 - puts function results in \$v0-\$v1
 - adjusts \$sp
 - returns to caller with "jr \$ra"

- Caller
 - restores any volatile register it had saved
 - examines \$v0-\$v1 if needed

Example of a call sequence

• Assume 2 arguments in \$t0 and \$t3 and we want to save the contents of \$t6 and \$t7

move	\$a0,\$t0	#1st argument in \$a0
move	\$a1,\$t3	#2nd argument in \$a1
subu	\$sp,\$sp,8	#room for 2 temps on stack
SW	\$t6,8(\$sp)	#save \$t6 on stack
SW	\$t7,4(\$sp)	#save \$t7 on stack
jal	target	

• Assume the callee does not need to save registers

target:	sw \$ra,0(\$sp)	#save return address
	subu \$sp,\$sp,4	# on stack

Return from the previous sequence

- The callee will have put the function results in \$v0-\$v1 addu \$sp,\$sp,4 #pop
 lw \$ra,0(\$sp) #return address in \$ra jr \$ra #to caller
- The caller will restore \$t6 and \$t7 and adjust stack

lw\$t6,8(\$sp)lw\$t7,4(\$sp)addu\$sp,\$sp,8