## Decisions

## CSE 410, Spring 2004 Computer Systems

http://www.cs.washington.edu/education/courses/410/04sp/

## Reading and References

- Sections 3.5, A.9, A. 10 through page A-54, Computer Organization and Design, Patterson and Hennessy


## goto considered harmful

- "Oh what a tangled web we weave, When first we practice to deceive!"
» Sir Walter Scott
- Branching in assembly language can turn your program into a rat's nest that cannot be debugged
- Keep control flow simple and logical
- Use comments describing the overall logic


## Conditional Branch



## Branch instructions

- Branch instructions are I-format instructions
» op code field
» two register fields
" 16-bit offset field
- Simplest branches check for equality
» beq \$t0, \$t1, address
» bne \$t0, \$t1, address


## Go to where?

- Calculating the destination address
" 4*(the 16-bit offset value)
» is added to the Program Counter (PC)
- The offset is a word offset in this case
- The base register is always the PC, so we don't need to specify it in the instruction
- Covers a range of $2^{16}$ words ( 64 KW )


## if (i==j) then $a=b ;$

- Assume all values are in registers
- Note that the test is inverted!
$\# \$ t 0=i, \$ t 1=j, \$ s 0=a, \$ s 1=b$
bne \$t0, \$t1, skip
move \$s0, \$s1
skip:


## while (s[i]==k) $i=i+j ;$

\# \$s0=addr(s), \$v1=i, \$a0=k, \$a1=j
loop:

| sll | \$v0,\$v1,2 \# v0 = 4*i |
| :---: | :---: |
| addu | \$v0,\$s0,\$v0 \# v0 = addr (s[i]) |
| 1w | \$v0,0 (\$v0) \# v0 = s[i] |
| addu | \$v1,\$v1,\$a1 \# i = i+j |
| beq | \$v0,\$a0,loop \# loop if equal |
| subu | \$v1,\$v1,\$a1 \# i = i-j |

for (i=0; i<10; i++) s[i] = i;

| \# \$s0=addr (s), \$t1=i |  |  |
| :--- | :--- | :--- |
| move | $\$ t 1, \$ z e r o$ | $\#$ i $=0$ |
| loop: |  |  |
| sll | $\$ t 0, \$ t 1,2$ | $\#$ t0 $=i * 4$ |
| addu | $\$ t 0, \$ s 0, \$ t 0$ | $\#$ to $=$ addr (s[i]) |
| sw | $\$ t 1,0(\$ t 0)$ | $\#$ s [i] $=i$ |
| addu | $\$ t 1, \$ t 1,1$ | $\#$ i++ |
| slt | $\$ t 0, \$ t 1,10$ | $\#$ if (i<10) $\$ t 0=1$ |
| bnez | $\$ t 0, l o o p$ | $\#$ loop if (i<10) |

\# \$s0=addr(s), \$t1=i
move
loop:
sll
addu $\$ t 0, \$ s 0, \$ t 0$
sw $\$ \mathrm{t} 1,0(\$ \mathrm{t} 0)$
addu $\$ \mathrm{t} 1, \$ \mathrm{t} 1,1$
slt $\$ t 0, \$ t 1,10$
bnez
\# i $=0$
\# to $=i * 4$
\# to = addr(s[i])
\# s[i] = i
\# i++
\# if (i<10) \$t0=1
\# loop if (i<10)

## Comparison instructions

- For comparisons other than equality

» slt : set less than<br>» sltu : set less than unsigned<br>» slti : set less than constant value<br>» sltiu : set less than unsigned constant

- set t0 to 1 if $\mathrm{t} 1<\mathrm{t} 2$
slt \$t0, \$t1, \$t2


## Pseudo-instructions

- The assembler is your friend and will build instruction sequences for you
- Original code:
bge $\$ a 0, \$ t 1$,end $\#$ if $a 0>=t 1$ skip
- Actual instructions:
slt $\$ a t, \$ a 0, \$ t 1$
beq $\$ a t, \$ 0$,end
\# if a0<t1 at=true
\# skip if at==false


## Jump Instructions

- Jump instructions provide longer range than branch instructions
- 26-bit word offset in J-format instructions
» j : jump
» jal : jump and link (store return address)
- 32-bit address in register jumps
» jr : jump through register
» jalr : jump through register and link


## J-format fields

| op code | word offset |
| :---: | :---: |
| 6 bits | 26 bits |

- The word offset value is multiplied by 4 to create a byte offset
" the result is 28 bits wide
- Then concatenated with top 4 bits of PC to make a 32 bit destination address


## Important Jumps

- Jump and link (jal)
» call procedure and store return address in $\$ \mathrm{ra}$
- Jump through register ( jr )
» return to caller using the address in \$ra
- We will talk about procedure calls in excruciating detail next lecture

