

# Database Systems CSE 414

## Lectures 8: Relational Algebra (Ch. 2.4, & 5.1)

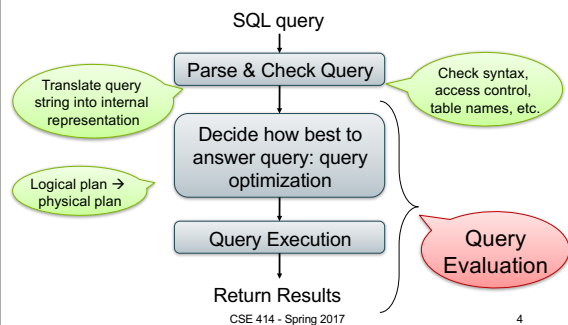
## Announcements

- WQ3 is due Sunday 11pm
- Azure codes will be sent out Wed/Thu
- Don't miss section tomorrow
  - will go through Azure setup and basic use
- HW3 will be posted by Thu night
  - due on Tuesday, 4/25 (in 13 days)

## Where We Are

- Motivation for using a DBMS for managing data
- SQL:
  - Declaring the schema for our data (CREATE TABLE)
  - Inserting data one row at a time or in bulk (INSERT/import)
  - Modifying the schema and updating the data (ALTER/UPDATE)
  - Querying the data (SELECT)
- **Next step: More knowledge of how DBMSs work**
  - Client-server architecture
  - Relational algebra and query execution

## Query Evaluation Steps



## The WHAT and the HOW

- SQL = **WHAT** we want to get from the data
- Relational Algebra = **HOW** to get the data we want
- Move from **WHAT** to **HOW** is **query optimization**
  - SQL ~> Relational Algebra ~> Physical Plan
  - Relational Algebra = Logical Plan

## Relational Algebra

## Sets v.s. Bags

- Sets: {a,b,c}, {a,d,e,f}, { }, . . .
- Bags: {a, a, b, c}, {b, b, b, b}, . . .

Relational Algebra has two semantics:

- Set semantics = standard Relational Algebra
- Bag semantics = extended Relational Algebra

DB systems implement bag semantics (Why?)

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## Relational Algebra Operators

- Union  $\cup$ , intersection  $\cap$ , difference  $-$
- Selection  $\sigma$
- Projection  $\pi$  ( $\Pi$ )
- Cartesian product  $\times$ , join  $\bowtie$
- Rename  $\rho$
- Duplicate elimination  $\delta$
- Grouping and aggregation  $\gamma$
- Sorting  $\tau$

RA

Extended RA

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## Union and Difference

$R1 \cup R2$   
 $R1 - R2$

What do they mean over bags ?

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## What about Intersection ?

- Derived operator using minus

$$R1 \cap R2 = R1 - (R1 - R2)$$

- Derived using join (will explain later)

$$R1 \cap R2 = R1 \bowtie R2$$

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## Selection

- Returns all tuples which satisfy a condition

$\sigma_c(R)$

- Examples

- $\sigma_{\text{Salary} > 40000}$ (Employee)
- $\sigma_{\text{name} = \text{"Smith"}}$ (Employee)

- The condition c can be =, <, <=, >, >=, <> combined with AND, OR, NOT

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Employee

SSN	Name	Salary
1234545	John	20000
5423341	Smith	60000
4352342	Fred	50000

$\sigma_{\text{Salary} > 40000}$ (Employee)

SSN	Name	Salary
5423341	Smith	60000
4352342	Fred	50000

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## Projection

- Eliminates columns

$$\pi_{A_1, \dots, A_n}(R)$$

- Example: project social-security number and names:
  - $\Pi_{SSN, Name}(Employee)$
  - Answer(SSN, Name)

Different semantics over sets or bags! Why?

Employee

SSN	Name	Salary
1234545	John	20000
5423341	John	60000
4352342	John	20000

$\pi_{Name, Salary}(Employee)$

Name	Salary
John	20000
John	60000
John	20000

Name	Salary
John	20000
John	60000

Bag semantics

Set semantics

Which is more efficient?

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## Composing RA Operators

Patient

no	name	zip	disease
1	p1	98125	flu
2	p2	98125	heart
3	p3	98120	lung
4	p4	98120	heart

$\pi_{zip, disease}(Patient)$

zip	disease
98125	flu
98125	heart
98120	lung
98120	heart

$\sigma_{disease='heart'}(Patient)$

no	name	zip	disease
2	p2	98125	heart
4	p4	98120	heart

$\pi_{zip, disease}(\sigma_{disease='heart'}(Patient))$

zip	disease
98125	heart
98120	heart

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## Cartesian Product

- Each tuple in R1 with each tuple in R2

$$R1 \times R2$$

- Rare in practice; mainly used to express joins

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## Cross-Product Example

Employee

Name	SSN
John	999999999
Tony	777777777

Dependent

EmpSSN	DepName
999999999	Emily
777777777	Joe

Employee X Dependent

Name	SSN	EmpSSN	DepName
John	999999999	999999999	Emily
John	999999999	777777777	Joe
Tony	777777777	999999999	Emily
Tony	777777777	777777777	Joe

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## Renaming

- Changes the schema, not the instance

$$\rho_{B_1, \dots, B_n}(R)$$

- Example:

$$\rho_{N, S}(Employee) \rightarrow Answer(N, S)$$

Not really used by systems, but needed on paper

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## Natural Join

$$R1 \bowtie R2$$

- Meaning:  $R1 \bowtie R2 = \pi_A(\sigma_\theta(R1 \times R2))$
- Where:
  - Selection  $\sigma$  checks equality of **all common attributes** (attributes with same names)
  - Projection  $\pi$  eliminates duplicate **common attributes**

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## Natural Join Example

R	A	B	S	B	C
	X	Y		Z	U
	X	Z		V	W
	Y	Z		Z	V
	Z	V			

$$R \bowtie S = \pi_{ABC}(\sigma_{R.B=S.B}(R \times S))$$

A	B	C
X	Z	U
X	Z	V
Y	Z	U
Y	Z	V
Z	V	W

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## Natural Join Example 2

AnonPatient P

age	zip	disease
54	98125	heart
20	98120	flu

Voters V

name	age	zip
p1	54	98125
p2	20	98120

P  $\bowtie$  V

age	zip	disease	name
54	98125	heart	p1
20	98120	flu	p2

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## Natural Join

- Given schemas  $R(A, B, C, D)$ ,  $S(A, C, E)$ , what is the schema of  $R \bowtie S$  ?
- Given  $R(A, B, C)$ ,  $S(D, E)$ , what is  $R \bowtie S$  ?
- Given  $R(A, B)$ ,  $S(A, B)$ , what is  $R \bowtie S$  ?

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AnonPatient (age, zip, disease)  
Voters (name, age, zip)

## Theta Join

- A join that involves a predicate

$$R1 \bowtie_{\theta} R2 = \sigma_{\theta}(R1 \times R2)$$

- Here  $\theta$  can be any condition
- For our voters/patients example:

$$P \bowtie_{P.zip = V.zip \text{ and } P.age \geq V.age - 1 \text{ and } P.age \leq V.age + 1} V$$

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## Equijoin

- A theta join where  $\theta$  is an equality predicate
- By far the most used variant of join in practice

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## Equijoin Example

AnonPatient P

age	zip	disease
54	98125	heart
20	98120	flu

Voters V

name	age	zip
p1	54	98125
p2	20	98120

$$P \bowtie_{P.age=V.age} V$$

P.age	P.zip	P.disease	P.name	V.zip	V.age
54	98125	heart	p1	98125	54
20	98120	flu	p2	98120	20

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## Join Summary

- **Theta-join:**  $R \bowtie_{\theta} S = \sigma_{\theta}(R \times S)$ 
  - Join of R and S with a join condition  $\theta$
  - Cross-product followed by selection  $\theta$
- **Equijoin:**  $R \bowtie_{\theta} S = \pi_{A}(\sigma_{\theta}(R \times S))$ 
  - Join condition  $\theta$  consists only of equalities
- **Natural join:**  $R \bowtie S = \pi_{A}(\sigma_{\theta}(R \times S))$ 
  - Equijoin
  - Equality on **all** fields with same name in R and in S
  - Projection  $\pi_{A}$  drops all redundant attributes

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## So Which Join Is It ?

When we write  $R \bowtie S$  we usually mean an equijoin, but we often omit the equality predicate when it is clear from the context

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## More Joins

- **Outer join**
  - Include tuples with no matches in the output
  - Use NULL values for missing attributes
  - Does not eliminate duplicate columns
- Variants
  - Left outer join
  - Right outer join
  - Full outer join

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## Outer Join Example

AnonPatient P

age	zip	disease
54	98125	heart
20	98120	flu
33	98120	lung

AnnonJob J

job	age	zip
lawyer	54	98125
cashier	20	98120

$$P \bowtie J$$

P.age	P.zip	disease	job	J.age	J.zip
54	98125	heart	lawyer	54	98125
20	98120	flu	cashier	20	98120
33	98120	lung	null	33	98120

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## More Examples

Supplier(sno, sname, scity, sstate)

Part(pno, pname, psize, pcolor)

Supply(sno, pno, qty, price)

Name of supplier of parts with size greater than 10

$\pi_{sname}(\text{Supplier} \bowtie \text{Supply} \bowtie (\sigma_{psize > 10}(\text{Part})))$

Name of supplier of red parts or parts with size greater than 10

$\pi_{sname}(\text{Supplier} \bowtie \text{Supply} \bowtie (\sigma_{psize > 10}(\text{Part}) \cup \sigma_{pcolor = \text{red}}(\text{Part})))$

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