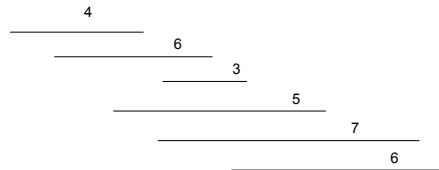


# CSE 421 Algorithms

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Lecture 16a  
Dynamic Programming

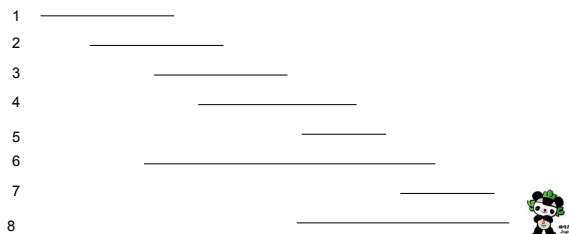
## Dynamic Programming

- Weighted Interval Scheduling
- Given a collection of intervals  $I_1, \dots, I_n$  with weights  $w_1, \dots, w_n$ , choose a maximum weight set of non-overlapping intervals



## Recursive Algorithm

Intervals sorted by finish time  
 $p[i]$  is the index of the last interval which finishes before  $i$  starts



## Optimality Condition

- $Opt[j]$  is the maximum weight independent set of intervals  $I_1, I_2, \dots, I_j$

## Algorithm

```
MaxValue(j) =  
  if j = 0 return 0  
  else  
    return max( MaxValue(j-1),  
               wj + MaxValue(p[ j ]))
```

## Run time

- What is the worst case run time of MaxValue
- Design a worst case input

## A better algorithm

$M[j]$  initialized to -1 before the first recursive call for all  $j$

```

MaxValue(j) =
  if j = 0 return 0;
  else if M[j] != -1 return M[j];
  else
    M[j] = max(MaxValue(j-1), wj + MaxValue(p[j]));
    return M[j];
  
```

## Iterative Algorithm

Express the MaxValue algorithm as an iterative algorithm

```

MaxValue {
  
```

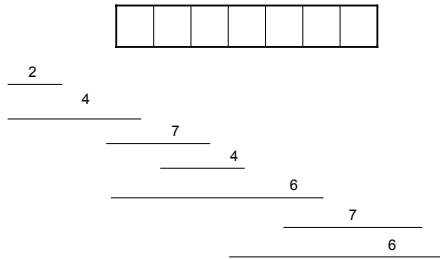
```

}
  
```



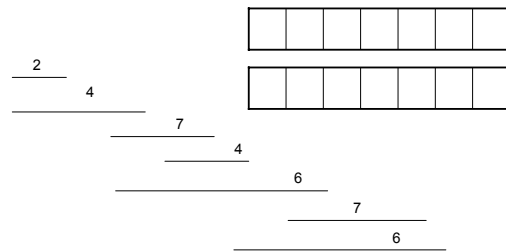
## Fill in the array with the Opt values

$Opt[j] = \max(Opt[j - 1], w_j + Opt[p[j]])$



## Computing the solution

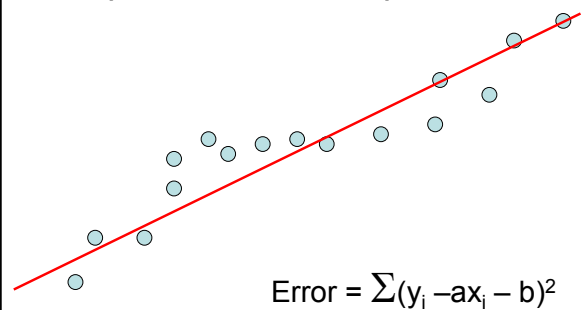
$Opt[j] = \max(Opt[j - 1], w_j + Opt[p[j]])$   
Record which case is used in Opt computation



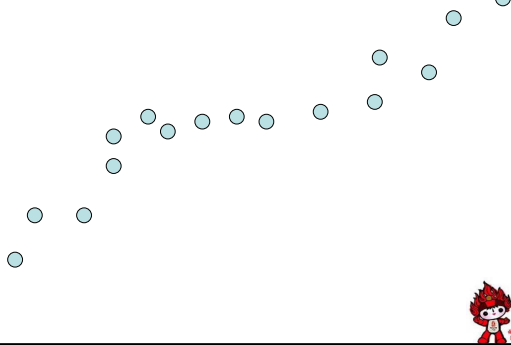
## Dynamic Programming

- The most important algorithmic technique covered in CSE 421
- Key ideas
  - Express solution in terms of a polynomial number of sub problems
  - Order sub problems to avoid recomputation

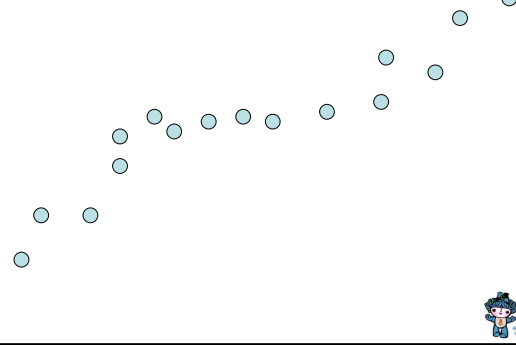
## Optimal linear interpolation



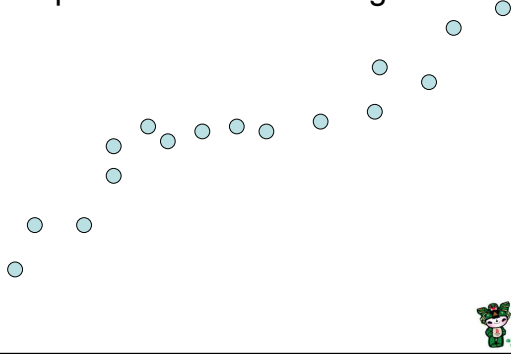
What is the optimal linear interpolation with three line segments



What is the optimal linear interpolation with two line segments

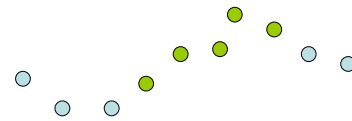


What is the optimal linear interpolation with n line segments



### Notation

- Points  $p_1, p_2, \dots, p_n$  ordered by x-coordinate ( $p_i = (x_i, y_i)$ )
- $E_{i,j}$  is the least squares error for the optimal line interpolating  $p_i, \dots, p_j$



### Optimal interpolation with two segments

- Give an equation for the optimal interpolation of  $p_1, \dots, p_n$  with two line segments
- $E_{i,j}$  is the least squares error for the optimal line interpolating  $p_i, \dots, p_j$



### Optimal interpolation with k segments

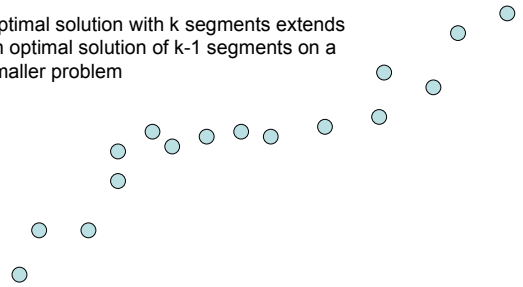
- Optimal segmentation with three segments
  - $\text{Min}_{i,j} \{E_{1,i} + E_{i,j} + E_{j,n}\}$
  - $O(n^2)$  combinations considered
- Generalization to k segments leads to considering  $O(n^{k-1})$  combinations

$\text{Opt}_k[j]$  : Minimum error approximating  $p_1 \dots p_j$  with  $k$  segments

How do you express  $\text{Opt}_k[j]$  in terms of  $\text{Opt}_{k-1}[1], \dots, \text{Opt}_{k-1}[j]$ ?

### Optimal sub-solution property

Optimal solution with  $k$  segments extends an optimal solution of  $k-1$  segments on a smaller problem



### Optimal multi-segment interpolation

Compute  $\text{Opt}[k, j]$  for  $0 < k < j < n$

for  $j := 1$  to  $n$

$\text{Opt}[1, j] = E_{1,j}$

for  $k := 2$  to  $n-1$

    for  $j := 2$  to  $n$

$t := E_{1,j}$

        for  $i := 1$  to  $j-1$

$t = \min(t, \text{Opt}[k-1, i] + E_{i,j})$

$\text{Opt}[k, j] = t$