

Extra Slides

Coin Changing

Greed is good. Greed is right. Greed works.
Greed clarifies, cuts through, and captures
the essence of the evolutionary spirit.
- Gordon Gecko (Michael Douglas)



Coin Changing

Goal. Given currency denominations: 1, 5, 10, 25, 100, devise a method to pay amount to customer using fewest number of coins.

Ex: 34¢.



Cashier's algorithm. At each iteration, add coin of the largest value that does not take us past the amount to be paid.

Ex: \$2.89.



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Coin-Changing: Greedy Algorithm

Cashier's algorithm. At each iteration, add coin of the largest value that does not take us past the amount to be paid.

```
Sort coins denominations by value:  $c_1 < c_2 < \dots < c_n$ .  
coins selected  
 $S \leftarrow \emptyset$   
while ( $x \neq 0$ ) {  
  let  $k$  be largest integer such that  $c_k \leq x$   
  if ( $k = 0$ )  
    return "no solution found"  
   $x \leftarrow x - c_k$   
   $S \leftarrow S \cup \{k\}$   
}  
return  $s$ 
```

Q. Is cashier's algorithm optimal?

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Coin-Changing: Analysis of Greedy Algorithm

Theorem. Greed is optimal for U.S. coinage: 1, 5, 10, 25, 100.

Pf. (by induction on x)

- Consider optimal way to change $c_i \leq x < c_{i+1}$: greedy takes coin k .
- We claim that any optimal solution must also take coin k .
 - if not, it needs enough coins of type c_1, \dots, c_{k-1} to add up to x
 - table below indicates no optimal solution can do this
- Problem reduces to coin-changing $x - c_k$ cents, which, by induction, is optimally solved by greedy algorithm. ■

k	c_k	All optimal solutions must satisfy	Max value of coins 1, 2, ..., $k-1$ in any OPT
1	1	$P \leq 4$	-
2	5	$N \leq 1$	4
3	10	$N + D \leq 2$	$4 + 5 = 9$
4	25	$Q \leq 3$	$20 + 4 = 24$
5	100	no limit	$75 + 24 = 99$

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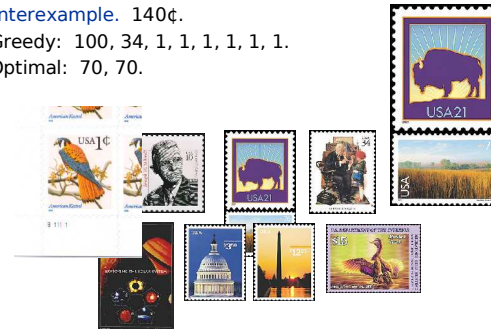
Selecting Breakpoints

Coin-Changing: Analysis of Greedy Algorithm

Observation. Greedy algorithm is sub-optimal for US postal denominations: 1, 10, 21, 34, 70, 100, 350, 1225, 1500.

Counterexample. 140¢.

- Greedy: 100, 34, 1, 1, 1, 1, 1, 1.
- Optimal: 70, 70.



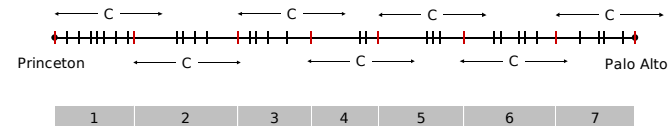
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Selecting Breakpoints

Selecting breakpoints.

- Road trip from Princeton to Palo Alto along fixed route.
- Refueling stations at certain points along the way.
- Fuel capacity = C .
- Goal: makes as few refueling stops as possible.

Greedy algorithm. Go as far as you can before refueling.



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Selecting Breakpoints: Greedy Algorithm

Truck driver's algorithm.

```

Sort breakpoints so that:  $0 = b_0 < b_1 < b_2 < \dots < b_n = L$ 

 $S \leftarrow \{0\}$  ← breakpoints selected
 $x \leftarrow 0$  ← current location

while ( $x \neq b_n$ )
  let  $p$  be largest integer such that  $b_p \leq x + C$ 
  if ( $b_p = x$ )
    return "no solution"
   $x \leftarrow b_p$ 
   $S \leftarrow S \cup \{p\}$ 
return  $S$ 
    
```

Implementation. $O(n \log n)$

- Use binary search to select each breakpoint p .

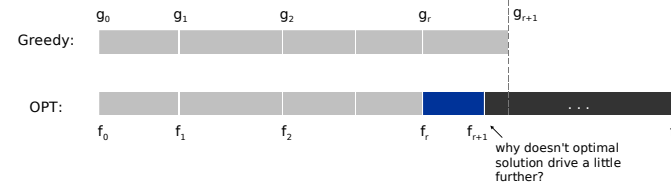
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Selecting Breakpoints: Correctness

Theorem. Greedy algorithm is optimal.

Pf. (by contradiction)

- Assume greedy is not optimal, and let's see what happens.
- Let $0 = g_0 < g_1 < \dots < g_r = L$ denote set of breakpoints chosen by greedy.
- Let $0 = f_0 < f_1 < \dots < f_q = L$ denote set of breakpoints in an optimal solution with $f_0 = g_0, f_1 = g_1, \dots, f_r = g_r$ for largest possible value of r .
- Note: $g_{r+1} > f_{r+1}$ by greedy choice of algorithm.



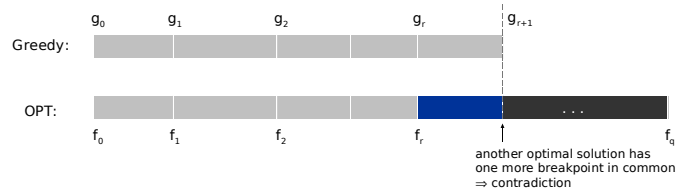
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Edsger W. Dijkstra

The question of whether computers can think is like the question of whether submarines can swim.

Do only what only you can do.

In their capacity as a tool, computers will be but a ripple on the surface of our culture. In their capacity as intellectual challenge, they are without precedent in the cultural history of mankind.

The use of COBOL cripples the mind; its teaching should, therefore, be regarded as a criminal offence.

APL is a mistake, carried through to perfection. It is the language of the future for the programming techniques of the past: it creates a new generation of coding bums.



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