

In class we discussed a pseudo-code of $BFS(s)$; Here I have modified the code to maintain the level of each vertex in the BFS tree, in other words, the array $L[]$ will have the shortest path distance from s to u for any vertex u in the connected component of s .

```

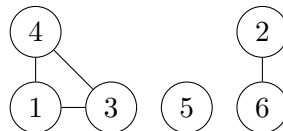
Function  $BFS(s)$ 
  Initialize: mark all vertices "undiscovered"
  mark  $s$  "discovered"
  queue = {  $s$  }
   $L[s]=0$ 
  while queue not empty do
     $u = \text{remove\_first}(\text{queue})$ 
    for each edge  $\{u, x\}$  do
      if  $x$  is undiscovered then
        mark  $x$  discovered
        append  $x$  on queue
         $L[x]=L[u]+1$ 
      end
    end
    mark  $u$  fully-explored
  end

```

Algorithm 1: Computes the shortest path distance from s

Next, we write a code to determine the connected components of a graph. When we call the function Connected-Components, it will construct an array A such that for all vertices v in the same connected component $A[v]$ is the same.

For example, consider the following graph; it has 3 connected components: $\{1, 3, 4\}$, $\{5\}$, $\{2, 6\}$. If we run the code on the following graph, we are going to make 3 BFS calls:



- 1) First we call $BFS(1)$ which visits vertices 1,3,4 and so we get $A[1] = A[3] = A[4] = 1$.
- 2) Then we call $BFS(2)$ which visits vertices 2,6 and so $A[2] = A[6] = 2$.
- 3) Then we call $BFS(5)$ which visits the vertex 5 and so we get $A[5] = 3$.

Note that we are not going to call $BFS(3)$, $BFS(4)$ and $BFS(6)$. Because by the time the main loop gets to vertices 3, 4, and 6 they are already fully-explored.

```

Function BFS(s,c)
  mark s "discovered"
  queue = { s }
  A[s]=c
  while queue not empty do
    | u = remove_first(queue)
    | for each edge {u,x} do
    | | if x is undiscovered then
    | | | mark x discovered
    | | | append x on queue
    | | | A[x]=c;
    | | end
    | end
    | mark u fully-explored
  end

```

```

Function Connected-Components
  Initialize: mark all vertices "undiscovered" and set c = 1
  for v = 1 → n do
    | if v is undiscovered then
    | | BFS(v,c)
    | | c=c+1
    | end
  end

```

Algorithm 2: Computes the Connected Components of a Graph

Also, observe that after running this code, for any pair of vertices u, v , there is a path connecting u to v in G if and only if $A[u] = A[v]$.