## Section 9

1. Consider the query $R(A, B)$ join $S(C, D)$ join $T(E, F)$ (the join condition is $B=C$ and $D=E)$. Suppose $M=100$, and $B(R)=30, B(S)=200, B(T)=60, B(R$ join $S)=80, B(S$ join $T)=50$. Design an optimal query plan that uses only main-memory hash join algorithms. Your plan may store intermediate results to disk if necessary.

Load R \& T into memory and create hash tables of them. Then read blocks of S one at a time, performing the joins in the following graph. All intermediate results are pipelined.

2. Consider the algebra plan below. Each of the joint operators is a main memory hash join algorithm, using the Open( ), GetNext( ), Close( ) interface. Assuming that all joins are pipelining, show the execution steps for computing the entire join.

## Answer(A,B,C,D,E,F,G,H)



Where R, S, T, U have the following content:

| R |  |
| :--- | :--- |
| A 1 | B |
| A2 | B |


| S |
| :--- |
| B |
| D1 |
| B | D 2 B

T

| D1 | F |
| :--- | :--- |
| D2 | F |

U

| $F$ | H1 |
| :--- | :--- |
| F | H2 |

T0.open
T2.open
U.open
U.getNext
U.getNext
U.getNext // got NULL
U.close
T.open

T2.getNext
T.getNext

T2.getNext
T2.getNext
T.getNext

T2.getNext
T2.getNext // got NULL
T2.close
T.close

T1.open
S.open
S.getNext
S.getNext
S.getNext // got NULL
S.close
R.open

T0.getNext
T1.getNext
R.getNext

T0.getNext
T0.getNext
T1.getNext
T0.getNext
T0.getNext
T1.getNext
R.getNext

T0.getNext
T0.getNext
T1.getNext
T0.getNext
T0.getNext
T1.getNext
R.getNext // got NULL

T0.close
T1.close R.close
(b) [10 points] Consider the following query, where $\bowtie$ denotes the natural join:

$$
R(A, B) \bowtie S(B, C) \bowtie T(C, D) \bowtie U(D, E)
$$

Here we only consider left linear plans
i. How many different left linear plans exist for this query ?

$$
\mathrm{n}!
$$

ii. Show two different left linear plans without cartesian products.
(((R join $S$ ) join $T)$ join $U)$
(((T join $S)$ join $U)$ join $R)$
iii. How many different plans without cartesian product exists for this query?

$$
2^{\wedge}(\mathrm{n}-1)
$$

