

# CSE/EE 461 – Lecture 11

## IP Addressing

---

David Wetherall  
djw@cs.washington.edu

## Last Two Times ...

---

- Focus
  - How do we calculate routes for packets?
- Routing Algorithms
  - Distance Vector routing (RIP)
  - Link State routing (OSPF)

Application
Presentation
Session
Transport
<b>Network</b>
Data Link
Physical

## This Lecture

---

- New focus
  - How do we make routing scale?
- IP Addressing
  - Hierarchy (class A, B, C, subnets)
  - Also allocation (DHCP, ARP)

Application
Presentation
Session
Transport
<b>Network</b>
Data Link
Physical

## Scalability Concerns

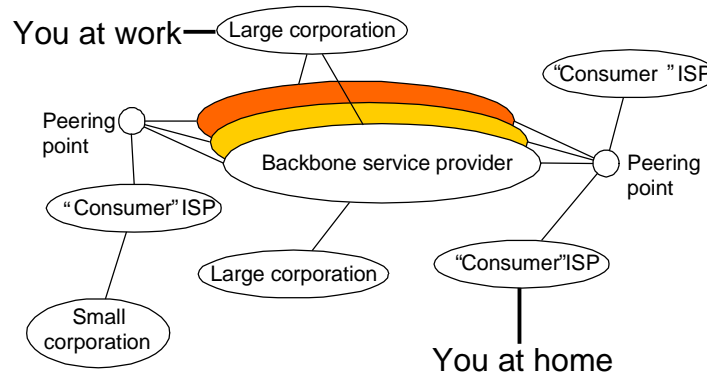
---

- Routing burden grows with size of an internetwork
  - Size of routing tables
  - Volume of routing messages
  - Amount of routing computation
- RIP/OSPF do not scale to the size of the Internet
- We must apply further techniques:
  - Hierarchical addressing ← Today
  - Use of structural hierarchy
  - Route aggregation

## Structure of the Internet

---

- Inter-domain versus intra-domain routing



djw // CSE/EE 461, Autumn 2001

L11.5

## IP Addresses

---

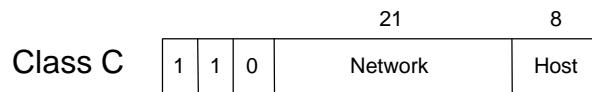
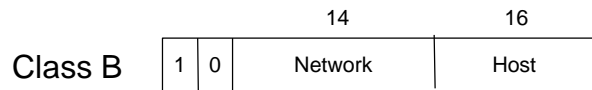
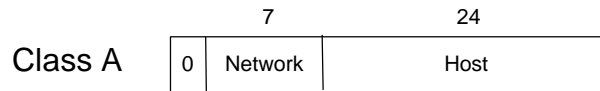
- Reflect location in topology; used for scalable routing
  - Unlike “flat” Ethernet addresses
- Interfaces on same network share prefix
  - Prefix administratively assigned (IANA or ISP)
  - Addresses globally unique
- Routing only advertises entire networks by prefix
  - Local delivery in a single “network” doesn’t involve router
  - (will make “network” precise later on)

djw // CSE/EE 461, Autumn 2001

L11.6

## IPv4 Address Formats

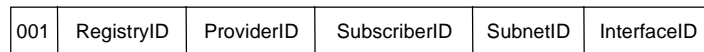
---



- 32 bits written in "dotted quad" notation, e.g., 18.31.0.135

## IPv6 Address Format

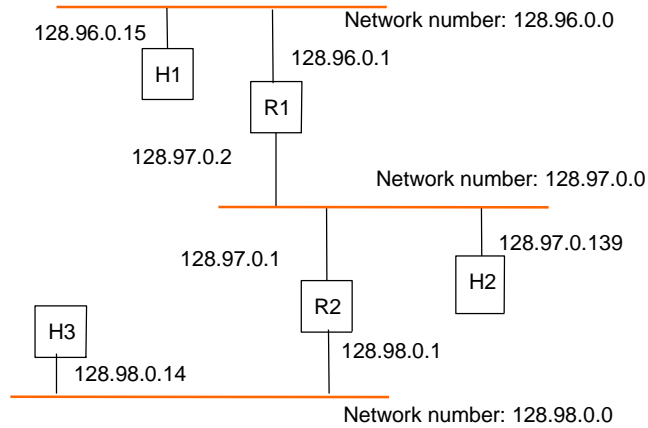
---



- 128 bits written in 16 bit hexadecimal chunks
- Still hierarchical, just more levels

## Network Example

---



djw // CSE/EE 461, Autumn 2001

L11.9

## Updated Forwarding Routine

---

- Used to be “look up destination address for next hop”
- Now addresses have network and host portions:
  - If host: if destination network is the same as the host network, then deliver locally (without router). Otherwise send to the router
  - If router: look up destination network in routing table to find next hop and send to next router. If destination network is directly attached then deliver locally.
- (Note that it will get a little more complicated later)

djw // CSE/EE 461, Autumn 2001

L11.10

## Getting an IP address

---

- Old fashioned way: sysadmin configured each machine
- Dynamic Host Configuration Protocol (DHCP)
  - One DHCP server with the bootstrap info
    - Host address, gateway address, subnet mask, ...
    - Find it using broadcast
  - Addresses may be leased; renew periodically
- “Stateless” Autoconfiguration (in IPv6)
  - Get rid of server – reuse Ethernet addresses for lower portion of address (uniqueness) and learn higher portion from routers

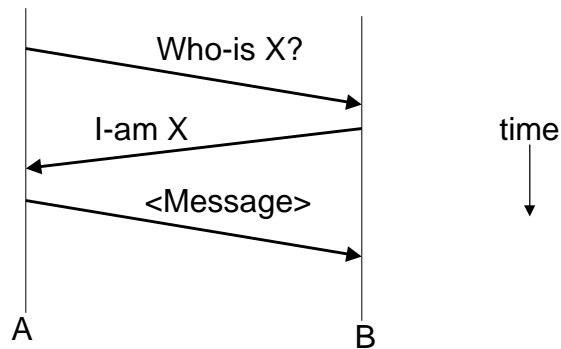
## Address Resolution Protocol (ARP)

---

- On a single link, need Ethernet addresses to send a frame ... source is a given, but what about destination?
  - Requires mapping from IP to MAC addresses
- ARP is a dynamic approach to learn mapping
  - Node A sends broadcast query for IP address X
  - Node B with IP address X replies with its MAC address M
  - A caches (X, M); old information is timed out (~15 mins)
  - Also: B caches A’s MAC and IP addresses, other nodes refresh

## ARP Example

- To send first message use ARP to learn MAC address
- For later messages (common case) don't need to ARP

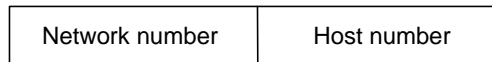


djw // CSE/EE 461, Autumn 2001

L11.13

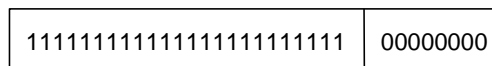
## Subnetting – More Hierarchy

- Split up one network number into multiple physical networks



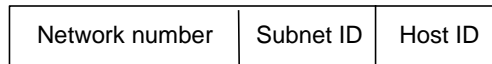
Class B address

- Internal structure isn't propagated



Subnet mask (255.255.255.0)

- Helps allocation efficiency



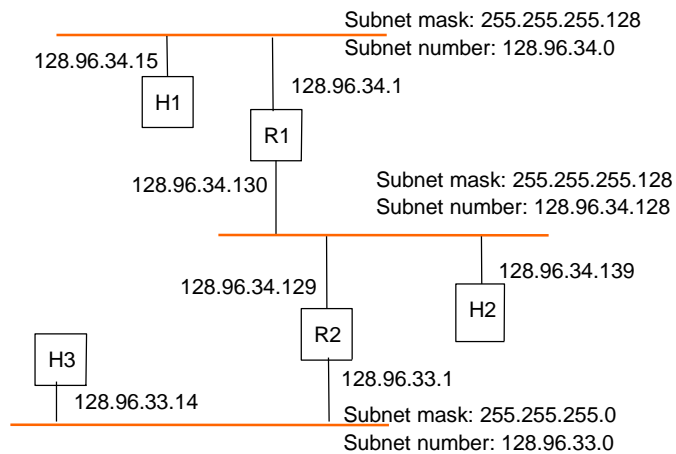
Subnetted address

djw // CSE/EE 461, Autumn 2001

L11.14

## Subnet Example

---



djw // CSE/EE 461, Autumn 2001

L11.15

## Updated Forwarding Routine

---

- Used to know network from address (class A, B, C)
- Now need to “search” routing table for right subnet
  - If host: easy, just substitute “subnet” for “network”
  - If router: search routing table for the subnet that the destination belongs to, and use that to forward as before
- (Note that it will get a little more complicated later :)

djw // CSE/EE 461, Autumn 2001

L11.16



## Key Concepts

---

- Hierarchical address allocation helps routing scale
  - Addresses are constrained by topology
  - Only need to advertise and compute routes for networks
  - Hide internal structure within a domain via subnets
  - Keep host simple and let routers worry about routing
- ARP learns the mapping from IP to MAC address