Why Multiprocessors?

Limits on the performance of a single processor: what are they?

Why Multiprocessors

Lots of opportunity

- Scientific computing/supercomputing
  - Examples: weather simulation, aerodynamics, protein folding
  - Each processor computes for a part of the grid
- Server workloads
  - Example: airline reservation database
  - Many concurrent updates, searches, lookups, queries
  - Processors handle different requests
- Media workloads
  - Processors compress/decompress different parts of image/frames
- Desktop workloads…
- Gaming workloads…

What would you do with 500 million transistors?

Issues in Multiprocessors

Which **programming model for interprocessor communication**

- shared memory
  - regular loads & stores
  - SPARCCenter, SGI Challenge, Cray T3D, Convex Exemplar, KSR-1&2, today's CMPs
- message passing
  - explicit sends & receives
  - TMC CM-5, Intel Paragon, IBM SP-2

Which **execution model**

- control parallel
  - identify & synchronize different asynchronous threads
- data parallel
  - same operation on different parts of the shared data space

Issues in Multiprocessors

How to **express parallelism**

- language support
  - HPF, ZPL
- runtime library constructs
  - coarse-grain, explicitly parallel C programs
- automatic (compiler) detection
  - implicitly parallel C & Fortran programs, e.g., SUIF & PTRANS compilers

Application development

- embarrassingly parallel programs could be easily parallelized
- development of different algorithms for same problem

Issues in Multiprocessors

How to get **good parallel performance**

- recognize parallelism
- transform programs to increase parallelism without decreasing processor locality
- decrease sharing costs

Flynn Classification

SISD: single instruction stream, single data stream
- single-context uniprocessors

SIMD: single instruction stream, multiple data streams
- exploits data parallelism
- example: Thinking Machines CM

MISD: multiple instruction streams, single data stream
- systolic arrays
- example: Intel iWarp, today's streaming processors

MIMD: multiple instruction streams, multiple data streams
- multiprocessors
- multithreaded processors
- parallel programming **&** multiprogramming
- relies on control parallelism: execute & synchronize different asynchronous threads of control
- example: most processor companies have CMP configurations

CM-1

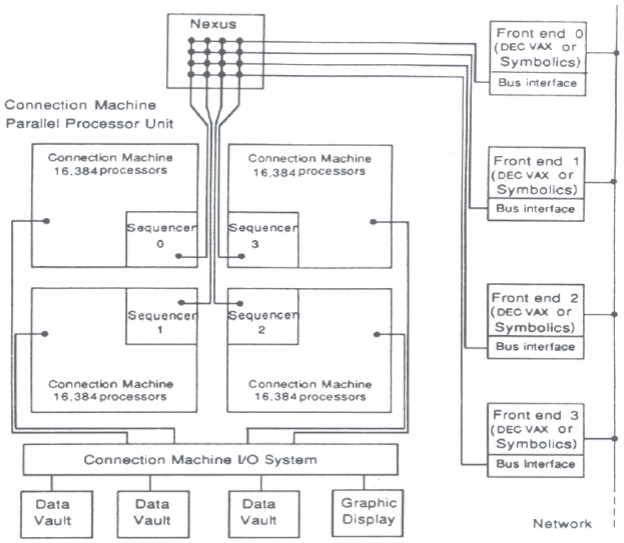
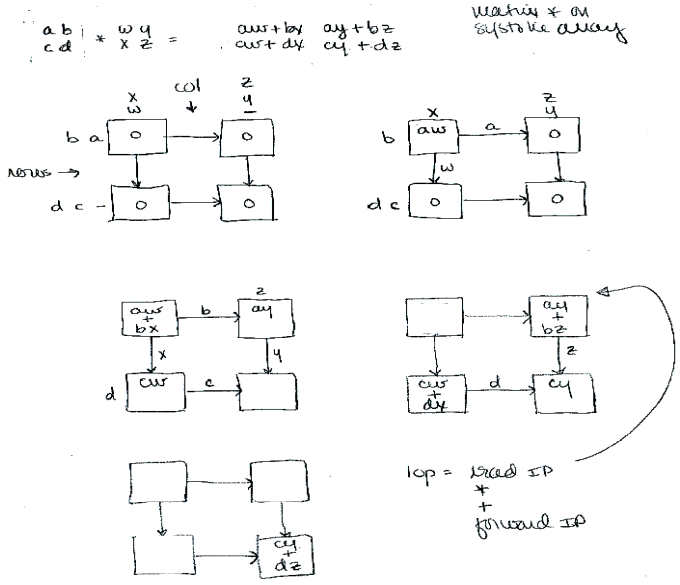


Figure 1. Connection Machine system organization.

Systolic Array

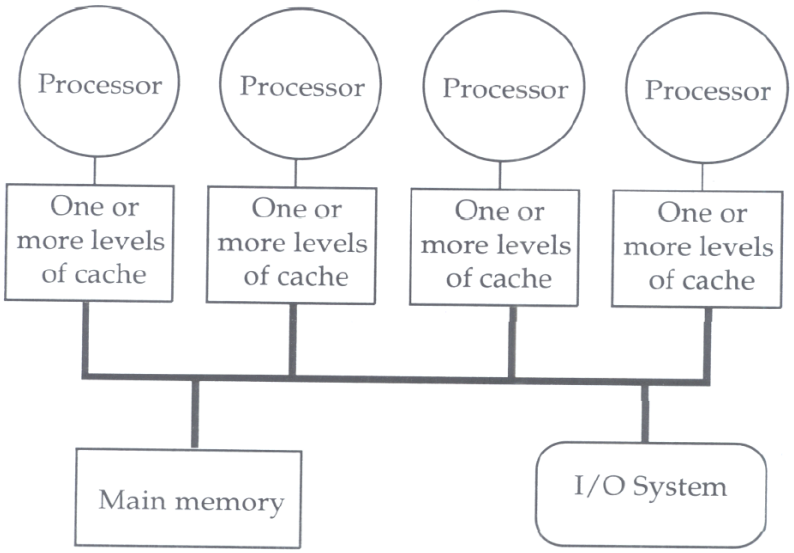


MIMD

Low-end

- bus-based
  - simple, but a bottleneck
  - simple cache coherency protocol
- physically centralized memory
- uniform memory access (UMA machine)
- Sequent Symmetry, SPARCCenter, Alpha-, PowerPC- or SPARC-based servers, most of today's CMPs

Low-end MP

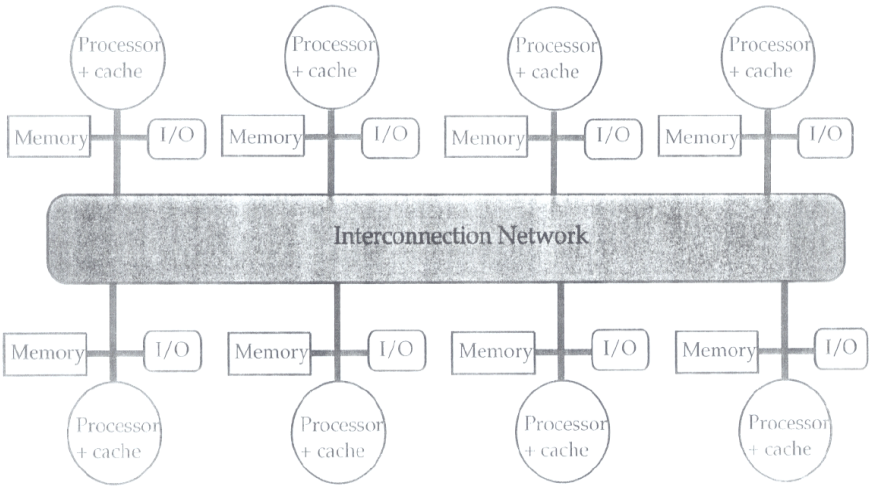


MIMD

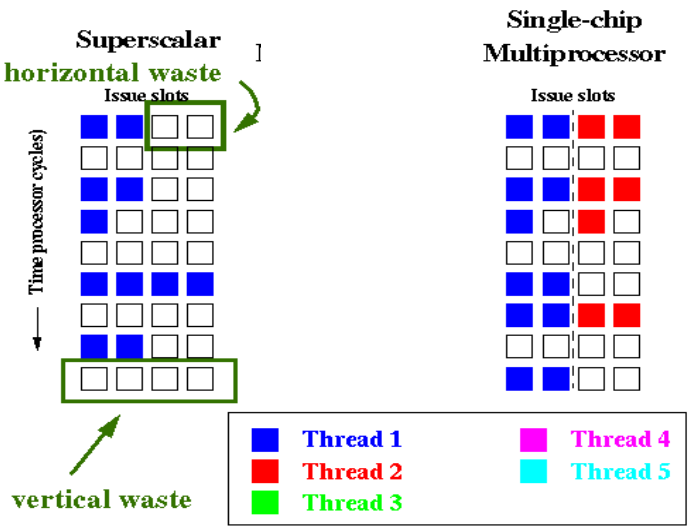
High-end

- higher bandwidth, multiple-path interconnect
  - more scalable
  - more complex cache coherency protocol (if shared memory)
  - longer latencies
- physically distributed memory
- non-uniform memory access (NUMA machine)
- could have processor clusters
- SGI Challenge, Convex Examplar, Cray T3D, IBM SP-2, Intel Paragon, Sun T1

High-end MP



Comparison of Issue Capabilities



Shared Memory vs. Message Passing

Shared memory

+ simple parallel programming model
  - global shared address space
  - not worry about data locality ***but***
    get better performance when program for data placement
    lower latency when data is local
      - ***but*** can do data placement if it is crucial, but don't have to
  - hardware maintains data coherence
    - synchronize to order processor's accesses to shared data
  - like uniprocessor code so parallelizing by programmer or compiler is easier

⇒ can focus on program semantics, not interprocessor communication

Shared Memory vs. Message Passing

Shared memory

- + low latency (no message passing software) ***but***
  - *overlap of communication & computation*
  - *latency-hiding techniques can be applied to message passing machines*
- + higher bandwidth for small transfers ***but***
  - *usually the only choice*

Shared Memory vs. Message Passing

Message passing

- + abstraction in the programming model encapsulates the communication costs ***but***
  - *more complex programming model*
  - *additional language constructs*
  - *need to program for nearest neighbor communication*
- + no coherency hardware
- + good throughput on large transfers ***but***
  - *what about small transfers?*
- + more scalable (memory latency for uniform memory doesn't scale with the number of processors) ***but***
  - *large-scale SM has distributed memory also*
    - ***hah!*** so you're going to adopt the message-passing model?

Shared Memory vs. Message Passing

Why there was a debate

- little experimental data
- not separate implementation from programming model
- can emulate one paradigm with the other
  - MP on SM machine
    message buffers in local (to each processor) memory
    copy messages by ld/st between buffers
  - SM on MP machine
    ld/st becomes a message copy
    slooooooooow

Who won?