

WaveScalar: the Executive Summary

A modern dataflow machine

- solves the language compatibility & memory ordering issues
- solves the scalability issue

The executive summary:

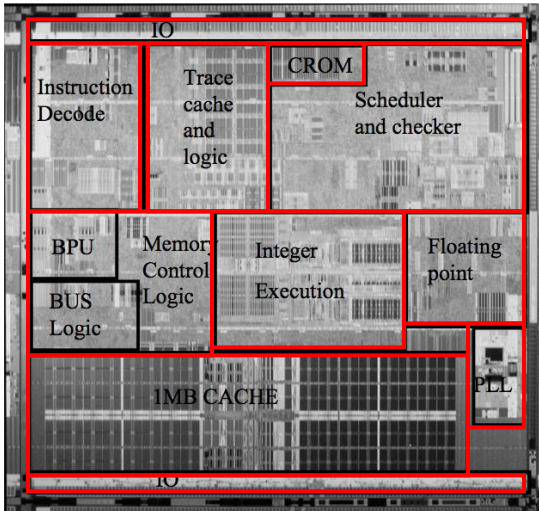
- good at exploiting ILP (dataflow parallelism)
- also traditional coarser-grain parallelism
 - cheap thread management
- low operand latency because of a hierarchical PE-interconnect organization
- memory ordering enforced through **wave-ordered memory**
 - can execute imperative language programs
 - no special dataflow languages

WaveScalar

Motivation stems from shrinking feature sizes:

- increasing disparity between computation (fast transistors) & communication (long wires)
- increasing circuit complexity
- decreasing fabrication reliability

Monolithic von Nuemann Processors



A success a few years ago.
But in 2016?

- ⊗ Performance
Centralized processing & control
Long wires
e.g., operand broadcast networks
- ⊗ Complexity
40-75% of “design” time is design verification
- ⊗ Defect tolerance
1 flaw -> tie pin, earrings, ...

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WaveScalar's Microarchitecture

Good performance via distributed microarchitecture 😊

- hundreds of PEs
- organized hierarchically for fast communication between neighboring PEs
- short point-to-point (producer to consumer) operand communication
- dataflow execution – no centralized control via a PC
- consequently scalable

Low design complexity through simple, identical PEs 😊

- design one & stamp out hundreds

Defect tolerance 😊

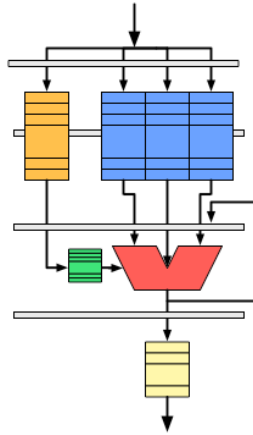
- route around a bad PE

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Processing Element



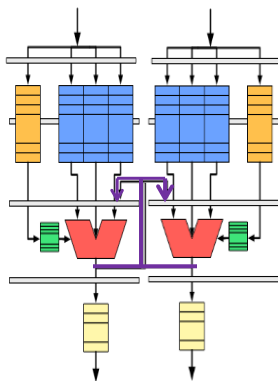
- Simple, small (.5M transistors)
- 5-stage pipeline (receive input operands, match tags, instruction issue, execute, send output)
- Holds 64 (decoded) instructions
- 128-entry token store
- 4-entry output buffer

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PEs in a Pod



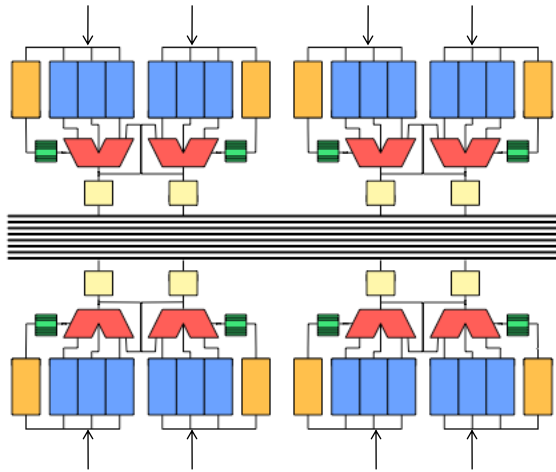
- Share operand bypass network
- Back-to-back producer-consumer execution across 2 PEs

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Domain

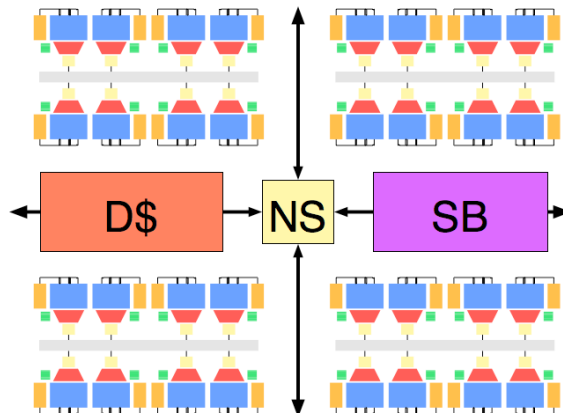


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Cluster



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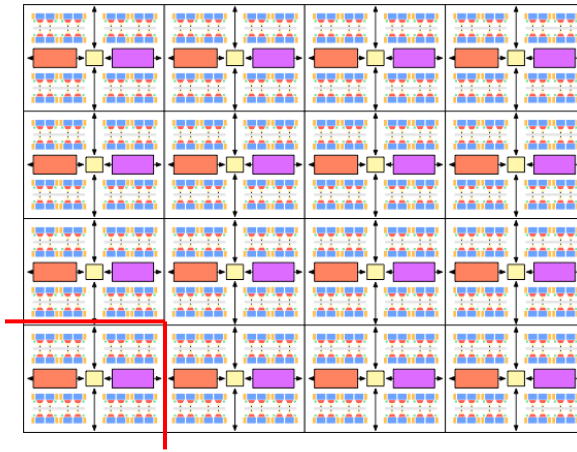
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WaveScalar Processor

Long distance

communication

- grid-based network
- 2-cycle hop/cluster
- dynamic routing



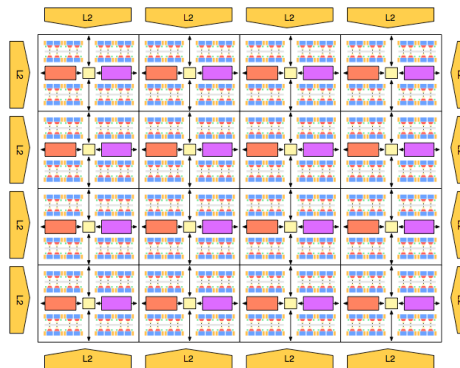
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Whole Chip

- Can hold 32K instructions
- Normal memory hierarchy
- Traditional directory-based cache coherence
- ~400 mm² in 90 nm technology
- 1GHz.
- ~85 watts



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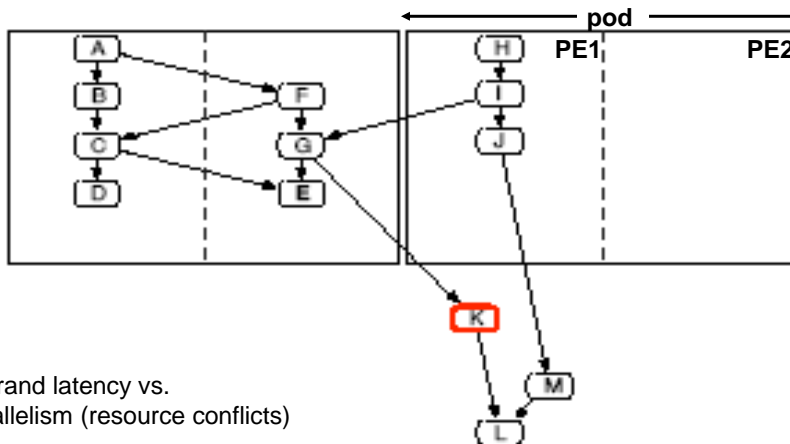
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WaveScalar Instruction Placement



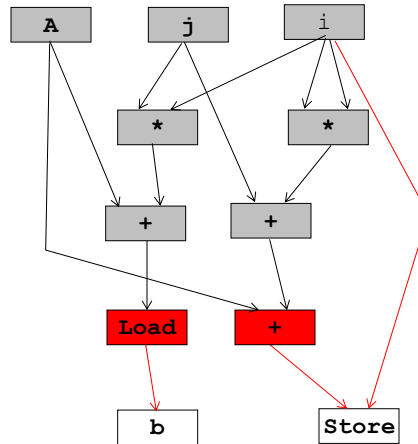
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Revisit Example to Illustrate the Memory Ordering Problem

```
A[j + i*i] = i;
b = A[i*j];
```



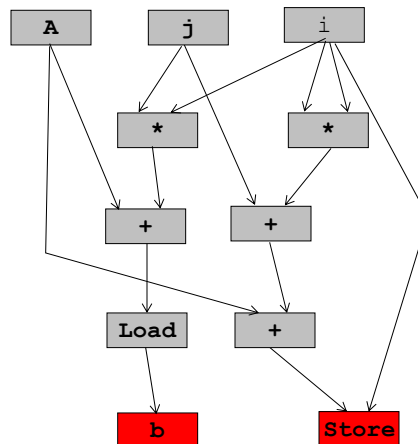
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Revisit Example to Illustrate the Memory Ordering Problem

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A[j + i*i] = i;
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```



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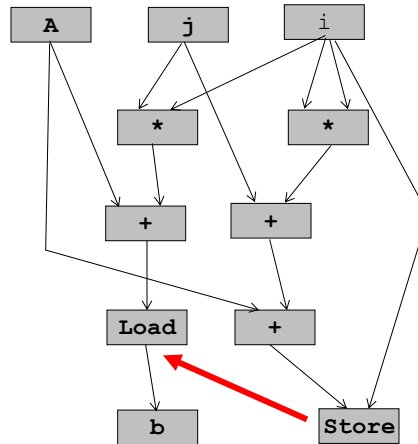
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Revisit Example to Illustrate the Memory Ordering Problem

`A[j + i*i] = i;`

`b = A[i*j];`



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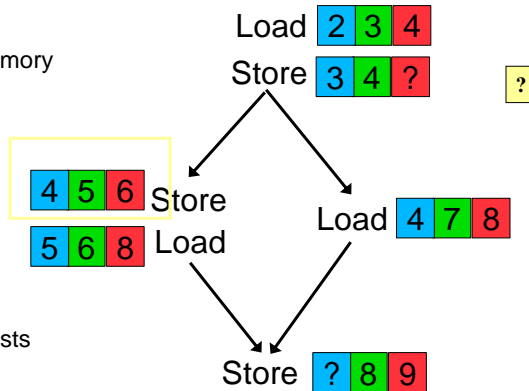
Wave-ordered Memory

- Compiler annotates memory operations

■ Sequence #

■ Successor

■ Predecessor

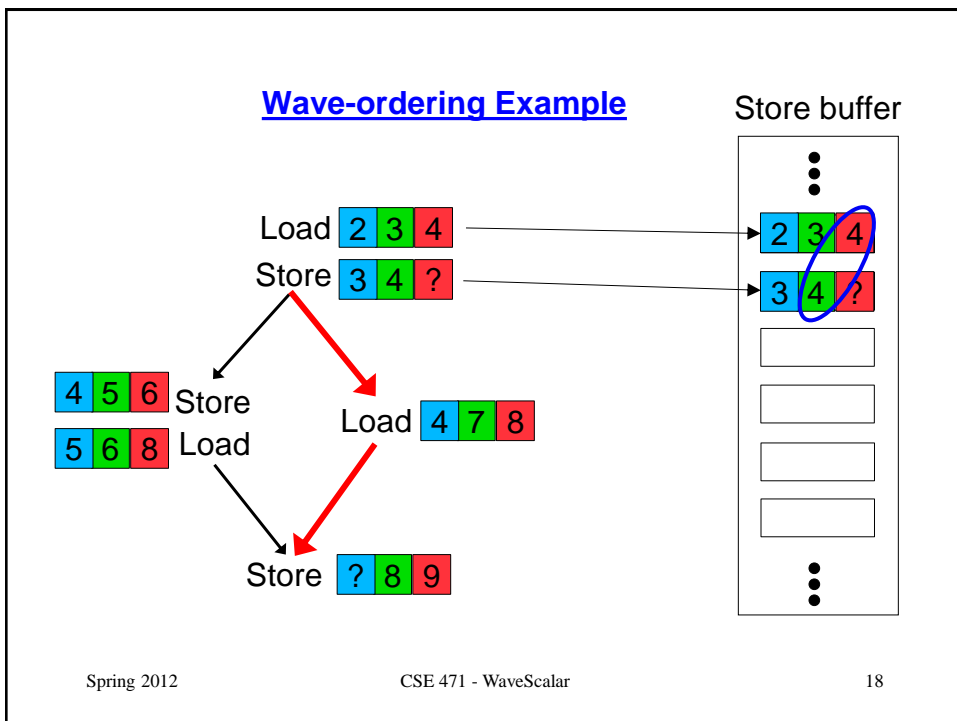
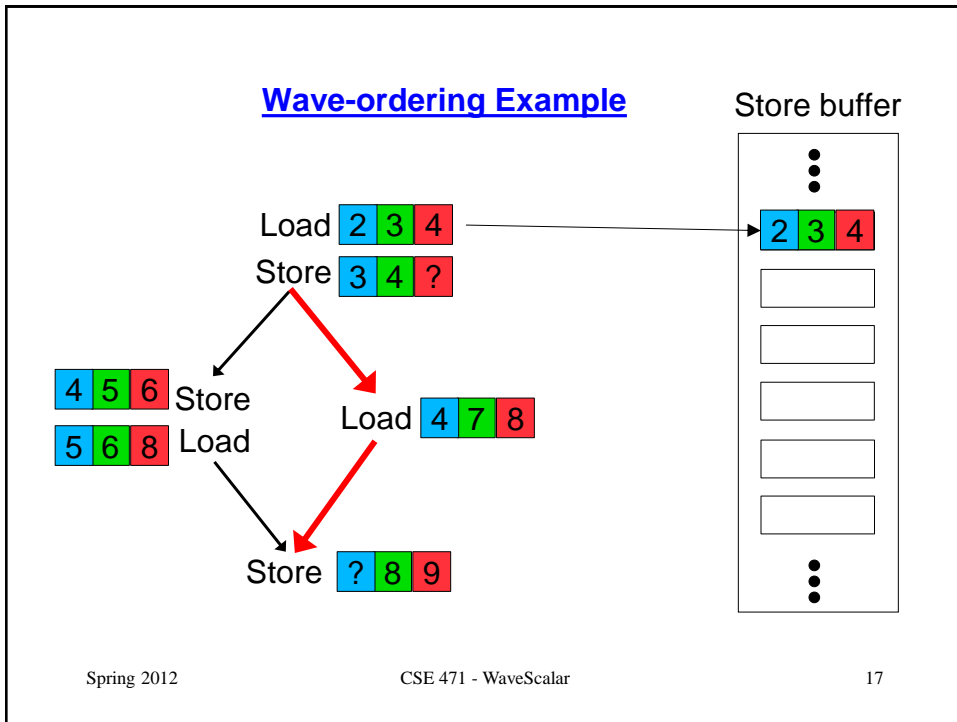


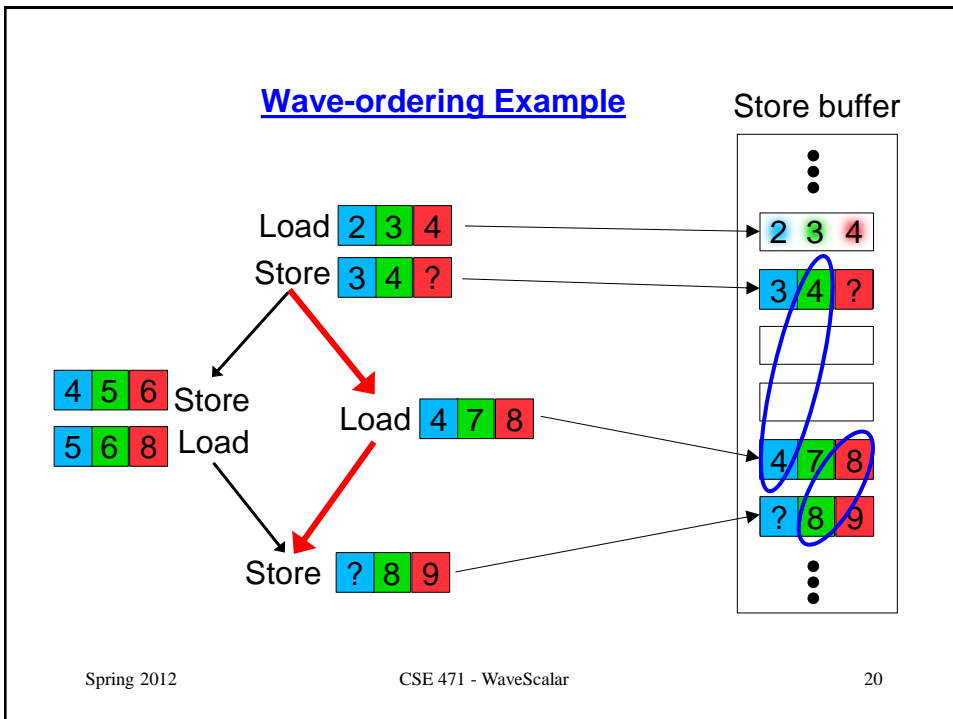
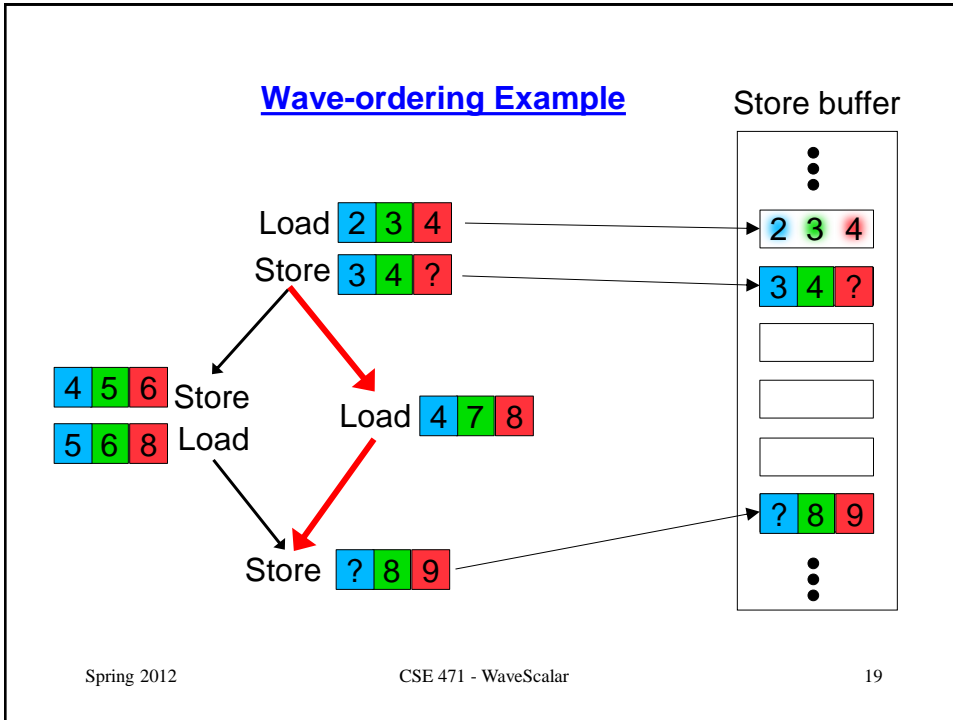
- Execute memory requests in any order
- Store buffer hardware reconstructs the correct order

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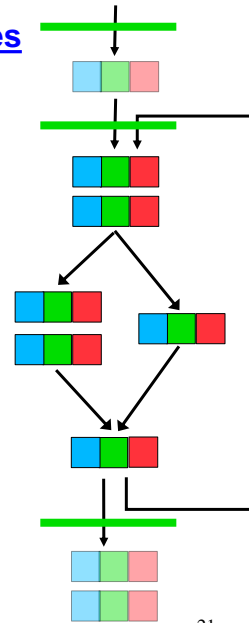
Wave-ordered Memory across Waves

Waves are loop-free sections of the dataflow graph

Each *dynamic* wave has a **wave number**
Wave number is incremented between waves

Ordering memory in a whole program:

- wave-numbers
- sequence numbers within a wave
- can execute imperative language programs



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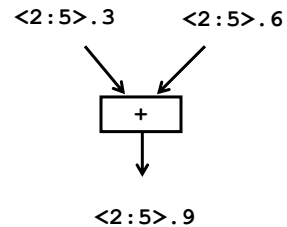
WaveScalar Tag-matching

WaveScalar tag

- thread identifier
- wave number

Token: **tag** & **value**

<ThreadID:Wave#>.value



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Multithreading the WaveCache

Architectural-support for WaveScalar threads

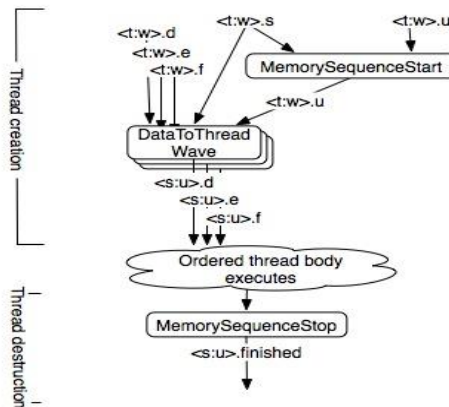
- instructions to start & stop memory orderings, i.e., threads
- memory-free synchronization to allow exclusive access to data (thread communicate instruction)
- “barrier” instruction to force all previous memory operations to fully execute (to allow other threads to see the results of this one’s memory ops)

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Creating & Terminating a Thread



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Multithreading the WaveCache

Combine to build threads with multiple granularities

- coarse-grain threads: 25-168X over a single thread; 2-16X over CMP, 5-11X over SMT
- fine-grain, dataflow-style threads: 18-242X over single thread
- a demonstration that one can combine the two in the same application (*equake*): 1.6X or 7.9X -> 9X

Important Issues

Original dataflow machines

- comparison to von Neumann architectures
- dataflow firing rule
- token
- handling branches
- problems

Modern dataflow machines, aka Wavescalar

- comparison to von Neumann microarchitecture
- how solve past dataflow problems
- hierarchical structure
- wave-ordered memory
- thread management