

## Control dependence

Must ensure side-effects occur in proper order  
Must ensure side-effects occur only under right conditions

CFG represents control dependence explicitly  
– but overspecifies control dependence requirements

## Control dependence graph

Program dependence graph (PDG):  
data dependence graph + control dependence graph (CDG)  
[Ferrante, Ottenstein, & Warren, TOPLAS 87]

Idea: represent controlling conditions directly

- complements data dependence representation

A node (basic block)  $Y$  is **control-dependent** on another  $X$  iff  $X$  determines whether  $Y$  executes, i.e.

- there exists a path from  $X$  to  $Y$  s.t. every node in the path other than  $X$  &  $Y$  is **post-dominated** by  $Y$
- $X$  is not post-dominated by  $Y$

Control dependence graph:

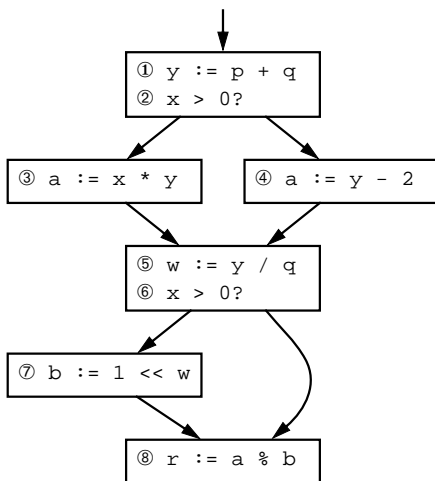
$Y$  proper descendant of  $X$  iff  $Y$  control-dependent on  $X$

- label each child edge with required branch condition
- group all children with same condition under **region** node

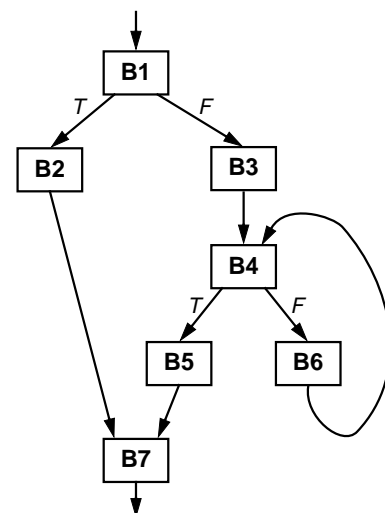
Two sibling nodes execute under same control conditions  $\Rightarrow$   
can be reordered or parallelized, as data dependences allow

Challenging to “sequentialize” back into CFG form

## Example



## An example with a loop



## Value dependence graphs

[Weise, Crew, Ernst, & Steensgaard, POPL 94]

- Idea: represent all dependences,  
including control dependences, as data dependences
- + simple, direct dataflow-based representation of all "interesting" relationships
  - analyses become easier to describe & reason about
  - harder to sequentialize into CFG

Control dependences as data dependences:

- control dependence on order of side-effects  
⇒ data dependence on reading & writing to global Store
- optimizations to break up accesses to single Store into separate independent chunks  
(e.g. a single variable, a single data structure)
- control dependence on outcome of branch  
⇒ a select node, taking test, then, and else inputs  
⇒ demand-driven evaluation model

Loops implemented as tail-recursive calls to local procedures

Apply CSE, folding, etc. as nodes are built/updated

## VDG for example, after store splitting

```
y := p + q
if x > 0 then a := x * y else a := y - 2
w := y / q
if x > 0 then b := 1 << w
r := a % b
```

