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CSE505: Programming Languages  
Lecture 18: Bounded polymorphism,  
Classless OOP

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# Revenge of Type Variables

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- Sorted lists in ML (partial):

```
make : ('a -> 'a -> int) -> 'a slist
cons : 'a slist -> 'a -> 'a slist
find : 'a slist -> ('a -> bool) -> 'a option
```

- Sorted lists with OO subtyping (null has any type)

```
interface Cmp { Int f(Object, Object); }
interface Pred { Bool g(Object); }
class SList {
  constructor (Cmp x) { ... }
  SList cons (Object x) { ... }
  Object find (Pred x) { ... }
}
```

# Still want generics

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- Will downcast args to f, arg to g, result of a find
- Not enforcing type equality of list elements
- OO subtyping no replacement for parametric polymorphism

So have both:

```
interface 'a Cmp { Int f('a, 'a); } // not a type
interface 'a Pred { Bool g('a); } // not a type
class 'a SList { // not a type (but Int SList is)
  constructor ('a Cmp x) { ... }
  SList cons ('a x) { ... }
  Object find ('a Pred x) { ... }
}
```

# Same Old Story

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- Interface and class declarations are parameterized; they produce types
- The constructor is polymorphic (for all  $T$ , given  $T$  `cmp`, it makes a  $T$  `sList`)
- If  $o$  has type  $T$  `sList`, its `cons` method takes  $T$  and returns a  $T$  `sList`

No more downcasts; the best of both worlds

# Complications

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“Interesting” interaction with overloading or multimethods

```
class B {  
  Int f(Int C x){1}  
  Int f(String C x){2}  
  Int g('a C x) { self.f(x) }  
}
```

Whether match is found depends on instantiation of 'a  
Cannot resolve static overloading at compile-time without  
code duplication

At run-time, need run-time type information

- Including instantiation of type constructors
- Or restrict overloading enough to avoid it

# Wanting bounds

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As expected, with subtyping and generics, want bounded polymorphism

Example:

```
interface I { unit print(); }  
class (<A:I) Logger {  
    'a item;  
    'a get() { item.print(); item }  
}
```

w/o polymorphism, `get` would return an `I` (not useful)

w/o the bound, `get` could not **send** print to `item`

# Fancy example

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With forethought, can use bounds to avoid some subtyping limitations

(Example lifted from Abadi/Cardelli text)

```
/* Herbivore1  $\leq$  Omnivore1 unsound */
interface Omnivore1 { unit eat(Food); }
interface Herbivore1 { unit eat(Veg); }
/* T Herbivore2  $\leq$  T Omnivore2 sound for any T */
interface ('a $\leq$ Food) Omnivore2 { unit eat('a); }
interface ('a $\leq$ Veg) Herbivore2 { unit eat('a); }
/* subtyping lets us pass herbivores to feed
   but only if food is a Veg */
unit feed('a food, 'a Omnivore animal) {
  animal.eat(food);
}
```

# Bounded Polymorphism

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Is useful in any language with universal types and subtyping.

Instead of  $\forall \alpha. \tau$  and  $\wedge \alpha. e$ , we have  $\forall \alpha < \tau'. \tau$  and  $\wedge \alpha < \tau'. e$ :

- Change  $\Delta$  to be a set of bounds ( $\alpha < \tau$ ) not just a set of type variables
- In  $e$  you can subsume from  $\alpha$  to  $\tau'$
- $e_1 [\tau_1]$  typechecks only if  $\tau_1$  “satisfies the bound” in the type of  $e_1$

One meta-theory drawback: When is  $\forall \alpha < \tau_1. \tau_2 \leq \forall \alpha < \tau_3. \tau_4$  ?

Contravariant bounds (and covariant bodies assuming bound) are sound, but makes subtyping undecidable.

Requiring invariant bounds (more restrictive) regains decidability.



# Classless OOP

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- OOP gave us code-reuse via inheritance and extensibility via late-binding
- But it also gave us a clunky, heavyweight class and named-type mechanism
- Can we throw out classes and still get OOP? Yes
- Can it have a type system that prevents “no match found” and “no best match” error?
  - yes, but we won’t get there.
- We will make up syntax as we go along
  - with subtitles in JavaScript
- This is mind-opening/bending stuff if you’ve never seen it

# Making objects directly

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- Everything is an object. You can make objects directly:

```
let p = [  
  field x = 7;  
  field y = 9;  
  right_quad() { x.gt(0) && y.gt(0) }  
]
```

- No classes. Constructors are easy to encode:

```
let make_p = [  
  doit(x0, y0) { [field x=x0; field y=y0; ... ] }  
]
```

# Making objects directly (JavaScript)

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- Everything is an object. You can make objects directly:

```
p = new Object;  
p.x = 7;  
p.y = 9;  
p.right_quad =  
function() { return (this.x > 0) && (this.y > 0); }
```

- No classes. Constructors are easy to encode:

```
function point(x, y) {  
  this.x = x;  
  this.y = y;  
  function right_quad() {return (this.x>0 && this.y>0);}  
}  
p = new point(7, 9);
```

# Inheritance and Override

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Building objects from scratch won't get us late-binding

and code reuse. Here is the trick:

- clone method produces a (shallow) copy of an object.
- method slots can be mutable

```
let o1 = [ // still have late binding
  odd = { if x.eq(0) then false else self.even(x-1) }
  even = { if x.eq(0) then true  else self.odd(x-1) }
]
let o2 = o1.clone()
o2.even := fun (x) { (x.mod(2)).eq(0) }
```

Language does not grow (just methods and mutable “slots”)

# Inheritance and Override (JavaScript)

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Building objects from scratch won't get us late-binding

and code reuse. Here is the trick:

- clone method produces a (shallow) copy of an object.
- method slots can be mutable

```
o1 = new Object;

o1.odd = function (x)
{if (x == 0) {return false;} else {return this.even(x-1);}}
o1.even = function (x)
{if (x == 0) {return true;} else {return this.odd(x-1);}}

o2 = o1.clone(); // not built-in in JavaScript
o2.even = function (x) { return x - (x / 2) == 0; }
```

Language does not grow (just methods and mutable “slots”)

# clone() in JavaScript

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```
Object.prototype.clone = function()
{
    var newObj = new this.constructor();
    newObj.__proto__ = this;    // not in IE
    return newObj;
};
```

# Extension

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- But that trick does not work to add slots to an object, a common use of subclassing.
- Having something like “`extend e1 (x=e2)`” that mutates `e1` to have a new slot is problematic semantically (what if `e1` has a slot named `x`) and for efficiency (may not be room where `e1` is allocated)
- Instead, we can build a new object with special parent slot:

```
[parent = e1; x = e2]
```

- parent is very special because definition of method lookup (the issue in OO) depends on it.

# Extension (in JavaScript)

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Extending an object:

```
function Parent() { ... }  
function Child(e1) { this.x = e1; }  
c = new Child(0);  
c.__proto__ = new Parent; // Not in IE
```

Another way of extending: via constructors

```
function Parent() { ... }  
function Child(e1) { this.x = e1; }  
Child.prototype = new Parent;  
c = new Child(0);
```



# Method Lookup

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- To find the **m** method of **o**
  - look for a slot named **m**
  - if not found, look in object held in parent slot
- But we still have late-binding: for method in parent slot, we still have **self** refer to the original **o**.
- Two inequivalent ways to define **parent = e1**:
  - delegation: **parent** refers to result of **e1**
  - embedding: **parent** refers to the result of **e1.clone**
- Mutation of result of **e1** (or its parent or grandparent or ...) exposes the difference.
- We'll assume delegation

# Oh So Flexible

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Delegation is more flexible (and simple) (and dangerous) than class-based OO: The object being delegated to is usually used like a class, but its slots may be mutable.

- Assigning to a slot in a delegated object changes every object that delegates to it (transitively)
  - clever change-propagation but as dangerous as globals (and more subtle?)
- Assigning to a parent slot is “dynamic inheritance” (changes where slots are inherited from)

Classes restrict what you can do and how you think (never thinking of clever run-time modifications of inheritance)

# Rarely What You Want

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We have the essence of OOP in a tiny language with more flexibility than we usually want.

Avoid it via careful coding idioms:

- create *trait/abstract* objects: just immutable methods (cf. abstract classes)
- extend with *prototype/template* objects: add mutable fields but don't mutate them (cf. classes)
- clone prototype to create *concrete/normal* objects (cf. constructors)

Traits can extend other traits and prototypes other prototypes (cf. subclasses)

# Coming Full Circle

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- Without separating first two roles, objects don't share method slots (wastes space), but immutably avoids danger.
- Late-binding still makes method-override work correctly
- This idiom is so important, it's worth having a type system that enforces it
- For example, a template object cannot have its members accessed (except clone).
- We end up getting close to classes, but from first principles and still allowing the full flexibility when you want it.

# A Word on Types

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- Untyped languages work (the OO Scheme) – may get a “no match found” exception at run-time. Very flexible.
- But we can develop type systems that restrict the language and prevent getting stuck without developing a class system
- Can base types on “derived from the same object”, which can form the basis for multimethods
- Summary: pure classless OO is a liberating way to think, especially if you learn workarounds in more restrictive languages.