## In-order vs. Out-of-order Execution

#### In-order instruction execution

- instructions are fetched, executed & completed in compilergenerated order
- one stalls, they all stall
- instructions are statically scheduled

### **Out-of-order instruction execution**

- instructions are fetched in compiler-generated order
- instruction completion may be in-order (today) or out-of-order (older computers)
- in between they may be executed in some other order
- independent instructions behind a stalled instruction can pass it
- instructions are dynamically scheduled

### **Out-of-order processors:**

- after instruction decode
  - check for structural hazards
    - an instruction can be issued when a functional unit is available
    - an instruction stalls if no appropriate functional unit
  - check for data hazards
    - an instruction can execute when its operands have been calculated or loaded from memory
    - · an instruction stalls if operands are not available

## **Out-of-order processors:**

- don't wait for previous instructions to execute if this instruction does not depend on them
- ready instructions can execute before earlier instructions that are stalled, e.g., waiting for their data to be loaded from memory
  - when go around a load instruction that is stalled for a cache miss:
    - use lockup-free caches that allow instruction issue to continue while a miss is being satisfied
    - the load-use instruction still stalls

## in-order processors

lw \$3, 100(\$4) add \$2, \$3, \$4 sub \$5, \$6, \$7 in execution, cache miss waits until the miss is satisfied waits for the add

### *out-of-order processors*

lw \$3, 100(\$4)
sub \$5, \$6, \$7
add \$2, \$3, \$4

in execution, cache miss in execution during the cache miss waits until the miss is satisfied

### **Out-of-order processors:**

- ready instructions can execute before earlier instructions that are stalled, e.g., waiting for their branch condition to be computed
  - when go around a branch instruction:
    - the instructions that are issued from the predicted path are issued speculatively, called speculative execution
    - speculative instructions can execute (but not commit) before the branch is resolved
    - if the prediction was wrong, speculative instructions are flushed from the pipeline
    - if prediction is right, instructions are no longer speculative

# **Speculation**

Instruction **speculation**: executing an instruction before it is known that it should be executed

- must be safe (no additional exceptions) or must handle the exceptions after the instruction is no longer speculative
- must generate the same results