

Foundations of Artificial Intelligence

CSE 573 — Fall 2001

Game Playing

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Game Playing – Why?

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## Game Playing

### An AI Favorite

- structured task
- not initially thought to require large amounts of knowledge
- focus on games of perfect information

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## Game Playing

Initial State

Operators

Terminal Test

Utility Function

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## Game Playing

**Initial State** is the initial board/position

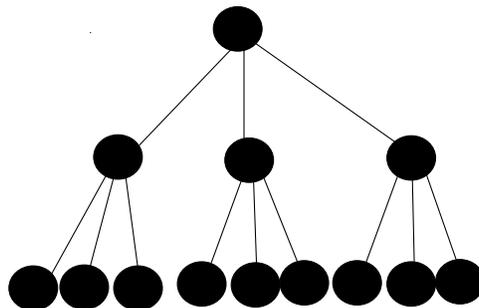
**Operators** define the set of legal moves from any position

**Terminal Test** determines when the game is over

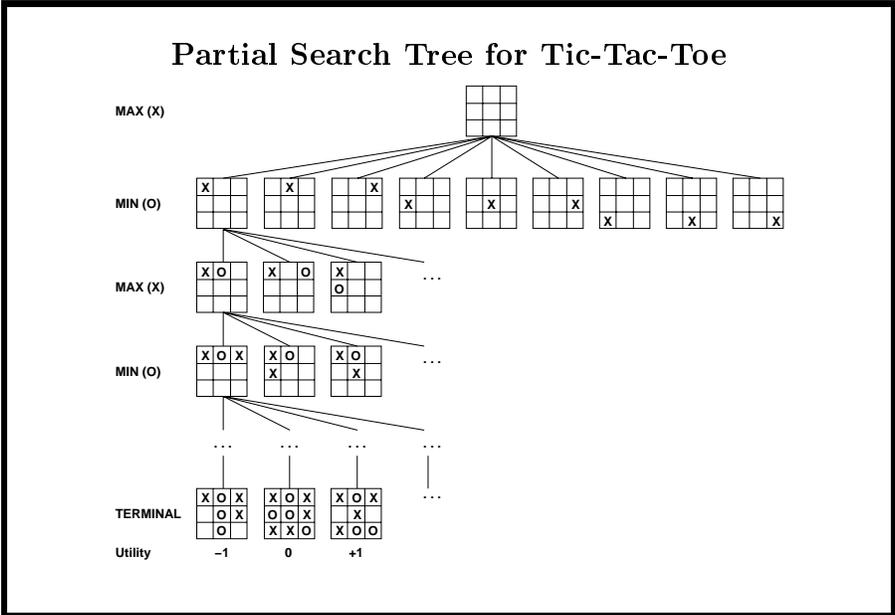
**Utility Function** gives a numeric outcome for the game

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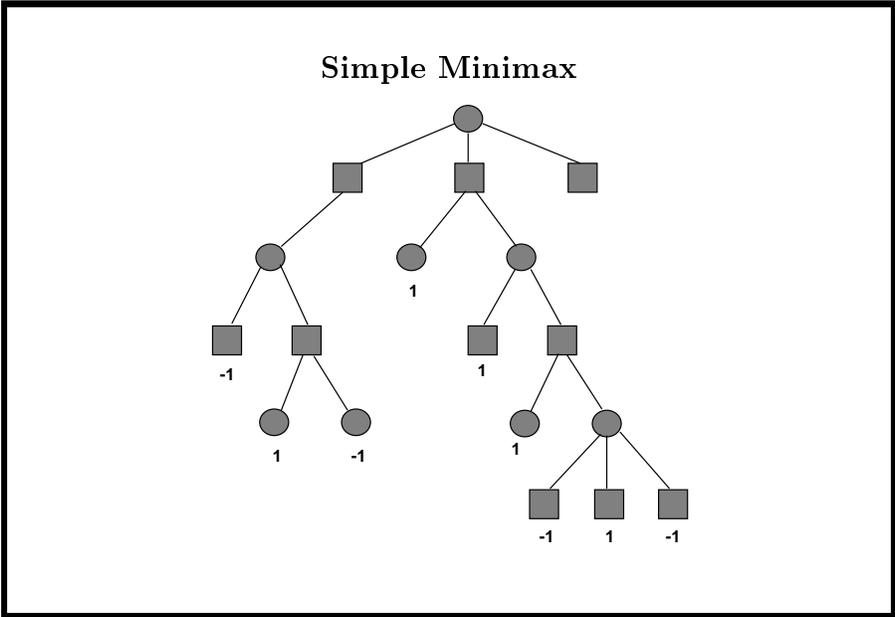
## Game Playing as Search



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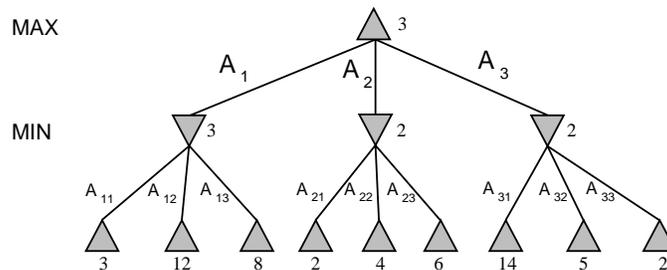
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### Simplified Minimax Algorithm

1. Expand the entire tree below the root.
2. Evaluate the terminal nodes as wins for the minimizer or maximizer.
3. Select an unlabeled node,  $n$ , all of whose children have been assigned values. If there is no such node, we're done — return the value assigned to the root.
4. If  $n$  is a minimizer move, assign it a value that is the minimum of the values of its children. If  $n$  is a maximizer move, assign it a value that is the maximum of the values of its children. Return to Step 3.

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### Another Example



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1. In game tree search, a *move* is a pair of actions. One player's action is a *ply*. 2-ply = one move.
2. Called a **minimax** decision because it maximizes the utility under the assumption that the opponent will play perfectly to minimize it.
3. Time complexity:  $O(b^m)$  ( $m$  plies and  $b$  branching.)  
Impractical for e.g. chess ( $b \approx 30$  to  $40$ ).

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### **The Need for Imperfect Decisions**

**Problem:** Minimax assumes the program has time to search to the terminal nodes.

**Solution:**

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### **The Need for Imperfect Decisions**

**Problem:** Minimax assumes the program has time to search to the terminal nodes.

**Solution:** Cut off search earlier and apply a **heuristic evaluation function** to the leaves.

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### **Static Evaluation Functions**

Minimax depends on the translation of board quality into a single, summarizing number. Can be difficult. Expensive.

- Add up values of pieces each player has (weighted by importance of piece). E.g. intro chess: pawn = 1pnt; knight or bishop = 3pnts; rook = 5 pnts; and queen = 9pnts.
- Isolated pawns are bad. How well protected is your king? How much maneuverability to you have? Do you control the center of the board?
- How many plies with perfect eval function?

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## Design Issues of Heuristic Minimax

**Evaluation Function:** What features should we evaluate and how should we use them? An evaluation function should:

- 1.
- 2.
- 3.

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## Linear Evaluation Functions

- $w_1 f_1 + w_2 f_2 + \dots + w_n f_n$
- $w$  — weight;  $f$  — feature
- Steps in designing an evaluation function:
  1. Pick informative features
  2. Tune the weights

Deep Blue: precision in eval (normalized, between 0 — 1) is  $10^{-3}$  to  $10^{-4}$  ! (lots of fine-tuning is important)

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### **Creating Evaluation Functions**

- Features / weights for chess?
- How tune weights?
- How find features?

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### **Design Issues of Heuristic Minimax**

**Search:** search to a constant depth

**Problems:**

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### Design Issues of Heuristic Minimax

- Some portions of the game tree may be “hotter” than others. Should search to *quiescence*. Continue along a path as long as one move’s static value stands out (indicating a likely capture).
- *Horizon effect*
- Secondary search. (*singular extension heuristic*)

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