

CSE P 501 – Compilers

Implementing ASTs

(in Java)

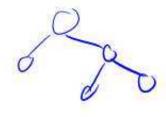
Hal Perkins Autumn 2009

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Agenda



- Representing ASTs as Java objects
- Parser actions
- Operations on ASTs
 - Modularity and encapsulation
- Visitor pattern
- This is a general sketch of the ideas more details and sample code online for MiniJava

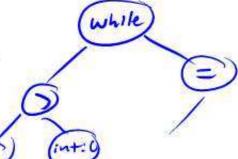
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Review: ASTs

 An Abstract Syntax Tree captures the essential structure of the program, without the extra concrete grammar details needed to guide the parser AST:



Example:

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Representation in Java

- Basic idea is simple: use small classes as records (or structs) to represent nodes in the AST
 - Simple data structures, not too smart
- But also use a bit of inheritance so we can treat related nodes polymorphically
- Following slides sketch the ideas not necessarily what you'll use in your project

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AST Nodes - Sketch

```
// Base class of AST node hierarchy
public abstract class ASTNode {
    // constructors (for convenience)
    // operations
    // string representation
    public abstract String toString();
    // visitor methods, etc.
}
```

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Some Statement Nodes

```
// Base class for all statements
public abstract class StmtNode extends ASTNode {(...)}
// while (exp) stmt
public class WhileNode extends StmtNode {
    public ExpNode exp;
    public StmtNode stmt;
    public WhileNode(ExpNode exp, StmtNode stmt) {
        this.exp = exp; this.stmt = stmt;
    }
    public String toString() {
        return "While(" + exp + ") " + stmt;
    }
}
(Note on toString: most of the time we'll want to print the tree in a separate traversal, so this is mostly useful for limited debugging)
```

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More Statement Nodes

```
// if (exp) stmt [else stmt]
[public class IfNode extends StmtNode {
    public ExpNode exp;
    public StmtNode thenStmt, elseStmt;
    public IfNode(ExpNode exp,StmtNode thenStmt,StmtNode elseStmt) {
        this.exp=exp; this.thenStmt=thenStmt;this.elseStmt=elseStmt;
    }
    public IfNode(ExpNode exp, StmtNode thenStmt) {
        this(exp, thenStmt, null);
    }
    public String toString() { ... }
}
```

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Expressions

8



More Expressions

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&c

- These examples are meant to get across the ideas, not necessarily to be used literally
 - E.g., you might find it much better to have a specific AST node for "argument list" that encapsulates the List of arguments
- You'll also need nodes for class and method declarations, parameter lists, and so forth
 - Starter code on the web for MiniJava

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Position Information in Nodes

- To produce useful error messages, it's helpful to record the source program location corresponding to a node in that node
 - Most scanner/parser generators have a hook for this, usually storing source position information in tokens
 - Included in the MiniJava starter code we distributed – useful to take advantage of it in your code

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AST Generation



- Idea: each time the parser recognizes a complete production, it produces as its result an AST node (with links to the subtrees that are the components of the production in its instance variables)
- When we finish parsing, the result of the goal symbol is the complete AST for the program

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Example: Recursive-Descent AST Generation

```
// parse while (exp) stmt
WhileNode whileStmt() {
    // skip "while ("
    getNextToken();
    getNextToken();

// parse exp
ExpNode condition = exp();
...
```

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AST Generation in YACC/CUP

- A result type can be specified for each item in the grammar specification
- Each parser rule can be annotated with a semantic action, which is just a piece of Java code that returns a value of the result type
- The semantic action is executed when the rule is reduced

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YACC/CUP Parser Specification

Specification

```
non terminal StmtNode stmt, whileStmt;
non terminal ExpNode exp;
...
stmt ::= ...
| WHILE LPAREN exp:e RPAREN stmt:s
{: RESULT = new WhileNode(e,s); :}
;
```

See the starter code for version with line numbers

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ANTLR/JavaCC/others

- Integrated tools like these provide tools to generate syntax trees automatically
 - Advantage: saves work, don't need to define AST classes and write semantic actions
 - Disadvantage: generated trees might not have the right level of abstraction for what you want to do
- For our project, do-it-yourself with CUP
 - The (revised) starter code contains the AST classes from the minijava web site

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Operations on ASTs

- Once we have the AST, we may want to
 - Print a readable dump of the tree (pretty printing)
 - Do static semantic analysis
 - Type checking
 - Verify that things are declared and initialized properly
 - Etc. etc. etc. etc.
 - Perform optimizing transformations on the tree
 - Generate code from the tree, or
 - Generate another IR from the tree for further processing

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Where do the Operations Go?

- Pure "object-oriented" style
 - Really smart AST nodes
 - Each node knows how to perform every operation on itself

```
public class WhileNode extends StmtNode {
  public WhileNode(...);
  public typeCheck(...);
  public StrengthReductionOptimize(...);
  public generateCode(...);
  public prettyPrint(...);
  ...
}
```

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Critique

- This is nicely encapsulated all details about a WhileNode are hidden in that class
- But it is poor modularity
- What happens if we want to add a new Optimize operation?
 - Have to open up every node class
- Furthermore, it means that the details of any particular operation (optimization, type checking) are scattered across the node classes

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Modularity Issues

- Smart nodes make sense if the set of operations is relatively fixed, but we expect to need flexibility to add new kinds of nodes
- Example: graphics system
 - Operations: draw, move, iconify, highlight
 - Objects: textbox, scrollbar, canvas, menu, dialog box, plus new objects defined as the system evolves

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Modularity in a Compiler

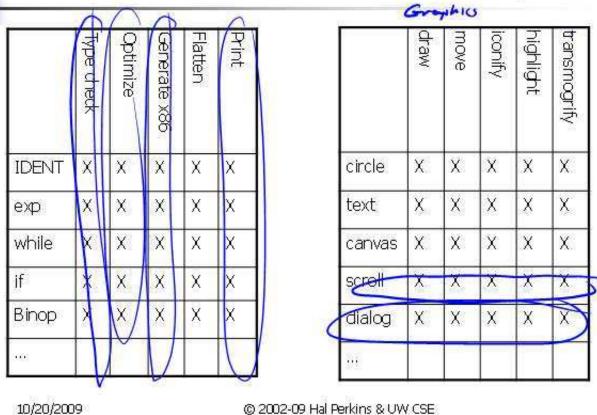
- Abstract syntax does not change frequently over time
 - ∴ Kinds of nodes are relatively fixed
- As a compiler evolves, it is common to modify or add operations on the AST nodes
 - Want to modularize each operation (type check, optimize, code gen) so its components are together
 - Want to avoid having to change node classes when we modify or add an operation on the tree

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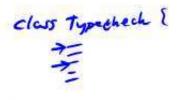
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Two Views of Modularity



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Visitor Pattern,

- Idea: Package each operation in a separate class
 - One operation method for each AST node kind
- Create one instance of this visitor class
 - Sometimes called a "function object"
- Include a generic "accept visitor" method in every node class
- To perform the operation, pass the "visitor object" around the AST during a traversal
 - This object contains separate methods to process each AST node type

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Avoiding instanceof

Next issue: we'd like to avoid huge if-elseif nests to check the node type in the visitor

```
void checkTypes(ASTNode p) {
   if (p instanceof WhileNode) { ... }
   else if (p instanceof IfNode) { ... }
   else if (p instanceof BinExp) { ... } ...
```

- Solution: Include an overloaded "visit" method in each AST node type and get the AST node to call back to the correct operation for that node(!)
 - "Double dispatch"

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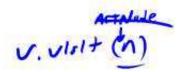


One More Issue

- We want to be able to add new operations easily, so the nodes shouldn't know anything specific about the actual visitor class(es)
- Solution: an abstract Visitor interface
 - AST nodes include "accept visitor" method for the interface
 - Specific operations (type check, code gen) are implementations of this interface

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Visitor Interface

```
interface Visitor {
    // overload visit for each AST node type
    public void visit(WhileNode s);
    public void visit(IfNode s);
    public void visit(BinExp e);
    ...
}
```

 Aside: The result type can be whatever is convenient, doesn't have to be void

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Specific class TypeCheckVisitor

```
// Perform type checks on the AST
public class TypeCheckVisitor implements Visitor {

// override operations for each node type
public void visit(BinExp e) {

// visit subexpressions – pass this visitor object
e.exp1.accept(this); e.exp2.accept(this);
// do additional processing on e before or after
}

public void visit(WhileNode s) { ... }

public void visit(IfNode s) { ... }

...
}

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```



Add Visitor Method to AST Nodes

 Add a new method to class ASTNode (base class or interface describing all AST nodes)

```
public abstract class ASTNode {
    ...
    // accept a visit from a Visitor object v
    public abstract void accept(Visitor v);
    ...
}
```

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arthode type visitor type





Override Accept Method in Each Specific AST Node Class

Example

```
public class WhileNode extends StmtNode {
    ...
    // accept a visit from a Visitor object v
    public void accept(Visitor v) {
        v.visit(this); // dynamic dispatch on "this" (WhileNode)
    }
    ...
}
```

- Key points
 - Visitor object passed as a parameter to WhileNode
 - WhileNode calls visit, which dispatches to visit(WhileNode) automatically – i.e., the correct method for this kind of node

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Encapsulation

- A visitor object often needs to be able to access state in the AST nodes
 - .:. May need to expose more state than we might do to otherwise
 - Overall a good tradeoff better modularity
 - (plus, the nodes are relatively simple data objects anyway – not hiding much of anything)

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Composite Objects

 If the node contains references to subnodes, we often visit them first (i.e., pass the visitor along in a depth-first traversal of the AST)
 public class WhileNode extends StmtNode {

```
// accept a visit from Visitor object v
public void accept(Visitor v) {
    this.exp.accept(v);
    this.stmt.accept(v);
    v.visit(this);
}
...
}
```

Other traversals can be added if needed

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Visitor Actions

- A visitor function has a reference to the node it is visiting (the parameter)
 - .: can access subtrees via that node
- It's also possible for the visitor object to contain local state (data), used to accumulate information during the traversal
 - Effectively "global data" shared by visit methods public class TypeCheckVisitor extends NodeVisitor { public void visit(WhileNode s) { ... } public void visit(IfNode s) { ... } ... private <local state>; }

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Responsibility for the Traversal

- Possible choices
 - The node objects (as done above)
 - The visitor object (the visitor has access to the node, so it can traverse any substructure it wishes)
 - Some sort of iterator object
- In a compiler, the first choice will handle many common cases

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References

- For Visitor pattern (and many others)
 - Design Patterns: Elements of Reusable
 Object-Oriented Software
 Gamma, Helm, Johnson, and Vlissides
 Addison-Wesley, 1995
- Specific information for MiniJava AST and visitors in Appel textbook & online

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Coming Attractions

- Static Analysis
 - Type checking & representation of types
 - Non-context-free rules (variables and types must be declared, etc.)
- Symbol Tables
- & more

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